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Department of Computer Science & Engineering

Operating Systems

(21CSE34)

OPERATING SYSTEM STRUCTURES

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Operating System Design and Implementation



- Design and Implementation of OS not “solvable”, but some approaches have proven successful.
- Internal structure of different Operating Systems can vary widely.
- Start the design by defining goals and specifications .
- Highest level: affected by choice of hardware, type of system.
- The requirements can be divided into User and System goals,
 - User goals – operating system should be convenient to use, easy to learn, reliable, safe, and fast.
 - System goals – operating system should be easy to design, implement, and maintain, as well as flexible, reliable, error-free, and efficient.

Operating System Design and Implementation



- Important principle to separate
- Policy: What will be done?
- Mechanism: How to do it?
- Mechanisms determine how to do something, policies decide what will be done
- The separation of policy from mechanism is a very important principle, it allows maximum flexibility if policy decisions are to be changed later (example – timer).
- Specifying and designing an OS is highly creative task of software engineering

Implementation



- Much variation.
- Early OSes in assembly language.
- Then system programming languages like Algol, PL/1.
- Now C, C++.
- Actually usually a mix of languages.
- Lowest levels in assembly.
- Main body in C.
- Systems programs in C, C++, scripting languages like PERL, Python, shell scripts.
- More high-level language easier to port to other hardware.
- But slower.
- Emulation can allow an OS to run on non-native hardware

Operating System Structure



- General-purpose OS is very large program.
- Various ways to structure ones.
- Simple structure – MS-DOS.
- More complex – UNIX.
- Layered – an abstraction.
- Microkernel –Mach
- Module Structure

Operating System Structure

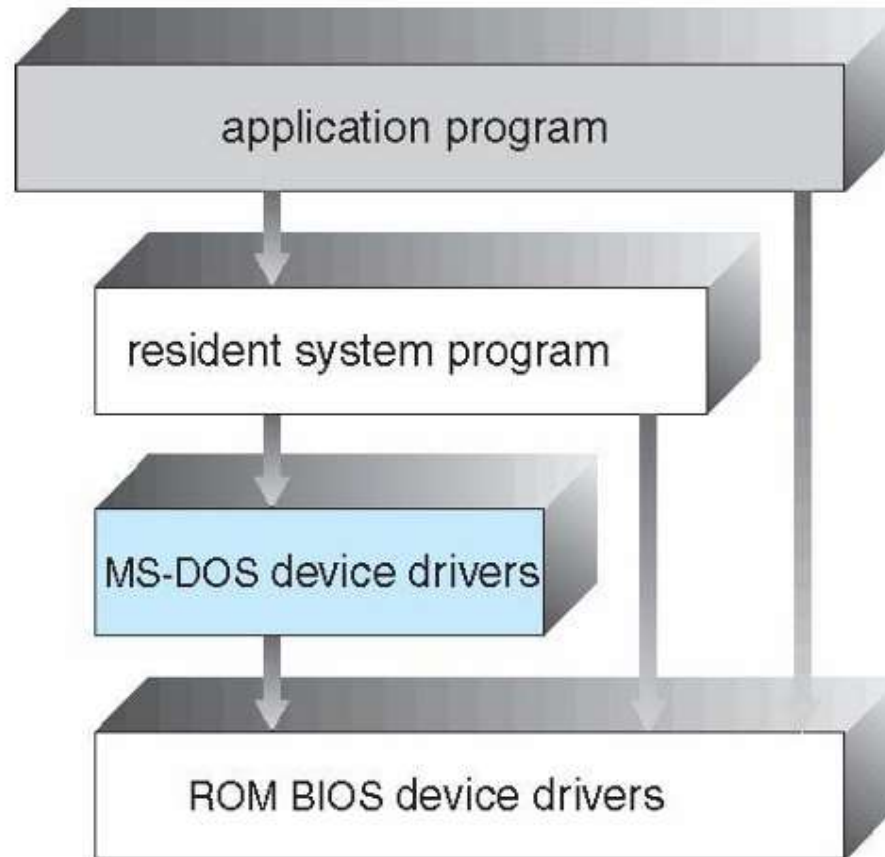


Fig 1: MS-DOS Layer Structure

Traditional Unix Structure



- UNIX – limited by hardware functionality, the original UNIX operating system had limited structuring.
- The UNIX OS consists of two separable parts .
- Systems programs The kernel
 - Consists of everything below the system-call interface and above the physical hardware
 - Provides the file system, CPU scheduling, memory management, and other operating-system functions; a large number of functions for one level

Traditional Unix Structure

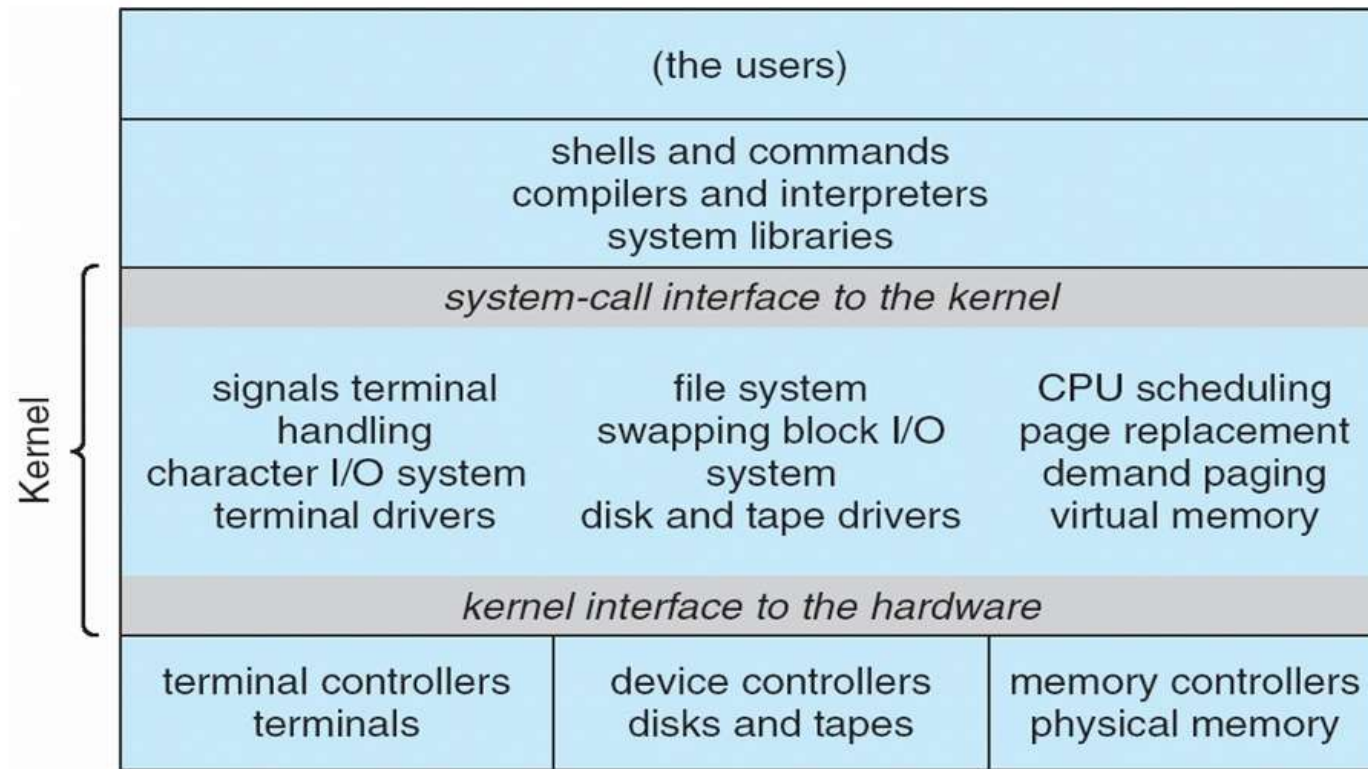


Fig 2: Traditional Unix Structure

Layered Structure



- The operating system is divided into a number of layers (levels), each built on top of lower layers. The bottom layer (layer 0), is the hardware; the highest (layer N) is the user interface.
- With modularity, layers are selected such that each uses functions (operations) and services of only lower-level layers
- Simplifies debugging and system verification

Layered Structure

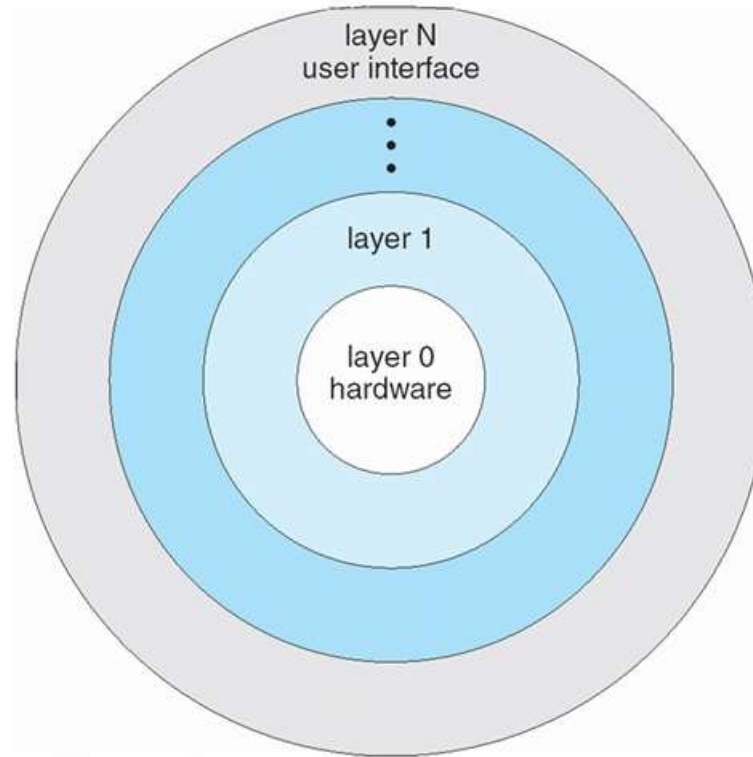


Fig 3: Layered Structure

Microkernel



- Moves as much from the kernel into user space Mach example of microkernel
- Mac OS X kernel (Darwin) partly based on Mach
- Communication takes place between user modules using message passing
- Benefits:
 - Easier to extend a microkernel
 - Easier to port the operating system to new architectures
 - More reliable (less code is running in kernel mode)
 - More secure
- Detriments:
 - Performance overhead of user space to kernel space communication

Microkernel

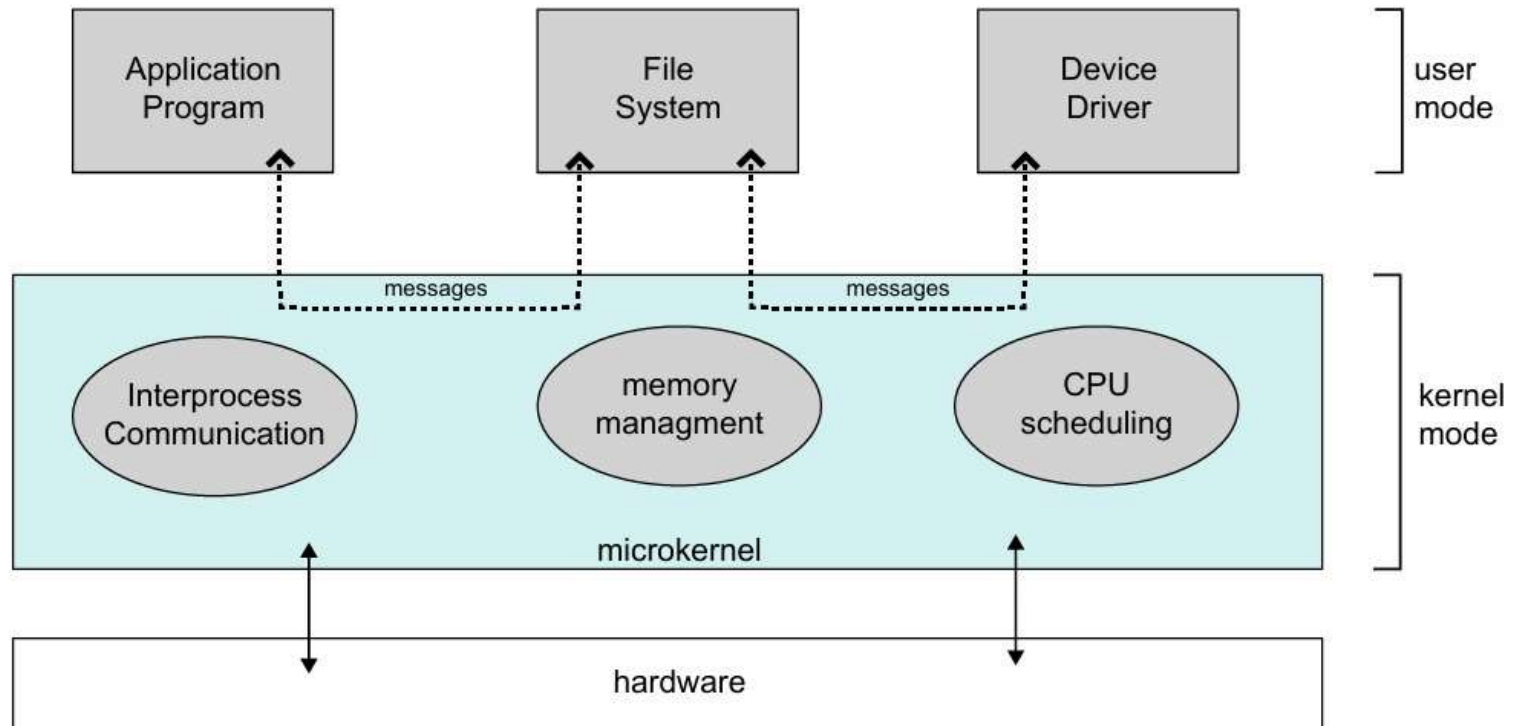


Fig 4: Microkernel

Module Structure



- Many modern operating systems implement loadable kernel modules
- Uses object-oriented approach
- Each core component is separate
- Each talks to the others over known interfaces
- Each is loadable as needed within the kernel
- Overall, similar to layers but with more flexible Linux, Solaris, etc

Module Structure

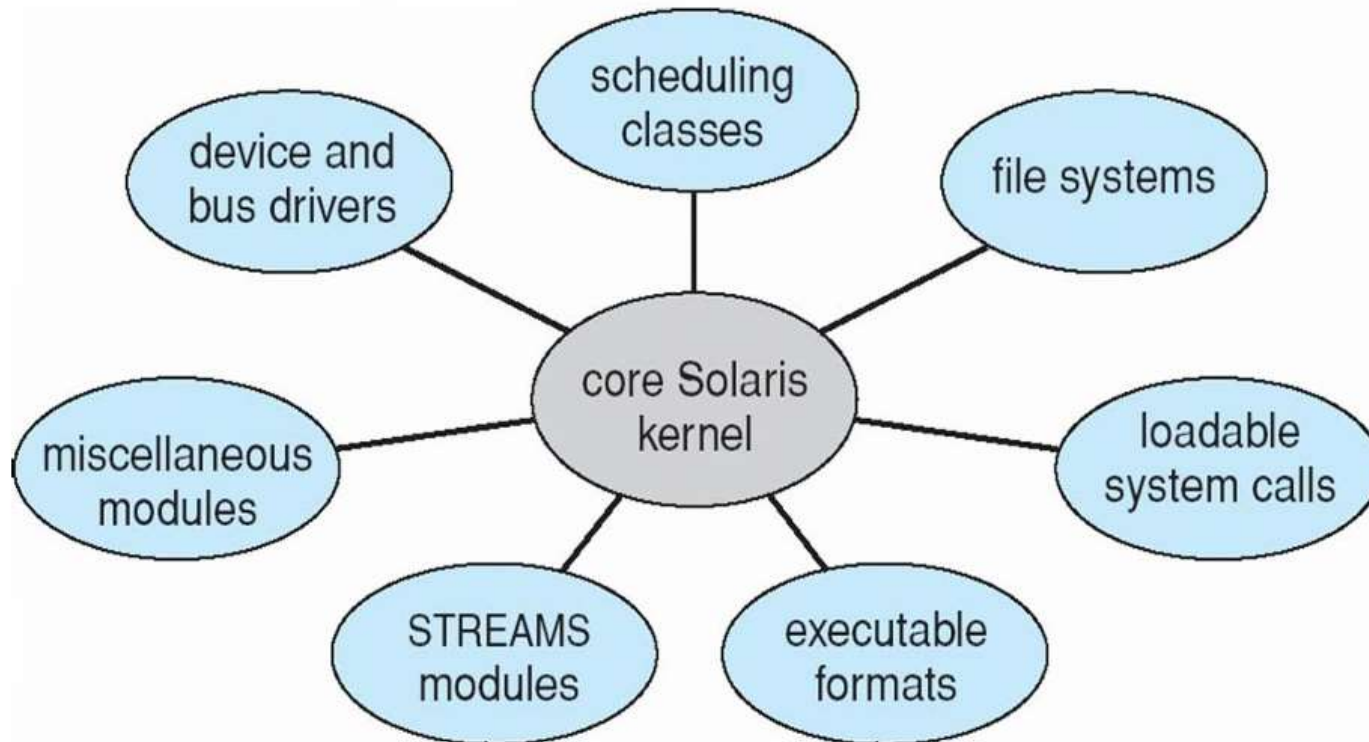


Fig 5: Module Structure