

Part I: Crypto

Chapter 2: Crypto Basics

MXDXBVTZWVMXNSPBQXLIMSCCSGXSCJXBOVQXCJZMOJZCVC
TVWJCZAAXZBCSSCJXBQCJZCOJZCNSPOXBXSBTWVWJC
JZDXGXXMOZQMSCSCJXBOVQXCJZMOJZCNSPJZHGXXMOSPLH
JZDXZAAXZBXHCSCJXTCSGXSCJXBOVQX

— plaintext from Lewis Carroll, *Alice in Wonderland*

The solution is by no means so difficult as you might be led to imagine from the first hasty inspection of the characters.

These characters, as any one might readily guess, form a cipher — that is to say, they convey a meaning...

— Edgar Allan Poe, *The Gold Bug*

Crypto

- ❑ **Cryptology** — The art and science of making and breaking “secret codes”
- ❑ **Cryptography** — making “secret codes”
- ❑ **Cryptanalysis** — breaking “secret codes”
- ❑ **Crypto** — all of the above (and more)

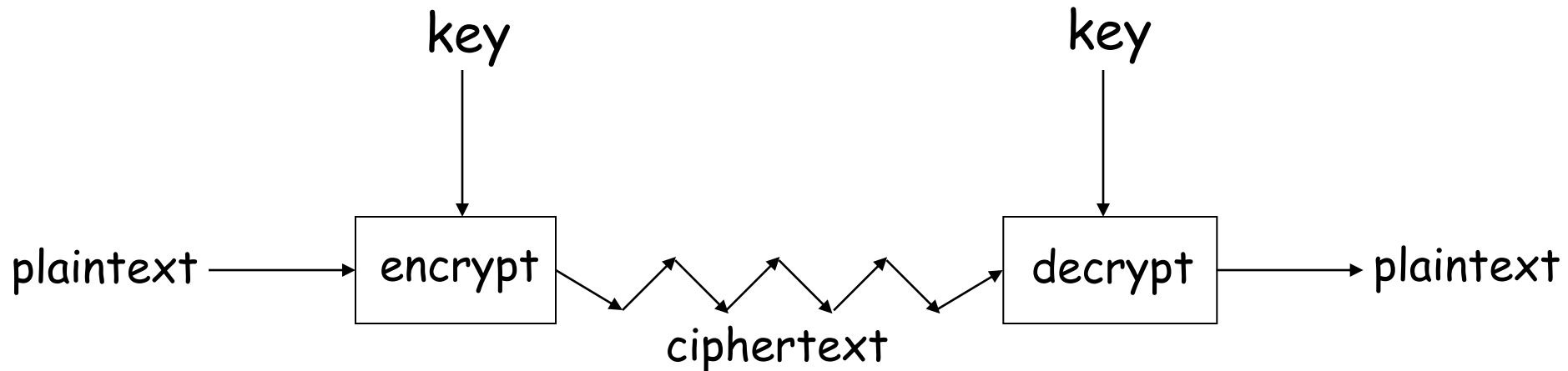
How to Speak Crypto

- ❑ A *cipher* or *cryptosystem* is used to *encrypt* the *plaintext*
- ❑ The result of encryption is *ciphertext*
- ❑ We *decrypt* ciphertext to recover plaintext
- ❑ A *key* is used to configure a cryptosystem
- ❑ A *symmetric key* cryptosystem uses the same key to encrypt as to decrypt
- ❑ A *public key* cryptosystem uses a *public key* to encrypt and a *private key* to decrypt

Crypto

- ❑ Basic assumptions
 - The system is completely known to the attacker
 - Only the key is secret
 - That is, crypto algorithms are not secret
- ❑ This is known as **Kerckhoffs' Principle**
- ❑ Why do we make such an assumption?
 - Experience has shown that secret algorithms tend to be weak when exposed
 - Secret algorithms never remain secret
 - Better to find weaknesses beforehand

Crypto as Black Box



A generic view of symmetric key crypto

Simple Substitution

❑ Plaintext: **fourscoreandsevenyearsago**

❑ Key:

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

❑ Ciphertext:

IRXUVFRUHDQGVHYHQBHDUVDJR

❑ Shift by 3 is "Caesar's cipher"

Caesar's Cipher Decryption

- Suppose we know a Caesar's cipher is being used:

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

- Given ciphertext:

VSRQJHEREVTXDUHSDQWV

- Plaintext: spongebobsquarepants

Caesar's Cipher

- Shift by n for some $n \in \{0, 1, 2, \dots, 25\}$
- Then key is n
- Example: key $n = 7$

Plaintext

a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G

Ciphertext

Cryptanalysis I: Try Them All

- ❑ We know Caesar's cipher (shift by n) used
 - But the specific key is unknown
- ❑ Given ciphertext: **CSYEVIXIVQMREXIH**
- ❑ How to determine the key?
- ❑ Only 26 possible keys — try them all!
- ❑ **Exhaustive key search**
- ❑ Solution: key is $n = 4$

Simple Substitution: General Case

- ❑ In general, simple substitution key can be any **permutation** of letters
 - Not necessarily a Caesar's cipher (shift)
- ❑ For example

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	J	I	C	A	X	S	E	Y	V	D	K	W	B	Q	T	Z	R	H	F	M	P	N	U	L	G	O

- ❑ In general, $26! > 2^{88}$ possible keys

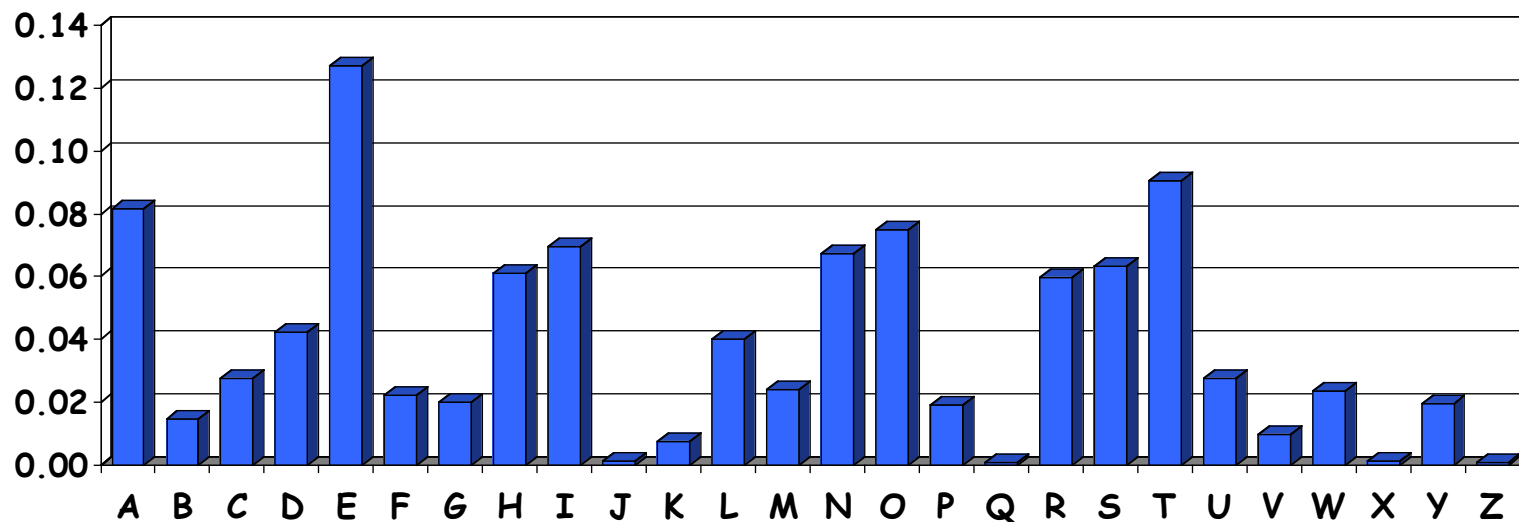
Cryptanalysis II: Be Clever

- ❑ We know that a simple substitution used
- ❑ But *not* necessarily a Caesar's cipher (shift)
- ❑ Find the key given the ciphertext:

PBFPVYFBQXZTYFPBFEQJHDXXQVAPTPQJKTOYQWIPBVWLXTOX
BTFXQWAXBVCXQWAXFQJWVLEQNTQZQGGQLFXQWAKVWLXQ
WAEBIPBFXFQVXGTVJVWLBTPQWAEBFPBFHCVLXBQUFEVWLXGD
PEQVPQGVPPBFTIXPFHXZHVFAGFOTHFEFBQUFTDHzBQPOTHXTY
FTODXQHFTDPTOGHFQPBQWAQJJTODXQHFOQPWTBDHHIXQV
APBFZQHCFWPFHPBFIPBQWKFABVYYDZBOTHBPBPQJTQOTOGHF
QAPBFEQJHDXXQVAVXEBQPEFZBVFOJIWFFACFCCFHQWAUVWF
LQHGFVAFXQHFUFHILTTAVWAFFAWTEVOITDHFHFQAITIXPFH
XAFQHEFZQWGFLVWPTOFFA

Cryptanalysis II

- ❑ Cannot try all 2^{88} simple substitution keys
- ❑ Can we be more clever?
- ❑ English letter frequency counts...



Cryptanalysis II

□ Ciphertext:

PBFPVYFBQXZTYFPBFEQJHDXQVAPTPQJKTOYQWIPBVWLXTOXBTFXQ
WAXBVCXQWAXFQJWVLEQNTQZQGGQLFXQWAKVWLXQWAEBIPBFXFQ
VXGTVJVWLBTPQWAEFBFBFHCVLXBQUFEVWLXGDPEQVPQGVPPBFTIXPFH
XZHVFAGFOTHFEBQUFTDHzBQPOTHXTYFTODXQHFTDPTOGHFQPBQW
AQJJTODXQHFOQPWTBDHHIXQVAPBFZQHCFWPFHPBFIQPBQWKFAVYY
DZBOTHQPBQPJTQTOGHFQAPBFEQJHDXQVAVXEBQPEFZBVFOJIWFF
ACFCCFHQWAVVWFLQHGFVAFXQHUFHILTTAVWAFFAWTEVOITDHFH
FQAITIXPFHXAFQHEFZQWGFLVWPTOFFA

□ Analyze this message using statistics below

Ciphertext frequency counts:

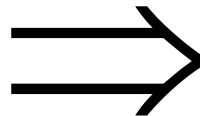
A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
21	26	6	10	12	51	10	25	10	9	3	10	0	1	15	28	42	0	0	27	4	24	22	28	6	8

Double Transposition

□ Plaintext: **attackxatxdawn**

	col 1	col 2	col 3
row 1	a	t	t
row 2	a	c	k
row 3	x	a	t
row 4	x	d	a
row 5	w	n	x

Permute rows
and columns



	col 1	col 3	col 2
row 3	x	t	a
row 5	w	x	n
row 1	a	t	t
row 4	x	a	d
row 2	a	k	c

□ Ciphertext: **xtawxnattxadakc**

□ Key is matrix size and permutations:
(3,5,1,4,2) and (1,3,2)

Cryptanalysis: Terminology

- ❑ Cryptosystem is **secure** if best know attack is to try all keys
 - Exhaustive key search, that is
- ❑ Cryptosystem is **insecure** if *any* shortcut attack is known
- ❑ But then insecure cipher might be harder to break than a secure cipher!
 - What the ... ?

One-Time Pad: Encryption

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

Encryption: Plaintext \oplus Key = Ciphertext

	h	e	i	l	h	i	t	l	e	r
Plaintext:	001	000	010	100	001	010	111	100	000	101
Key:	111	101	110	101	111	100	000	101	110	000
Ciphertext:	110	101	100	001	110	110	111	001	110	101
	s	r	l	h	s	s	t	h	s	r

One-Time Pad: Decryption

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

Decryption: Ciphertext \oplus Key = Plaintext

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
Key:	111	101	110	101	111	100	000	101	110	000
Plaintext:	001	000	010	100	001	010	111	100	000	101
	h	e	i	l	h	i	t	l	e	r

One-Time Pad

Double agent claims following “**key**” was used:

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
“ key ”:	101	111	000	101	111	100	000	101	110	000
“Plaintext”:	011	010	100	100	001	010	111	100	000	101
	k	i	l	l	h	i	t	l	e	r

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

One-Time Pad

Or, might claim the key is...

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
"key":	111	101	000	011	101	110	001	011	101	101
<hr/>										
"Plaintext":	001	000	100	010	011	000	110	010	011	000
	h	e	l	i	k	e	s	i	k	e

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

One-Time Pad Summary

- ❑ **Provably** secure
 - Ciphertext gives **no** useful info about plaintext
 - All plaintexts are *equally likely*
- ❑ BUT, only when be used correctly
 - Pad must be random, used only once
 - Pad is known only to sender and receiver
- ❑ Note: pad (key) is same size as message
- ❑ So, why not distribute message itself, instead of the pad?

Codebook Cipher

- ❑ Literally, a book filled with “codewords”
- ❑ Encryption via codebook

Februar	13605
fest	13732
finanzielle	13850
folgender	13918
Frieden	17142
Friedenschluss	17149
:	:

- ❑ Modern block ciphers are codebooks!
- ❑ More about this later...

Zimmerman Telegram

- Perhaps most famous codebook ciphertext ever

WESTERN UNION TELEGRAM

Send the following telegram, subject to the terms on back hereof, which are hereby agreed to

GERMAN LEGATION
MEXICO CITY

via Galveston

JAN 29 1917

130	13042	13401	8501	115	3528	416	17214	8491	11310
18147	18222	21560	10247	11518	23677	13605	3494	14936	
98092	5905	11311	10392	10371	0302	21290	5161	39695	
23571	17504	11269	18276	18101	0317	0228	17694	4473	
22284	22200	19452	21589	87893	5569	13918	8958	12137	
1333	4725	4458	5905	17166	13851	4458	17149	14471	6708
13850	12224	6929	14991	7382	15857	67893	14218	36477	
5870	17553	67893	5870	5454	16102	15217	22801	17138	
21001	17388	7446	23638	18222	6719	14331	15021	23845	
3156	23552	22096	21604	4797	9497	22464	20855	4377	
23610	18140	22260	5905	13347	20420	39689	13732	20667	
6929	5275	18507	52262	1340	22049	13339	11265	22295	
10439	14814	4178	6992	8784	7032	7357	6926	52262	11267
21100	21272	9346	9559	22464	15874	18502	18500	15857	
2188	5376	7381	98092	16127	13486	9350	9220	76036	14219
5144	2831	17920	11347	17142	11264	7667	7762	15099	9110
10482	97556	3569	3670						

BEPNSTORFF.

Charge German Embassy.

Early 20th Century

- ❑ WWI — Zimmerman Telegram
- ❑ “Gentlemen do not read each other's mail”
 - Henry L. Stimson, Secretary of State, 1929
- ❑ WWII — **golden** age of cryptanalysis
 - Midway/Coral Sea
 - Japanese **Purple** (codename MAGIC)
 - German **Enigma** (codename ULTRA)

Post-WWII History

- ❑ Claude Shannon — father of the science of information theory
- ❑ Computer revolution — lots of data to protect
- ❑ Data Encryption Standard (DES), 70's
- ❑ Public Key cryptography, 70's
- ❑ CRYPTO conferences, 80's
- ❑ Advanced Encryption Standard (AES), 90's
- ❑ The crypto genie is out of the bottle...

Claude Shannon

- ❑ Founded field of information theory
- ❑ His 1949 paper: [Comm. Thy. of Secrecy Systems](#)
- ❑ Fundamental concepts
 - **Confusion** — obscure relationship between plaintext and ciphertext
 - **Diffusion** — spread plaintext statistics through the ciphertext
- ❑ Proved one-time pad is secure
- ❑ One-time pad is confusion-only, while double transposition is diffusion-only

Taxonomy of Cryptography

□ Symmetric Key

- Same key for encryption and decryption
- Modern types: Stream ciphers, Block ciphers

□ Public Key (or “asymmetric” crypto)

- Two keys, one for encryption (public), and one for decryption (private)
- And digital signatures — nothing comparable in symmetric key crypto

□ Hash algorithms

- Can be viewed as “one way” crypto

Taxonomy of Cryptanalysis

- ❑ From perspective of info available to Trudy...
 - Ciphertext only — Trudy's worst case scenario
 - Known plaintext
 - Chosen plaintext
 - "Lunchtime attack"
 - Some protocols will encrypt chosen data
 - Adaptively chosen plaintext
 - Related key
 - Forward search (public key crypto)
 - And others...