## CS & IT ENGINEERING



F

Flow Control

Lecture No-9



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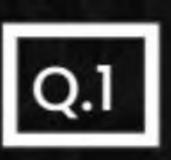


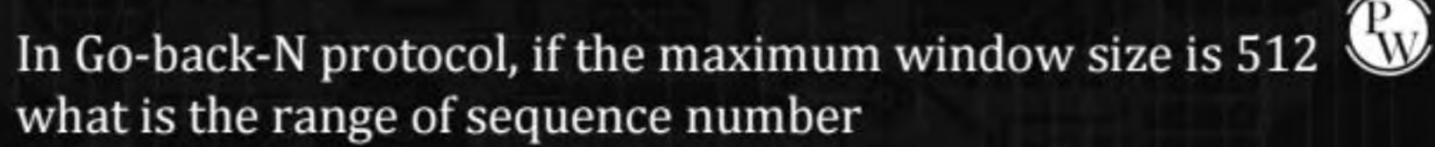
## TOPICS TO BE COVERED

- Problem solving on GB-N



## Problem Solving on GB-N Protocol





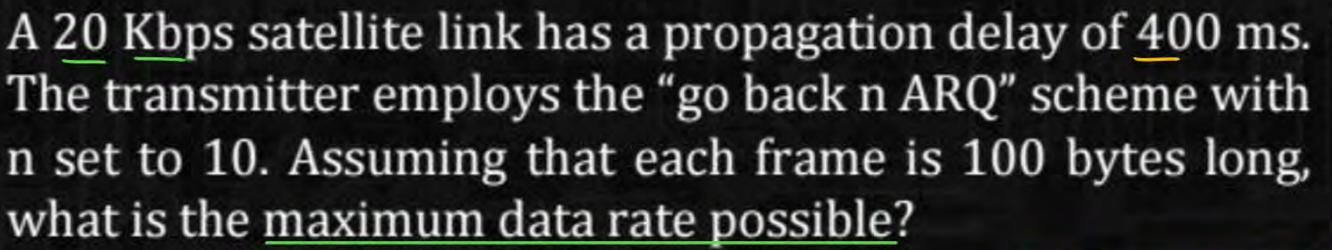


- 0 to 513
- 1 to 513
- 0 to 512
  - 1 to 512

9n GB-N 9F sender window size = N then minimum sequence No required in GB-N  $= W_S + W_R = N + 1 (0 - N)$ 

Window sindy size = 512 min. sezuence No required in GBN= 519+1=513[0-518]

## Q.2





| A )    | - | Kb           | ne |
|--------|---|--------------|----|
| $\sim$ | J | $\mathbf{n}$ | D2 |
|        |   |              |    |

$$B = 20 \text{KbPs} = 20 \times 10^3 \text{ bits/suc}$$
 $R = 400 \text{ msuc}$ 
 $N = 10$ 

Throughput = n \*B

= 0.4761 \* 20KbPs

= 9.52 KbPs

= 10KbPs

efficiency = Usefultime total time

$$= 10 \times 40 \text{ mssc}$$

$$40 \text{ mssc} + 2 \times 400 \text{ mssc}$$





Assume we need to design Go-back-N sliding window protocol for a network in which bandwidth is 1 Mbps and average distance between sender and receiver is 5000 Km. Assume that average packet size is 5000 bits. Propagation speed in the media is  $2 \times 10^8$  m/sec. In GB-10 If process delay is 0.5 Msec and queuing delay is 2 msec then what is the efficiency.

A 99%

B= 1mbPs=106 bits/sec

U=2×105 km/ SQC

B 57%

d= 5000 km Packet size or Frame size = 5000 bits

N=10 > Pod=0.5 m sec

C

87%

Gd=amsec

D 67%

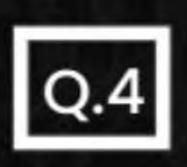
$$T_{d}(F) = \frac{F69ml \, Size}{89ndwidth}$$

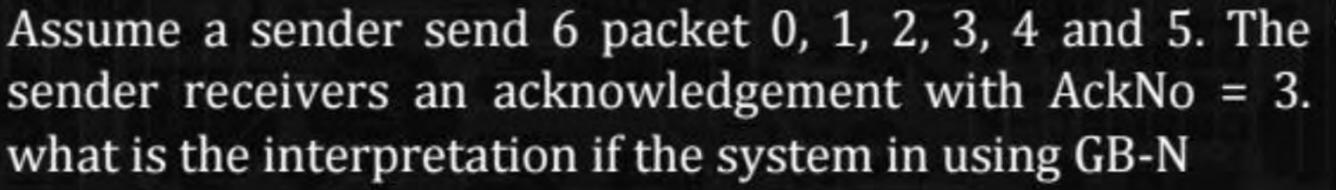
$$= \frac{5000 \, bits}{10^{5} \, bits/sec}$$

$$= 5 \times 10^{-3} \, sec = 5 \, msec$$

$$\frac{1}{10} = \frac{1}{2000} \times \frac{1}{1000} = \frac{1}{2000} \times \frac{1}{1000} = \frac{1}{2000} \times \frac{1}{1000} = \frac{1}{2000} \times \frac{1}{1000} = \frac{1}{20000} \times \frac{1}{1000} = \frac{1}{2000} \times \frac{1}{1000} = \frac{1}{1000} \times \frac{1}{1000} = \frac{1}$$











It means that packet 3 has been received uncorrupted



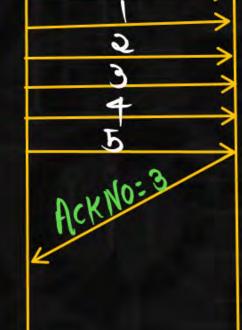
It means packet 0, 1, 2 have received uncorrupted and receiver is expected packet 3

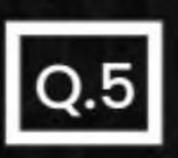


Ack does not say anything about other packet

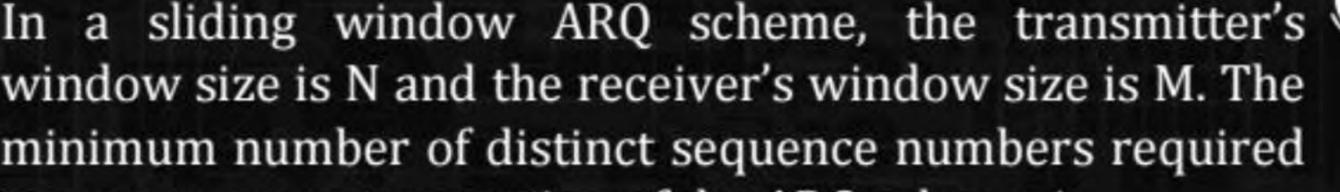


All the above





In a sliding window ARQ scheme, the transmitter's window size is N and the receiver's window size is M. The minimum number of distinct sequence numbers required to ensure correct operation of the ARQ scheme is



min (M, N)

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| В | , | max | (M, | N) |
|---|---|-----|-----|----|
|   |   |     |     |    |

| (C) | M - | + N |
|-----|-----|-----|

| Ws | WR | min seewince No required |
|----|----|--------------------------|
| N  | M  | N+M                      |
|    |    |                          |



Consider packet size is 1000 bits and distance between two hosts is 5 km, 1 Mbps link with signal speed 2 ms/km (ms per km) is used, the efficiency in percentage if GB-N protocol is used and N is set to 7 333.

```
Packet size or Framesize = 1000 bits
                                               N=7
 d= 5km > B= 1mbPs=106 bits/sec
                                              Td(F) = Frames 12e
                                                      Bandwidth
 Propagation time For one km= amsec
                                                   = 1000 bits
 Propagation dulay For 5km = 5x2msec = 10msec
                                                       108 byts / SOC
                                                   = 103 sqc = 1 M sqc
```

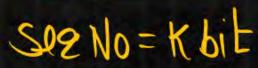




In GB-N Protocol the packet size is 1000 bytes transmission time for one packet is 1ms. If distance between hosts is 10km and signal speed is 5ms per km (5ms/km) and frame sequence number are 6 bit long in frame format then the throughput (in Mbps) is 439.

$$T_{d}(F) = 1 \text{ marc} = 10^{3} \text{ sac}$$
  
 $d = 10 \text{ km}$ 





$$\frac{W_S}{a^{6}-1} \qquad \frac{W_R}{1}$$



$$\frac{10^{-3} \text{sec} = 8000 \text{ bits}}{8}$$

$$\beta = \frac{8000 \text{ bits}}{10^{-3} \text{ sec}}$$





