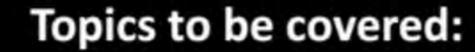


# CS & IT

Engineering







R

Lypractice on undecidability

**Topics Covered in Previous Session:** 

Pw

-> Undecidability problems Languages



1. Which of the following is decidable?

A. Finite Language → Decidable language (Recorsive)

B. Regular Language → T

C. Context Free Language

D. All of the above



- 2. Which of the following is decidable?
- A. Complement of a Regular Language → Regular → Decidable

  B. Complement of a Context Free Language → CSL
- C. Both A and B
  - D. None of these



3. Which of the following is decidable?

A. Set of regular languages over alphabet. =  $\{R_1, R_2, R_3, R_4, \dots \}$ B. Set of non regular languages over alphabet. =  $\{R_1, R_2, R_3, R_4, \dots \}$ 

C. Both A and B

D. None

Here there are some

non regulars which Can be not REL



#### 4. Which of the following is uncountable ?

Countain  $\Delta$ .  $\Sigma = \{a, b\}$ . finite Regular B.  $\Sigma^*$  over  $\Sigma = \{a, b\}$ .

C.  $2^{\Sigma^*}$  over  $\Sigma = \{a, b\}$ .

D. None



## 5. Which of the following is undecidable ?

A. 
$$\Sigma = \{a, b\}$$
.  $\Sigma^* \text{ over } \Sigma = \{a, b\}$ .

B. 
$$\Sigma^*$$
 over  $\Sigma = \{a, b\}$ .

D. None



6. Which of the following is decidable?

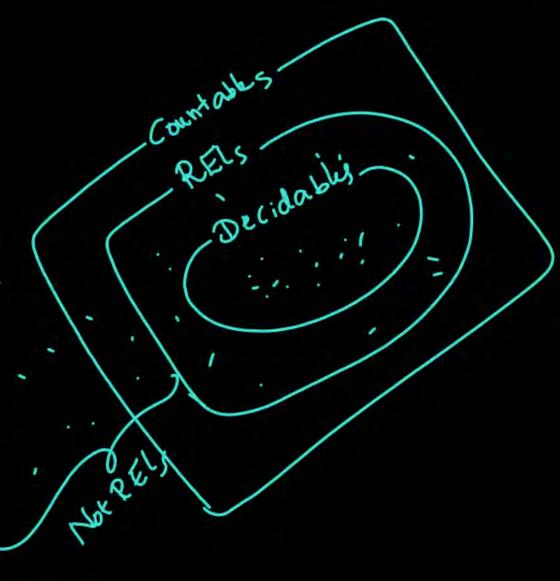
B. Set of all strings over finite alphabet. 2 C. Set of all languages over finite alphabet. — Where Reinstein B. Both A and B

D. Both A and B



### 7. Which of the following is FALSE?

- A. Every decidable set is countable set.
- B. Every REL is countable set.
- C. Every Non-REL is uncountable set. False
  - D. None of these





# 8. Let L = { M | M is a DFA that accepts ab }. Then L is \_\_\_\_\_

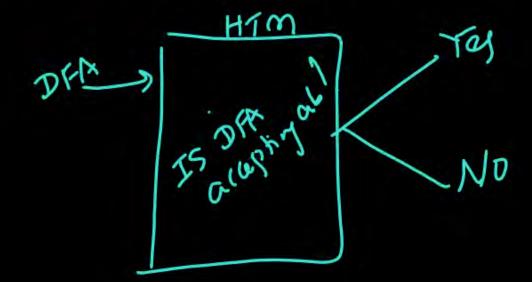
- A. Decidable Language
- B. REL but not recursive
- C. Not REL
- D. None of these



9. "Whether given DFA accepts ab" is \_\_\_\_\_

A. Decidable Language

- B. REL but not recursive
- C. Not REL
- D. None of these





#### 10. Let L = { M | M is a DFA that halts on ab }. Then L is

Halfing

A. Decidable Language

- B. Undecidable Language
- C. Regular Language
- D. Both A and C



Let L = { R | R is regular language over alphabet}.

Then L is \_\_\_\_\_

L- Set of all regular languages

- A. Decidable Language
- B. Undecidable Language
  - C. Regular Language
  - D. Both A and C



12. Let L = { R | R is a regular expression that generates string w}.

L = Set of all regular expression generating w

- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C

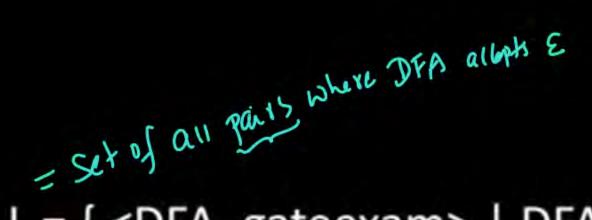


13. Let  $L = \{R \mid R \text{ is a regular grammar that generates string } w\}$ .

- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C



- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C





15. Let L = { <DFA, gateexam> | DFA accepts string epsilon}.

Then L is

- A. Decidable Language
  - B. Undecidable Language
  - C. Regular Language
  - D. Both A and C





16. Let L = { DFA | DFA accepts finite language}.
Then L is \_\_\_\_

- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C



- A. Decidable Language
  - B. Undecidable Language
  - C. Regular Language
  - D. Both A and C



To any State of the state of th

18. Let L = { <F, R> | F is FA and R is Regular Expression, L(F)=L(R)}.

- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C



A. Decidable Language

- B. Undecidable Language
- C. Regular Language
- D. Both A and C



20. Let L = { w | w belongs to {a, b}\*, w starts with a}. = 
$$a(a+b)^*$$
  
= { $a, aa, ab, ...$ } =  $a(a+b)^*$   
Then L is \_\_\_\_\_\_ Regular

- A. Decidable Language
- B. Undecidable Language
- C. Regular Language
- D. Both A and C



21 Let I = SHTM

21. Let L = {HTM | HTM accepts epsilon}. = Set if an Him a caph &

Then L is \_\_\_\_\_

Members lip Decidable

- A. Decidable Language
  - B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



22. Let L = { TM | TM accepts ab}. Membership for TM

Membership for TM

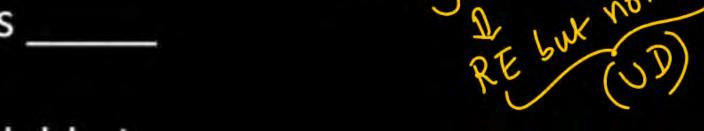
ORE but not rec

(UD)

- B. Undecidable Language but partially decidable RE Language but not even partially decidable decidable
  - D. None



23. Let L = { TM | TM halts on ab}.



- A. Decidable Language
- B. Undecidable Language but partially decidable -
  - C. Undecidable language but not even partially decidable
  - D. None



24. Let L = { TM | TM accepts a particular string w}.

Then L is \_\_\_\_\_

Membership for Im

RE but not rec

- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



25. Let L = { TM | TM accepts some string }.

Then L is \_\_\_\_\_ Non emphase set of the last not be a l

- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially NA REL decidable
  - D. None



- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



29. Let L = { TM | TM accepts non-empty language}.



- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



- A. Decidable Language
  - B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



31. Let L = { TM | TM accepts recursive set}. → Not REL
Then L is \_\_\_\_

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



32. Let L = { TM | TM accepts regular set}.

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



33. Let L = { TM | TM has at least 3 states}.
Then L is

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



## 34. Let L = { TM | TM accepts at least 3 strings}. $\rightarrow$ RE like rook with the lis

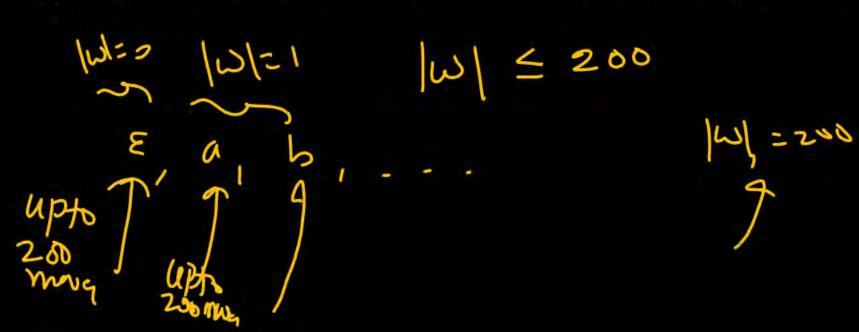
- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None

TM IM



35. Let L = { TM | TM reaches state q within 200 moves}. Then L is

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None





- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None

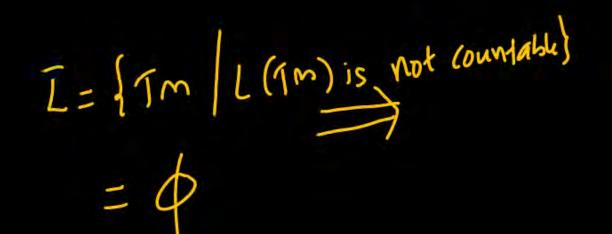


- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



38. Let  $L = \{ TM \mid L(TM) \text{ is countable} \}$ .

Trivial



- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



40. Let  $L = \{ < M, w, k > | M \text{ is a TM that halts on w within k steps} \}$ .

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



41. Let L = { < M, w, k > | M is a TM that does not halt on w within k steps}.

Then L is

The it hatts with x moves eximely to the hatts with x moves eximely the states with the moves eximely the states with the stat

A. Decidable Language

- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None



42. Let L = { M | M is a TM, L(TM) is subset of {a, bb}}.

Then L is \_\_\_\_\_

L(1m) = fa,663

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



43. Let L = { M | M is a TM, L(TM) is not subset of {a, bb}}.  $L(1m) = \{\alpha, bb\}$ 

Then L is

Find some string accepted by TM

Other Kan

a 4 66

- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



44. Let L = { TM | TM halts after 100 moves}.

Then L is \_\_\_\_\_

Whether In halts after boomones.

(In will not halt within too moves)

A. Decidable Language

- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None

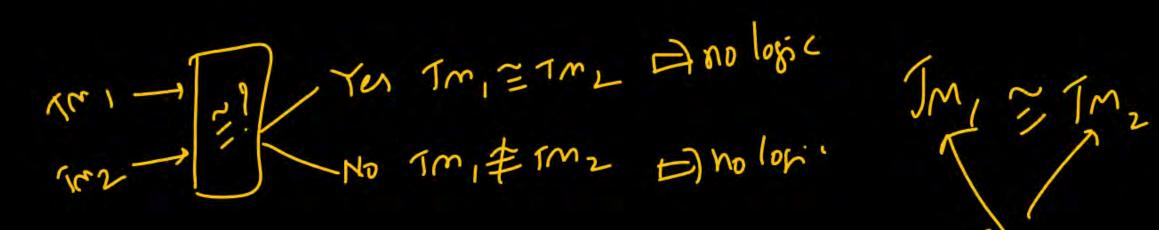
|W|=0=) 101st move upto 100 length strings



45. Let L = { TM | TM has 2 transitions}.

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
- D. None





46. Let  $L = \{ \langle M1, M2 \rangle \mid M1 \text{ is equivalent to } M2 \}$ .

Then L is

Equivalence for Tons

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None

Pw

- set of an undicidable languages

47. Let L = { R | R is not decidable language}.

Then L is

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



48. Let L = { R | R is decidable language}.

- A. Decidable Language
- B. Undecidable Language but partially decidable
  - C. Undecidable language but not even partially decidable
  - D. None



49. Let  $L = \{ R \mid R \text{ is not REL} \}$ .

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially decidable
  - D. None



50. Let  $L = \{ R \mid R \text{ is countable language} \}$ .

Then L is

- A. Decidable Language
- B. Undecidable Language but partially decidable
- C. Undecidable language but not even partially
- decidable
- D. None

UG (unrestricted Grammer) LHS -> RHS I) (VUT)\*-->(VUT)\* I) (VUT)\* V(VUT)\* -> (VUT)\*/
RHJ

CSG



7 But | LHS | = | RHS |



$$S \rightarrow \alpha AB$$
 $\alpha A \rightarrow \alpha A$ 
 $\alpha B \rightarrow b$ 



-> practice

