

CONSTRAINED OPTIMIZATION

A constrained optimization problem has the form

$$\begin{aligned} & \min_x f(x) \\ & \text{s.t. } g_i(x) \leq 0, \quad i=1, \dots, m \\ & \quad h_i(x) = 0, \quad i=1, \dots, n \end{aligned}$$

where $x \in \mathbb{R}^d$. If x satisfies all the constraints, it is said to be feasible. Assume f is defined at all feasible points.

The Lagrangian

The Lagrangian function is

$$L(x, \lambda, \nu) := f(x) + \sum_{i=1}^m \lambda_i g_i(x) + \sum_{j=1}^n \nu_j h_j(x)$$

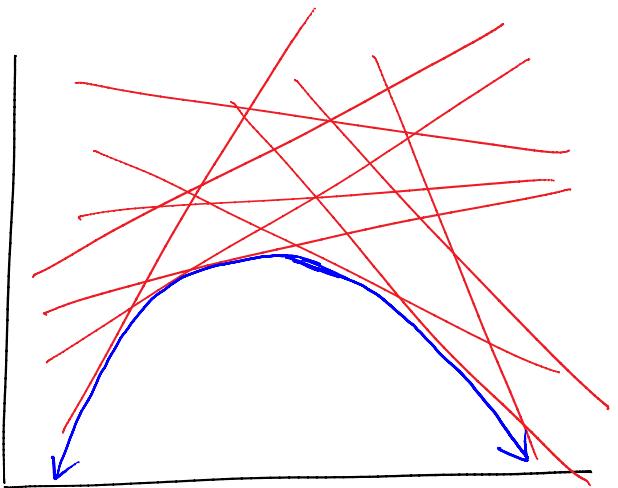
and $\lambda = [\lambda_1, \dots, \lambda_m]^T$ and $\nu = [\nu_1, \dots, \nu_n]^T$ are called Lagrange multipliers or dual variables.

The Lagrange dual function is

$$L_d(\lambda, \nu) := \min L(x, \lambda, \nu)$$

$L_D(\lambda, v)$

Note L_D is concave,
being the point-wise
minimum of a family
of affine functions



The dual optimization problem is

$$\max_{\lambda, v : \lambda_i \geq 0} L_D(\lambda, v)$$

Similarly, the primal function is

$$L_p(x) := \max_{\lambda, v : \lambda_i \geq 0} L(x, \lambda, v)$$

and the primal optimization problem is

$$\min_x L_p(x)$$

Notice that

$$L_p(x) = \begin{cases} f(x) & \text{if } x \text{ is feasible} \\ \infty & \text{otherwise} \end{cases}$$

Therefore, the primal problem and the original

problem have the same solution, yet the primal problem is unconstrained.

Weak Duality

Denote the optimal objective function values of the primal and dual

$$p^* = \min_x L_p(x) = \min_x \max_{\lambda, \nu: \lambda_i \geq 0} L(x, \lambda, \nu)$$

$$d^* = \max_{\lambda, \nu: \lambda_i \geq 0} L_D(x) = \max_{\lambda, \nu: \lambda_i \geq 0} \min_x L(x, \lambda, \nu).$$

Weak duality refers to the following fact, which always holds

Proposition $d^* \leq p^*$

Proof | Let \tilde{x} be feasible. Then for any λ, ν with $\lambda_i \geq 0$

$$L(\tilde{x}, \lambda, \nu) = f(\tilde{x}) + \sum \lambda_i g_i(\tilde{x}) + \sum \nu_j h_j(\tilde{x}) \leq f(\tilde{x})$$

Hence

$$L_D(\lambda, \nu) = \min_x L(x, \lambda, \nu) \leq f(\tilde{x})$$

This is true for any feasible \tilde{x} , so

$$L_D(\lambda, \nu) \leq \min_{\tilde{x} \text{ feasible}} f(\tilde{x}) = p^*.$$

Taking the max over $\lambda, \nu : \lambda_i \geq 0$, we have

$$d^* = \max_{\lambda, \nu : \lambda_i \geq 0} L_D(\lambda, \nu) \leq p^*. \quad \square$$

Strong Duality

If $p^* = d^*$, we say strong duality holds.

The original unconstrained optimization problem is said to be convex if f and g_1, \dots, g_m are convex functions and h_1, \dots, h_n are affine.

We state the following without proof.

Theorem] If the original problem is convex and a constraint qualification holds, then $p^* = d^*$.

Examples] of constraint qualifications

- All g_i are affine
- (Strict feasibility) $\exists x$ s.t. $h_j(x) = 0 \forall j$ and $g_i(x) < 0 \forall i$.

KKT Conditions

From now on, assume $f, g_1, \dots, g_m, h_1, \dots, h_n$ are differentiable. For unconstrained optimization, we know $\nabla f(x^*) = 0$ is necessary for x^* to be a global minimizer, and sufficient if f is additionally convex. The following two results generalize these properties to constrained optimization.

Theorem (Necessity) If $p^* = d^*$, x^* is primal optimal, and (λ^*, ν^*) is dual optimal, then the Karush-Kuhn-Tucker (KKT) conditions hold:

$$(1) \quad \nabla f(x^*) + \sum_i \lambda_i^* \nabla g_i(x^*) + \sum_i \nu_i^* \nabla h_i(x^*) = 0$$

$$(2) \quad g_i(x^*) \leq 0$$

$$(3) \quad h_i(x^*) = 0$$

$$(4) \quad \lambda_i^* \geq 0$$

$$(5) \quad \lambda_i^* g_i(x^*) = 0 \quad \forall i \quad (\text{complementary slackness})$$

Proof (2)-(3) hold since x^* is feasible. (4) holds by definition of the dual problem. To prove (5) and (1):

$$f(x^*) = L_D(\lambda^*, \nu^*) \quad [\text{by strong duality}]$$

$$\begin{aligned}
&= \min_x \left(f(x) + \sum \lambda_i^* g_i(x) + \sum \nu_i^* h_i(x) \right) \\
&\leq f(x^*) + \sum \lambda_i^* g_i(x^*) + \sum \nu_i^* h_i(x^*) \\
&\leq f(x^*) \quad [\text{by (2)-(4)}]
\end{aligned}$$

and therefore the two inequalities are equalities. Equality of the last two lines implies $\lambda_i^* g_i(x^*) = 0 \ \forall i$.

Equality of the 2nd and 3rd lines implies x^* is a minimizer of $L(x, \lambda^*, \nu^*)$ w.r.t. x . Therefore

$$\nabla_x L(x^*, \lambda^*, \nu^*) = 0,$$

which is (i). □

Theorem | (Sufficiency) If the original problem is convex and $\tilde{x}, \tilde{\lambda}, \tilde{\nu}$ satisfy the KKT conditions, then \tilde{x} is primal optimal, $(\tilde{\lambda}, \tilde{\nu})$ is dual optimal, and strong duality holds.

Proof | By (2) and (3), \tilde{x} is feasible. By (4), $L(x, \tilde{\lambda}, \tilde{\nu})$ is convex in x . By C.1, \tilde{x} is a minimizer of $L(x, \tilde{\lambda}, \tilde{\nu})$. Then

$$\begin{aligned}
 L_D(\tilde{x}, \tilde{\nu}) &= L(\tilde{x}, \tilde{\lambda}, \tilde{\nu}) \\
 &= f(\tilde{x}) + \underbrace{\sum \tilde{\lambda}_i g_i(\tilde{x}) + \sum \tilde{\nu}_i h_i(\tilde{x})}_{=0 \text{ by (5) and (3)}} \\
 &= f(\tilde{x}).
 \end{aligned}$$

Therefore $p^* = d^*$ and the result follows. \square

In conclusion, if a constrained optimization problem is differentiable and convex, then the KKT conditions are necessary and sufficient for primal/dual optimality (with zero duality gap). Thus, the KKT conditions can be used to solve such problems.

Saddle Point Property

If \tilde{x} is primal optimal, $(\tilde{\lambda}, \tilde{\nu})$ is dual optimal, and strong duality holds, then $(\tilde{x}, \tilde{\lambda}, \tilde{\nu})$ is a saddle point of L , i.e.,

$$L(x, \lambda, \nu) \leq L(\tilde{x}, \tilde{\lambda}, \tilde{\nu}) \leq L(x, \lambda, \nu)$$

for all $x \in \mathbb{R}^d$, $\lambda \in \mathbb{R}^m$ with $\lambda_i \geq 0$, and $\nu \in \mathbb{R}^n$.

The proof is left as an exercise.

