1 Materials and Methods

This section describes the theory for the encoding and Meggitt decoding of a cyclic code. Both encoding and decoding is explained with concrete examples which are also testable with the enclosed MATLAB scripts. Knowledge is obtained through TI-INCO course lessons, the "Essentials of Eror-Control Coding" text book by Jorge Castiñeira Moreira and Patrick Guy Farrell.

1.1 Encoding

Encoding for the code vector in systematic form:

$$X^{n-k}m(X) = q(X)g(X) + p(X)$$
(1)

By adding the redundancy polynomial to the shifted message polynomial, the encoded vector in systematic form is obtained.

$$c(X) = X^{n-k}m(X) + p(X)$$
(2)

A concrete example with the generator polynomial $g(X) = 1 + X^4 + X^6 + X^7 + X^8$ is used to encode the following message polynomial: $m(X) = 1 + X^2 + X^4 + X^6$, corresponding to m = [1010101]. The code is a cyclic code, $C_{cyc}(15,7)$.

First the message is right shifted n-k times.

$$X^{15-7}m(X) = X^8 + X^{10} + X^{12} + X^{14}$$
(3)

Find p(X) by taking the remainder from $X^{n-k}m(x)$ divided by g(X).

$$X^{8} + X^{7} + X^{6} + X^{4} + 1$$

$$X^{14} + X^{12} + X^{10} + X^{8}$$

$$X^{14} + X^{13} + X^{12} + X^{10} + X^{6}$$

$$X^{13} + X^{8} + X^{6}$$

$$X^{13} + X^{12} + X^{11} + X^{9} + X^{5}$$

$$X^{12} + X^{11} + X^{9} + X^{8} + X^{6} + X^{5}$$

$$X^{10} + X^{9} + X^{6} + X^{5} + X^{4}$$

$$X^{10} + X^{9} + X^{8} + X^{6} + X^{2}$$

$$X^{8} + X^{5} + X^{4} + X^{2}$$

$$X^{8} + X^{5} + X^{4} + X^{2}$$

$$X^{8} + X^{7} + X^{6} + X^{5} + X^{4} + 1$$

$$y(X) = X^{10} + X^{$$

Now c(X) can be calculated:

$$c(X) = X^{8}m(X) + p(X) = 1 + X^{2} + X^{5} + X^{6} + X^{7} + X^{8} + X^{10} + X^{12} + X^{14}$$

$$(4)$$

$$c = [101001111010101] (5)$$

It is easily seen that the encoded vector is in systematic form, hence the message vector corresponds to the last 7 bits.

1.2 Decoding