

CAE-II OSPART-13

- ⑦ → Consider 5 philosophers who spend their lives thinking and eating. The philosophers share a circular table surrounded by 5 chairs; each belonging to one philosopher.
- In center of the table is a bowl of rice, and the table is laid with 5 single chop sticks. When a philosopher thinks, he/she does not interact with her colleagues.
- From time-to-time, a philosopher gets hungry & ~~tries to~~ tries to pick up the chop stick ~~closest~~ closest to her.
- ~~A philosopher~~ They can pick only one chop stick at a time, he/she eats without realising his/her chop sticks. When she's finished eating, puts down both of chop sticks ~~or he's~~ and starts ~~the~~ thinking again.

→ The problem is considered as a classic example synchronization problem as it an example of a large class of concurrency control problems. One simple solution is to represent each chop stick with a semaphore.

→ Problems with Semaphores

① Incorrect use of semaphore operations.

① signal (mutex) ... wait (mutex)

② wait (mutex) ... wait (mutex)

③ Omitting of wait (mutex) or signal (mutex) (or both)

② Deadlock and starvation are possible.

Monitor Solution for Dining Philosopher

Monitor

Monitor Dining Phil

```
{
    enum (thinking; hungry, eating) state[5];
    condition self[5];
    void pickup (int i) {
        state[i] = hungry;
        test[i];
        if state
            if (state[i] == thinking)
                if (state[i] != eating)
                    self[i].wait;
    }
}
```

```

void putdown (int i)
{
    state[i] = thinking;
    // test left and right neighbors.
test
    test((i+4) % 5);
    test((i+1) % 5);
}

```

```

void test (int i)
{
    if ((state[(i+4) % 5] != eating) &&
        (state[i] == hungry) &&
        (state[(i+1) % 5] != eating))
    {
        state[i] = eating;
        self[i].signal();
    }
}

```

```

initialization-code()

```

```

{
    for (int i = 0; i < 5; i++)
        state[i] = thinking;
}

```


→ ~~Each~~ Each Philosopher 'i' invokes the operations pickup() and putdown() in the following sequence:

Dining Phil. ~~pick~~ pickup(i);

eat

Dining Phil. putdown(i);

→ No deadlock but starvation is possible

9(i) The following are Page Replacement 'algorithms.

(1) LIFO (First In First Out)

→ The 1st page has to be replaced first and which page has to be deleted first

(2) → Replace the pages that has been been in the RAM / ROM for the longest time per period

(left side)
(left side)

for Eg.

String: 7 0 1 2 0 3 0 4 2 3 0 3 1 2 0

Page frames															
F ₁	⑦	7	7	②	2	2	2	④	4	4	0	0	0	0	①
F ₂		①	0	0	①	③	3	3	②	2	2	2	1	1	1
F ₃			①	1	1	1	①	0	0	③	3	3	3	3	2
	*	*	*	*	hit	*	*	*	*	*	*	Present hit	*	*	hit

Page hit: Hits $\rightarrow 3$

Page faults: $\rightarrow 12$

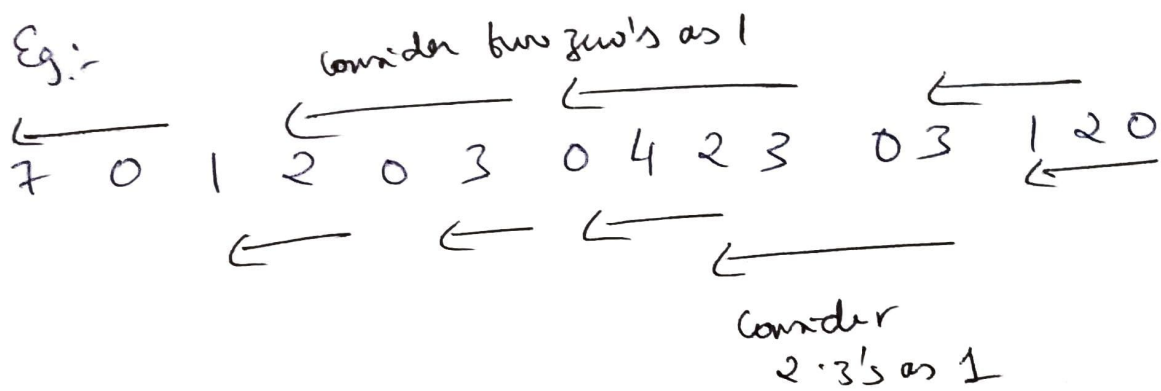
Ref string: $\rightarrow 15$

② LRU:- Last Recently Used.

\rightarrow 3 frames = 3 page no. (compare)

\rightarrow left side page no. are replaced with another

Eg:-

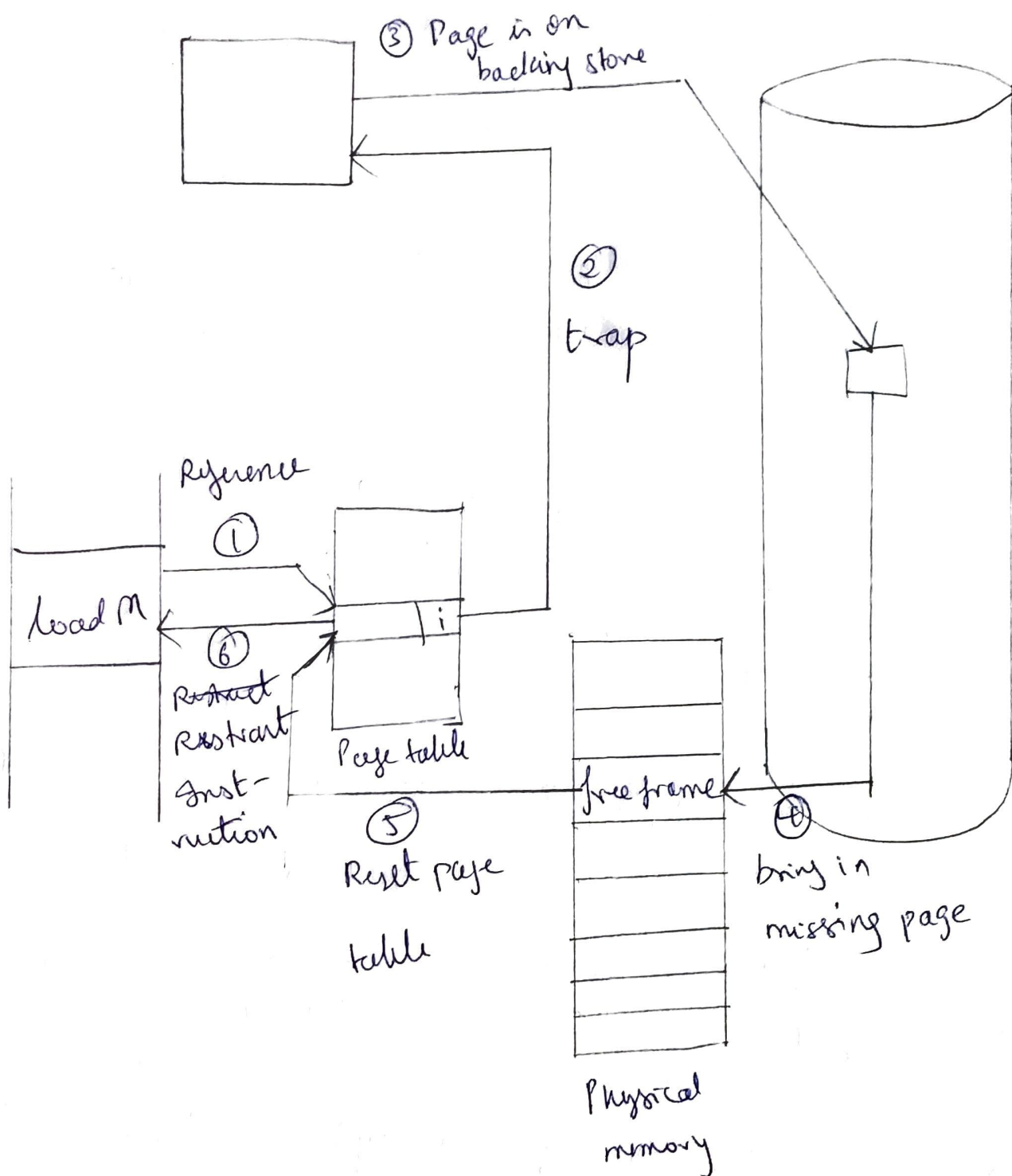


F ₁	7	7	⑦	2	2	2	②	4	4	④	0	0	①	2	2
F ₂		0	①	0	0	0	0	①	3	3	3	3	③	0	
F ₃			1	1	①	3	3	③	2	2	2	②	1	1	1
	*	*	*	*	hit	*	hit	hit	*	*	*	hit	*	*	*

page hit = 3

" fault = 12

Steps in handling a page fault are:



PART-A

① These are the process that can affect / are affected by other processes running on the system

② It is a synchronisation object that controls access by multiple processes to a common resource.

types:

① Binary Semaphore

② Counting "

③ Mutex "

③ The request for any resource will be granted if the resulting state of the system doesn't cause deadlock in the system. is called deadlock avoidance

④ It is a type of interrupt, raised by the HW when a running prog. accesses a memory page.

Control fragmentation

⑧ Internal Fragmentation

① But fit Block search is the solution

② Occurs when paging is employed.

Compaction is the solution

Occurs when segmentation is employed