



SATHYABAMA

INSTITUTE OF SCIENCE AND TECHNOLOGY

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SLR (Simple LR) Parser

Dr.Y.Bevish Jinila

Department of Information Technology



Augmented Grammar

- If G is a grammar with start symbol, then G' is the augmented grammar for G with a new start symbol S' and production $S' \rightarrow S$

PURPOSE : To indicate the parser when it should stop parsing and announce acceptance of the input



Example

$S \rightarrow AA$

$A \rightarrow aA \mid b$

Soln:

The augmented grammar is

$S' \rightarrow S$

$S \rightarrow AA$

$A \rightarrow aA \mid b$



LR(0) items

- An LR (0) item is a **production G with dot at some position on the right side of the production.**
- LR (0) items is useful to indicate that how much of the input has been scanned up to a given point in the process of parsing.



Example

$S \rightarrow AA$

LR(0) items

$S \rightarrow \cdot AA$ (Not seen anything on right)

$S \rightarrow A \cdot A$ (Seen A, yet to see another A)

$S \rightarrow AA \cdot$ (Seen everything)



Construction of SLR Parsing Table

- Compute two functions
 - Closure
 - Goto

Closure operation

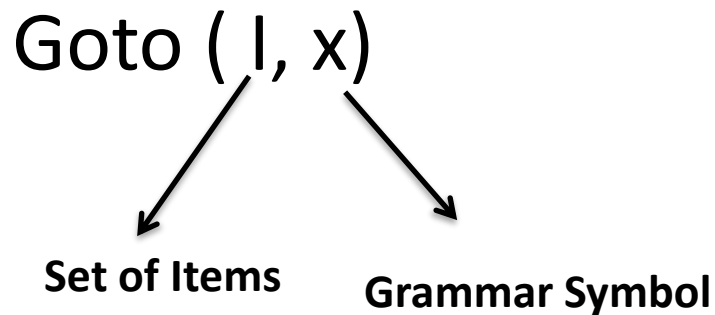
If I is a set of items for a grammar G , then $\text{closure}(I)$ is the set of items constructed from I by the two rules.

- (1) Initially, every item in I is added to $\text{closure}(I)$
- (2) If $A \rightarrow \alpha.B\beta$ is in $\text{closure}(I)$ and $B \rightarrow \gamma$ is a production, then add the item $B \rightarrow \cdot\gamma$ to I , if it is not already there.

(We apply this rule until no more new items can be added to $\text{closure}(I)$)



Goto operation



It is defined to be the closure of the set of all items $[A \rightarrow \alpha x. \beta]$ such that $[A \rightarrow \alpha. x \beta]$ is in I



Problem

Construct the SLR parsing table and parse the string abab for the following grammar

$S \rightarrow AA$

$A \rightarrow aA \mid b$

Soln:

1. Augmented Grammar

$S' \rightarrow S$

$S \rightarrow AA$

$A \rightarrow aA \mid b$

LR (0) Items

I_0

$S' \rightarrow \cdot S$
 $S \rightarrow \cdot AA$
 $A \rightarrow \cdot aA$
 $A \rightarrow \cdot b$

Goto (I_0 , S)

I_1

$S' \rightarrow S \cdot$

Goto (I_0 , A)

I_2

$S \rightarrow A \cdot A$
 $A \rightarrow \cdot aA$
 $A \rightarrow \cdot b$

Goto (I_0 , a)

I_3

$A \rightarrow a \cdot A$
 $A \rightarrow \cdot aA$
 $A \rightarrow \cdot b$

Goto (I_0 , b)

I_4

$A \rightarrow b \cdot$

Goto (I_2 , A)

I_5

$S \rightarrow AA \cdot$

Goto (I_2 , a)

I_3

$A \rightarrow a \cdot A$
 $A \rightarrow \cdot aA$
 $A \rightarrow \cdot b$

Goto (I_2 , b)

I_4

$A \rightarrow b \cdot$

Goto (I_3 , A)

I_6

$A \rightarrow aA \cdot$

Goto (I_3 , a)

I_3

$A \rightarrow a \cdot A$
 $A \rightarrow \cdot aA$
 $A \rightarrow \cdot b$

Goto (I_3 , b)

I_4

$A \rightarrow b \cdot$

SLR Parsing Table

$S \rightarrow AA$ -----(1)

$A \rightarrow aA$ -----(2)

$A \rightarrow b$ -----(3)



State	Action			Goto	
	a	b	\$	S	A
0	S3	S4		1	2
1			Accept		
2	S3	S4			5
3	S3	S4			6
4	R3	R3	R3		
5			R1		
6	R2	R2	R2		

Reduce Action

I4 :A -> b.

Follow(A)=
FIRST(A)= {a,b,\$}

I5: S->AA.

Follow(S)= {\$}

I6: A -> aA .

Follow(A)=
FIRST(A)= {a,b,\$}



Parsing

Actions : Shift, Reduce, Accept, Error

Stack	Input	Action
0	abab \$	Shift (S3)
0a3	bab \$	Shift (S4)
0a3b4	ab \$	Reduce (R3) A -> b
0a3A6	ab \$	Reduce (R2) A -> aA
0A2	ab \$	Shift (S3)
0A2a3	b\$	Shift (S4)
0A2a3b4	\$	Reduce (R3) A -> b
0A2a3A6	\$	Reduce (R2) A -> aA
0A2A5	\$	Reduce (R1) S -> AA
0S1	\$	Accept

State	Action			Goto	
	a	b	\$	S	A
0	S3	S4		1	2
1			Accept		
2	S3	S4			5
3	S3	S4			6
4	R3	R3	R3		
5			R1		
6	R2	R2	R2		



Thank you