

CSE 6140/ CX 4140
Computational Science and Engineering
ALGORITHMS

Coping with NP-completeness - 4
Local Search

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Based on slides by Prof. Ümit V. Çatalyürek and Bistra Dilkina

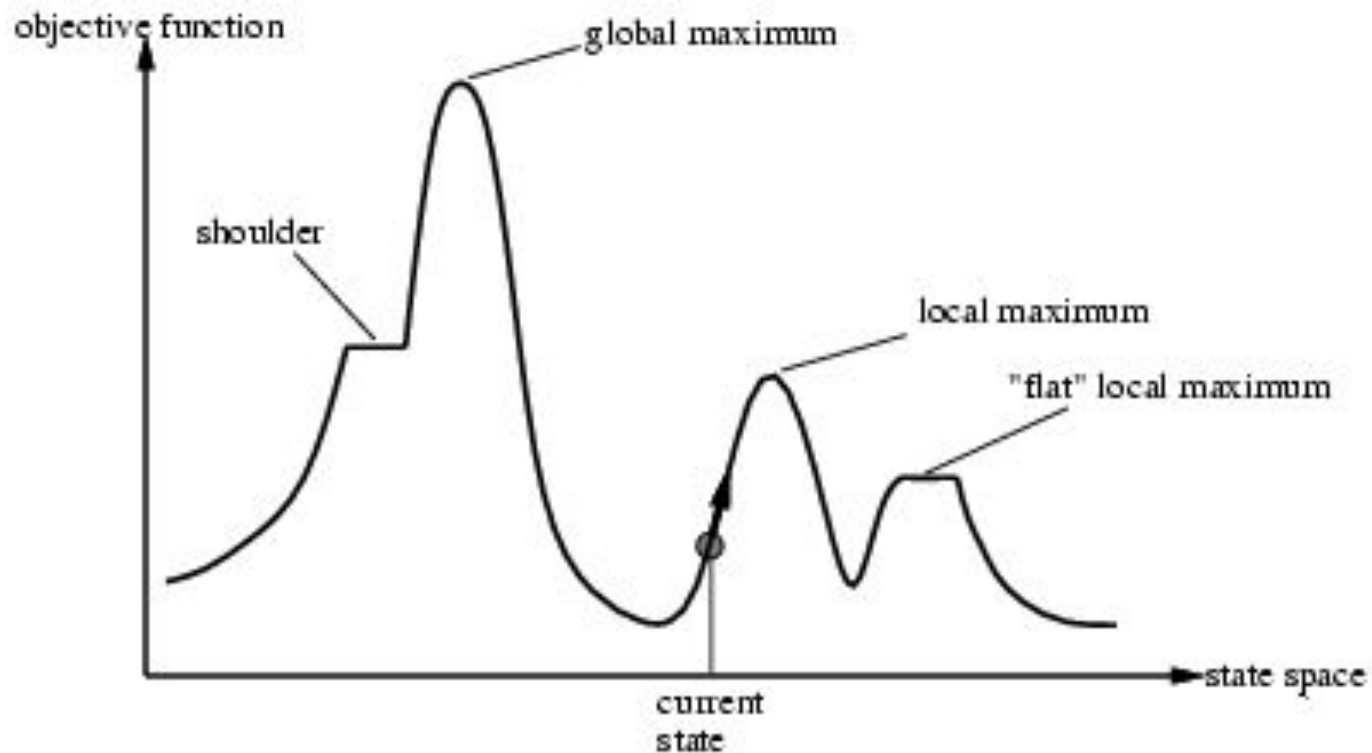
HEURISTICS/LOCAL SEARCH

Local Search (LS)

- Start from initial position
- Iteratively move from current position to one of neighboring positions
- Use evaluation function to choose among neighboring positions

State space landscape

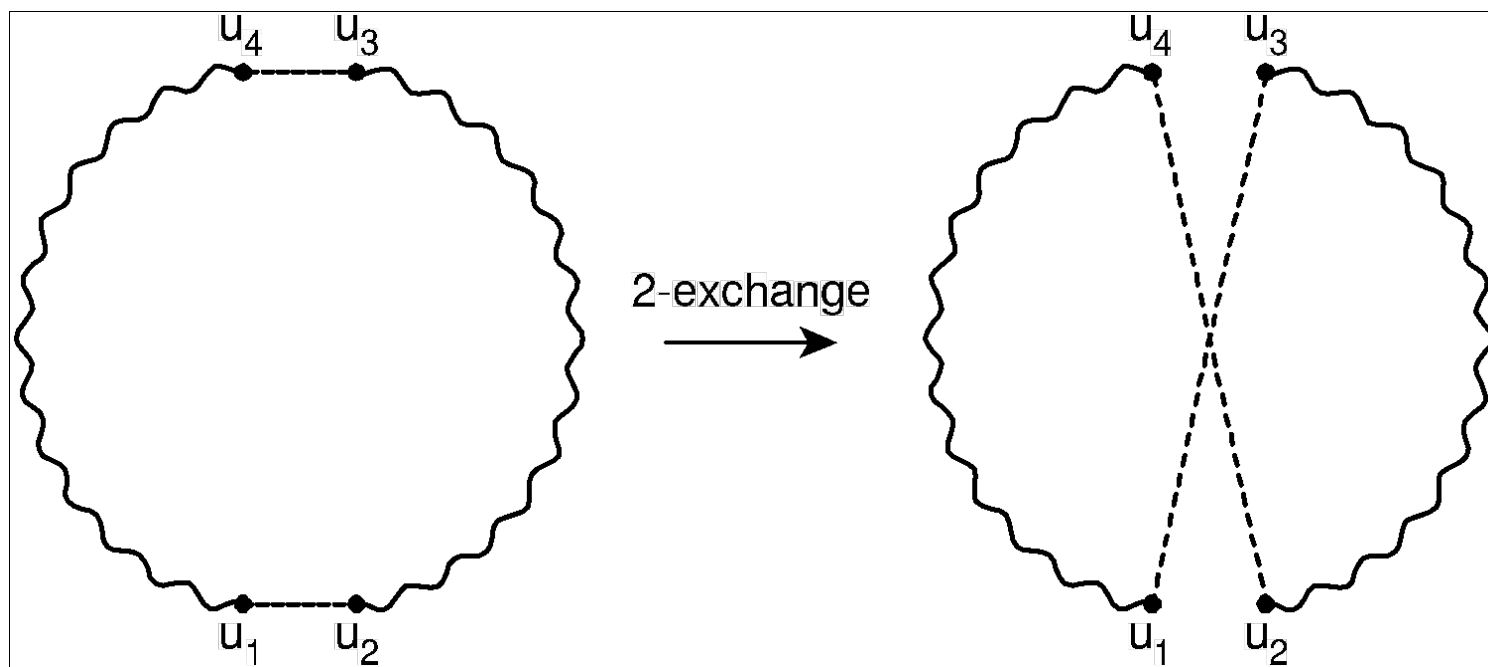
- Objective function defines *state space landscape*



Local Search (LS) Algorithms

- search space S
 - TSP: set of all permutations of vertices (all “potential solutions”)
- solution set $S' \subseteq S$
 - TSP: the tours of minimum length
- neighborhood relation $N \subseteq S \times S$
 - A way to move from one potential solution to another
 - TSP: neighboring tour differ in several edges
- evaluation function $g : S \rightarrow \mathbb{R}^+$
 - TSP: length of tour

Symm. TSP --- search neighborhood



Search Space: all permutations of the cities (each defines a cycle)

3-opt – delete 3 edges and reconnect fragments into 1 cycle

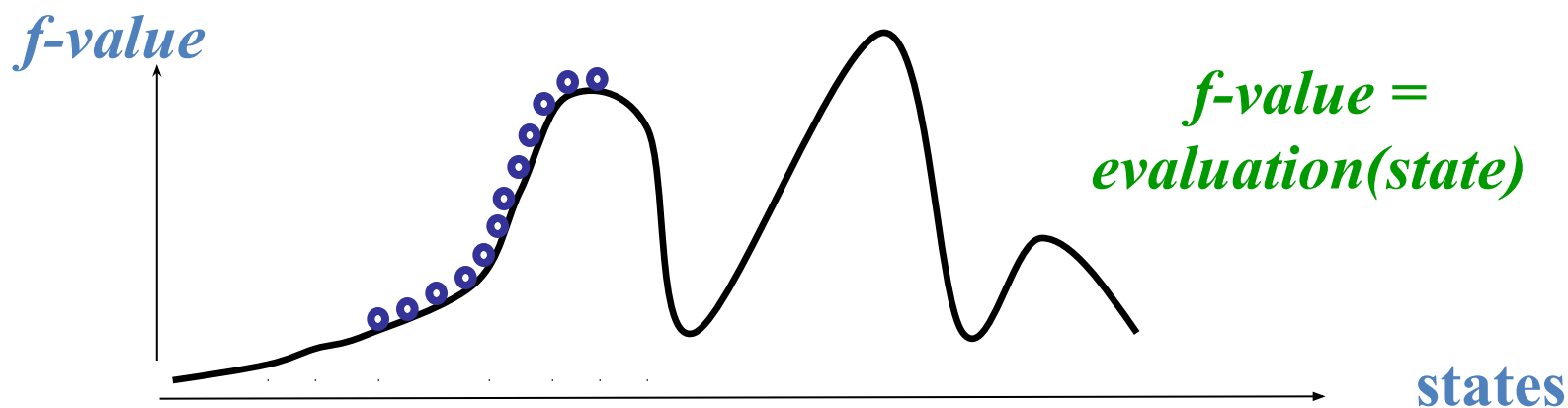
k-opt – delete k edges and reconnect fragments into 1 cycle

Iterative Improvement (Greedy Search)

- Initialize search at some point of search space
- In each step, move from the current search position to a neighboring position with better evaluation function value

Hill climbing (Best Improvement Search)

- Choose the neighbor with the largest improvement as the next state



```
while f-value(state) < f-value(next-best(state))  
    state := next-best(state)
```


Hill climbing

function Hill-Climbing(*problem*) **returns** a *solution state*

current \leftarrow Make-Node(Initial-State[*problem*])

loop do

next \leftarrow a highest-valued successor of *current*

if Value[*next*] < Value[*current*] **then return** *current*

current \leftarrow *next*

end

Problems with iterative improvement

- Advantages:
 - Very fast, works well for certain problems

- Disadvantages:
 - What if there are multiple peaks?
 - Hill climbing gets stuck at all peaks, known as local maxima
 - Optimal solution is highest peak – global maximum
 - May result in extremely suboptimal solution if many peaks
 - Being misguided by evaluation/objective function

Stochastic Local Search

- Randomize initialization step
- Randomize search steps such that suboptimal/worsening steps are allowed
- Improved performance & robustness
- Typically, degree of randomization controlled by noise parameter

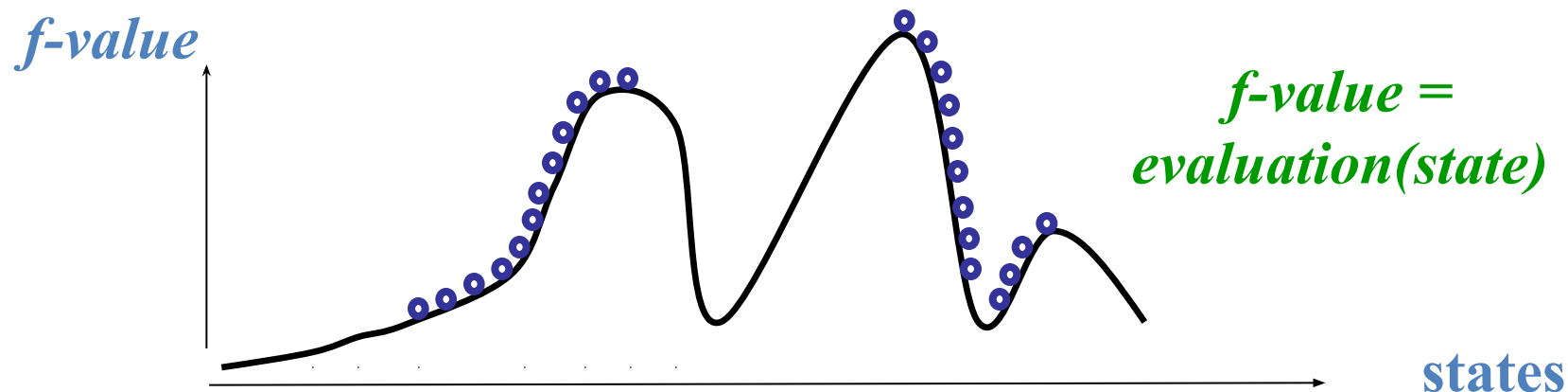
Stochastic Local Search

- Pros:
 - for many combinatorial problems more efficient than systematic search
 - easy to implement
 - easy to parallelize

- Cons:
 - often incomplete (no guarantees for finding existing solutions)
 - highly stochastic behavior
 - often difficult to analyze theoretically/empirically

Random restart hill climbing

- Start at random solution
- Hill-climb until local optima
- Start at another random position



Randomized Iterative Improvement:

- *Idea*: escape local maxima by allowing some "bad" moves
- *Initialize* search at some point of search space
- *Search steps*:
 - With probability p , move from current search position to a randomly selected neighboring position
 - Otherwise, move from current search position to neighboring position with better evaluation function value.

The WalkSAT Algorithm

```
WalkSAT(CNF, max-tries, max-flips, p) {  
    for i ← 1 to max-tries do  
        solution = random truth assignment  
        for j ← 1 to max-flips do  
            if all clauses in CNF satisfied then  
                return solution  
            c ← random unsatisfied clause in CNF  
            with probability p  
                flip a random variable in c  
            else  
                flip variable in c that maximizes  
                    number of satisfied clauses  
        return failure  
}
```

Simulated annealing (SA)

- Combinatorial search technique inspired by the physical process of annealing [Kirkpatrick et al. 1983, Cerny 1985]
- Outline
 - Select a neighbor at random.
 - If better than current state, go there (improving move).
 - Otherwise, go there with some probability (worsening move).
 - **Probability goes down** with time (similar to temperature cooling)
- When probability is high -> diversify (many worsening moves)
- When probability is low -> intensify (focus on improving moves)

SA analogy

- Annealing is process of heating and cooling metals in order to improve strength
- Idea: Controlled heating and cooling of metal
 - When hot, atoms move around
 - When cooled, atoms find configuration with lower internal energy (i.e. makes metal stronger)
- Analogy:
 - Temperature = probability of accepting worse neighboring solution
 - When temperature is high, likely to accept worse neighboring solutions (but may lead to better overall solution)
 - Cooling represents shrinking probability of accepting worse solutions

SA Pseudo code

function Simulated-Annealing(*problem, schedule*) **returns**
solution state

current \leftarrow Make-Node(Initial-State[*problem*])

for *t* \leftarrow 1 **to** *infinity* (*Iters, Time cutoff*)

T \leftarrow *schedule*[*t*] // *T* (a function of time *t*) goes downwards.

if *T* = 0 **then** **return** *greedy from current*

next \leftarrow Random-Successor(*current*)

ΔE \leftarrow f-Value[*next*] - f-Value[*current*]

if $\Delta E > 0$ **then** *current* \leftarrow *next*

else *current* \leftarrow *next* with probability $e^{\Delta E/T}$

end

Simulated annealing (SA)

- **Acceptance criterion (Metropolis condition):** choose new solution s' over old solution s with probability (maximization)

$$\Pr(s', s) = \begin{cases} 1 & \text{if } f(s') > f(s) \\ \exp\left\{\frac{f(s') - f(s)}{T}\right\} & \text{otherwise} \end{cases}$$

- **Initial temperature** T_0
- **Annealing (cooling) schedule:** how to update the temperature
 - E.g., $T = a T$ with $a = 0.95$ (geometric schedule)
 - Number of iterations, neighbourhood size
- **Stopping criterion**
 - E.g., no improved solution found for a number of iterations (or number of temperature values)

SA for TSP [Johnson & McGeoch 1997]

- Baseline implementation:
 - start with random initial solution
 - use 2-exchange neighborhood
 - simple annealing schedule;
- -> relatively poor performance

- Improvements:
 - look-up table for acceptance probabilities
 - neighborhood pruning
 - low-temperature starts

SA with restarts

- **RESTARTS:**

Sometimes it is better to move back to a solution that was significantly better rather than always moving from the current state.

- The decision to restart could be based on several criteria.
 - based on a fixed number of steps
 - the current solution is much worse than the best so far
 - too many iterations without improvement
 - restarting randomly, etc.

Summary of Simulated Annealing

- Historically important
- Easy to implement
- Has interesting theoretical properties (convergence), but these are of very limited practical relevance
- Achieves good performance (often at the cost of substantial run-times)

Tabu Search

- Combinatorial search technique which heavily relies on the use of an explicit memory of the search process to guide search process [Glover 1989, 1990]
- Memory typically contains only specific attributes of previously seen solutions
- Simple tabu search strategies exploit only short term memory
- More complex tabu search strategies exploit long term memory

Next lecture

- More on Tabu search
- Iterated local search
- Constructing initial solutions

- Approximation algorithms

- Hw4 deadline extended:
Nov. 18, 11:59pm (instead of Nov. 13)
hard deadline, no grace period!