

DP from Midterm 2016

Pb States  $\{1 \dots n\}$ , state  $i$ :  $p_i$  (pop)  
 $v_i$  (votes)

$$V = \underbrace{\left( \sum_{i=1}^n v_i \right)}_{V_{\text{tot}}} / 2 + 1$$

required to win

Goal: Find set of states  $S$  that minimizes

$$\sum_{i \in S} p_i \quad \text{s.t.} \quad \sum_{i \in S} v_i \geq V$$

$\text{MinPop}(i, v)$ : Min pop. of a subset of states  $\{1, \dots, i\}$  s.t. their votes sum to  $v$  (exactly)

a) Prove opt. substructure.

Suppose we have opt. sol. for obtaining  $v$  using states  $\{1..i\}$ , a value  $\text{MinPop}(i, v)$

→  $S(i, v)$  is the set of states included in the vote for this sol.

In a solution, either  $i \in S(i, v)$  or  $i \notin S(i, v)$

• If  $i \in S(i, v)$ ,  $S(i, v) \setminus \{i\}$  is a set of states

with exactly  $v - v_i$  votes. If this set is not

the opt. (for states  $\{1..i-1\}$ ), then we replace it with the opt.  $S(i-1, v - v_i)$ , add  $i$  to this set

→ another sol with exactly  $v$  votes but less population

• If  $i \notin S(i, v)$ , the  $S(i, v)$

must be opt. way to get  $v$  votes

from  $1..i-1$  (otherwise build better sol)

b/



19

e/

For  $j=1$  to  $n$ ,  $\{ V_{tot} += V_j; P_{tot} += P_j \}$

**// base cases**

For  $i = 0$  to  $n$ ,  $s[i, 0] = 0$ ;

For  $v = 1$  to  $V_{tot}$ ,  $S[0, v] = P_{tot} + 1$ ; // (+∞)

```
// loop
```

For  $i = 1$  to  $n$

For  $v = 1$  to  $V_{tot}$

$$S[i, v] = \min(S[i-1, v], S[i-1, v-v_i] + p_i)$$

1 (check  $v \geq v_i$  to add)

// sol.

$$Min = Phot + 1$$

For  $v = V$  to  $V_{hot}$  {  $min = S[n, v]$  }