

# Self-Affine Sets with Positive Lebesgue Measure

Karma Dajani, Kan Jiang and Tom Kempton

February 14, 2014

## Abstract

Using techniques introduced by C. Güntürk, we prove that the attractors of a family of overlapping self-affine iterated function systems contain a neighbourhood of zero for all parameters in a certain range. This corresponds to giving conditions under which a single sequence may serve as a ‘simultaneous  $\beta$ -expansion’ of different numbers in different bases.

## 1 Introduction

Given real numbers  $1 < \beta_1 < \beta_2$ , we define contractions  $T_i : \mathbb{R}^2 \rightarrow \mathbb{R}^2$  by

$$T_i(x, y) = \left( \frac{x + i}{\beta_1}, \frac{y + i}{\beta_2} \right).$$

A classical result of Hutchinson [4] asserts that there exists a unique non-empty compact set  $A_{\beta_1, \beta_2}$  satisfying

$$A_{\beta_1, \beta_2} = T_{-1}(A_{\beta_1, \beta_2}) \cup T_1(A_{\beta_1, \beta_2}).$$

If  $\beta_1 \neq \beta_2$  then the contractions  $T_i$  are affine contractions and  $A_{\beta_1, \beta_2}$  is termed a self-affine set. Since  $\beta_1, \beta_2 < 2$ , the two contracted copies  $T_{-1}(A_{\beta_1, \beta_2})$  and  $T_1(A_{\beta_1, \beta_2})$  overlap. There are many fundamental open questions about the structure of overlapping self-affine sets, see for example [5, 6, 7].

The family  $A_{\beta_1, \beta_2}$  of sets was studied in [7], where Shmerkin proved that there exists an open set  $K \subset (1, 2)^2$  such that for almost every pair  $(\beta_1, \beta_2) \in K$  the corresponding set  $A_{\beta_1, \beta_2}$  has positive Lebesgue measure. This was done by studying the absolute continuity of a certain measure defined on  $A_{\beta_1, \beta_2}$ . In this article we prove that  $A_{\beta_1, \beta_2}$  contains a neighbourhood of  $(0, 0)$  for all  $(\beta_1, \beta_2) \in (1, 1 + C)^2$  for some positive constant  $C$  which is explicitly defined later.

In fact this problem is closely related to the problem of ‘simultaneous  $\beta$ -expansions’ studied by Güntürk in [3]. Given  $\beta \in (1, 2)$  and  $x \in [\frac{-1}{\beta-1}, \frac{1}{\beta-1}]$ , a  $\beta$ -expansion of  $x$  is a sequence  $\underline{a} \in \{-1, 1\}^{\mathbb{N}}$  for which

$$\sum_{n=1}^{\infty} a_n \beta^{-n} = x.$$

This definition can be extended to  $\beta > 2$  by letting the digits  $a_n$  come from a larger digit set.

For typical  $x$  the  $\beta$ -expansion of  $x$  is not unique, indeed almost every  $x \in [\frac{-1}{\beta-1}, \frac{1}{\beta-1}]$  has uncountably many  $\beta$ -expansions, see [8]. This allows one, given  $x$ , to search for  $\beta$ -expansions of  $x$  with interesting properties, such as a given digit frequency or that the sequence is a  $\beta$ -expansion of  $x$  for more than one  $\beta$ .

In [3], Güntürk proved that given  $\beta_1, \beta_2 > 1$  and  $(x_1, x_2) \in \mathbb{R}^2$  there exists a sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  satisfying

$$\sum_{n=1}^{\infty} a_n \beta_k^{-n} = x_k$$

for each  $k \in \{1, 2\}$  whenever a certain algorithm can be implemented, see Proposition 2.1. It was claimed without proof<sup>1</sup> that there exist constants  $C, \delta > 0$  such that the algorithm can be implemented whenever  $\beta_1, \beta_2 \in (1, 1 + C)$  and  $(x_1, x_2) \in (-\delta, \delta)^2$ . We prove this fact and provide suitable constants  $C$  and  $\delta$  explicitly. We also prove a number of related results including results on finding  $\beta$ -expansions with given digit frequency and finding sequences which serve as multiple expansions for a range of  $\beta_1, \beta_2$ . An interesting facet of our work is that the techniques of Güntürk which we use are quite distinct from the usual fractal geometry techniques for studying self-affine sets.

The following is our main theorem.

**Theorem 1.1.** *There exists a constant  $C \approx 0.05$  such that for any  $1 < \beta_1 < \beta_2 < 1 + C$ , there exists  $\delta = \delta(\beta_1, \beta_2)$  such that for any pair  $(x_1, x_2) \in [-\delta, \delta]^2$ , there exists a sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  such that*

$$\left( \sum_{n=1}^{\infty} a_n \beta_1^{-n}, \sum_{n=1}^{\infty} a_n \beta_2^{-n} \right) = (x_1, x_2). \quad (1)$$

In the self-affine setting, this theorem corresponds to saying that the sequence  $(a_n)$  is a coding of the pair  $(x_1, x_2)$  in  $A_{\beta_1, \beta_2}$ , and in particular that  $(x_1, x_2) \in A_{\beta_1, \beta_2}$ . This leads immediately to the following corollary.

**Corollary 1.1.** *For all any  $1 < \beta_1 < \beta_2 < 1 + C$  we have that the self-affine fractal  $A_{\beta_1, \beta_2}$  contains a neighbourhood of  $(0, 0)$ .*

The constant  $\delta$  is explicitly computable. If  $\beta_1$  tends to  $\beta_2$  the constant  $\delta$  tends to zero.

**Remark 1.1.** *An important special case of Corollary 1.1 is the case  $x_1 = x_2$ . This was the main motivation of Güntürk for his original article because of its relevance to analogue digital conversion, see [3]. While in general the constant  $\delta$  depends on  $\beta_1, \beta_2$ , in the case that  $x_1 = x_2$  we can choose  $\delta = 0.16$  independently of  $\beta_1, \beta_2$  to give that for all  $1 < \beta_1 < \beta_2 < 1 + C$  and  $x \in [-0.16, 0.16]$  there exists a sequence  $(a_i) \in \{-1, 1\}^{\mathbb{N}}$  such that*

$$x = \sum_{n=1}^{\infty} \frac{a_n}{\beta_1^n} = \sum_{n=1}^{\infty} \frac{a_n}{\beta_2^n}.$$

---

<sup>1</sup>Güntürk stated in [3] that details would be provided in a later publication, but has confirmed to us that, due to other commitments, no such publication will be forthcoming. Since the techniques of [3] are rather different from the standard techniques for analysing self-affine sets, and the results are interesting, we take the liberty of providing a proof of the stated results of Güntürk in this article.

Using the same techniques, we can also find  $\beta$ -expansions of real numbers which have certain given digit frequencies. It was stated in [3] that the following theorem should follow by suitably adapting the proof of Theorem 1.1, we provide the appropriate adaptation and prove the result giving explicit constants.

**Theorem 1.2.** *Let  $C_1 > 0$  satisfy  $(1+C_1)+2(1+C_1)^3 = 6$ . Then for all  $1 < \beta < 1+C_1$ , there exists  $\delta = \delta(\beta)$  such that for any  $x \in [-\delta, \delta]$  there exists a sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  satisfying*

$$\sum_{n=1}^{\infty} a_n \beta^{-n} = x \quad (2)$$

and

$$x = \lim_{n \rightarrow \infty} \frac{a_1 + a_2 + \cdots + a_n}{n}. \quad (3)$$

One can read off limiting digit frequencies of the sequence  $(a_n)$  from equation 3 by noting that

$$\frac{1 - \frac{a_1 + a_2 + \cdots + a_n}{n}}{2} = \frac{|\{k \in \{1, \dots, n\} : a_k = -1\}|}{n}.$$

Proofs of Theorems 1.1 and 1.2 are given in the next two sections. In the final section we state some further corollaries and remarks.

## 2 Proof of Theorem 1.1

As stated in the introduction, we are using many of the ideas of [3]. For clarity, we have amalgamated these ideas to form the following proposition, which was proved in [3]. The remainder of our proof of Theorem 1.1, which gives conditions under which the algorithm in Proposition 2.1 can be implemented, is new.

**Proposition 2.1.** *Given  $1 < \beta_1 < \beta_2 < 2$  and  $(x_1, x_2) \in \mathbb{R}^2$  suppose that one can implement the following algorithm.*

1. *For  $L > 2$  pick real numbers  $h_1, \dots, h_L$  with  $h_L \neq 0$  and*

$$h_{L-1} = h_{L-2} = 0,$$

*such that  $\beta_1, \beta_2$  are roots of the polynomial  $P(z) = z^L - \sum_{k=1}^L h_k z^{L-k}$ .*

2. *Pick real numbers  $u_{-L+1}, u_{-L+2}$  which satisfy the equation*

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = h_L \begin{pmatrix} \beta_1^{-1} & \beta_1^{-2} \\ \beta_2^{-1} & \beta_2^{-2} \end{pmatrix} \begin{pmatrix} u_{-L+1} \\ u_{-L+2} \end{pmatrix}$$

*Set  $u_{-L+3} = \cdots = u_0 = 0$ .*

3. *Find a sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  such that*

$$u_n := \sum_{k=1}^L h_k u_{n-k} - a_n$$

*satisfies  $u_n \in [-1, 1]$  for each  $n \in \mathbb{N}$ .*

Then the sequence  $(a_n)_{n=1}^{\infty}$  will satisfy equation (1).

In this article we give rigorous conditions under which the algorithm of Güntürk can be implemented leading to a proof of Theorem 1.1. For completeness we also give the proof of Proposition 2.1. We begin by introducing the polynomial  $P$ ,  $P$  was chosen because it has relatively low degree and satisfies the conditions of Proposition 2.1, but it is likely that better bounds on  $C$  and  $\delta$  can be obtained by choosing a better polynomial  $P$ .

**Definition 2.1.** Given  $\beta_1, \beta_2 > 1$ , we define the polynomial  $P$  by

$$P(x) = x^4 - h_1x^3 - h_2x^2 - h_3x - h_4$$

where

$$\begin{aligned} h_1 &= \frac{(\beta_1 + \beta_2)(\beta_1^2 + \beta_2^2)}{\beta_1^2 + \beta_1\beta_2 + \beta_2^2} \\ h_2 &= 0 \\ h_3 &= 0 \\ h_4 &= \frac{-(\beta_1\beta_2)^3}{\beta_1^2 + \beta_1\beta_2 + \beta_2^2}. \end{aligned}$$

We further define the constant  $C$  by

$$C := \sqrt[3]{\sqrt{10} - 2} \approx 0.05.$$

**Lemma 2.1.** The polynomial  $P$  satisfies  $P(\beta_1) = P(\beta_2) = 0$ .

*Proof.* Defining,

$$b = \frac{\beta_1\beta_2(\beta_1 + \beta_2)}{\beta_1^2 + \beta_1\beta_2 + \beta_2^2},$$

and

$$c = \frac{(\beta_1\beta_2)^2}{\beta_1^2 + \beta_1\beta_2 + \beta_2^2}$$

gives us that

$$(x - \beta_1)(x - \beta_2)(x^2 + bx + c) = x^4 - h_1x^3 - h_2x^2 - h_3x - h_4 = P(x).$$

Then  $\beta_1$  and  $\beta_2$  are roots of  $P$ . □

**Lemma 2.2.** For  $\beta_1, \beta_2 \in (1, 1 + C)$  we have that

$$\sum_{n=1}^4 |h_n| = |h_1| + |h_4| \leq 2.$$

*Proof.* Expanding out, we see that

$$\begin{aligned} \sum_{n=1}^4 |h_n| &= |h_1| + |h_4| \\ &= \frac{(\beta_1 + \beta_2)(\beta_1^2 + \beta_2^2) + \beta_1^3\beta_2^3}{\beta_1^2 + \beta_1\beta_2 + \beta_2^2} \\ &\leq \frac{(2 + 2C)2(1 + C)^2 + (1 + C)^6}{3} \\ &\leq 2 \end{aligned}$$

whenever  $\beta_1, \beta_2 \in (1, 1+C)$ , as required. Indeed,  $C$  was chosen to be the largest constant such that the above inequalities hold.  $\square$

We now prove Theorem 1.1 using Proposition 2.1.

*Proof.* We set

$$\begin{aligned} u_{-3} &= \frac{\beta_1^2 + \beta_1\beta_2 + \beta_2^2}{(\beta_2 - \beta_1)\beta_1\beta_2} \left( \frac{x_1}{\beta_2^2} - \frac{x_2}{\beta_1^2} \right), \\ u_{-2} &= \frac{\beta_1^2 + \beta_1\beta_2 + \beta_2^2}{(\beta_2 - \beta_1)\beta_1\beta_2} \left( \frac{x_2}{\beta_1} - \frac{x_1}{\beta_2} \right), \\ u_{-1} &= u_0 = 0. \end{aligned}$$

These choices of  $u_i$  ensure that condition (2) of Proposition 2.1 is satisfied, i.e.

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = h_4 \begin{pmatrix} \beta_1^{-1} & \beta_1^{-2} \\ \beta_2^{-1} & \beta_2^{-2} \end{pmatrix} \begin{pmatrix} u_{-3} \\ u_{-2} \end{pmatrix}.$$

Condition (1) has already been shown to hold for our choice of  $P$  by Lemma 2.1. It remains to show that condition (3) holds, i.e. that one can choose some sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  such that defining  $u_n$  for  $n \in \mathbb{N}$  by

$$u_n := \sum_{k=1}^L h_k u_{n-k} - a_n \quad (4)$$

gives  $u_n \in [-1, 1]$  for each  $n \in \mathbb{N}$ . Since  $h_2 = h_3 = 0$  the above equation for  $u_n$  becomes

$$u_n = h_1 u_{n-1} + h_4 u_{n-4} - a_n.$$

We set

$$a_n = \begin{cases} -1 & h_1 u_{n-1} + h_4 u_{n-4} < 0 \\ +1 & h_1 u_{n-1} + h_4 u_{n-4} \geq 0 \end{cases}. \quad (5)$$

Now we observe that, if for some  $k \in \mathbb{N}$  one has that  $u_{k-1}, u_{k-4} \in [-1, 1]$ , then it follows from Lemma 2.2 that

$$h_1 u_{k-1} + h_4 u_{k-4} \in [-2, 2].$$

Hence it follows that

$$u_k := h_1 u_{k-1} + h_4 u_{k-4} - a_k \in [-1, 1],$$

and hence by induction that  $u_n \in [-1, 1]$  for each  $n \in \mathbb{N}$ .

Now we define

$$\delta = \frac{\beta_1^2 \beta_2^2 (\beta_2 - \beta_1)}{(\beta_1^2 + \beta_1 \beta_2 + \beta_2^2)(\beta_1 + \beta_2)}.$$

We see that  $\delta > 0$  whenever  $\beta_2 > \beta_1$ , but that  $\delta \rightarrow 0$  as  $\beta_2 - \beta_1 \rightarrow 0$ . From the definition of  $u_{-3}, u_{-2}$  we see that for  $x_1, x_2 \in [-\delta, \delta]^2$  and  $\beta_1, \beta_2 \in (1, 1+C)$  we have that  $u_{-3}, u_{-2} \in [-1, 1]$ . Since  $u_{-1} = u_0 = 0$  it follows by induction that  $u_n \in [-1, 1]$  for each  $n \in \mathbb{N}$ . Hence conditions (1), (2) and (3) of Proposition 2.1 are satisfied, and so the sequence  $(a_n)$  satisfies equation 1 and Theorem 1.1 is proved.  $\square$

It remains only to give a formal proof of Proposition 2.1.

*Proof.* We give a proof for the case  $L = 4$ , which is the case that we have used. From condition (2), we have that

$$x_i = h_4(u_{-3}\beta_i^{-1} + u_{-2}\beta_i^{-2})$$

for  $i = 1, 2$ . Rewriting condition (3) gives us that  $a_n = \sum_{k=0}^4 h_k u_{n-k}$ . Then summing gives us that

$$\sum_{n=1}^{\infty} a_n \beta_i^{-n} = \sum_{n=1}^{\infty} \sum_{k=0}^4 h_k u_{n-k} \beta_i^{-n} \quad (6)$$

where  $h_0 = -1$ . Since the sequence  $(u_n)$  is bounded, we have by Fubini's theorem that

$$\sum_{n=1}^{\infty} a_n \beta_i^{-n} = \sum_{k=0}^4 \sum_{n=1}^{\infty} h_k u_{n-k} \beta_i^{-n} = \sum_{k=0}^4 h_k \beta_i^{-k} \sum_{n=-k+1}^{\infty} u_n \beta_i^{-n}$$

Here the first equality involved using equation (6) and swapping the order of summation by Fubini. The second equality is just a change of variables. Now, by separating the terms for positive and negative  $n$  in the right hand side of the above equation, we have that

$$\begin{aligned} \sum_{n=1}^{\infty} a_n \beta_i^{-n} &= \left( \sum_{k=0}^4 h_k \beta_i^{-k} \right) \left( \sum_{n=1}^{\infty} \frac{u_n}{\beta_i^n} \right) + h_1 \beta_i^{-1} u_0 + h_2 \beta_i^{-2} (u_0 + u_{-1} \beta_i) \\ &+ h_3 \beta_i^{-3} (u_0 + u_{-1} \beta_i + u_{-2} \beta_i^2) + h_4 \beta_i^{-4} (u_0 + u_{-1} \beta_i + u_{-2} \beta_i^2 + u_{-3} \beta_i^3). \end{aligned}$$

Since  $\beta_i$  is the root of  $P(x)$  we have  $\sum_{k=0}^4 h_k \beta_i^{-k} = 0$  and so the first term vanishes. From conditions (1) and (2), we have  $u_{-1} = u_0 = h_2 = h_3 = 0$ . Then, removing the zero terms, the right hand side of the above equation becomes  $h_4(u_{-3}\beta_i^{-1} + u_{-2}\beta_i^{-2})$ , which by condition (2) is equal to  $x_i$ . We conclude that

$$\sum_{n=1}^{\infty} a_n \beta_i^{-n} = h_4(u_{-3}\beta_i^{-1} + u_{-2}\beta_i^{-2}) = x_i$$

as required. This completes the proof of Proposition 2.1. □

Finally we comment that in the case that  $x_1 = x_2$  we can give values of  $\delta$  which are independent of  $\beta_1, \beta_2 \in (1, 1+C)$ . Our bound on  $\delta$  was to ensure that  $u_{-2}, u_{-3} \in [-1, 1]$ . If  $x_1 = x_2$  then

$$\begin{aligned} |u_{-2}| \leq |u_{-3}| &= \left| x_1 \left( \frac{\beta_1^2 + \beta_1 \beta_2 + \beta_2^2}{(\beta_1^3 \beta_2^3)} \right) (\beta_1 + \beta_2) \right| \\ &\leq |x_1| 6(1+C)^3 \leq 1 \end{aligned}$$

whenever  $|x_1| \leq \delta = \frac{1}{6(1+C)^3} \approx 0.16$ .

### 3 $\beta$ -expansions with a given digit frequency.

With some modifications, the algorithm used in the proof of Proposition 2.1 can also be utilized to prove Theorem 1.2. The following is analagous to Proposition 2.1.

**Proposition 3.1.** *Given  $1 < \beta < 2$  and  $x \in [-\delta, \delta]$  for some  $\delta$  which will be set in the process of proof, suppose that one can implement the following algorithm.*

1. For  $L > 2$  pick real numbers  $h_1, \dots, h_L$  with  $h_L \neq 0$  and

$$h_{L-1} = h_{L-2} = 0,$$

such that  $1, \beta$  are roots of the polynomial  $P(z) = z^L - \sum_{k=1}^L h_k z^{L-k}$ .

2. Pick a real number  $u_{-L+1}$ , which satisfies the equation

$$x = \frac{h_L(\beta - 1)u_{-L+1}}{\beta(\beta - 2)}.$$

Set  $u_{-L+2} = \dots = u_0 = 0$ .

3. Find a sequence  $(a_n) \in \{-1, 1\}^{\mathbb{N}}$  such that

$$u_n := \left( \sum_{k=1}^L h_k u_{n-k} \right) + x - a_n$$

satisfies  $u_n \in [-1, 1]$  for each  $n \in \mathbb{N}$ .

Then the sequence  $(a_n)_{n=1}^{\infty}$  satisfies equations (2) and (3).

Such sequences are known as ‘hybrid encoders’. We begin by proving Proposition 3.1, this is similar to the proof of Proposition 2.1.

*Proof.* We begin by rearranging condition (3) of Proposition 3.1 to give

$$a_n = \left( \sum_{k=1}^4 h_k u_{n-k} \right) + x - u_n = \left( \sum_{k=0}^4 h_k u_{n-k} \right) + x.$$

Then we have that

$$\sum_{n=1}^{\infty} a_n \beta^{-n} = \sum_{n=1}^{\infty} \sum_{k=0}^4 h_k u_{n-k} \beta^{-n} + x \sum_{n=1}^{\infty} \frac{1}{\beta^n} \quad (7)$$

where  $h_0 = -1$ . We now follow the reasoning of the proof of Proposition 2.1 exactly, to yield that

$$\sum_{n=1}^{\infty} a_n \beta^{-n} = h_4(u_{-3}\beta^{-1} + u_{-2}\beta^{-2}) + \frac{x}{\beta - 1}.$$

Unlike in Proposition 2.1, we also have that  $u_{-2} = 0$ , so we conclude that

$$\sum_{n=1}^{\infty} a_n \beta^{-n} = h_4(u_{-3}\beta^{-1}) + \frac{x}{\beta - 1},$$

and picking  $u_{-3} = \frac{\beta(\beta-2)x}{h_4(\beta-1)}$  yields that

$$x = \sum_{n=1}^{\infty} a_n \beta^{-n}.$$

It remains to prove part two of the theorem, that

$$x = \lim_{n \rightarrow \infty} \frac{a_1 + a_2 + \cdots + a_n}{n}.$$

Now we have from the condition (3) of Proposition 3.1 that

$$\left( \sum_{k=1}^L h_k u_{n-k} \right) + x - a_n - u_n = \left( \sum_{k=0}^L h_k u_{n-k} \right) + x - a_n = 0.$$

Then

$$\left( \frac{1}{N} \sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k} \right) + x - \frac{1}{N} \sum_{n=1}^N a_n = 0.$$

We shall prove that  $\sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k}$  is bounded by some constant independent of  $N$ , which will give

$$\lim_{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k} = 0$$

and hence that

$$x = \lim_{n \rightarrow \infty} \frac{a_1 + a_2 + \cdots + a_n}{n}.$$

Now we have that

$$\sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k} = (u_1 + u_2 + \cdots + u_{N-L})(h_0 + h_1 + \cdots + h_L) + \text{extra terms},$$

where there are  $N(L+1) - ((N-L)(L+1)) = L(L+1)$  extra terms, each of which are bounded in absolute value by

$$\left( \max_{k \in \{0, \dots, L\}} |h_k| \right) (\sup_{n \in \mathbb{N}} u_n) \leq \max_{k \in \{0, \dots, L\}} |h_k| \leq M$$

for some constant  $M$ . But  $h_0 + h_1 + \cdots + h_L = 0$ , and so we see that

$$\left| \sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k} \right| \leq 0 + L(L+1)M$$

which is independent of  $N$ , and so

$$\lim_{N \rightarrow \infty} \frac{1}{N} \sum_{n=1}^N \sum_{k=0}^L h_k u_{n-k} = 0$$

as required. □



Now we prove Theorem 1.2.

*Proof.* The proof of  $\sum_{n=1}^{\infty} a_n \beta^{-n} = x$  is almost the same as the proof of Theorem 1.1. We follow the construction of Lemma 2.1 replacing  $\beta_1$  by 1 and  $\beta_2$  by  $\beta$ . This gives

$$P(z) = z^L - \sum_{k=1}^L h_k z^{L-k} = (z-1)(z-\beta)(z^2 + az + b),$$

where  $b = \frac{\beta(1+\beta)}{1+\beta+\beta^2}$  and  $c = \frac{\beta^2}{1+\beta+\beta^2}$ .

Then we have  $h_1 = \frac{(1+\beta)(1+\beta^2)}{1+\beta+\beta^2}$  and  $h_4 = \frac{-\beta^3}{1+\beta+\beta^2}$ . We choose  $C_1$  such that

$$\begin{aligned} \sum_{k=1}^4 |h_k| &= |h_1| + |h_4| \\ &= \frac{(1+\beta)(1+\beta^2) + \beta^3}{1+\beta+\beta^2} \\ &< \frac{(1+C_1)((1+C_1)^2+1) + (1+C_1)^3}{3} \\ &= 2 \end{aligned}$$

where  $C_1$  is the real root of  $\frac{(1+x)((1+x)^2+1)+(1+x)^3}{3} = 2$ .

Since  $\sum_{k=1}^4 |h_k| = |h_1| + |h_4| < 2$ , we can choose  $\delta_0 > 0$  such that  $\sum_{k=1}^4 |h_k| = |h_1| + |h_4| \leq 2 - \delta_0 < 2$ , thus for any  $x$  satisfying  $|x| \in [0, \delta_0]$  we have

$$\sum_{n=1}^4 |h_k| = |h_1| + |h_4| \leq 2 - \delta_0 \leq 2 - |x| < 2$$

The next step is to prove the boundness of  $u_n$ . Choosing  $\delta_1 = \frac{h_4(\beta-1)}{\beta(\beta-2)}$  we have that

$$|u_{-3}| = \left| \frac{\beta(\beta-2)x}{h_4(\beta-1)} \right| \leq 1.$$

Finally, if we take  $\delta = \min\{\delta_0, \delta_1\}$ , then this choice can ensure that

$$\sum_{n=1}^4 |h_k| = |h_1| + |h_4| \leq 2 - x$$

and  $|u_{-3}| \leq 1$  hold simultaneously. We also have that  $u_{-2} = u_{-1} = u_0 = 0$ . We let the sequence  $(a_n)$  be chosen as follows:

$$a_n = \begin{cases} -1 & \sum_{k=1}^L h_k u_{n-k} + x < 0 \\ +1 & \sum_{k=1}^L h_k u_{n-k} + x \geq 0 \end{cases}. \quad (8)$$

Then by induction we have that  $u_n \in [-1, 1]$  for all  $n \in \mathbb{N}$ , and hence the conditions of Proposition 3.1 are fulfilled and Theorem 1.2 is proved.  $\square$

## 4 Further Remarks

We have the following further remarks.

- (i) We have proved that if  $\beta_1$  and  $\beta_2$  are very close to 1 then  $A_{\beta_1, \beta_2}$  has an interior, but it is unlikely that our bounds are optimal, see for example the diagrams in [3]. Our proof was based on choosing an expansion  $(a_n)_{n=1}^\infty$  of pairs  $(x_1, x_2)$  using equation (3). Perhaps by using a more sophisticated algorithm one may hope to gain a truer picture of the conditions under which our technique can be made to work.
- (ii) The IFS which we study is a little different to that studied by Shmerkin in [7], since we use digit set  $\{-1, 1\}$  rather than  $(-\frac{1}{\gamma}, -\frac{1}{\lambda})$  and  $(\frac{1}{\gamma}, \frac{1}{\lambda})$ . However such changes of digit set do not affect whether the attractor of the corresponding IFS has an interior.
- (iii) We note that if  $\beta_1\beta_2 > 2$  then  $A_{\beta_1, \beta_2}$  cannot have an interior. Güntürk gave a volume covering argument to prove this. In fact one can say more, the sets  $A_{\beta_1, \beta_2}$  fall into the setting of ‘self-affine sets of Keakeya type’ studied in [5], and so by Theorem 3.3 of that paper we have that

$$\dim_B(A_{\beta_1, \beta_2}) = 1 + \frac{\log \frac{2}{\beta_1}}{\log \beta_2} < 2$$

whenever  $\beta_1\beta_2 > 2$  and  $1 < \beta_1 < \beta_2 < 2$ .

- (iv) Our approach to generating sequences  $\underline{a}$  which satisfy the conditions of Theorem 1.1 is in some sense dynamical, we have an algorithm which chooses a value of  $(a_n)$  based on the vector  $(u_n, u_{n-1}, u_{n-2}, u_{n-3})$ , and then maps this vector to the vector  $(u_{n+1}, u_n, u_{n-1}, u_{n-2})$  and repeats the operation. This system is reminiscent of shift radix systems, see [1], except that we have a displacement by  $a_n$ .

Our algorithm is far less simple than corresponding algorithms for generating expansions in the one dimensional case, such as the random  $\beta$ -transformation of [2]. It would be nice to have an analogue of the random  $\beta$ -transformation for the higher dimensional case which produces expansions of pairs  $(x_1, x_2)$  in a more direct and understandable way.

- (v) One can use Remark 1.1 to consider when, for a specific sequence  $(a_n)$  and real number  $x$ , there exist  $\beta_1, \beta_2$  such that

$$x = \sum_{i=1}^{\infty} a_i \beta_1^{-i} = \sum_{i=1}^{\infty} a_i \beta_2^{-i}.$$

Given  $\underline{a} \in \{-1, 1\}^\mathbb{N}$  we define the function  $f_{\underline{a}} : (1, 2] \rightarrow \mathbb{R}$  by

$$f_{\underline{a}}(\beta) = \sum_{n=1}^{\infty} \frac{a_n}{\beta^n}.$$

The function  $f_{\underline{a}}$  is continuous and differentiable. We call a sequence  $\underline{a} = (a_n)$  a *simultaneous encoder* of  $x$  if there exist  $1 < \beta_1 < \beta_2 < 2$  such

that  $x = f_{\underline{a}}(\beta_1) = f_{\underline{a}}(\beta_2)$ . By Remark 1.1, for any  $1 < \beta_1 < \beta_2 < 1 + C$  and any  $x \in [-0.16, 0.16]$  one can find a simultaneous encoder  $\underline{a}$  of  $x$  satisfying  $x = f_{\underline{a}}(\beta_1) = f_{\underline{a}}(\beta_2)$ . By the extreme value theorem, the function  $f_{\underline{a}}$  has global extrema in  $[\beta_1, \beta_2]$ . We let  $\beta, \beta_0 \in [\beta_1, \beta_2]$  be the values where the global minimum and global maximum take place. Let  $y_1 = f_{\underline{a}}(\beta)$  and  $y_2 = f_{\underline{a}}(\beta_0)$ . Then by the intermediate value theorem, the sequence  $\underline{a}$  is a simultaneous encoder for all  $z \in (y_1, y_2)$ . Thus, if we define the set

$$E_{\underline{a}} := \{x \in \mathbb{R} : \underline{a} \text{ is a simultaneous encoder of } x\}.$$

then the above argument shows that either  $E_{\underline{a}}$  is empty, or is a single point or contains an interval.

## Acknowledgements

The second author was supported by the China Scholarship Council. The third author was supported by the Dutch Organisation for Scientific Research (NWO) grant number 613.001.022.

## References

- [1] Shigeki Akiyama and Klaus Scheicher. From number systems to shift radix systems. *Nihonkai Math. J.*, 16(2):95–106, 2005.
- [2] Karma Dajani and Cor Kraaikamp. Random  $\beta$ -expansions. *Ergodic Theory Dynam. Systems*, 23(2):461–479, 2003.
- [3] C.S. Güntürk. Simultaneous and hybrid beta-encodings. In *Information Sciences and Systems, 2008. CISS 2008. 42nd Annual Conference on*, pages 743–748, 2008.
- [4] John E. Hutchinson. Fractals and self-similarity. *Indiana Univ. Math. J.*, 30(5):713–747, 1981.
- [5] Antti Käenmäki and Pablo Shmerkin. Overlapping self-affine sets of Keakeya type. *Ergodic Theory Dynam. Systems*, 29(3):941–965, 2009.
- [6] Yuval Peres and Boris Solomyak. Problems on self-similar sets and self-affine sets: an update. In *Fractal geometry and stochasticity, II (Greifswald/Koserow, 1998)*, volume 46 of *Progr. Probab.*, pages 95–106. Birkhäuser, Basel, 2000.
- [7] Pablo Shmerkin. Overlapping self-affine sets. *Indiana Univ. Math. J.*, 55(4):1291–1331, 2006.
- [8] Nikita Sidorov. Almost every number has a continuum of  $\beta$ -expansions. *Amer. Math. Monthly*, 110(9):838–842, 2003.