

# A Unified Analysis of Dynamic Interactive Learning

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# Some Motivation: Recommender Systems

- Recommender systems are integral to many online services today (Amazon, Google, Spotify,...)
- These systems are not static, users preferences change due to external factors
- The system itself may also induce changes in the user
- What is a good mathematical model for this scenario, and can we obtain convergence results?

# Previous Work on Static Models

Emamjomeh-Zadeh and Kempe introduced a (static) model for learning with user feedback [1]. It is assumed the user has a finite set of possible preferences modelled as a graph, and a single target preference ranked as ‘the best’.

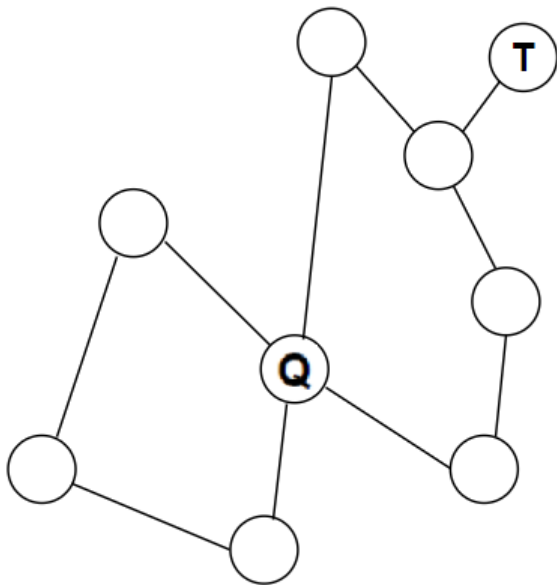
The learner queries vertices in  $v \in V(G)$ , and receives feedback from the user as an edge  $(v, u) \in E(G)$  on the shortest path from  $v$  to  $t$

# Static Models (cont.)

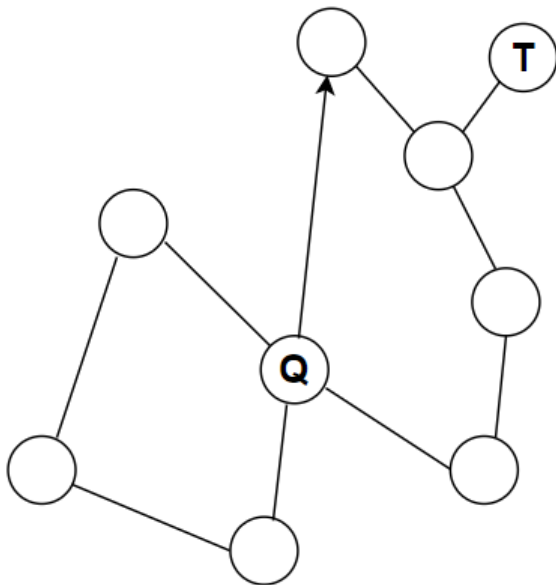
A specification can be given as:

- The learner has access to a weighted graph  $G = (V, E, w)$ , nodes represent concepts and edges are possible transitions between them
- A target concept  $t \in V(G)$  is a fixed node that the learner wishes to discover
- The learner proposes a vertex  $v$  as the target, and receives feedback as an edge out of  $v$
- Edge weights satisfy a key property: if the learner proposes  $q \neq t$ , correct user feedback  $z$  lies on a shortest path from  $q$  to  $t$
- user feedback may be noisy (and adversarial) with probability  $p < \frac{1}{2}$

# A Simple Example



# A Simple Example



# Quick Recap on Weighted Majority

**parameter:**  $\epsilon \in [0, 1]$

Initialize the weights  $w_i = 1$  for all experts.

For each round  $t$ :

    Make predictions using weighted majority vote based on  $w$ .

    For each expert  $i$ :

        If the  $i$ -th expert's prediction is correct,  $w_i$  stays the same.

        Otherwise,  $w_i \leftarrow w_i(1 - \epsilon)$ .

- Setting:  $N$  experts, each makes predictions that are correct some of the time
- Updates bias towards experts which have a history of being correct
- Theorem [2] : MW makes at most  $\frac{\log N + m_T^* \log \frac{1}{\epsilon}}{\log \frac{2}{1+\epsilon}}$  mistakes after  $T$  rounds, where  $m_T^*$  is the number of mistakes made by the best expert

# A Multiplicative Weights algorithm for Graph Feedback

Emamjomeh-Zadeh and Kempe [1] gave a multiplicative weight update algorithm for the digraph problem:

## Noiseless case

- Assign initial likelihoods  $\lambda(v) = 1$  for each node  $v \in G$ .
- Query the 'median' node with shortest weighted distance to all other nodes:  $q = \operatorname{argmin}_v \sum_{v'} d(v, v') \cdot \lambda(v')$ .
- Feedback  $z$  is a neighbor of the median.  $\forall v \in G$  :
  - if  $v$  is *consistent* with  $z$ ,  $\lambda(v) \times 1$
  - if  $v$  is *inconsistent* with  $z$ ,  $\lambda(v) \times 0$

$v$  is *consistent* with  $z$  means  $z$  is on a shortest  $q - v$  path, ie.,  $v$  could be the target
- $O(\log n)$  queries suffice [1]



# A Multiplicative Weights algorithm for Graph Feedback, Noisy Setting

## Noisy case

- Feedback is incorrect with probability  $p < 1/2$  known to the learner
- Similar weight update idea: after querying median  $q$  and receiving feedback  $z$ ,  $\forall v \in G$  :  
if  $v$  is *consistent* with  $z$ ,  $\lambda(v) \times (1 - p)$   
if  $v$  is *inconsistent* with  $z$ ,  $\lambda(v) \times p$
- with probability  $1 - \delta$ , target  $t$  can be found using

$$\frac{(1 - \delta)}{1 - H(p)} \cdot (\log n + o(\log n) + O(\log^2(1/\delta)))$$

queries. [1]

# Interactive Learning of a Dynamic Structure

In a follow-up paper, Emamjomeh-Zadeh and Kempe [3] proposed an extension to the previous paper, where the target is allowed to move  $B$  times in the digraph.

- **Shifting-target** model: the target moves in a pre-determined subset of  $k$  nodes. The subset is unknown, but the parameter  $k$  is known.

Eg: multiple users

- **Drifting-target** model: there's a transition graph  $G'$  with maximum degree  $\Delta$ , and the target moves following the transition graph ( $G'$  is known to the learner, and not necessarily the same as the feedback graph  $G$ )

Eg: user preference changes over time

# Generic Algorithm for Dynamic Interactive Learning

## Sketch of the algorithm

- Consider all  $V^R$  possible sequences of targets throughout the  $R$  rounds, and treat these as Experts in the Multiplicative Weights Framework
- Query the 'median node' at each round, where this depends on the weight of sequences.
- On feedback, update the weight of *sequences* of nodes, as well as individual nodes.
- This gives polynomial mistake bounds, but at the cost of exponential computation. In the special cases of Drifting and Shifting, [3] give (quasi) polynomial-time algorithms.
- The key observation for speed-up is that we don't need to keep track of sequence weights, just node weights

# Previous Bounds [3]

## Shifting Target

$$\begin{aligned} \frac{1}{1 - H(p)} \cdot (k \cdot \log n + B \cdot \log k) \\ \leq M(A) \leq \\ \frac{1}{1 - H(p)} \cdot (k \log n + (B + 1) \log k + R \cdot H(B/R)) \end{aligned}$$

## Drifting Target

$$\begin{aligned} \frac{1}{1 - H(p)} \cdot (\log n + B \cdot \log \Delta) \\ \leq M(A) \leq \\ \frac{1}{1 - H(p)} \cdot (\log n + B \cdot \log \Delta + R \cdot H(B/R)) \end{aligned}$$

# Our Contributions

- ① We propose a generalized transition model that includes the Shifting and Drifting models as special cases
- ② We improve the lower bound for the Drifting target model to match the upper bound given in [3]
- ③ We analyze algorithms for low diameter graphs that require  $O(1)$  computation per round

# A General Model

- Our model is inspired by the Drifting target model over a larger graph  $G'$
- Feedback graph  $G = (V, E, w)$   
Transition graph  $G' = (V', E', \pi)$ , in general  $G'$  is a bigger graph
- Allow multiple vertices in  $G'$  to represent each node in  $G$

## Sketch of the algorithm

- Apply a multiplicative weights style algorithm, with ‘experts’ given as sequences of vertices in the transition graph:  
 $(V(G))^R$
- Transition graph  $G'$  bounds the number of possible sequences there are as  $n' \cdot \Delta'^B \cdot \binom{R}{B}$
- Likelihoods in the  $r$ 'th round are aggregated from duplicate vertices in the transition graph
- Compute median  $q_r$  using likelihoods in feedback graph
- update weights in  $G'$  based on feedback and transition information, then aggregate into  $G$ .

# Interactive Learning Likelihood Upper Bound

Let  $n', \Delta'$  be the number of vertices and maximum degree of the feedback graph  $G'$ .

## Theorem

*[GMR '23] The Likelihood Update Algorithm runs in time  $O(\Delta' \cdot n' + \text{poly}(n))$ , uses space  $O(n')$ , and makes no more than*

$$\frac{1}{1 - H(p)} \cdot (\log n' + B \cdot \log \Delta' + R \cdot H(B/R))$$

*mistakes in expectation.*



# Simplified Shifting Target Analysis

## Corollary

([3], [GMR '23]) In the Shifting Target model, The Likelihood Update Algorithm runs in time  $O(k^2 \cdot n^k)$ , uses space  $O(k \cdot n^k)$ , and makes at most

$$\frac{1}{1 - H(p)} \cdot \left( k \cdot \log n + (B + 1) \cdot \log k + R \cdot H(B/R) \right)$$

mistakes in expectation.

## Proof.

The transition graph  $G'$  consists of  $\binom{n}{k}$  disconnected sub-graphs, where each sub-graph is a clique of size  $k$ , corresponding to a subset of  $k$  vertices in  $V$ . Apply the Theorem from the previous slide. □

# Improved Drifting Target Lower Bound

## Theorem

*For every  $n$  and  $\Delta'$ , there exists a Drifting Target problem such that every algorithm makes at least*

$$\frac{1}{1 - H(1 - p)} \cdot \left( \log n + B \cdot \log \Delta' + R \cdot H(B/R) \right) \\ - o \left( \log n + B \cdot \log \Delta' + R \cdot H(B/R) \right)$$

*mistakes in expectation.*

# Constant Computation Algorithms

Consider the following simple algorithm:

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**Algorithm 1** ‘Follow the Feedback’ Procedure for Interactive Learning

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$q_1 \leftarrow \operatorname{argmin}_{i \in V} \sum_{j \in V} w(i, j)$  {Start with a ‘center’ vertex}  
**for**  $1 \leq r \leq R$  **do**  
     $z_r \leftarrow$  feedback from adversary after querying  $q_r$   
     $q_{r+1} \leftarrow z_r$  {Follow the feedback for next round}  
**end for**

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A natural question is, for what classes of graphs and under what transition models does this give polynomial mistake bounds?

# A Diameter Bound

## Theorem

*Suppose  $G'$  is as before, and let  $d$  be the diameter of  $G'$ . Then the 'Follow the feedback' algorithm makes at most*

$$\frac{1}{1-p} \cdot \left( dB - \frac{pB}{1-2p} + pR \right)$$

*mistakes in expectation.*

## Proof Idea

- 1 Reduce  $G'$  to  $P_d$ , a path on  $d$  vertices
- 2 Algorithm 1 makes a mistake anytime the feedback is not the target vertex
- 3 Apply Markov chain hitting time bounds on  $P_d$