**Title: Sorting and Algorithm Efficiency** 

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**Section: 1** 

**Assignment: 1** 

**Description: HW1 Report** 

#### **Question 1:**

(a) [5 points] Show that  $f(n) = 4n^5 + 2n^3 + 3n$  is  $O(n^5)$  by specifying appropriate c and  $n_0$  values in Big-O definition.

a) show that 
$$f(n) = 4n^5 + 2n^3 + 3n$$
 is  $O(n^5)$  by specifying  $c$  and  $n_0$  values  $4n^5 + 2n^3 + 3n \le c \cdot n^5$   $(4n^4 + 2n^2 + 3) \le c \cdot n^5$   $(n > 0)$   $4n^4 + 2n^2 + 3 \le c \cdot n^4$  If  $c = 5$  and  $n_0 = 2 \Rightarrow 4n^4 + 2n^2 + 3 \le 5n^4$   $2n^2 + 3 \le n^4$  for all  $n \ge n_0 = 2$   $11 \le 16$   $0 \le n^4 - 2n^2 - 3$   $0 \le (n^2 - 1)^2 - 4$  Hence, since there exist constants  $c$  and  $n_0$  such that  $f(n) \le c \cdot n^5 f_{n,n}$   $n \ge n_0$ ,  $f(n) = 4n^5 + 2n^3 + 3n$  is  $O(n^5)$ .

(b) [10 points] Trace the following sorting algorithms to sort the array [40, 25, 65, 45, 50, 35, 55, 38, 30, 42] in ascending order. Use the array implementation of the algorithms as described in the textbook and show all major steps.

#### **Selection Sort:**

Some 65,42 [40, 25, 65, 45, 50, 35, 55, 38, 30, 42] initial array swap 55,30 [40, 25, 42, 45, 50, 35, 55, 38, 30, 65] after 1st swap swap 50,38 [40,25,42,45,50,35,30,38,55,65] after 2nd swap swap 45,30 [40,25,42,45,50,35,30,50,55,65] after 3nd swap swap 45,30 [40,25,42,45,38,35,30,50,55,65] after 3nd swap swap 42,35 [40,25,42,30,38,35,45,50,55,65] after 4th swap swap 40,38 [40,25,42,30,38,35,45,50,55,65] after 5th swap swap 38,30 [38,25,35,30,40,42,45,50,55,65] after 5th swap swap 38,30 [38,25,35,36,40,42,45,50,55,65] after 9th swap swap 35,35 [30,25,35,38,40,42,45,50,55,65] after 9th swap swap 30,25 [30,25,35,38,40,42,45,50,55,65] after 9th swap swap 30,25 [30,25,35,38,40,42,45,50,55,65] after 9th swap [25,30,35,38,40,42,45,50,55,65] after 9th swap swap is swited

#### **Insertion Sort:**

```
[10] 25, 65, 45, 50, 35, 55, 38, 40, 42]

1st iteration | Sorted | Unsorted |
```

### **Question 2:**

#### **Screenshot of the output in part c:**

#### han.arslan@diikstra:~

```
Last login: Wed Jul 12 15:31:13 2023 from 10.201.182.133
[han.arslan@dijkstra ~]$ ls
hw1 main.cpp Makefile sorting.cpp sorting.h
[han.arslan@dijkstra ~]$ ./hw1
Initial array:
[10 5 9 16 17 7 4 12 19 1 15 18 3 11 13 6]
BUBBLE SORT
Number of key comparisons: 114
Number of data moves: 180
Array after bubble sort:
[1 3 4 5 6 7 9 10 11 12 13 15 16 17 18 19]
Initial array:
[10 5 9 16 17 7 4 12 19 1 15 18 3 11 13 6]
MERGE SORT
Number of key comparisons: 46
Number of data moves: 128
Array after merge sort:
[1 3 4 5 6 7 9 10 11 12 13 15 16 17 18 19]
Initial array:
[10 5 9 16 17 7 4 12 19 1 15 18 3 11 13 6]
QUICK SORT
Number of key comparisons: 45
Number of data moves: 102
Array after quick sort:
[1 3 4 5 6 7 9 10 11 12 13 15 16 17 18 19]
```

# Screenshots of the output in part d:

#### **For Random Arrays:**

RANDOM ARRA	vc		
KANDOM AKKA	15		
Analysis of	Bubble Sort		
Array Size		compCount	moveCount
4000	99.642 ms	7997054	11952156
8000	437.515 ms	39535622	48018132
12000	1032.85 ms	102346163	107814255
16000	1882.47 ms	196326647	193199748
20000	2978.37 ms	321512279	300311103
24000	4335.72 ms	477655537	430238673
28000	5967.2 ms	665235900	591769419
32000	7819.77 ms	882993806	763237938
36000	9973.78 ms	1132514864	974440695
40000	12379.8 ms	1416437184	1201930278
	123/9.0 1113	141043/104	1201930278
Analysis of	Merge Sort		
	Elapsed time	compCount	moveCount
4000	1.4289 ms	42812	95808
8000	3.08963 ms	93647	207616
12000	4.56718 ms	147588	327232
16000	6.76977 ms	203230	447232
20000	8.48942 ms	260908	574464
24000	10.2195 ms	319254	702464
28000	11.9259 ms	378797	830464
32000	13.7025 ms	438639	958464
36000	15.432 ms	499754	1092928
40000	17.1405 ms	561599	1228928
Analysis of	Quick Sort		
Array Size	Elapsed time	compCount	moveCount
4000	0.938943 ms	53334	88603
8000	1.94086 ms	117510	184120
12000	3.25356 ms	209046	345101
16000	4.30396 ms	260165	431155
20000	5.4572 ms	347626	514646
24000	6.75027 ms	422828	670272
28000	7.94667 ms	474753	794117
32000	8.96964 ms	547338	857466
36000	10.731 ms	656366	1171966
40000	11.8474 ms	717328	1181766

Output 1: Algorithm comparison with randomly generated arrays

# For Ascending Arrays:

10000		,1,520	1101700
ASCENDING A	RRAYS		
Analysis of	Bubble Sort		
Array Size	Elapsed ti	me compCount	moveCount
4000	0.023328 ms	3999	0
8000	0.046222 ms	7999	0
12000	0.069264 ms	11999	0
16000	0.092283 ms	15999	0
20000	0.11909 ms	19999	0
24000	0.138332 ms	23999	0
28000	0.161285 ms	27999	0
32000	0.184784 ms	31999	0
36000	0.207585 ms	35999	0
40000	0.234549 ms	39999	0
Analysis of			
Array Size			
4000	0.911124 ms	24372	95808
8000	2.07275 ms	52767	207616
12000	3.04355 ms	84741	327232
16000	4.57007 ms	113538	447232
20000	5.73533 ms	148870	574464
24000	6.8372 ms	181453	702464
28000	7.95839 ms	213974	830464
32000	9.04467 ms	243026	958464
36000	10.2301 ms	280923	1092928
40000	11.3377 ms	317667	1228928
	0.1-1.5-1		
Analysis of		mo compCount	moveCount
Array Size 4000	Elapsed ti	ne compCount 7998000	moveCount 15996
8000	134.444 ms	31996000	31996
12000	302.75 ms	71994000	47996
16000	538.462 ms	127992000	63996
20000	841.462 ms	199990000	79996
24000	1211.7 ms	287988000	95996
28000	1649.28 ms	391986000	111996
32000 32000	2154.24 ms	511984000	127996
36000	2726.35 ms	647982000	143996
40000	3365.76 ms	799980000	159996
40000	3303.70 IIIS	799980000	139990

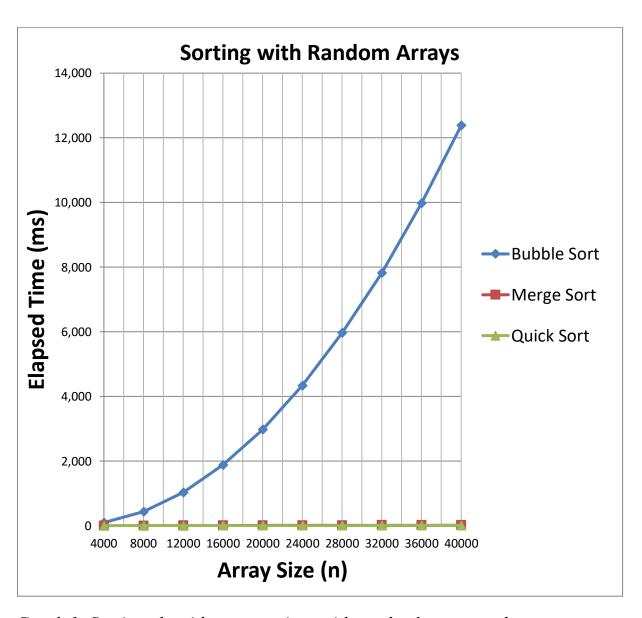
Output 2: Algorithm comparison with ascending arrays

# For Descending Arrays:

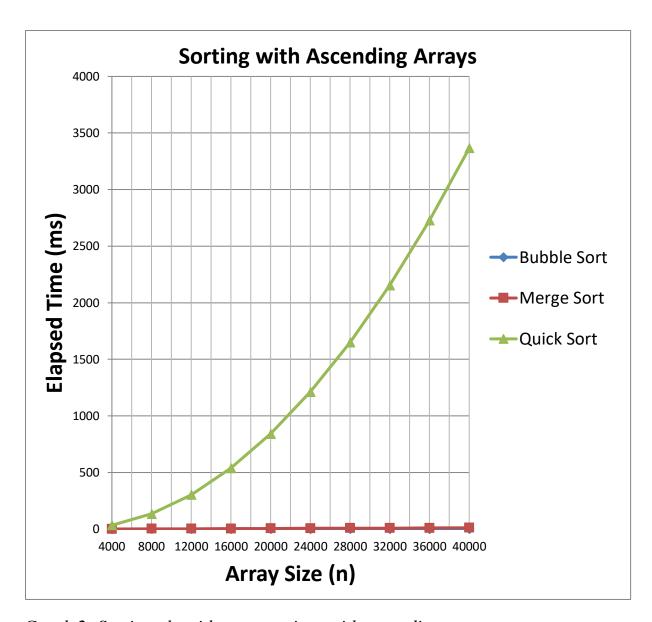
DESCENDING /	ΔRRΔVS		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	133330
Analysis of	Bubble Son	rt		
Array Size			compCount	moveCount
4000	114.093 r	ns	7998000	23992644
8000	456.62 r	ns	31996000	95985285
12000	1029.39 r		71994000	215978010
16000	1832.93 r	ns	127992000	383970864
20000	2862.91 r	ns	199990000	599963340
24000	4123.23 r	ns	287988000	863956194
28000	5611.63 r	ns	391986000	1175948886
32000	7328.86 r	ns	511984000	1535941224
36000	9288.7 r	ns	647982000	1943934519
10000	11453.1 r	ns	799980000	-1895040220
Analysis of	Merge Sort	t		
Array Size	Elapsed		compCount	moveCount
4000	0.889467 r		23728	95808
8000	2.02065 r		51456	207616
12000	3.00394 r		79312	327232
16000	4.47285 r		110912	447232
20000	5.59931 r		139216	574464
24000	6.70066 r		170624	702464
28000	7.81423 r	ns	202512	830464
32000	8.90455 r		237824	958464
36000	10.0108 r		267280	1092928
10000	11.1395 r	ns	298432	1228928
Analysis of				
Array Size			compCount	moveCount
4000	63.5496 r		7566650	11367726
8000	255.172 r		30362316	45578996
12000	573.656 r		68341787	102565916
16000	1018.76 r		121525205	182358853
20000	1593.16 r		189680257	284609044
24000	2296.54 r		273255678	409990068
28000	3115.13 r		371028721	556667097
32000	4068.87 r		484564742	726988773
36000	5181.45 r		616523979	924946204
10000	6364.14 r		758434476	1137829253
[han.arslan(	@dijkstra /	~]\$ <b>_</b>		

Output 3: Algorithm comparison with descending arrays

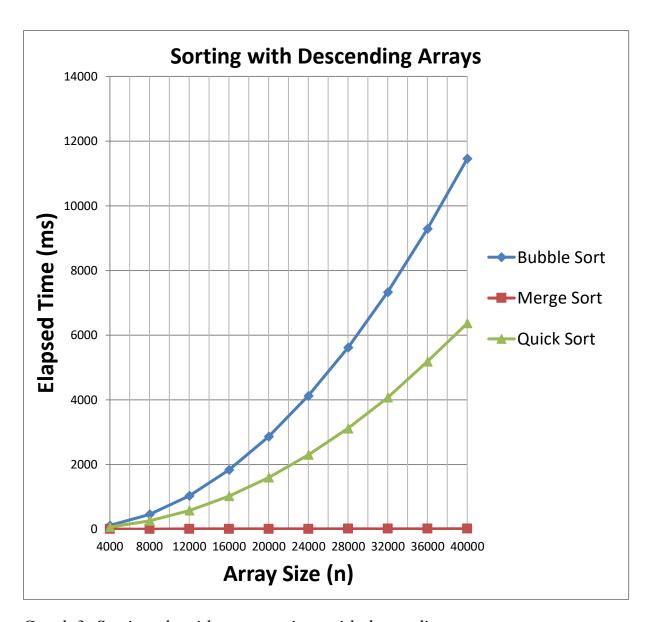
# **Question 3:**



Graph 1: Sorting algorithm comparison with randomly generated arrays



Graph 2: Sorting algorithm comparison with ascending arrays



Graph 3: Sorting algorithm comparison with descending arrays

#### **Comments:**

- ❖ Bubble sort worst case behaviour is O(n²) and it happens when the array is in reverse order. By looking at the graph 3, experimental results with descending arrays are consistent with theoretical results.
- ❖ Bubble sort average case behaviour is O(n²), by looking at the graph 1 experimental results are consistent with theoretical results since the arrays are randomly generated.

- \* Bubble sort makes  $O(n^2)$  key comparisons and moves in worst case and average case, as seen in output 3 and output 1.
- ❖ Bubble sort best case behaviour is O(n) and it happens when the array is already sorted. This case happens in graph 2, when the array is in ascending order. The moveCount variable is 0 in all array samples, but bubble sort algorithm stil do (n-1) key comparisons as seen in output 2.
- ❖ There is integer overflow in descending array analysis with bubble sort of array size 40000, hence moveCount variable is negative.
- ❖ Merge sort is O(n\*logn) in all cases, independent of the array configuration.
  Graph 1, 2 and 3 are consistent with this theoretical behaviour.
- ❖ Merge sort makes O(n\*logn) key comparisons in average case and worst case, hence all three outputs are consistent with this theoretical behaviour.
- ❖ Quick sort is O(n\*logn) algorithm in best case and average case, results in graph 1 is consistent with this behaviour.
- ❖ In worst case, when the array is already sorted or in reverse order, quick sort is O(n²). Graph 2 and 3 shows this behaviour of quick sort algorithm's worst case.
- ❖ Although merge sort is stable and O(n\*logn) in every case, quick sort is faster in terms of elapsed time in average case, as seen in output 1. This behaviour is most likely caused by merge sort's requirement of an extra array in its execution.
- ❖ Bubble sort's best case behavior O(n) is actually faster than merge sort and quick sort in all cases, as seen in the outputs, however this case rarely occurs in real-life implementations.