

# Chapter 6: Image Compression

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# Introduction

## Redundant information

Three types of redundancies to be explored by compressors:

- coding: minimise the length of the representation
  - e.g., if you have AAAAAAABB, it could be compressed as 8A2B
- inter-pixel: minimise the repetitions (redundancies)
  - e.g., represent regions (or blobs) as vectors
- psycho-visual: lossy, disregard information not “visible”
  - Humans do not see subtle changes in illuminat

# Introduction

## Compression Ratio

$$C_r = \frac{d_1}{d_2} \quad (1)$$

where:  $d_1$  and  $d_2$  are the amounts of data in the original and compressed images Average compression ratios rarely go over 30 for normal coding.

## Lossy x Loss-less compression

- Loss-less (non-lossy): recover the exact input
- Lossy: the image recovered looks the same to users, but it is not

# Shannon's Information Theory and Entropy

## Entropy

$$H = - \sum_{i=0}^N p_i \log(p_i) \quad (2)$$

Where:  $p_i$  is the probability of a symbol to appear in a stream.

## Example

8 symbols (000, 001, 010, 011, 100, 101, 110, and 111) with equal probabilities would yield maximum Entropy:

$$H = - \sum_{i=0}^N p_i \log_2(p_i) = -8 \frac{1}{8} \log_2\left(\frac{1}{8}\right) = 3$$

## More examples on Entropy

For example, suppose that in the 3 bits system we have symbols A, B, C, D, E, F, G and H. The probabilities are: A(0.4) B(0.3) C(0.1) D(0.1) E(0.06) F(0.04) G(0) and H(0).

$$H = -(0.4 \log_2(0.4) + 0.3 \log_2(0.3) + 0.1 \log_2(0.1) + 0.06 \log_2(0.06) + 0.04 \log_2(0.04))$$

$$H = 2.1435$$

Another example, A and B in equal quantities (probabilities are 0.5). All other symbols have probabilities equals zero.

$$H = -2 (0.5) \log_2(0.5) = 1$$

The entropy expresses the average number of bits that would take to represent the same stream using a *variable-length code*.

# Loss-less compression

## Huffman

- Huffman coding (1952)
- So-called Entropy coder, it converts the probabilities into a sequence of bits
- Entropy:  $H = -\sum p_i \log_2(p_i)$
- Max Entropy for a 1 byte code: 8
- When a file has Entropy 8, it is not compressible using Huffman's code

## LZW

- LZW (Lempel-Ziv-Welch, 1984)
- Builds a dynamic dictionary of repeated sequences
- Plagued by patents (Unisys and IBM)

# Huffman

## The method

- Create an ordered list for each symbol
- Symbol can be of any length
- Each arbitrary size symbol is represented by an increasing number of bits
- *Compression Ratio* is calculated based on the total number of bits



## Huffman pseudo-code

**Require:** Heap (minHeap, or a priority queue)  $P$

**Require:** Tree node  $T_{huf}$  and  $N$  tree nodes  $T_i$ , temporary tree nodes  $T_a$ ,  $T_b$  and  $T$

- 1: Create a leaf  $T_i$  for each symbol (using the probability as a key)
- 2: Insert all  $T_i$  in the heap  $P$
- 3: **while** there is more than one node in the Heap  $P$  **do**
- 4:   delete the two nodes from  $P$  (the two with lowest probability) and copy to  $T_a$  and  $T_b$
- 5:   create a **new** tree node  $T$
- 6:   make the probability of  $T$  equals to the sum of the probabilities of  $T_a + T_b$
- 7:   make  $T_a$  and  $T_b$  children (smallest on left) of the new node
- 8:   insert the new tree node back into the heap  $P$
- 9: **end while**
- 10: the remaining node in  $P$  (with probability 1) is the root of tree  $T_{huf}$ .
- 11: Traverse  $T_{huf}$ , marking 0 if going left and 1 if going right.
- 12: When reaching a leaf, output the code.

# Huffman

## Initialisation and step 1

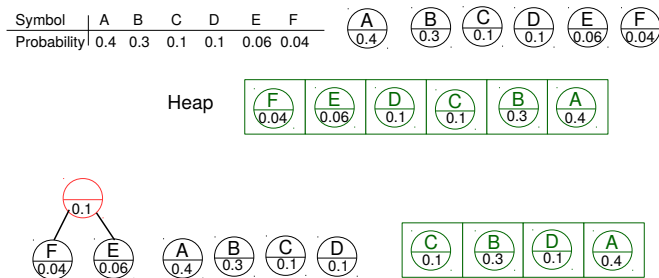


Figure 1: Huffman Example.

# Huffman

steps 2 and 3

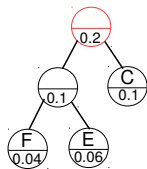
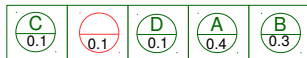
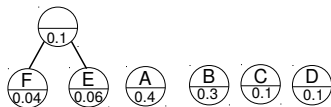


Figure 2: Huffman Example.

# Huffman

step 4

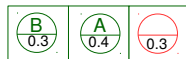
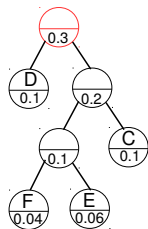
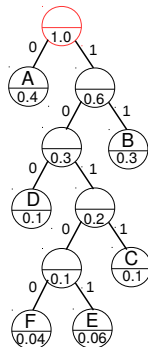
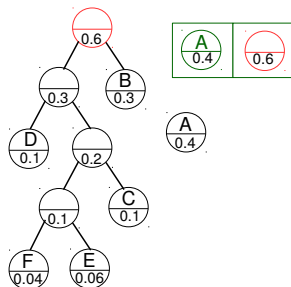


Figure 3: Huffman Example.

# Huffman

final steps



Coding

A 0  
 B 11  
 D 100  
 C 1011  
 E 10101  
 F 10100

Figure 4: Huffman Example.

# Huffman

## Compression ratio

$$C_r = \frac{3}{(0.4 * 1 + 0.3 * 2 + 0.1 * 3 + 0.1 * 4 + 0.06 * 5 + 0.04 * 5)} = 1.36$$

The average number of bits predicted by the Entropy of the uncompressed file is:

$$H = - \sum p_i \log_2(p_i) = 2.1435$$

so the compression predicted by the number of bits is  $\frac{3}{2.1435} = 1.39$ , very close to the real result.

# LZW

## Introduction

- GIF and TIF formats can use this method
- Patents created controversy: PNG format was created
- Patents expired in 2003 and 2004
- Many variants: LZ77, LZ78, LZMW (1985), LZAP (1988), LZWL etc

# LZW

## The method

- Create a dictionary
- Uses inter-code redundancies: i.e. repeated 'letters' are used once
- In images, repeated sequences of pixels are encoded this in this way



# LZW

## The algorithms

- Input: file1 (or image1)
- Output: file2 (or image2)
- Strings S, O and character C.
- Array T[N] (N=number of entries)
- + means append

# LZW

## Compression

1. Initialise a table with a single value strings (say from 0 to 255).
2. Initialise  $S = \text{pixel}[0]$
3. WHILE there are pixels in *image* DO
  - $C = \text{get input character}$
  - IF  $S+C$  is in the string table then
    - $S = S+C$
  - ELSE
    - output the code for  $S$
    - add  $S+C$  to the string table
    - $S = C$
  - END of IF
4. END of WHILE
5. output the code for  $S$

# LZW

## Decompression

1. Initialise a table with single value strings (say from 0 to 255).
2. Read  $O = \text{pixel}[0]$  of *image2*
3. output  $O$
4. WHILE there are pixels in *image2* DO  
    Read  $N = \text{pixel}[i]$   
     $S = \text{get translation of } N$   
    output  $S$   
     $C = \text{first character in } S$   
    add  $O + C$  to the translation table  $T$   
     $O = N$
5. END of WHILE

# LZW

## Example

Let us use a 16 byte sequence:

39 39 126 126 39 39 126 126 39 39 126 126 39 39  
126 126

# LZW

## Example

Table 1: LZW compression.

S	C	Output	Code	Dictionary String
39				
39	39	39	256	39-39
39	126	39	257	39-126
126	126	126	258	126-126
126	39	126	259	126-39
39	39			
39-39	126	256	260	39-39-126
126	126			
126-126	39	258	261	126-126-39
39	39			
39-39	126			
39-39-126	126	260	262	39-39-126-126
126	39			
126-39	39	259	263	126-39-39
39	126			
39-126	126	257	264	39-126-126
126		126		

# LZW

## Compression ratio

- Original sequence with 8 bits: 39 39 126 126 39 39 126 126  
39 39 126 126 39 39 126 126
- Compressed with 9 bits: 39 39 126 126 256 258 260 259 257  
126
- $C_r = (16 * 8bits)/(10 * 9bits) = 1.42$

# LZW

## Example

Table 2: LZW decomposition.

<i>N</i>	<i>O</i>	<i>C</i>	<i>S</i>	Dictionary String	Code
	39				
39	39	39	39	39-39	256
126	39	126	126	39-126	257
126	126	126	126	126-126	258
256	126	39	39-39	126-39	259
258	256	126	126-126	39-39-126	260
260	258	39	39-39-126	126-126-39	261
259	260	126	126-39	39-39-126-126	262
257	259	39	39-126	126-39-39	263
126	257	126	126	39-126-126	264

# Lossy compression

## Modelling the data to increase redundancies

- Cannot recover the exact original image
- However, “good enough” quality
- Use of a model for the data, i.e., use specific properties of the data:
  - Quantisation
  - Finite coefficients of frequency domain (DCT, FFT etc)
  - Colour space



# Quantisation

- Staircase most commonly used
- Humans do not perceive subtle changes in intensities, specially if not neighbours
- Can be applied to colour spaces (e.g., remember YUV411)
- Non-linear quantisation (e.g., Lloyd-Max quantisers, with different quantisation for different intervals)

# DCT

## Introduction

- Remember Fourier Series?
- Use Cosines only
- Discrete Cosine Transform uses  $N \times N$  number of coefficients
- Redundancies yield coefficients with low or zero values.
- On decompressing, we delete a number of low value coefficients.

# DCT

## Equations

A Discrete Cosine Transform for a 2D square sized image ( $N \times N$ )  $i(x, y)$  is:

$$C(u, v) = \alpha(u)\alpha(v) \sum_{x=0}^{N-1} \sum_{y=0}^{N-1} i(x, y) \cdot \cos \left[ \frac{(2x+1)u\pi}{2N} \right] \cdot \cos \left[ \frac{(2y+1)v\pi}{2N} \right] \quad (3)$$

Where:

$$\alpha(0) = \sqrt{1/N}$$

$$\alpha(u) = \sqrt{2/N} \text{ for } u = 1, 2, \dots, N-1$$

# DCT

## Equations

To recover the image given the coefficients  $C(u, v)$ , the following equation is used:

$$i(x, y) = \sum_{u=0}^{N-1} \sum_{v=0}^{N-1} \alpha(u) \alpha(v) C(u, v) \cdot \cos \left[ \frac{(2x+1)u\pi}{2N} \right] \cdot \cos \left[ \frac{(2y+1)v\pi}{2N} \right] \quad (4)$$

# DCT

## Basis functions visualised

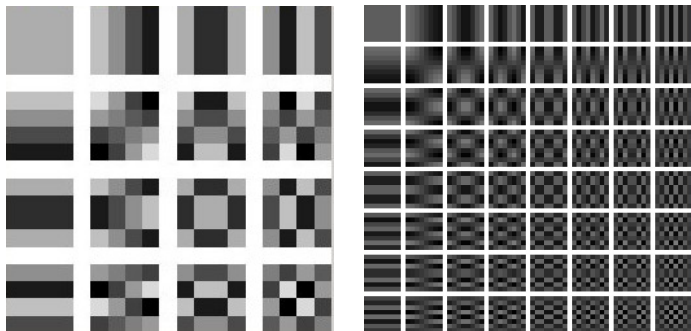


Figure 5: Basis functions for a DCT transform. a) 4x4 transform b) 8x8 transform.

# DCT

## Example

Image 8x8

1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23
1	127	127	120	120	123	124	23

DCT result

765	-24	-281	-32	-237	-20	-118	-3
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0

Figure 6: DCT example.

# DCT

## Observations:

- Loss-less DCT compression: encode with all coefficients
- Lossy DCT compression: encode with the largest coefficients
- e.g., JPEG (DCT, quantise, encode)

# Video Compression

## Introduction

- In video, we can compress every frame.
- Are there Inter-frame redundancies?
- How to explore them?



# Video Compression

## Mpeg

- “Infinite” ways to write codecs (coder-decoder)
- Attempts to standardise them: MPEG-1, MPEG-2, MPEG-4, MPEG-7, MPEG-21...
- The ultimate goal is to compress and decompress efficiently

# Video Compression

## Mpeg

- If only objects move, we could keep the background in one frame...
- However, in practise there are problems with that.
- E.g., occlusion, too many things moving etc.
- noisy videos take more space...

# Video Compression

## Mpeg-2

- Three frame types
- I frames: intra-coded (jpeg).
- P frames: predicted frames, use previous I or other P frames with few blocks changes.
- B frames: Bi-directional frames, predicted frames using past and future frames.

# Video Compression

## Mpeg-2

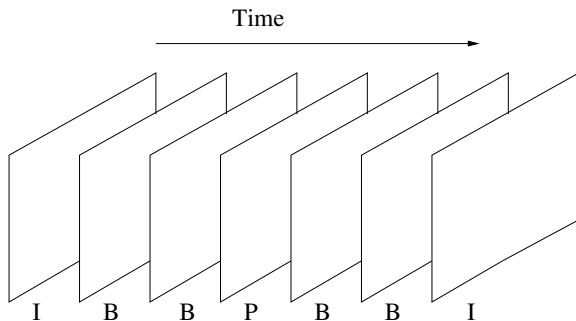


Figure 7: MPEG frame types.

## Exercises 1 and 2

- 1) Write a simple version of Huffman code that is able to compress a grey-scale image in OpenCV. How does your compression rate fare against jpg? (to check that, you can export a jpg image in GIMP, after transforming it into grey-scale.)
- 2) Using the DCT transform above, write a program to compute the cosine transform coefficients for an 8x8 block. Compress it using Huffman code.

## Exercises 3 and 4

- 3) Using the DCT transform above, write a program to compute the cosine transform coefficients for an  $16 \times 16$  block. Compress it using Huffman code. Any difference in the compression rate?
- 4) Modify the program to compute the inverse cosine transform. Analyse the errors in the resulting images by simple rounding (to integers) and quantisation of the coefficients compared to the use of floating point numbers to represent the coefficients.

# Bibliography

(partial)



R. C. Gonzalez and R. E. Woods, *Digital Image Processing*, Prentice Hall, 2002.



M. Sonka, V. Hlavac, and R. Boyle, *Image Processing, Analysis and Machine Vision*, PWS Publishing, 1999