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#### Outline

- Basic concepts from Information
   Theory
- Some thoughts on compression of scientific data
- Simple example
- · Conclusions





## Basic Quantities in Information Theory

· Data stream S and for x ∈ S let

· Shannon Information Content

$$h(x) = -\log_2 p_x$$

· Entropy

H(S)= - ∑ px log 2 px

⇒ Noisy/random data has HIGH ENTROPY





### Relevance to Data Storage

#### SOURCE CODING THEOREM

N items of data from S require at Least NH(S) bits of storage.

- · Information content related to likelihood.
- · Noisy/random data has LARGE H(S) and hence ALMOST INCOMPRESSIBLE.
- E.G. Gomez & Cappello 2013 show only 15% compression (lossless).





#### Data Compression

DATA STREAM

..., Xn, Xn+1) ...

RAW DATA

STEP 1

PREPROCESSOR

e.g. XOR in FCP

..., yn , yn+1) ---

PROCESSED DATA

STEP 2

REDUNDANCY REMOVAL

e.g. Duplicates

FINAL COMPRESSED DATA

STEP 3

ENTROPY ENCODING

e.g. Huffman Tree

..., Zn, Zn+1, ...



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#### Data Compression

· Noisy/random data has HIGH H(S)

=> limited scope for lossless compression

· Applying hierarchy of thresholds to data to reduce entropy H(S)

ENTROPY
ENCODING ..., Zn, Zn+1)...

Huffman Tree

· Map onto storage hierarchy

HIGHER FIDELITY; LESS COMPRESSION

٧S

LOWER THRESHOLD

· LOW ENTROPY

· HIGH COMPRESSIBLE

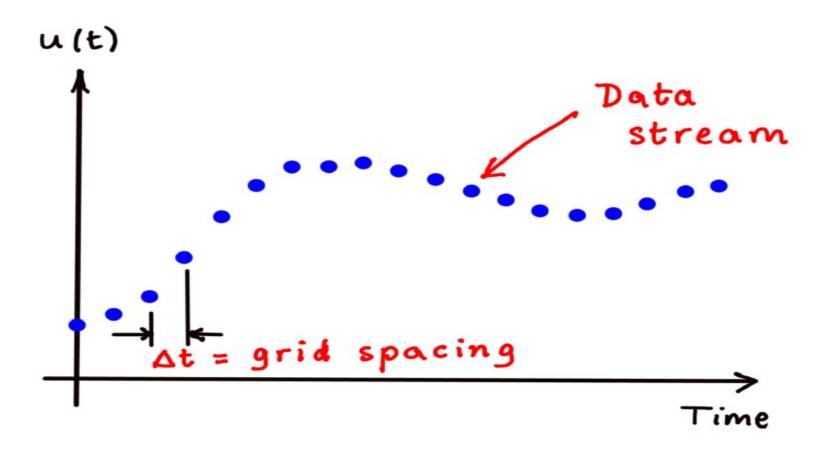
HIGHER THRESHOLD

- · ALMOST LOSS LESS
- · LOW COMPRESSIBLE

LOWER FIDELITY; HIGH COMPRESSION,

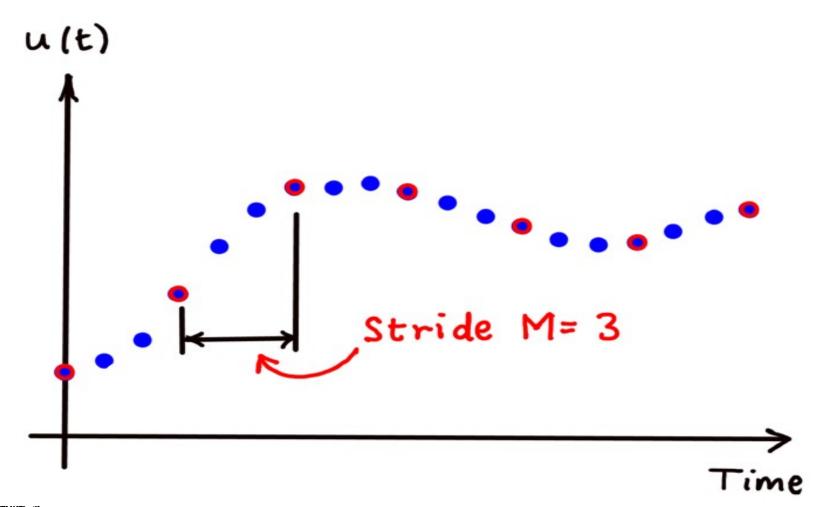






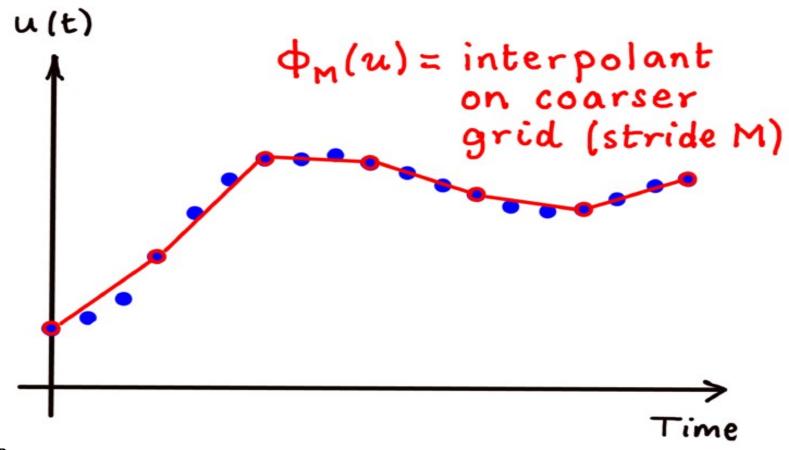






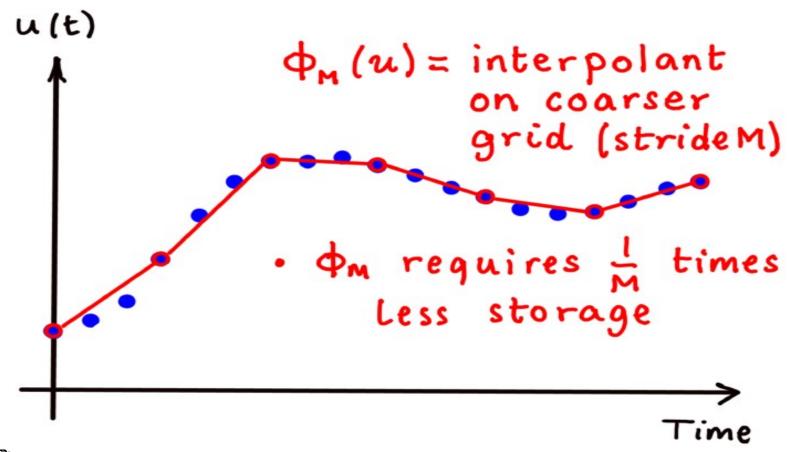






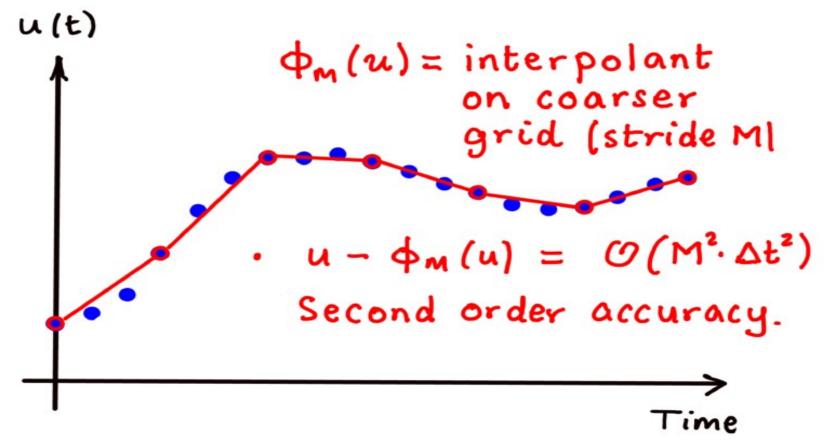








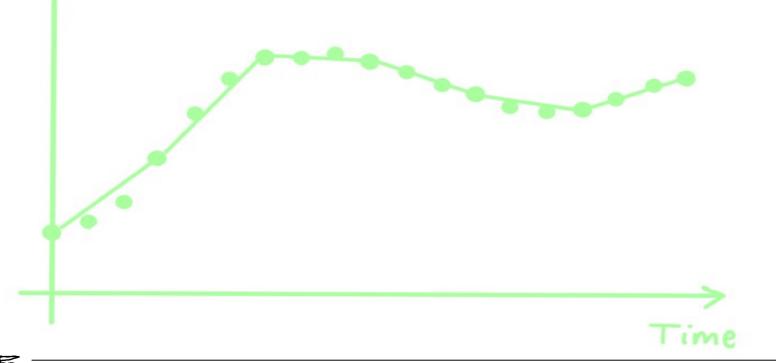








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- New storage cost =  $\frac{32 \,\text{N}}{\text{M}}$  bits  $+ \left\{23 - \log_2\left(\text{C·M}^2 \cdot \Delta t^2\right)\right\} \text{N}$

• Ratio = 
$$\frac{1}{M} + \frac{23 - \log_2 (C \cdot M^2 \cdot \Delta t^2)}{32}$$
  
=  $(\frac{1}{M} - \frac{1}{16} \log_2 M) - \frac{1}{16} \log_2 \Delta t + const.$ 





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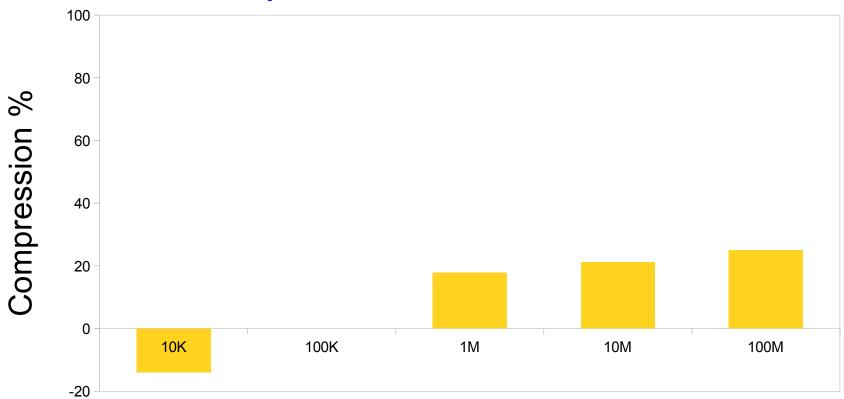
$$= \left(\frac{1}{M} - \frac{1}{16} \log_2 M\right) - \frac{1}{16} \log_2 \Delta t + const.$$

Min. independent of N. At

Benign



## Compression (No auditor)

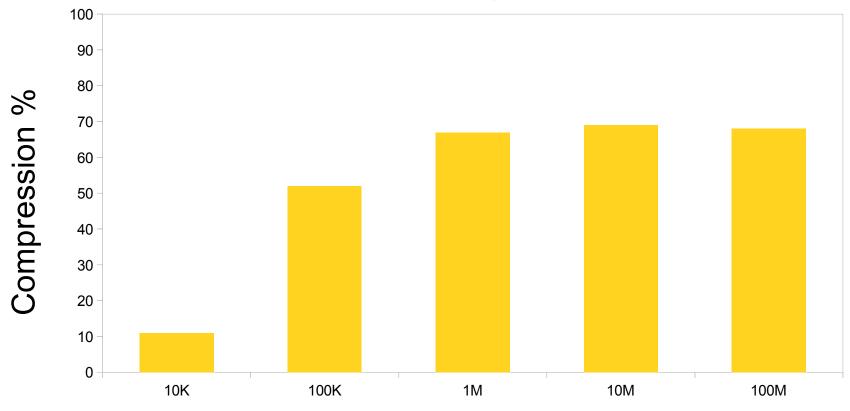


Size of Data Set





# Compression Interpolation (Stride=5)

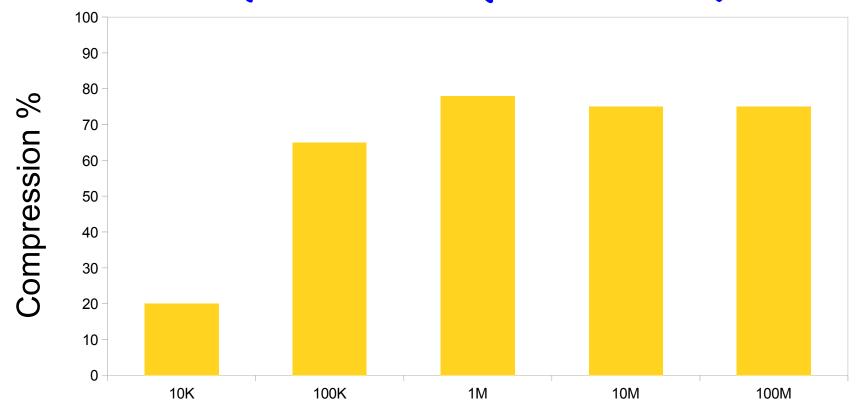


Size of Data Set





# Compression Interpolation (Stride=10)

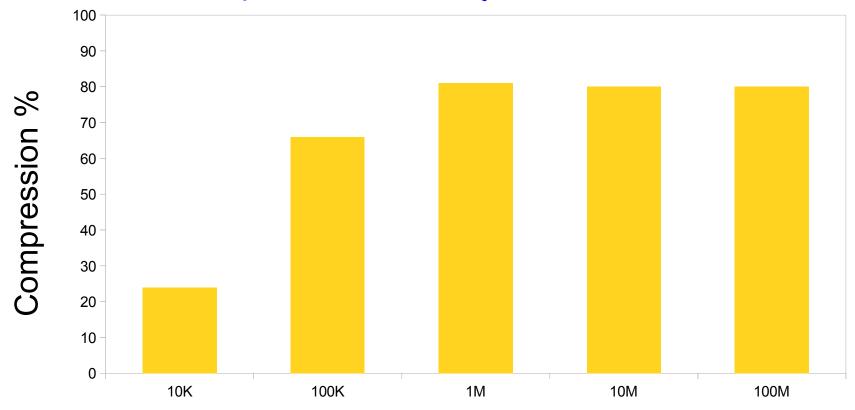


Size of Data Set





# Compression Interpolation (Stride=20)

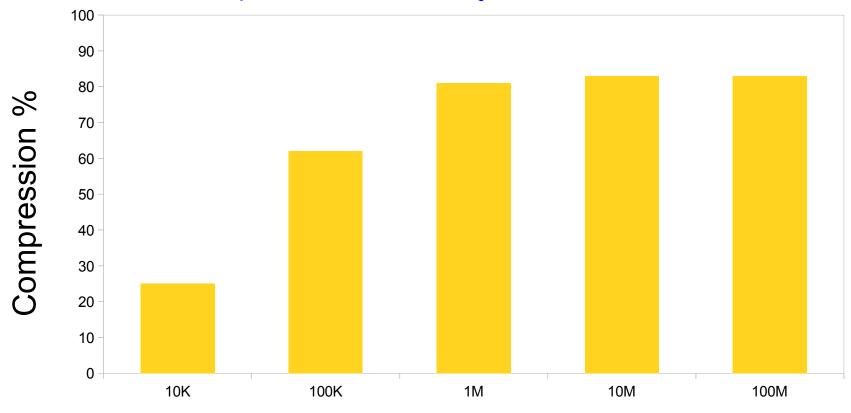


Size of Data Set





# Compression Interpolation (Stride=50)

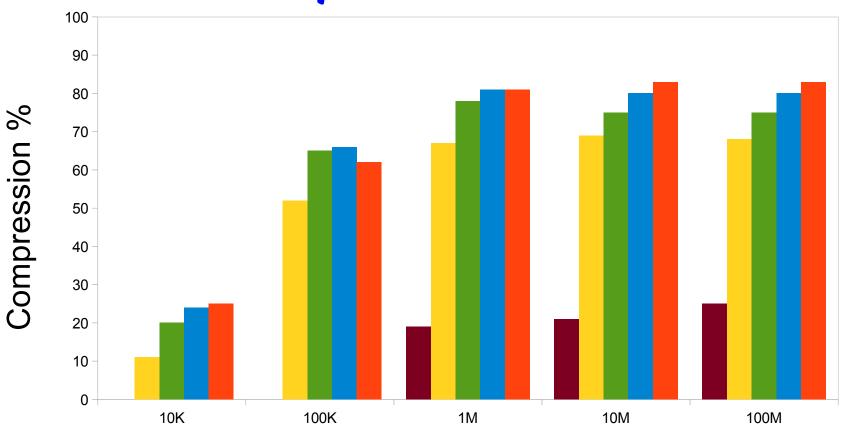


Size of Data Set





### Compression Interpolation Auditor



Size of Data Set





