

# Technical Session 3

## The CLEARSY Safety Platform

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11th March 2025



Engineering and  
Physical Sciences  
Research Council

# Agenda

- ▶ Introduction to proven software
- ▶ Introduction to the CLEARSY Safety Platform
- ▶ Development process (demo video)
- ▶ Bits of B
- ▶ Using the modelling interface



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# Introduction to Proven Software

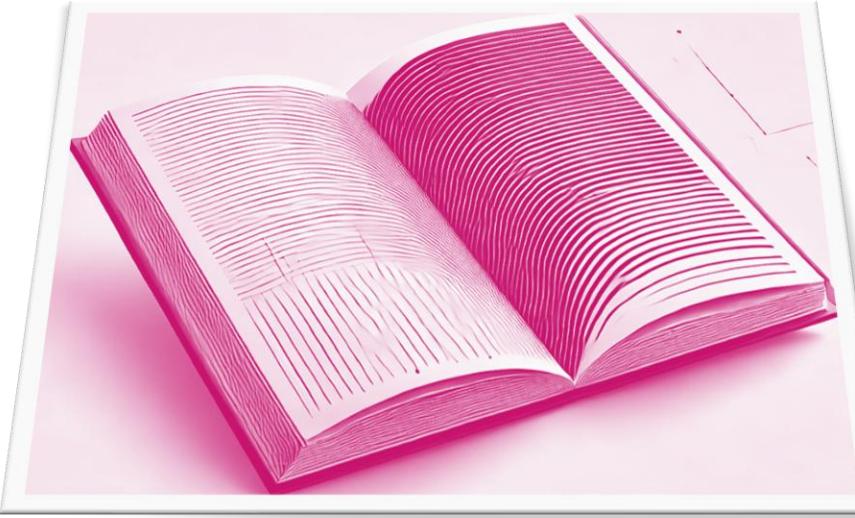


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# Introduction to proven software: code manually produced and tested

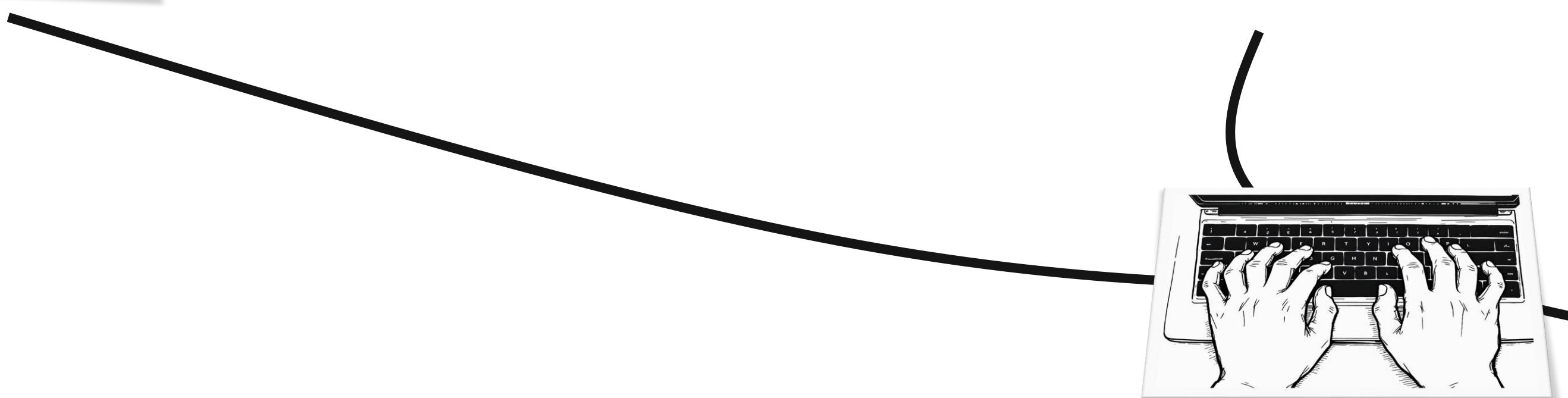
requirements



code

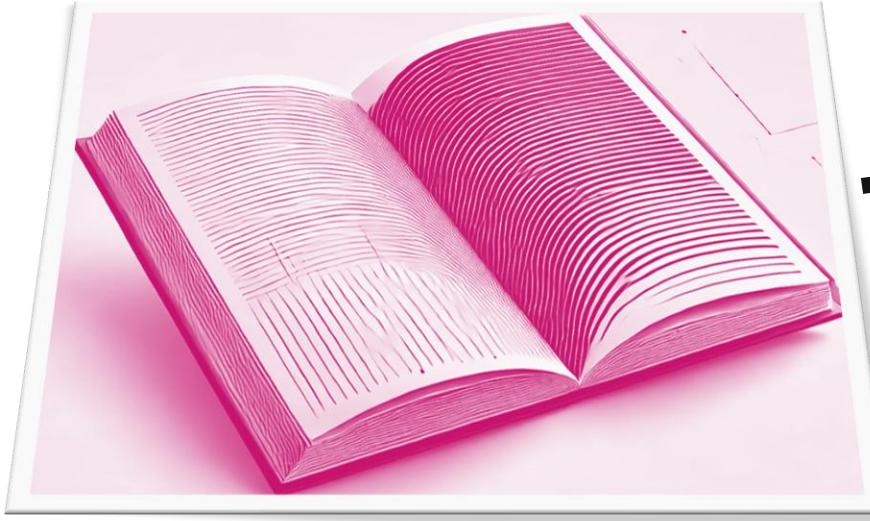


tests



# Introduction to proven software: code manually produced and **formally verified**

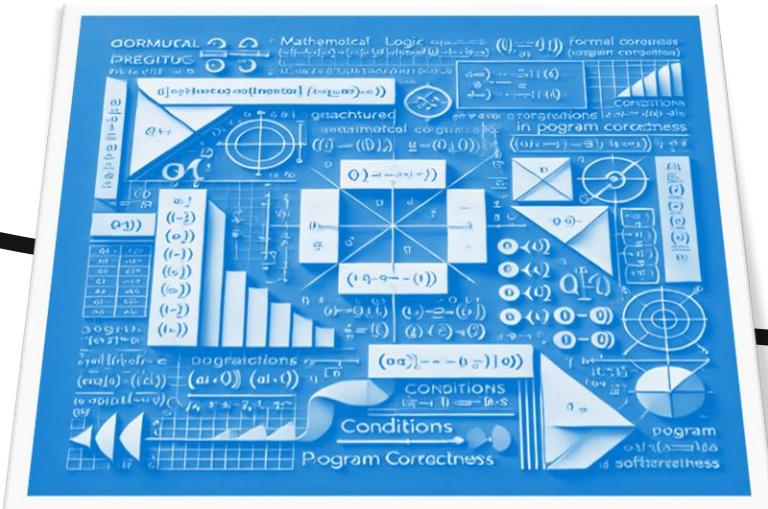
**requirements**



**code**

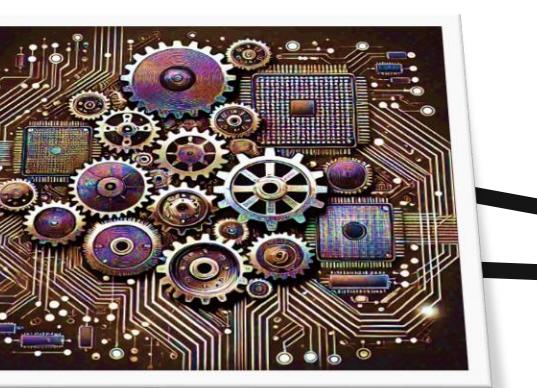


**assertions**

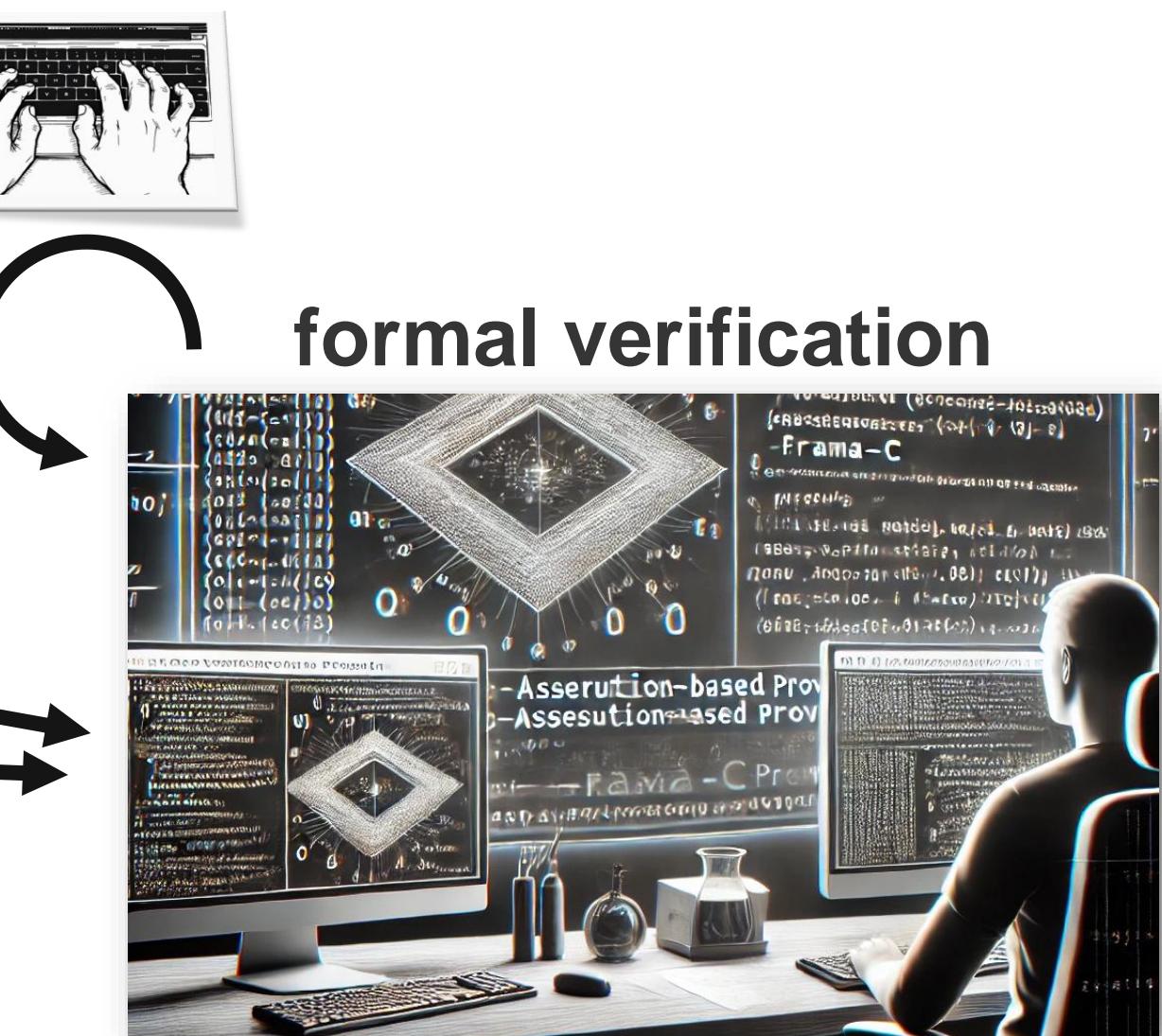


Assertions are mathematical predicates  
for hypotheses and properties to verify

**verifier**

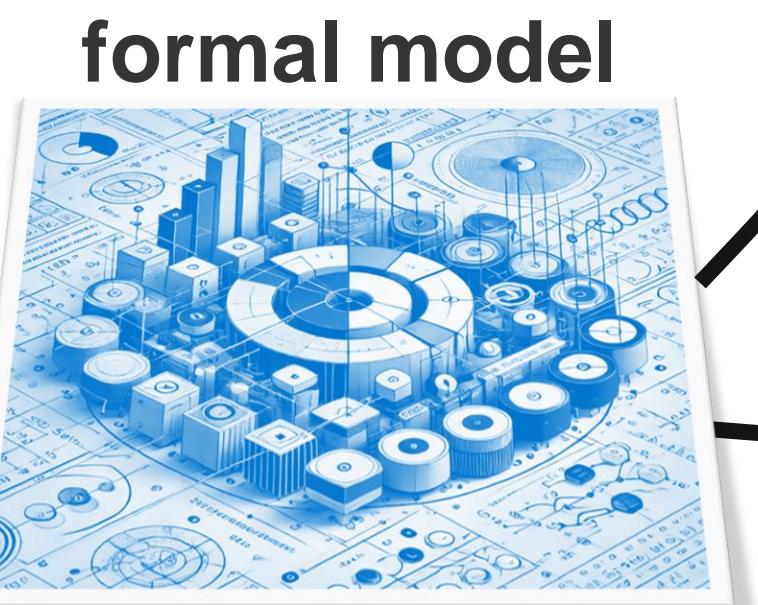
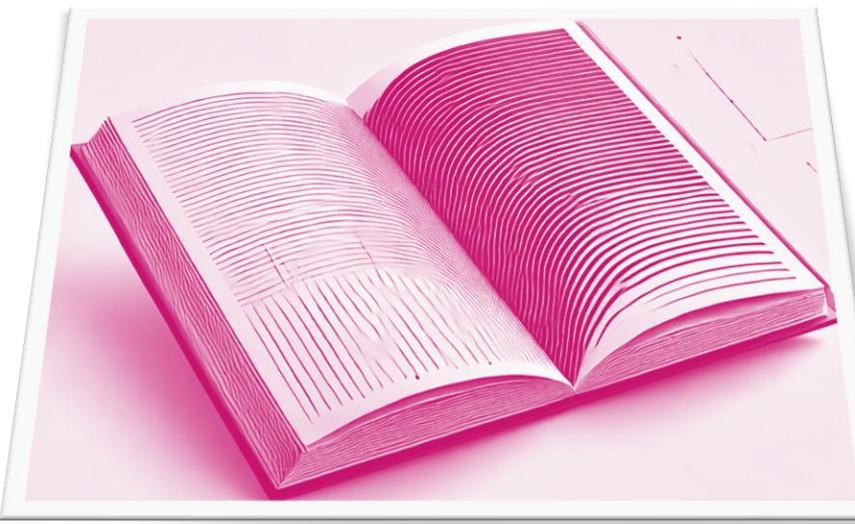


**formal verification**



# Introduction to proven software: code formally specified and verified

requirements



Formal model use mathematical notation  
to specify and implement behaviour

code



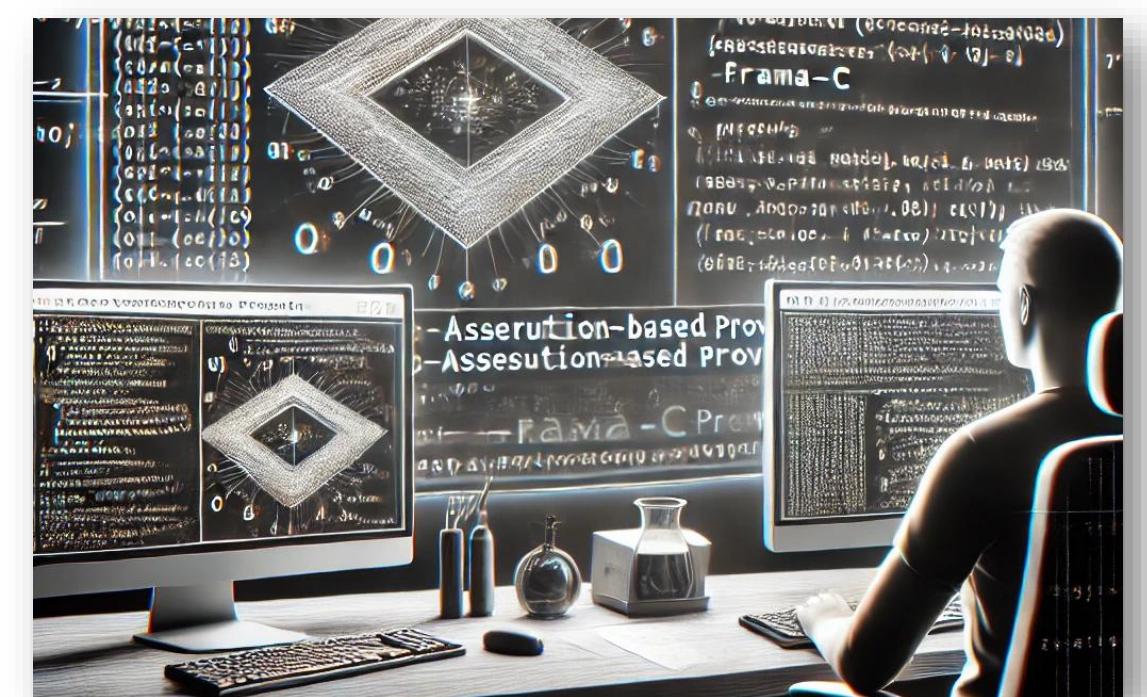
code generator



prover



formal proof



# Introduction to proven software: formal proof

## Requirement

*V1* is a variable of type Integer

Its initial value is 0

The function *increment* adds one to *V1*

## Specification

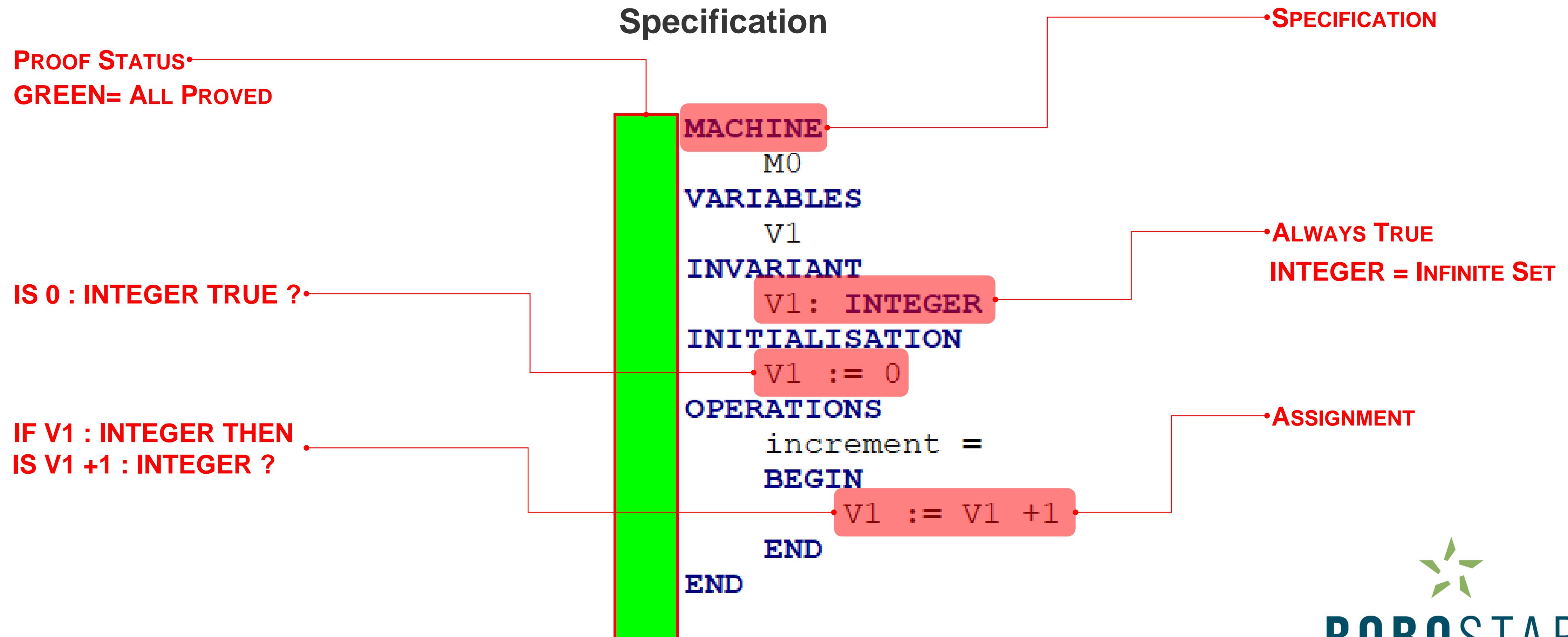
```
MACHINE  
M0  
VARIABLES  
    V1  
INVARIANT  
    V1: INTEGER  
INITIALISATION  
    V1 := 0  
OPERATIONS  
    increment =  
    BEGIN  
        V1 := V1 +1  
    END  
END
```

•SPECIFICATION

•TYPING  
INTEGER = INFINITE SET

•ASSIGNMENT

# Introduction to proven software: formal proof



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# Introduction to proven software: formal proof

Specification	Implementation
<p><b>SAME VARIABLE</b></p> <p><b>CONCRETE MEANS TO BE TRANSLATED IN C</b></p> <p><b>INT = FINITE SET</b> <math>0..2^{16}-1</math></p> <p><b>FALSE IF <math>V1 = 2^{16}-1</math></b></p> <pre>MACHINE M0 VARIABLES V1 INVARIANT V1: INTEGER INITIALISATION V1 := 0 OPERATIONS     increment =     BEGIN         V1 := V1 + 1     END END</pre>	<p><b>IMPLEMENTATION M0_i</b></p> <p><b>REFINES M0</b></p> <p><b>CONCRETE_VARIABLES</b></p> <p><b>V1</b></p> <p><b>INVARIANT</b></p> <p><b>V1 : INT</b></p> <p><b>INITIALISATION</b></p> <p><b>V1 := 0</b></p> <p><b>OPERATIONS</b></p> <p><b>increment =</b></p> <p><b>BEGIN</b></p> <p><b>V1 := V1 + 1</b></p> <p><b>END</b></p> <p><b>END</b></p>

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# Introduction to proven software: formal proof

**IF GREEN, NO TESTING REQUIRED**

**PROOF STATUS**

**NOT GREEN = NOT PROVED AUTOMATICALLY.  
EITHER WRONG OR NOT PROVABLE BY THE TOOL**

**Implementation**

```
IMPLEMENTATION M0_i
REFINES M0
CONCRETE_VARIABLES
    V1
INVARIANT
    V1 : INT
INITIALISATION
    V1 := 0
OPERATIONS
    increment =
BEGIN
    V1 := V1 + 1
END
END
```

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# Introduction to the **CLEARSY Safety Platform**



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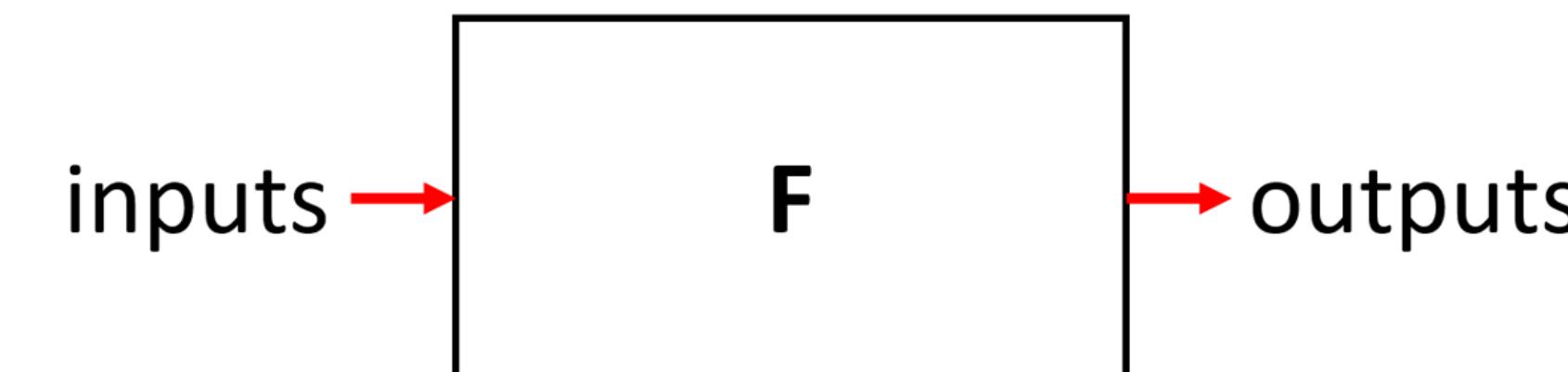
# Introduction to the CLEARSY Safety Platform

Summary of  
<https://youtu.be/A8uemvcclcU>

## What a Safety Computer is

### Safety computer

- $F == (\text{read inputs}, \text{compute}, \text{set outputs})^*$
- $F$  could harm / kill people
- Ability to check if able to execute  $F$  properly



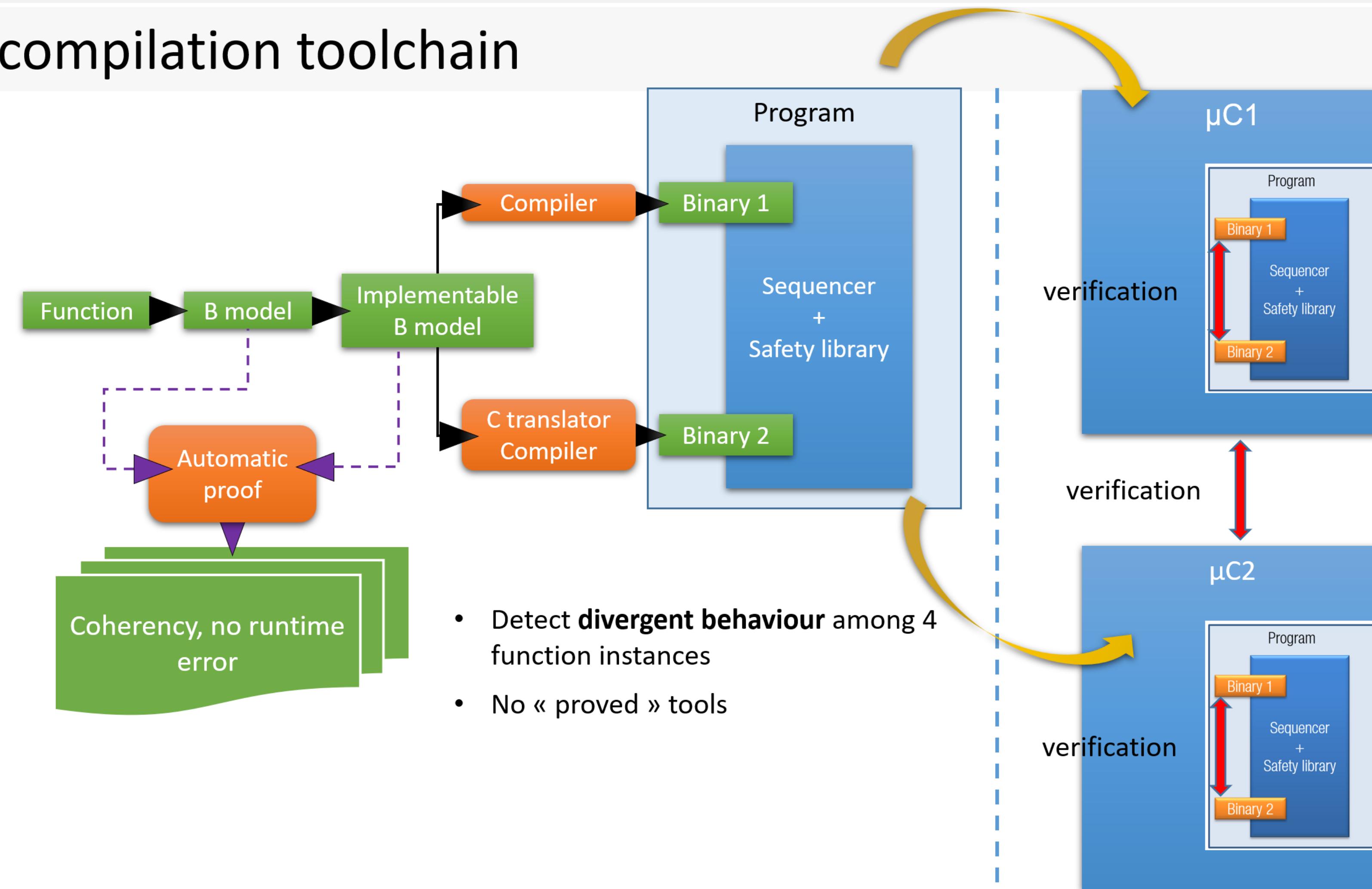
- $F$  is not safe just because a safety computer is used  
**« execute the right  $F$  and execute the  $F$  right »**

\*: for space applications, could be a reset when computer hit by high energy particle

# Introduction to the CLEARY Safety Platform

Summary of  
<https://youtu.be/A8uemvcclcU>

## Double compilation toolchain



# Introduction to the CLEARSY Safety Platform

Summary of  
<https://youtu.be/A8uemvcclcU>

## Verification

- Starter kits for education:
  - Programmed with B
  - SK0 available since Q1 2019: 5 digital I/O
  - SK0 software simulator (no safety)



Handle failures:

- **Systematic** (buggy code generator and compiler, etc.)
- **Random** (memory corruption, failing transistor, degrading clock, etc.)

**Safety is built-in, out of reach of the developer who cannot alter it**

# Introduction to the CLEARY Safety Platform

Summary of  
<https://youtu.be/A8uemvcclcU>

## Summary

- CLEARY Safety Platform
  - Safety computer
    - Execute 4 instances of the same function on 2 processors
    - Verify health regularly
    - Stop application execution and deactivate outputs if problem
  - IDE
    - Application developed with B formal language
    - Model mathematically proved



# CLEARSY Safety Platform Programming Model

- ▶ The execution is cyclic
- ▶ The function is executed regularly as often as possible similar to arduino programming (`setup()`, `loop()`)
- ▶ No underlying operating system
- ▶ No `interrupt()`
- ▶ No predefined cycle time (if outputs are not set and cross read every **50ms**, board enters panic mode)
- ▶ No `delay()`
- ▶ Inputs are values captured at the beginning of a cycle (digital I/O)
- ▶ Outputs are maintained from one cycle to another (digital I/O)
- ▶ Project skeleton is generated from board description (I/O used, naming)
- ▶ Programming is specifying and implementing the function *user\_logic*

```
init();  
while (1) {  
    instance1();  
    instance2();  
}
```



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# CLEARSY Safety Platform Applications (Industrial Version)

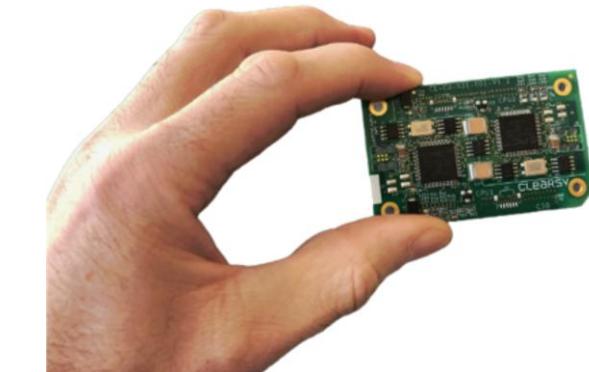
## Platform Screen Doors Controller



São Paulo, Monorail line 15  
Stockholm, City Line  
Brisbane, Cross River Rail



## SIL4 Certified



Design examination type certificate  
Certificat de type par examen de la conception

N° 9594/0262 edition 2

Attributed to

Délivré à

CLEARSY

320 Av. Archimède - Pléiades III  
F-13100 Aix-en-Provence

by

par

CERTIFER SA

18 Rue Edmond Membrey  
F-59300 VALENCIENNES

Which certifies that the design of the following product:  
Qui certifie que la conception du produit suivant :

GENERIC PRODUCT

CLEARSY SAFETY PLATFORM

BASELINE 1.0.2

Meets SIL4 requirements of the standards  
Est conforme aux exigences SIL4 des normes

CENELEC EN 50126:2017, EN 50129:2018, EN 50128:2011

## Autonomous Train Localization



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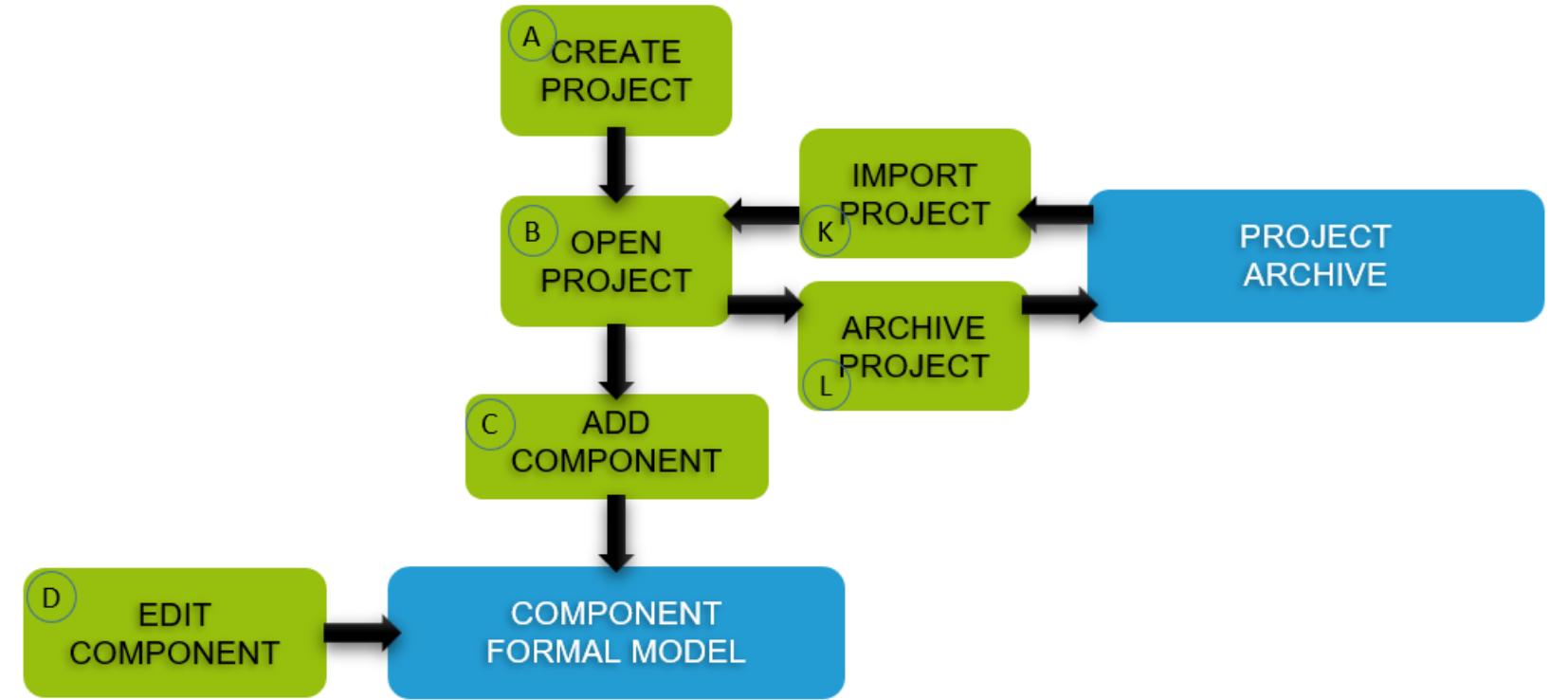
# Development Process



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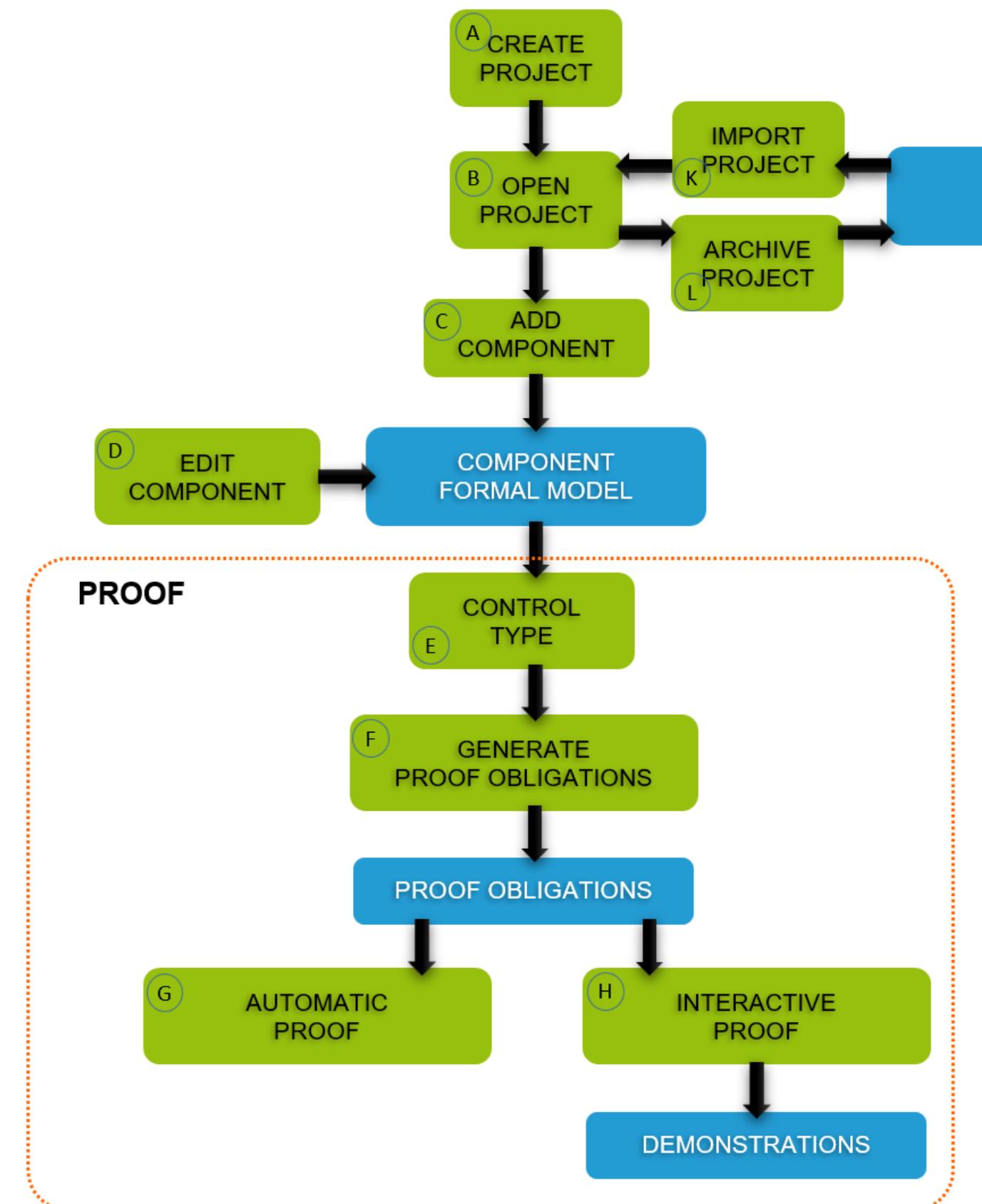
# Development process



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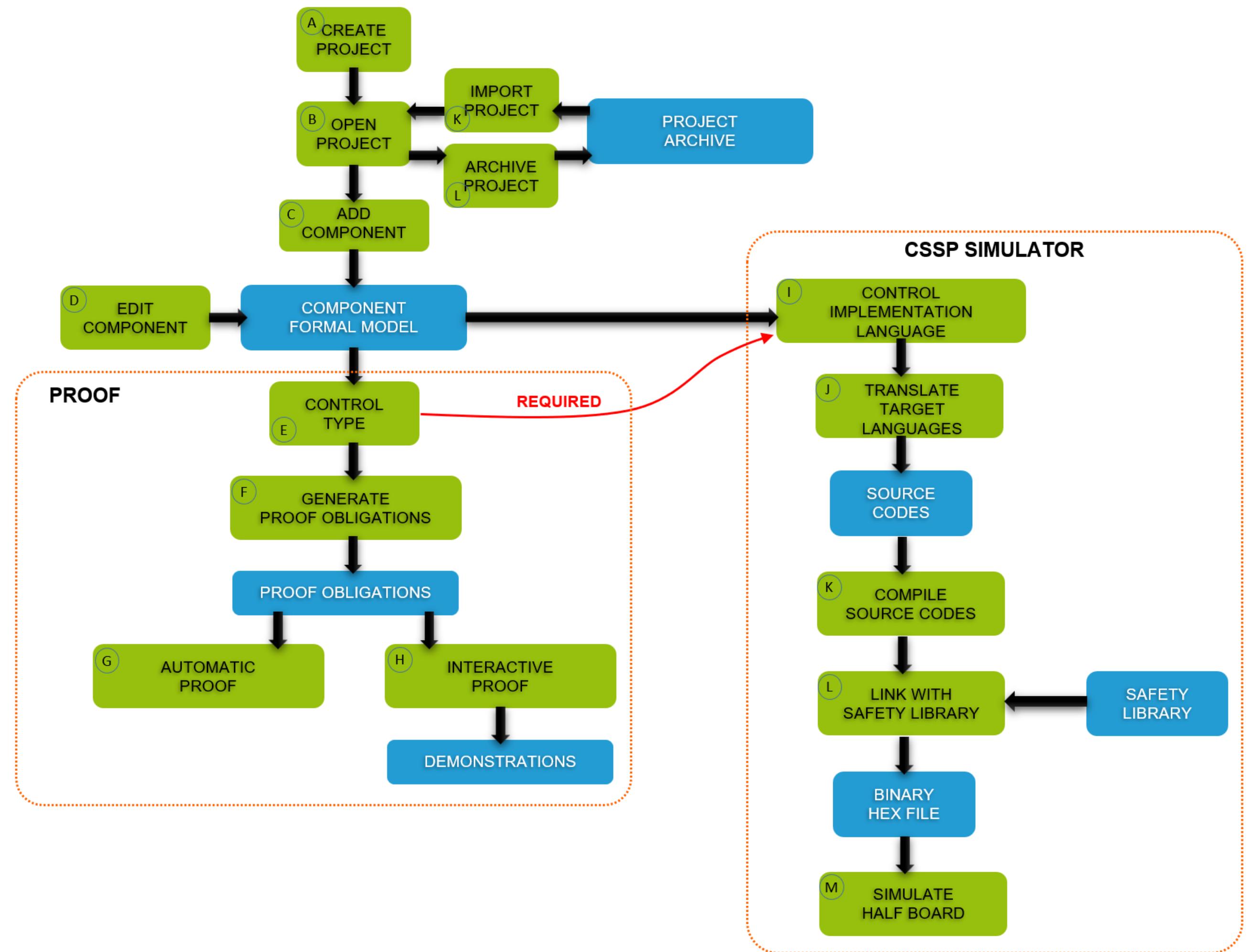
# Development process



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# Development process



# ROBOSTAR

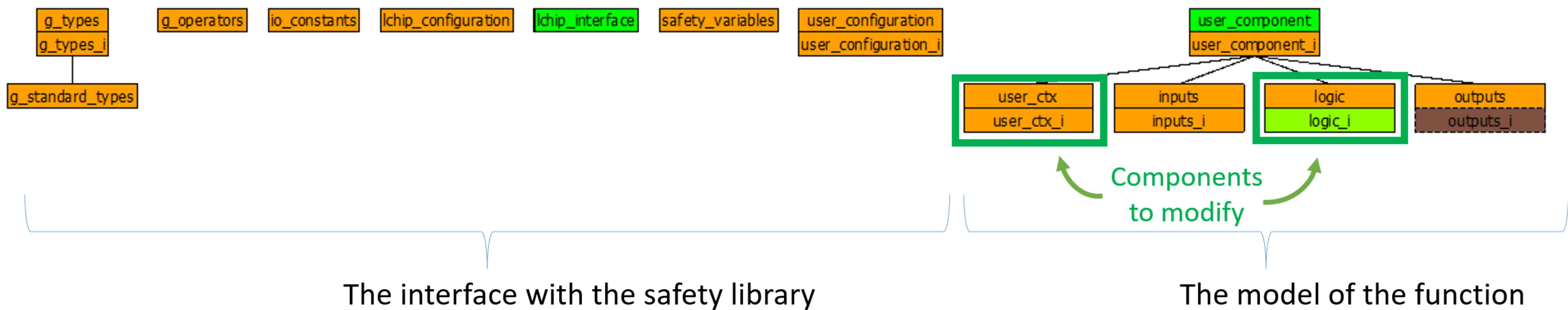
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# CLEARSY Safety Platform Project

## A CSSP project is a B project

- is generated automatically from board configuration (# IOs, naming)

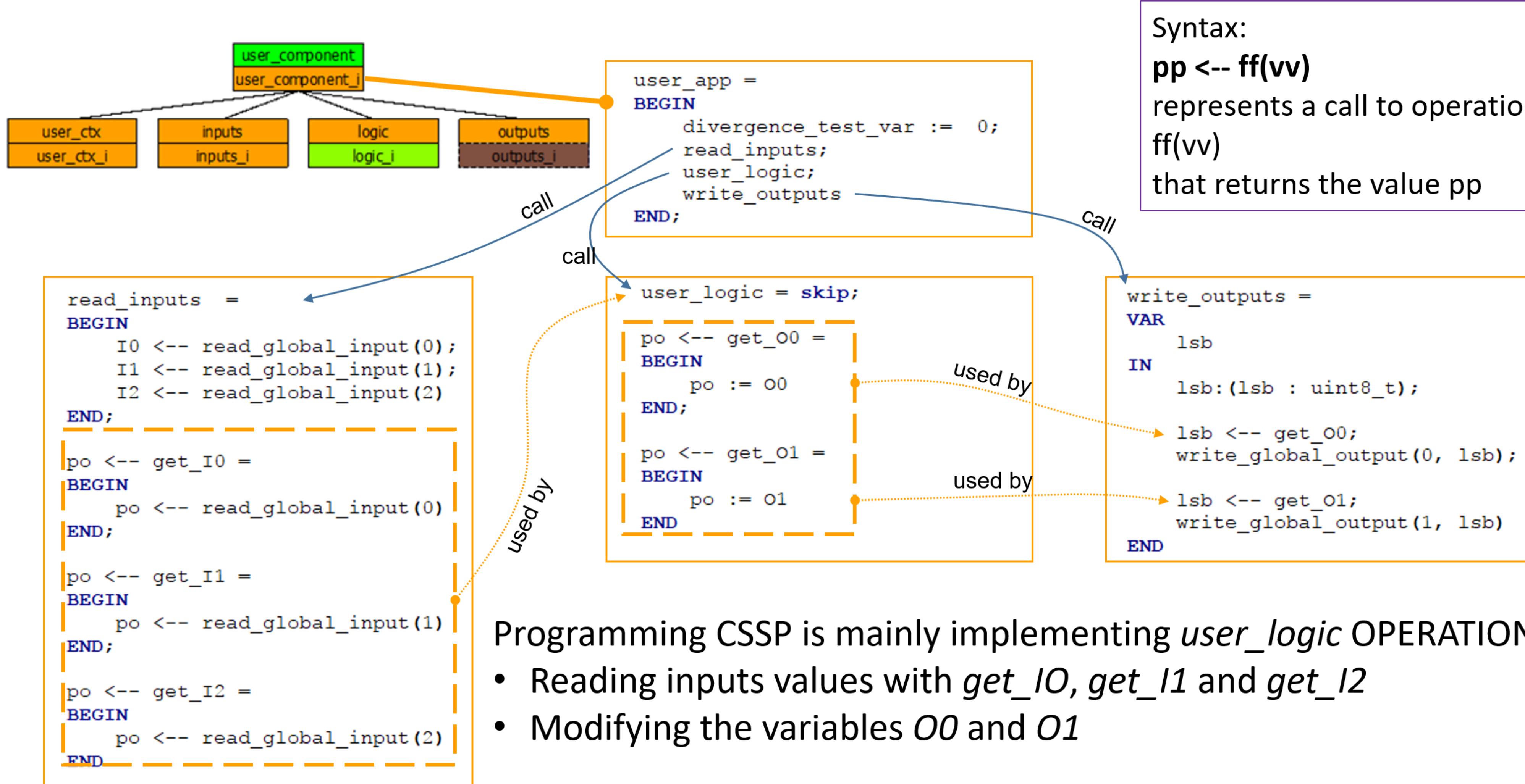
## It contains



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# Generated Models



# Bits of B



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# Bits of B : B Variables Declaration

## specification

```
ABSTRACT_VARIABLES  
 00,  
 01
```

: means « belongs

```
INVARIANT  
 00 : uint8_t &  
 01 : uint8_t
```

|| means « in parallel », « at the same time »

```
INITIALISATION  
 00 :: uint8_t  
 01 :: uint8_t
```

:: means « any value within »

## implementation

// pragma SAFETY\_VARS ————— Contains variables that will be verified

```
CONCRETE_VARIABLES  
 00,  
 01,  
 TIME_A,  
 STATUS
```

} Variables local to implementation

```
INVARIANT  
 00 : uint8_t &  
 01 : uint8_t &  
 TIME_A : uint32_t &  
 STATUS : uint8_t
```

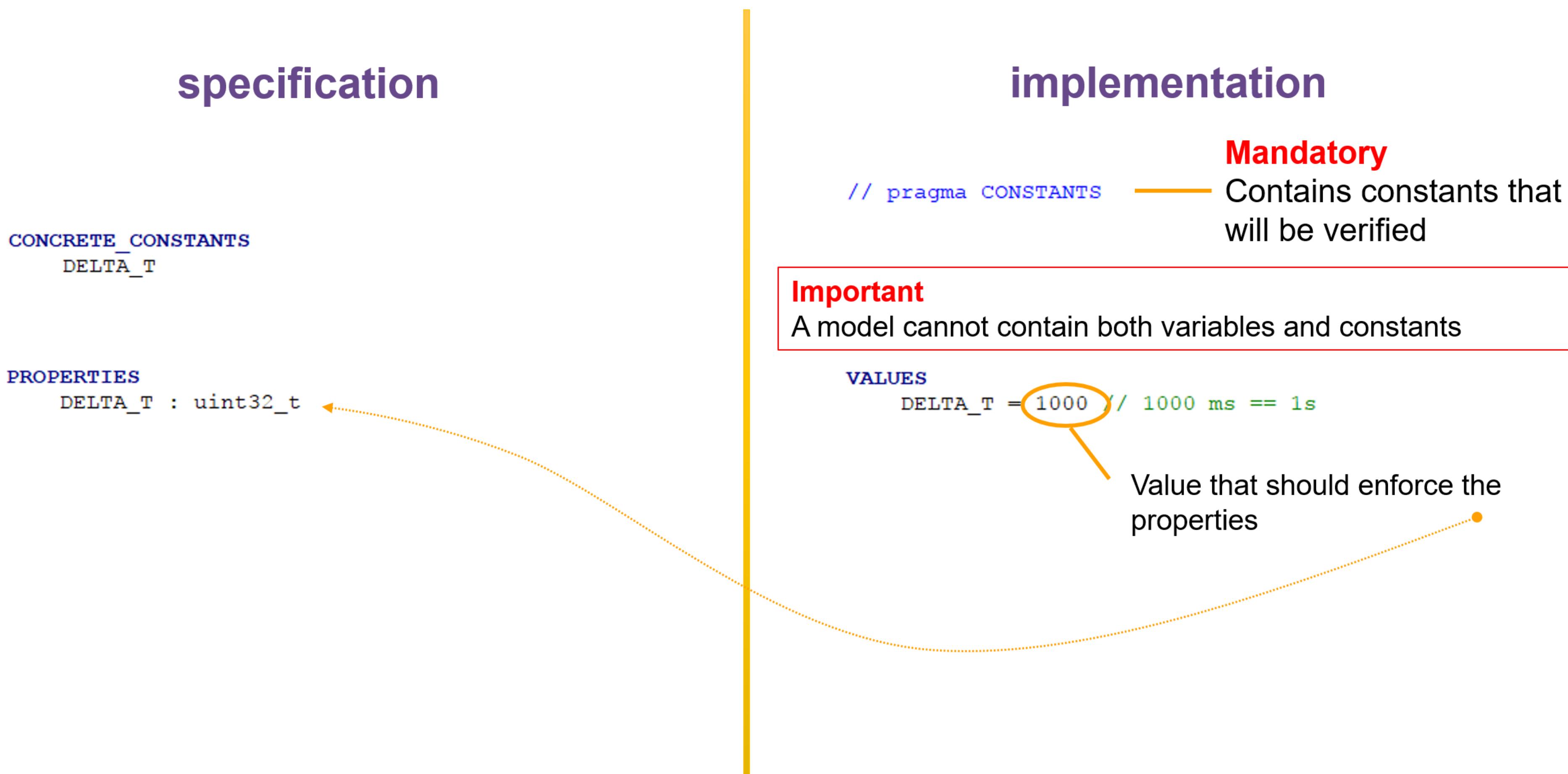
```
INITIALISATION  
 00 := IO_OFF;  
 01 := IO_OFF;  
 TIME_A := 0;  
 STATUS := SFALSE
```



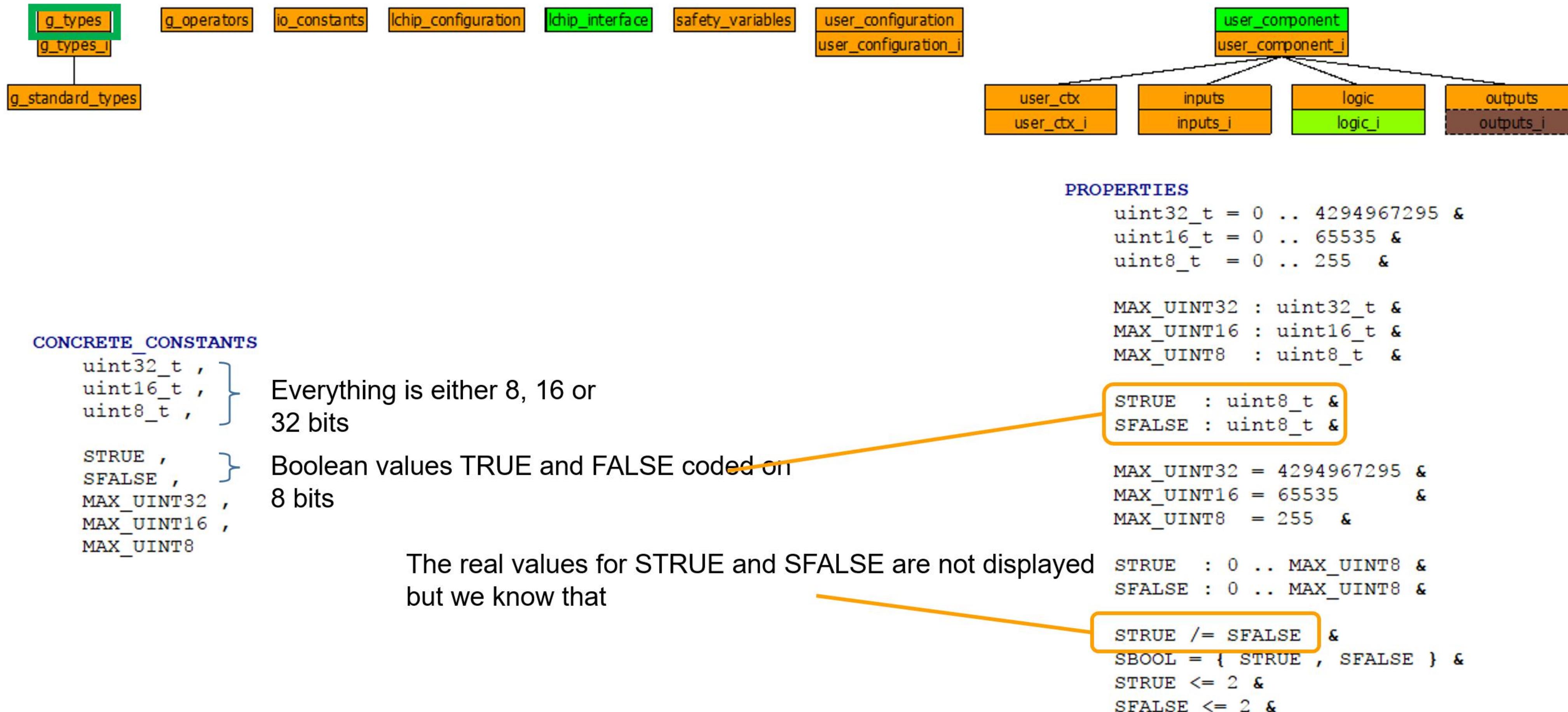
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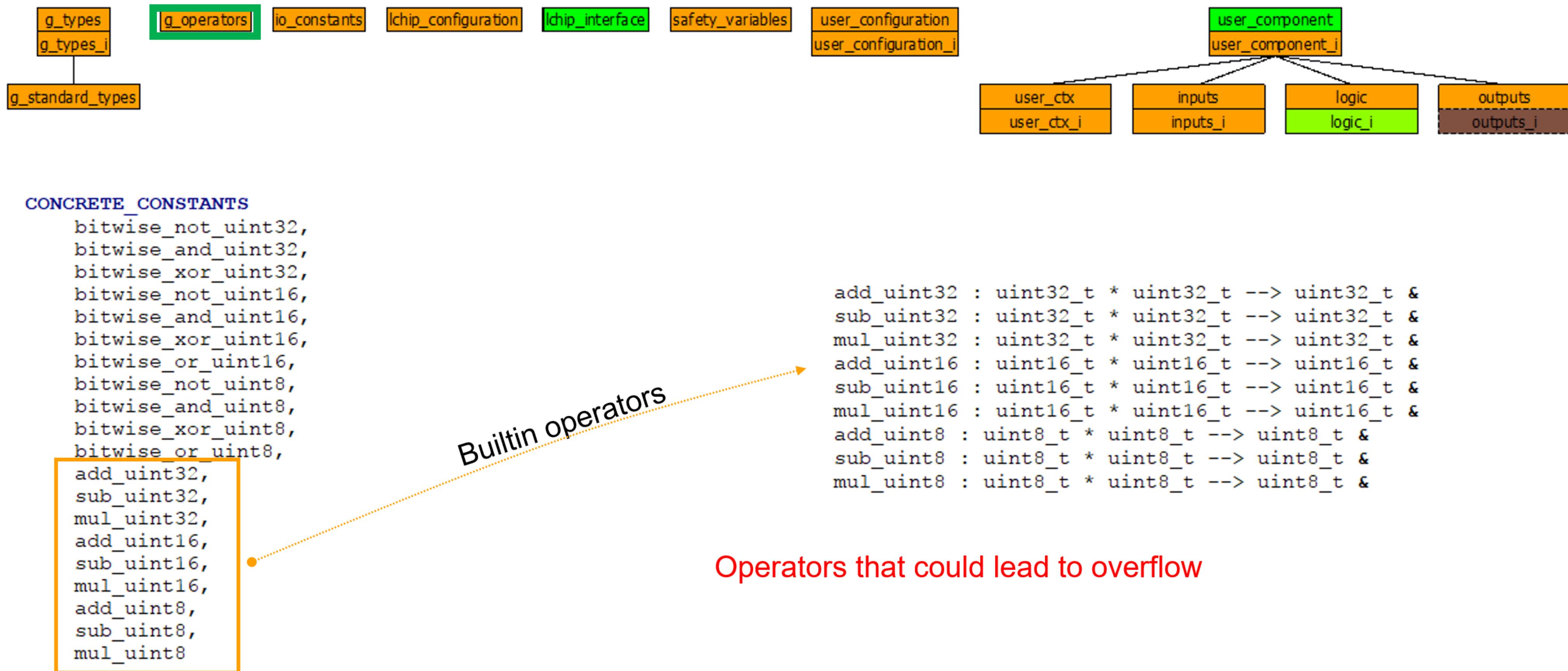
# Bits of B : B Constants Declaration



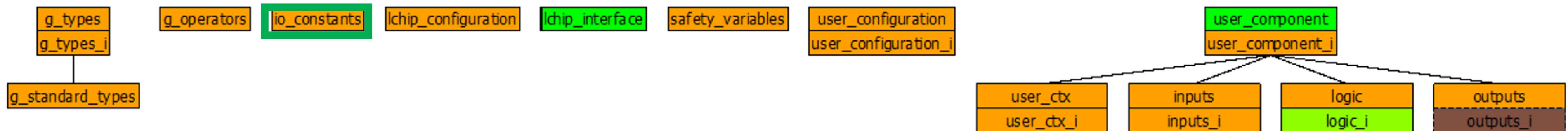
# Bits of B : CLEARSY Safety Platform Supported Types



# Bits of B : Unsigned INT Operators



# Bits of B : Inputs/outputs



**ABSTRACT\_CONSTANTS**  
TIME,  
IO\_STATE

**CONCRETE\_CONSTANTS**  
IO\_ON,  
IO\_OFF

**PROPERTIES**

```
TIME = uint32_t &
IO_STATE = uint8_t &

IO_ON : uint8_t &
IO_OFF : uint8_t &
IO_ON /= IO_OFF &
IO_ON : IO_STATE &
IO_OFF : IO_STATE
```

inputs and outputs state

values used by digital inputs and outputs

coded on 8 bits

## Verification

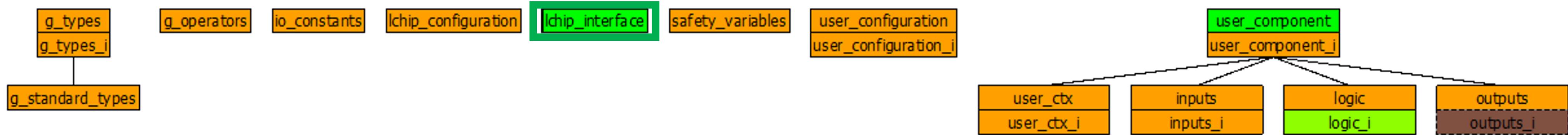
If a digital output is valued with a value different from `IO_ON` or `IO_OFF` then  $SK_0$  stops in error mode



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# Bits of B : Time



```
out <-- get_ms_tick =  
PRE  
    out : uint32_t  
THEN  
    out := ms_tick  
END
```

Returns the number of milliseconds elapsed since last reboot

# Bits of B: B Operations

Operations are populated with substitutions

Available substitutions in specification are different from the ones available in implementation

## specification

Express the properties that the variables comply with when the operation is completed independently from the algorithm implemented (*post-condition*)

To simplify, always use « becomes such that substitutions »

```
user_logic =
  BEGIN
    oo, o1 : (
      oo : uint8_t &
      o1 : uint8_t &
      not(oo = o1) )
  END;
```

Typing (mandatory)  
Constraints (optional)

# Bits of B: Implementation

## implementation

```
user_logic = skip; — do nothing
```

```
user_logic =
BEGIN
  o0 := IO_ON;
  o1 := IO_OFF
END;
```

```
user_logic =
BEGIN
  IF Var8 = 0 THEN
    o0 := IO_ON
  ELSE
    o1 := IO_ON
  END
END;
```

### Important

Only single condition (no conjunction nor disjunction)  
= < <= operators only

```
user_logic =
BEGIN
  VAR time_ IN
    time_ : (time_ : uint32_t);
    time_ <-- get_ms_tick;
  IF 2000 <= time_ THEN
    o1 := IO_ON
  END
END;
```

Local variables declaration  
Operation call

### Important

Local variables have to be typed first using  
« becomes such that » substitution

Constraints on the language to simplify the compiler

# Bits of B: user\_logic

## specification

```
user_logic = skip;
```

skip means « do no alter the variables of the model »

```
MACHINE logic
  OPERATIONS
    user_logic = skip;
    po <-- get_o1 =
      PRE
        po : uint8_t
      THEN
        po := o1
      END;
    po <-- get_o2 =
      PRE
        po : uint8_t
      THEN
        po := o2
      END
  ABSTRACT_VARIABLES
    o1,
    o2
  INVARIANT
    o1 : uint8_t &
    o2 : uint8_t
  INITIALISATION
    o1 :: uint8_t ||
    o2 :: uint8_t
  END
```

## implementation

```
user_logic = skip;
```

Minimum example:

- do nothing; outputs remain in their initial state (INITIALISATION)

```
IMPLEMENTATION logic_i
  REFINES logic
  SEES
    g_types,
    g_operators,
    io_constants,
    lchip_interface,
    inputs
    // pragma SAFETY_VARS
  CONCRETE_VARIABLES
    o1,
    o2
  INVARIANT
    o1 : uint8_t &
    o2 : uint8_t
  INITIALISATION
    o1 := IO_OFF;
    o2 := IO_OFF
  OPERATIONS
    user_logic = skip;
    po <-- get_o1 =
      BEGIN
        po := o1
      END;
    po <-- get_o2 =
      BEGIN
        po := o2
      END
  END
```



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# Bits of B: user\_logic

## specification

```
user_logic =
BEGIN
  o0 :: uint8_t ||  O0 and O1 belong to
  o1 :: uint8_t      their type
END
```

```
user_logic =   :( ) means « becomes such that »
BEGIN
  o0, o1 : (
    o0 : uint8_t &
    o1 : uint8_t &
    not(o0 = o1)  O0 and O1 belong to
  )                  their type
END               and O0 is different from O1
```

```
user_logic =
BEGIN
  o0 := IO_ON ||  Set O0 and reset O1
  o1 := IO_OFF
END
```

## implementation

```
user_logic =
BEGIN
  o0 := IO_ON;    Set O0 then
  o1 := IO_OFF;  reset O1
END
```

« then » is related to the valuation of O0 regarding O1  
O0 and O1 will be positioned at the same time at the end of the cycle



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# Using the Modelling Interface



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# Using the modelling interface



- ▶  $I_1, I_2, O_1, O_2$  belongs to  $\{IO\_OFF, IO\_ON\}$
- ▶  $O_1$  is  $IO\_ON$  iff both  $I_1$  and  $I_2$  are  $IO\_ON$
- ▶  $O_2$  is the complement of  $O_1$



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# Using the modelling interface



- ▶ I1, I2, O1, O2 belongs to `{IO_OFF, IO_ON}` • is unsigned 8 bit integer enumeration
- ▶ O1 is IO\_ON iff both I1 and I2 are IO\_ON
- ▶ O2 is the complement of O1

```
user_logic =  
BEGIN  
    o1, o2: (  
        o1 : uint8_t &  
        o2 : uint8_t  
    )  
END;
```

« 01, 02 become such that

- Type of O1 is unsigned 8 bit integer
- Type of O2 is unsigned 8 bit integer »

Minimum specification : « 01 and 02 are modified in accordance with their type »



# Using the modelling interface



- ▶ I1, I2, O1, O2 belongs to {IO\_OFF, IO\_ON}
- ▶ O1 is IO\_ON iff both I1 and I2 are IO\_ON
- ▶ O2 is the complement of O1

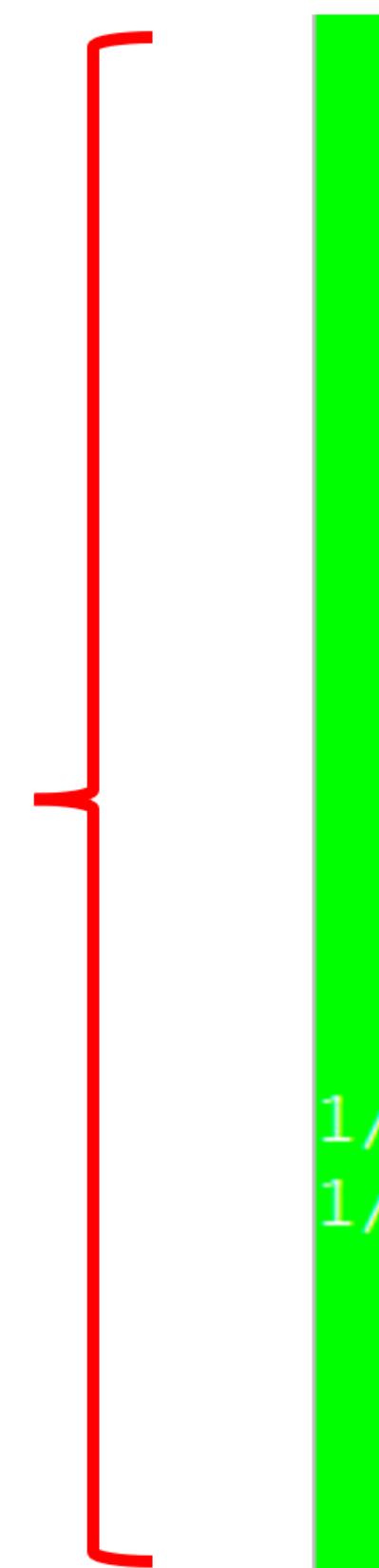
```
user_logic =  
BEGIN  
    o1, o2: (  
        o1 : uint8_t &  
        o2 : uint8_t &  
        (o1=IO_ON <=> (I1=IO_ON & I2=IO_ON)) &  
        not(o1 = o2)  
    )  
END
```



# Using the modelling interface



Green means  
« Implementation  
Fully proved  
Against  
Specification »

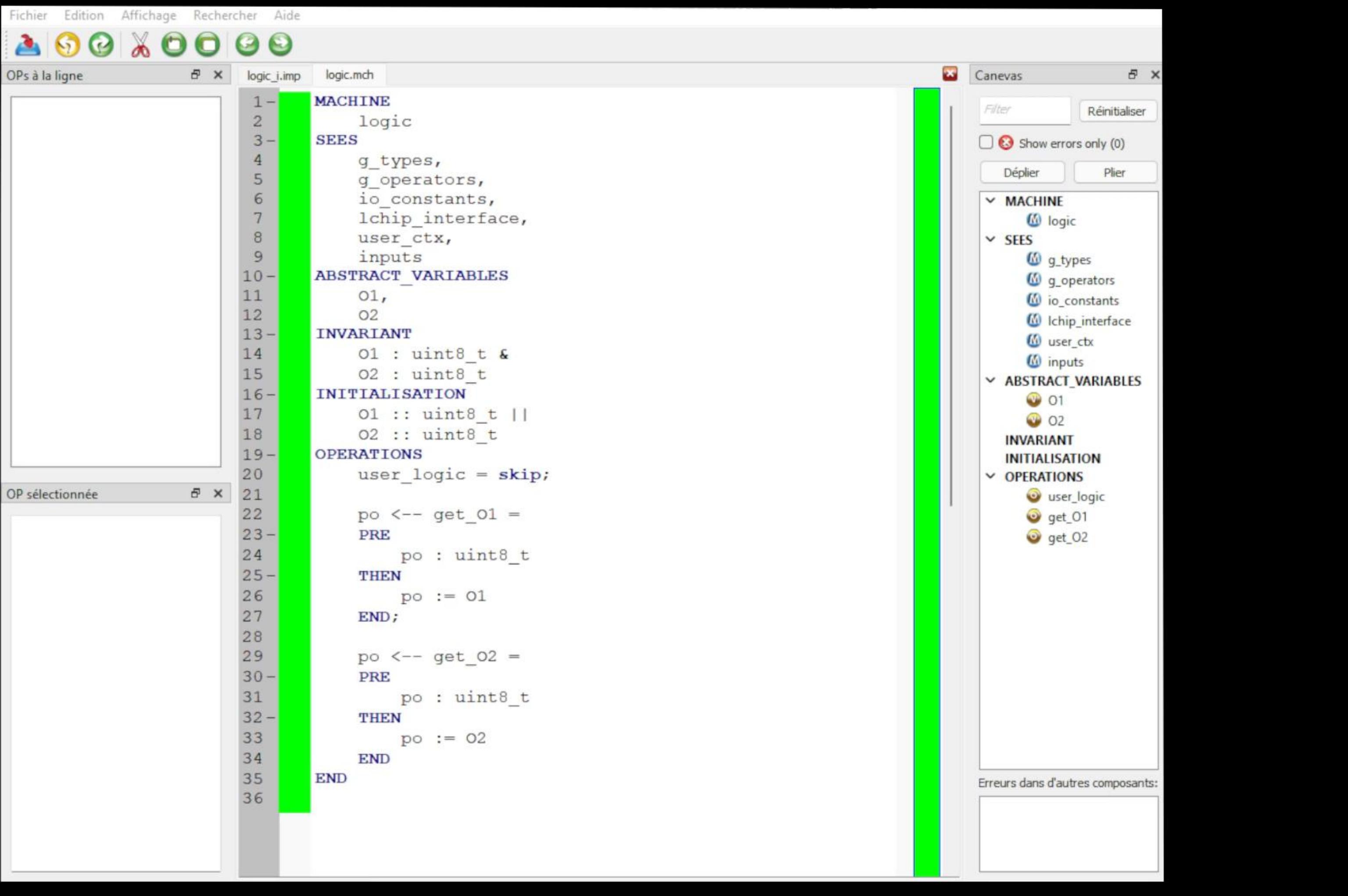


```
user_logic =
BEGIN
  VAR i1_, i2_ IN
    i1_ : (i1_ : uint8_t);
    i2_ : (i2_ : uint8_t);

    i1_ <- get_I1;
    i2_ <- get_I2; ] Get I1 and I2 values

    o1 := IO_OFF;
    o2 := IO_ON;
    IF i1_ = IO_ON THEN
      IF i2_ = IO_ON THEN
        o1 := IO_ON;
        o2 := IO_OFF
      END
    END
  END
END
```

# Using the modelling interface



The screenshot shows the Robostar modelling interface with the following components:

- Top Bar:** Fichier, Edition, Affichage, Rechercher, Aide.
- Toolbar:** Standard icons for file operations.
- Left Panel:** "OPs à la ligne" tab (selected) and "logic.mch" tab.
- Code Editor:** Displays the following UML-like code:

```
1 - MACHINE
2 -   logic
3 - SEES
4 -   g_types,
5 -   g_operators,
6 -   io_constants,
7 -   lchip_interface,
8 -   user_ctx,
9 -   inputs
10 - ABSTRACT_VARIABLES
11 -   O1,
12 -   O2
13 - INVARIANT
14 -   O1 : uint8_t &
15 -   O2 : uint8_t
16 - INITIALISATION
17 -   O1 ::= uint8_t ||| O2 ::= uint8_t
18 -
19 - OPERATIONS
20 -   user_logic = skip;
21 -
22 -   po <-- get_O1 =
23 -     PRE
24 -       po : uint8_t
25 -     THEN
26 -       po := O1
27 -     END;
28 -
29 -   po <-- get_O2 =
30 -     PRE
31 -       po : uint8_t
32 -     THEN
33 -       po := O2
34 -     END
35 -
36 - END
```
- Right Panel:** "Canevas" tab (selected). It contains a tree view of the model elements and a "Filter" button.
- Bottom Panel:** Shows error messages: "Erreurs dans d'autres composants:" followed by a redacted box.

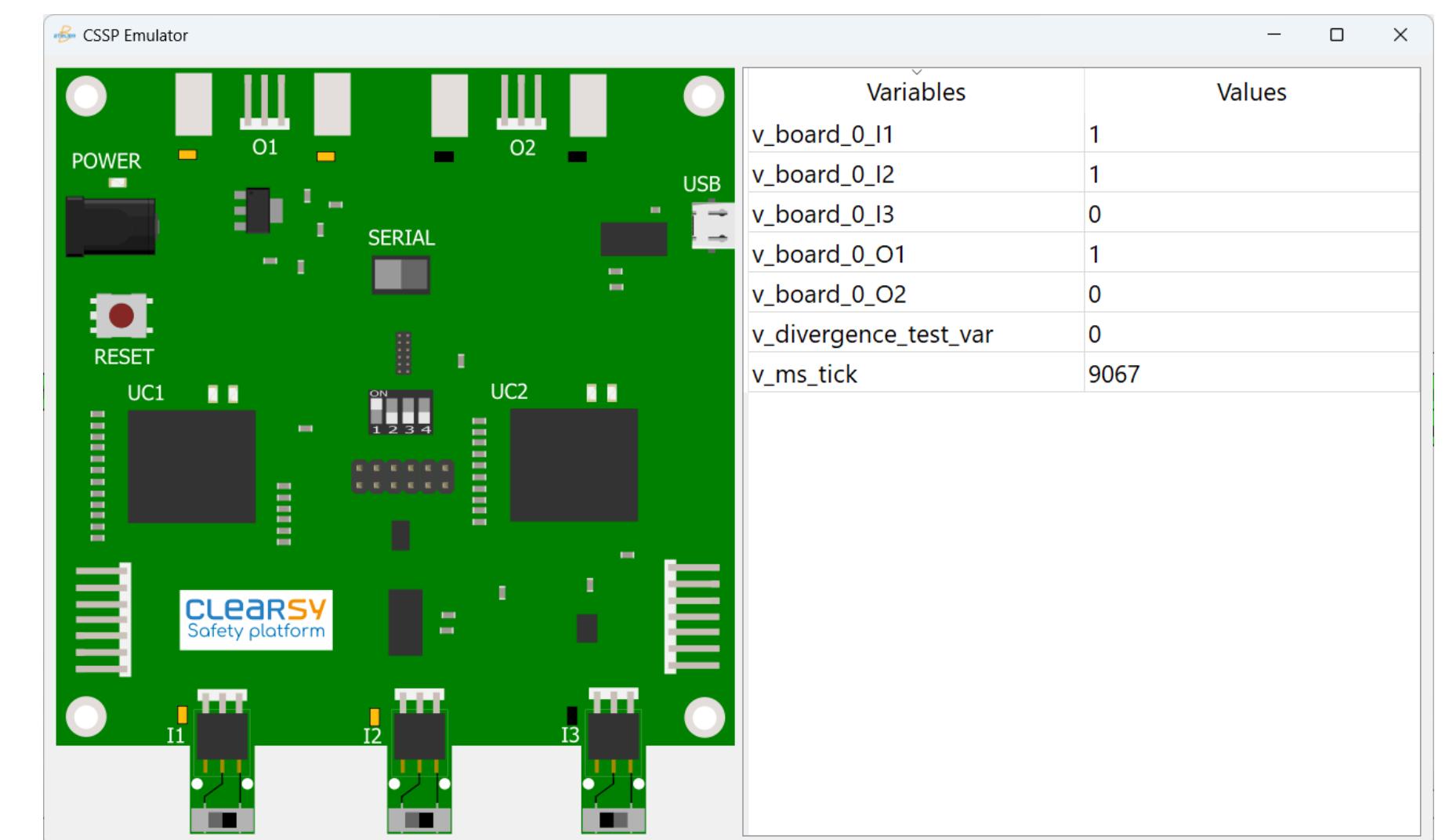
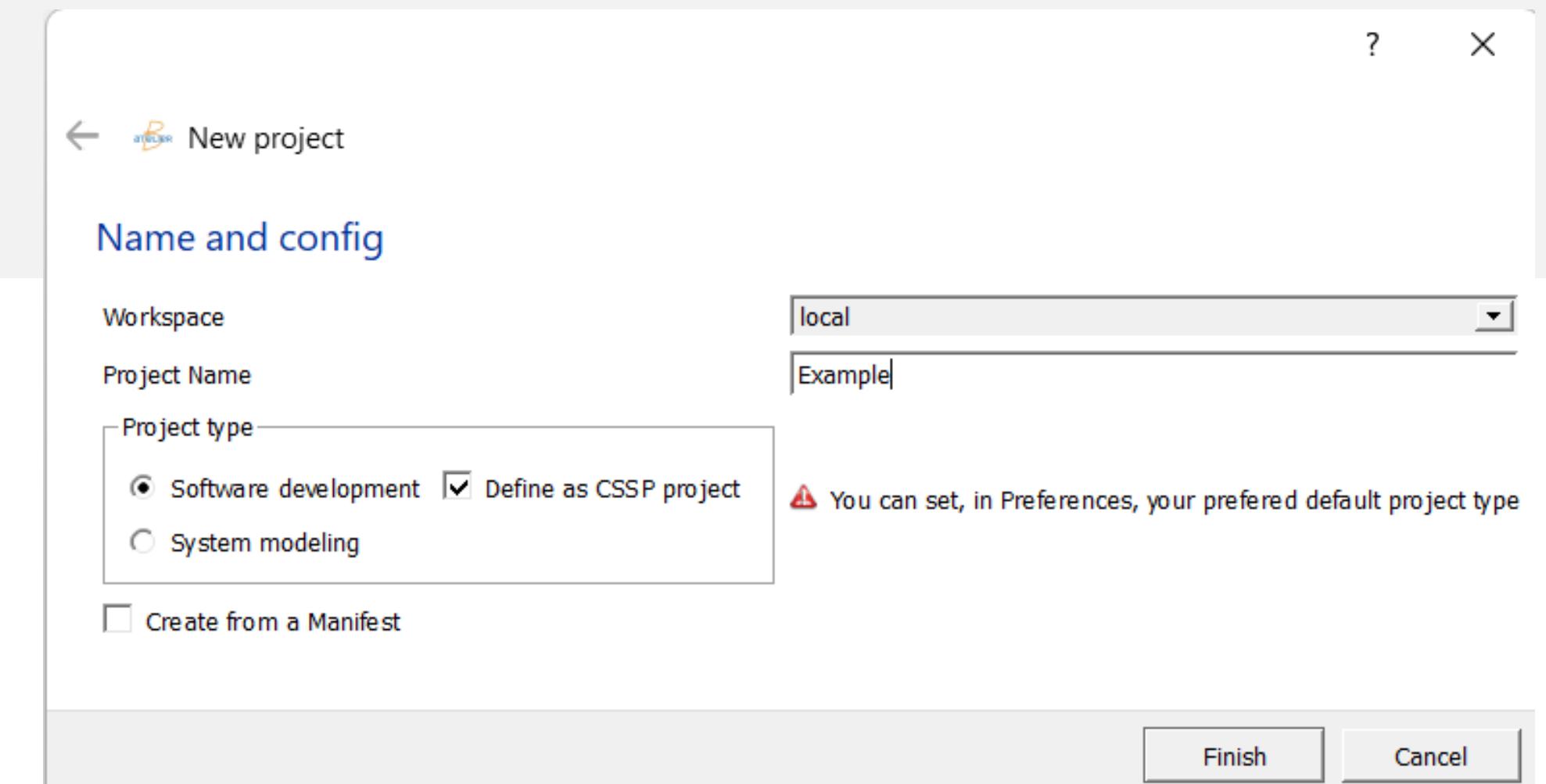


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# Using the modelling interface

- Edit logic.mch
- Edit logic\_i.imp
- Save to check if the model
  - has a correct type
  - is proved
- Start the simulator (Ctrl+D)
  - Once the simulator has started, click on the inputs to observe the expected behaviour



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Next

# Programming The CLEARY Safety Platform

Paulo Bezerra, Thierry Lecomte



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