

Lab03-Recursive Function

CS363-Computability Theory, Xiaofeng Gao, Spring 2016

* Please upload your assignment to FTP or submit a paper version on the next class

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1. Show that the following functions are primitive recursive:

$$(a) \text{ half}(n) = \begin{cases} \frac{n}{2}, & \text{if } n \text{ is even,} \\ \frac{n-1}{2}, & \text{if } n \text{ is odd.} \end{cases}$$

Proof:

$$\text{half}(0) = 0$$

.

$$\text{half}(x+1) = \begin{cases} \frac{x}{2}, & \text{if } x \text{ is even,} \\ \frac{x+1}{2}, & \text{if } x \text{ is odd.} \end{cases} = \begin{cases} \text{half}(x), & \text{if } x \text{ is even,} \\ \text{half}(x) + 1, & \text{if } x \text{ is odd.} \end{cases}$$

So, the function is primitive recursive.

(b) $\max\{x_1, x_2, \dots, x_n\}$ = the maximum of x_1, x_2, \dots, x_n .

(c) $f(x)$ = the sum of all prime divisors of x .

(d) $g(x) = x^x$.

2. Show the computability of the following functions by minimalisation.

(a) $f^{-1}(x)$, if $f(x)$ is a total injective computable function.

(b) $f(a) = \begin{cases} \text{the least non-negative integral root of } p(x) - a & (a \in \mathbb{N}), \\ \text{undefined if there's no such root,} \end{cases}$

where $p(x)$ is a polynomial with integer coefficients.

(c) $f(x, y) = \begin{cases} x/y & \text{if } y \neq 0 \text{ and } y|x, \\ \text{undefined} & \text{otherwise.} \end{cases}$

3. Let $\pi(x, y) = 2^x(2y + 1) - 1$. Show that π is a computable bijection from \mathbb{N}^2 to \mathbb{N} , and that the functions π_1, π_2 such that $\pi(\pi_1(z), \pi_2(z)) = z$ for all z are computable.

4. Show that the following function is primitive recursive (with the help of $\pi(x, y)$, perhaps):

$$\begin{aligned} f(0) &= 1, \\ f(1) &= 1, \\ f(n+2) &= f(n) + f(n+1). \end{aligned}$$

5. Coding Technology.

Any number $x \in \mathbb{N}$ has a unique expression as

$$(1) \ x = \sum_{i=0}^{\infty} \alpha_i 2^i, \text{ with } \alpha_i = 0 \text{ or } 1, \text{ for all } i.$$

Hence, if $x > 0$, there are unique expressions for x in the forms

$$(2) \ x = 2^{b_1} + 2^{b_2} + \dots + 2^{b_l}, \text{ with } 0 \leq b_1 < b_2 < \dots < b_l \text{ and } l \geq 1, \text{ and}$$

(3) $x = 2^{a_1} + 2^{a_1+a_2+1} + \dots + 2^{a_1+a_2+\dots+a_k+k-1}$. (The expression (3) is a way of regarding x as coding the sequence (a_1, a_2, \dots, a_l) of numbers)

Show that each of the functions α , l , b , a defined below is computable.

(a) $\alpha(i, x) = \alpha_i$ as in the expression (1);

(b) $l(x) = \begin{cases} l \text{ as in (2),} & \text{if } x > 0, \\ 0 & \text{otherwise;} \end{cases}$

(c) $b(x) = \begin{cases} b_i \text{ as in (2),} & \text{if } x > 0 \text{ and } 1 \leq i \leq l, \\ 0 & \text{otherwise;} \end{cases}$

(d) $a(i, x) = \begin{cases} a_i \text{ as in (3),} & \text{if } x > 0 \text{ and } 1 \leq i \leq l, \\ 0 & \text{otherwise;} \end{cases}$