

# Lab05-Numbering Programs

CS363-Computability Theory, Xiaofeng Gao, Spring 2015

\* Please upload your assignment to FTP or submit a paper version on the next class.

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1. Show that there is a total computable function  $k$  such that for each  $n$ ,
  - (a)  $k(n)$  is an index of the function  $\lfloor \sqrt[n]{x} \rfloor$ .
  - (b)  $W_{k(n)}^{(m)} = \{(y_1, \dots, y_m) : y_1 + y_2 + \dots + y_m = n\}$  ( $m \geq 1$ ).
  - (c)  $E_{k(n)} = W_n$ .
2. (a) Find  $P_{2057}$ . Distinguish what are  $\phi_{2057}(x)$  and  $\phi_{2057}^{(n)}(x_1, \dots, x_n)$  and their corresponding  $W_{2057}(x)$ ,  $E_{2057}(x)$  and  $W_{2057}^{(n)}(x)$ ,  $E_{2057}^{(n)}(x)$ ;  
 (b) Let  $P$  be the program  $J(1,2,4)$ ,  $Z(1)$ ,  $S(1)$ . Calculate  $\gamma(P)$ .
3. Use Cantor's Diagonal Method to complete the following tasks:
  - (a) Let  $f_0, f_1, \dots$  be an enumeration of partial functions from  $\mathbb{N}$  to  $\mathbb{N}$ . Construct a function  $g$  from  $\mathbb{N}$  to  $\mathbb{N}$  such that  $Dom(g) \neq Dom(f_i)$  for each  $i$ .
  - (b)  $f$  is a partial function from  $\mathbb{N}$  to  $\mathbb{N}$ . Construct a non-computable function  $g$  such that  $g(x) \simeq f(x)$  for  $x \leq m$ , here  $m \in \mathbb{N}$ .
4. (a) Show that the functions  $s_n^m$  defined in s-m-n theorem are all primitive recursive.  
 (b) Show that for each  $m \in \mathbb{N}$  there is a total  $(m+1)$ -ary computable function  $s^m$  such that for all  $n$ ,  $\phi_e^{(m+n)}(\mathbf{x}, \mathbf{y}) \simeq \phi_{s^m(e, \mathbf{x})}^{(n)}(\mathbf{y})$ , where  $\mathbf{x}, \mathbf{y}$  are  $m$ - and  $n$ -tuples respectively.  
 (*Hint. Consider the definition of  $s_n^m(e, \mathbf{x})$  given in the proof of s-m-n Theorem. The only way in which  $n$  was used was in determining how many of the  $r_1, r_2, \dots$  to transfer to  $R_{m+1}, R_{m+2}, \dots$ . Now recall that the effect of  $P_e$  depends only on the original contents of  $R_1, \dots, R_{\rho(P_e)}$ , where  $\rho$  is the function defined as the smallest register number such that none of the registers larger than  $\rho(P_e)$  will be affected during the computation of  $P_e$ ;  $\rho(P_e)$  is independent of  $n$ .)*)
5. Alternative Selection of  $\pi$

The  $\pi$  function where  $\pi(x, y) = 2^x(2y + 1) - 1$  can enumerate linearly all pairs of natural numbers  $(x, y) \in \mathbb{N} \times \mathbb{N}$ . However, it does not generate a completely diagonal-zigzag trace in the first quadrant of the plane. Correspondingly, instead of applying this  $\pi$  function, we can define an alternative bijection  $\pi'$ , such that  $\pi' : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  and it grows diagonally according to the right figure. Thus we have:

$$\begin{aligned} \pi'(0, 0) &= 0, \pi'(1, 0) = 1, \pi'(0, 1) = 2, \\ \pi'(2, 0) &= 0, \pi'(1, 1) = 4, \pi'(0, 2) = 5, \text{ etc.} \end{aligned}$$

Now please develop a mathematical formula for  $\pi'$ , (like the notation of original  $\pi$ ), and prove the correctness of your design.

**Hint:**

- (a) For the pairs in the diagonal we have the property that  $x + y$  is a constant.
- (b) The diagonal given by  $x + y = n$  consists of  $n + 1$  pairs.
- (c) You can count the elements occurring before the pair  $(x_0, y_0)$ .

$x \backslash y$	0	1	2	3	4
0	0	2	5	9	14
1	1	4	8	13	19
2	3	7	12	18	25
3	6	11	17	24	32
4	10	16	23	31	40