

CS381 Exercise 5

Name: Zhang Yupeng

Student ID: 5130309468

1. Prove the low-high algorithm for computing \odot

Proof:

We define that $a \% b$ is equal to $a \bmod b$.

The low-high algorithm for computing \odot is:

$$ab \% (2^n + 1) = \begin{cases} ab \% 2^n - ab / 2^n & , ab \% 2^n \geq ab / 2^n \\ ab \% 2^n - ab / 2^n + 2^n + 1 & , ab \% 2^n < ab / 2^n \end{cases}$$

The $ab \% 2^n$ corresponds the lower n bits of ab .

The $ab / 2^n$ corresponds the higher n bits of ab .

$$\text{Let } ab = q(2^n + 1) + r = \begin{cases} q2^n + (q + r) & , q + r < 2^n \\ (q + 1)2^n + (q + r - 2^n) & , q + r \geq 2^n \end{cases}$$

First, when $q + r < 2^n$, $ab / 2^n = q$, $r + q = ab - q2^n$, $r = ab - q2^n - q$

That is, $r = ab \% 2^n - ab / 2^n$

Next, when

$$q + r \geq 2^n, ab / 2^n = q + 1, q + r - 2^n = ab - (q + 1)2^n, r = ab - (q + 1)2^n - (q + 1) + 2^n + 1$$

That is, $r = ab \% 2^n - ab / 2^n + 2^n + 1$

So, we prove the low-high algorithm for computing \odot

2. Prove that the In-structure in IDEA is an involution.

Proof:

From the figure following we can see that the input of In-structure is K_1, \dots, K_4 , and the input of MA-structure is $a = K_1 \oplus K_3$ and $b = K_2 \oplus K_4$, and the output of MA-structure is c and d .

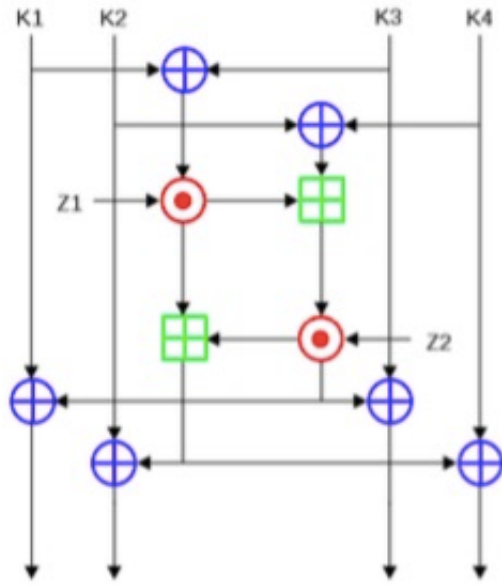


Figure 1: In-structure

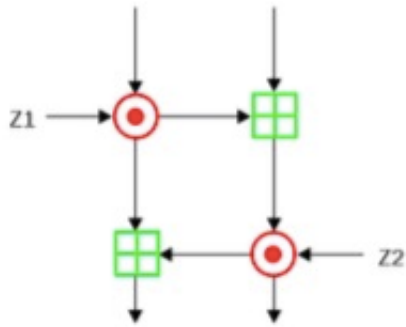


Figure 2: MA-structure

Then, we can see that after one In-structure operation, the output is:

$$K'_1 = K_1 \oplus d$$

$$K'_2 = K_2 \oplus c$$

$$K'_3 = K_3 \oplus d$$

$$K'_4 = K_4 \oplus c$$

In the second round of In-structure operation, the input of the second MA-structure will be:

$$d' = K'_1 \oplus K'_3 = K_1 \oplus d \oplus K_3 \oplus d = K_1 \oplus K_3$$

$$b' = K_2 \oplus K_4$$

So, the output of the second MA-structure is still c and d .

Therefore, the output of the second In-structure will be:

$$K_1'' = K_1' \oplus d = K_1 \oplus d \oplus d = K_1$$

$$K_2'' = K_2$$

$$K_3'' = K_3$$

$$K_4'' = K_4$$

So, we've proved that the In-structure is involution.