

ADS

2. Red-Black tree

- (1) every node is black or red
 - (2) root is black
 - (3) every leaf is black (NIL)
 - (4) if a node is red, both child are black
 - (5) all path contain the same number of black nodes
- black-height of any node x , denoted by $bh(x)$ is the number of black nodes on any simple path from x (x not included) down to a leaf $bh(Tree) = bh(w)$
- N internal nodes at most $2(\ln N + 1)$ height

Insert:

Color red
- can be down iteratively

Case: 1

Case: 2

Case: 3

Symmetric

$$T = O(h) = O(\ln N)$$

3. B+ tree (order M)

- (1) the root is either a leaf or has 2 and M children
 - (2) All non-leaf nodes (except the root) have between $\lceil M/2 \rceil$ and M children.
 - (3) All leaves are the same depth
- a general B+ tree of order M
Depth(M, N) = $O(\log_{\lceil M/2 \rceil} N)$

4. Inverted Index FILE

Term - Document Incidence Matrix $T(M, N)$
Compact-Version Inverted File Index $T_{find}(M, N)$
 $= O(M \log M \log N)$
 $= O(\log N)$

Inverted file contains a list of pointers

NO.	Term	Times, Documents
1	truck	<3; 1, 3, 4>

$\Rightarrow <3; (1;3); (3;3); (4;7)>$

```

while ( read a document D ) {
  while ( read a term T in D ) {
    if ( Find ( Dictionary, T ) == False )
      Insert ( Dictionary, T );
    Get T's posting list;
    Insert a node to T's posting list;
  }
  write the inverted index to disk.
}

```

word streaming stop words

Access a term: 1. Search trees
Hashing

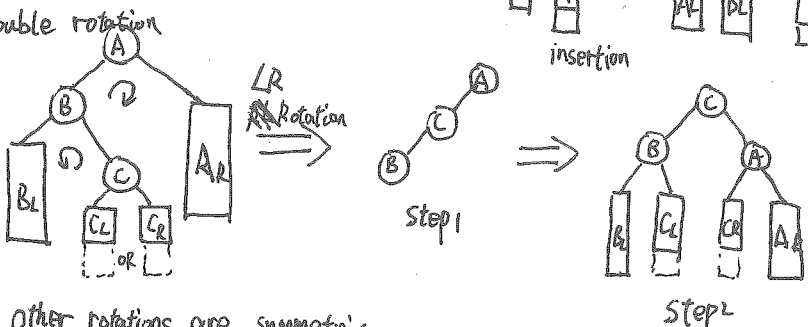
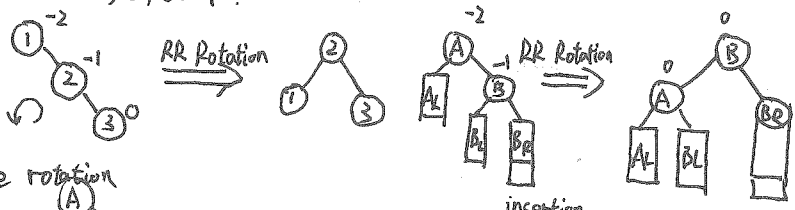
Relevant R_R Irrelevant I_R
Retrieved R_N Not Retrieved I_N

Precision $P = \frac{R_R}{R_R + I_R}$
Recall $R = \frac{R_R}{R_R + R_N}$
Precision $= \frac{R_R}{R_R + I_R}$
Recall $R = \frac{R_R}{R_R + R_N}$

1. AVL & Splay tree

The height of an empty tree is defined to be -1.

The balance factor $BF(node) = h_L - h_R$. In an AVL tree, $BF(node) = -1, 0, \text{ or } 1$.



Other rotations are symmetric

n_h be the minimum number of nodes in the height balanced tree of height h .

$$n_h = n_{h-1} + n_{h-2} + 1$$

$$n_h = F_{h+2} - 1 \quad F_i \approx \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^i \Rightarrow n_h \approx \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^{h+2} - 1 \Rightarrow h = O(\ln n)$$

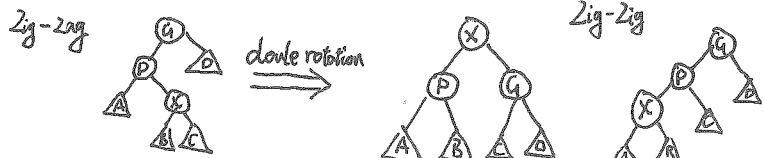
splay tree:

any M consecutive tree operations starting from an empty tree take at most $O(M \log N)$ time.

After a node is accessed, it is pushed to the root by a series of AVL tree rotations.

Case 1: P is the root \rightarrow Rotate X and P

Case 2: P is not the root



deletions:

Step 1: Find X;

Step 2: Remove X;

Step 3: FindMax(T_L);

Step 4: Make T_R the right child of the root of T_L

Delete: case 1: x is the point to be delete

case 2: continue to the path of x

case 3: add 1 black to x

case 4: Retrieved

Relevant R_R Irrelevant I_R
Not Retrieved R_N Irrelevant I_N

Precision $P = \frac{R_R}{R_R + I_R}$

Recall $R = \frac{R_R}{R_R + R_N}$

how fast does it index
how fast does it search
expressiveness of query language