CMPT 295

Unit - Machine-Level Programming

Lecture 20 – Assembly language – Array – 1D

Recursive Function - countOnesR(...)

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

What does this function do?

Recursive Function – Example – Base Case

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl %eax, %eax
 testq
        %rdi, %rdi
 je
         done
         %rbx
 pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call countOnesR
 addq
         %rbx, %rax
         %rbx
 popq
done:
 ret
```

Recursive Function – Example - Saving registers

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl
        %eax, %eax
         %rdi, %rdi
 testq
 jе
         done
         %rbx
 pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call countOnesR
 addq
         %rbx, %rax
         %rbx
  popq
done:
  ret
```

Recursive Function – Example - Call Setup

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl
        %eax, %eax
         %rdi, %rdi
 testq
 je
      done
         %rbx
  pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call countOnesR
 addq
         %rbx, %rax
         %rbx
 popq
done:
  ret
```

Recursive Function – Example – Recursive Call

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl
        %eax, %eax
         %rdi, %rdi
 testq
 je
         done
         %rbx
 pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call
         countOnesR
 addq
         %rbx, %rax
         %rbx
 popq
done:
  ret
```

Recursive Function – Example – Result

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl
        %eax, %eax
         %rdi, %rdi
 testq
 je
       done
        %rbx
  pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call countOnesR
 addq
         %rbx, %rax
         %rbx
 popq
done:
 ret
```

Recursive Function – Example – Clean-up and return

```
/* Recursive counter of 1's */
long countOnesR(unsigned long x) {
  if (x == 0)
    return 0;
  else
    return (x & 1) + countOnesR(x >> 1);
}
```

```
countOnesR:
 xorl
        %eax, %eax
        %rdi, %rdi
 testq
 jе
      done
         %rbx
 pushq
         %rdi, %rbx
 movq
 andl
         $1, %ebx
 shrq
         %rdi
 call countOnesR
 addq
         %rbx, %rax
         %rbx
 popq
done:
 ret
```

Recursive Function – Example – Test Cases

Test Case 1

Input: x = 0

Expected result: 0

Test Case 2

Input: x = 7

Expected result: 3

base + displacement Stack Variables

countOnesR: xorl %eax, %eax %rdi, %rdi testq done jе %rbx pushq %rdi, %rbx movq andl **\$1,** %ebx %rdi shrq call countOnesR addq %rbx, %rax %rbx popq done: ret

Register Table:

Last Lecture - x86-64 "register saving" convention

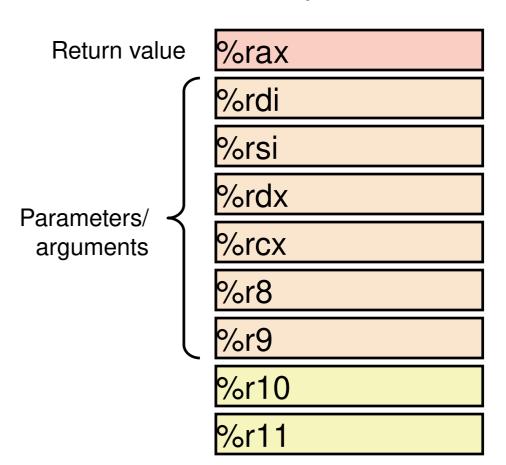
callee saved registers:

Callee must save & restore

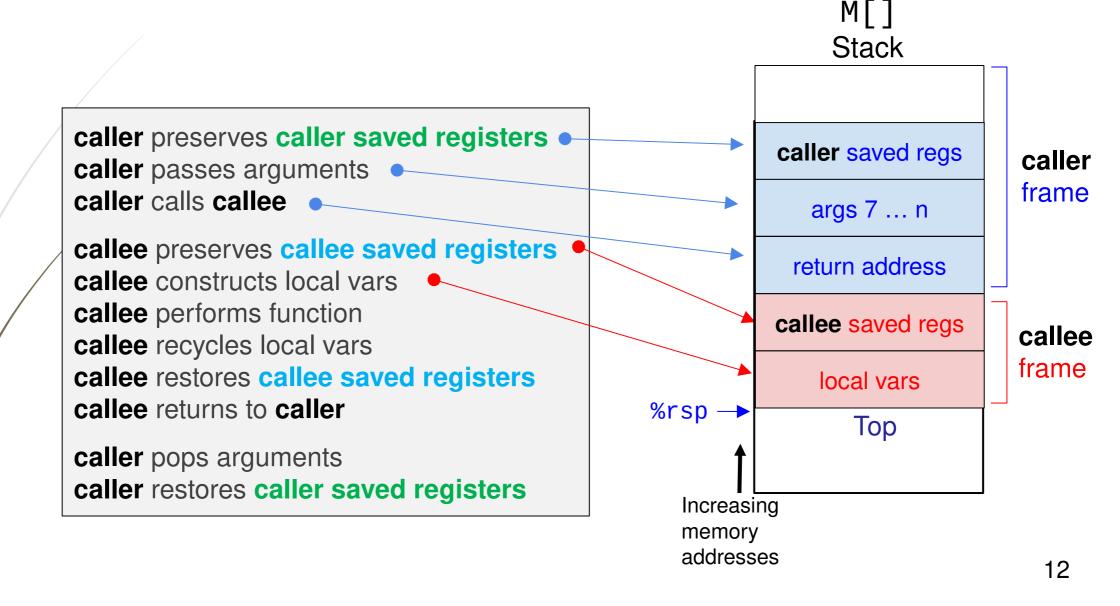
%rbx %r12 %r13 %r14 %r15 %rbp %rsp

caller saved registers:

- Caller must save & restore
- Can be modified by callee



Last Lecture - x86-64 conventions and stack frame



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Last Lecture

- Recursion
 - Handled without special consideration using ...
 - Stack frames
 - 1 x86-64 Function call and Register saving conventions

Today's Menu

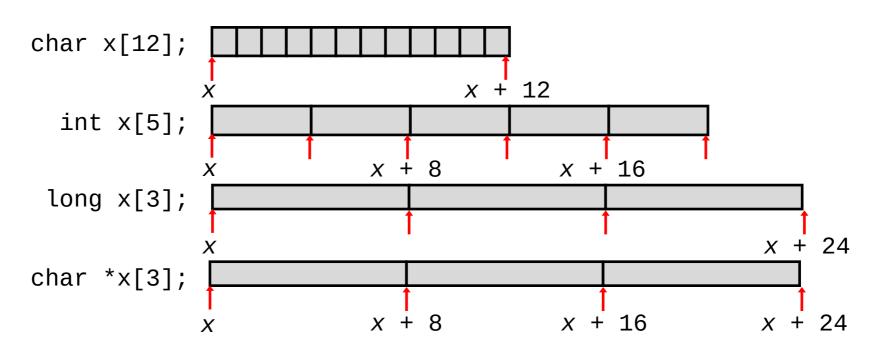
- Introduction
 - C program -> assembly code -> machine level code
- Assembly language basics: data, move operation
 - Memory addressing modes
- Operation leaq and Arithmetic & logical operations
- ☐ Conditional Statement Condition Code + cmovX
- Loops
- Function call Stack
 - Overview of Function Call
 - Memory Layout and Stack x86-64 instructions and registers
 - Passing control
 - Passing data Calling Conventions
 - Managing local data
 - Recursion
- Array
- Buffer Overflow
- Floating-point operations



1D Array

T **A**[*n*];

- Array of data type T and length n
- Contiguously allocated region of n * L bytes in memory where L = **sizeof**(*T*)



Accessing 1D Array

- Address of A[i] = base address + i * L
- A can be used as a pointer to array element 0
- \Box Can increment a pointer **A** by adding L to the address

```
_____C Type Value
    x [4] int
    x int *
    x + 1 int *
    &x[2] nt *
    x[5] nt
    *(x+1) int
    x + i int *
```

Manipulating 1D array – Example - main.c

```
#include <stdio.h>
                                          This program defines 4 arrays:
#include <stdlib.h>

    an array of char's,

    an array of short's,

char sumChar(char *, int);
                                           an array of int's,
short sumShort(short *, int);

    an array of long's

int sumInt(int *, int);
                                          then it calls the appropriate sum function, i.e., the one
long sumLong(long *, int);
                                          that sums elements of the correct data type.
#define N 6
                      Test cases
                                                           Expected results
char AC[N] = \{-58, 22, 101, -15, 72, 27\}; // Expected results: sum = -107
short AS[N] = \{-58, 22, 101, -15, 72, 27\}; // Expected results: sum = 149
int AI[N] = \{258, 522, 1010, -15, -3372, 27\}; // Expected results: sum = -1570
long AL[N] = \{258, 522, 1010, -15, -3372, 27\}; // Expected results: sum = -1570
void main () {
    printf("The total of AC is %hhi.\n", sumChar(AC, N));
    printf("The total of AS is %hi.\n", sumShort(AS, N));
    printf("The total of AI is %d.\n", sumInt(AI, N));
    printf("The total of AL is %ld.\n", sumLong(AL, N));
    return;
```

Manipulating 1D array – Example - sums.s - Part 1

- Register %rdi contains starting address of array (base address of array)
- Register %esi contains size of array (N)
- Register %ecx contains array index
- Register %al or %ax contains the running sum

```
.qlobl
          sumChar
sumChar:
   movl $0, %eax
          $0,
              %ecx
   movl
loopChar:
          %ecx, %esi
   cmpl
   jle
          endloop1
   addb (%rdi,%rcx,1), %al
   incl
          %ecx
          loopChar
   jmp
endloop1:
   ret
```

```
.globl
           sumShort
sumShort:
   xorl %eax, %eax
   xorl %ecx, %ecx
           cond1
   jmp
loopShort:
   addw
           (%rdi, %rcx, 2), %ax
   incl
           %ecx
cond1:
   cmpl %esi, %ecx
   jne
           loopShort
   ret
```

Manipulating 1D array – Example - sums.s - Part 2

- Register %rdi contains starting address of array (base address of array)
- Register %esi contains size of array (N)
- Register %ecx contains array index
- Register %eax or %rax contains the running sum

```
.qlobl
            sumInt
sumInt:
    xorl %eax, %eax
    xorl %ecx, %ecx
            cond2
    jmp
loopInt:
           (%rdi, %rcx, 4), %eax
    addl
    incl
            %ecx
cond2:
           %esi, %ecx
    cmpl
    jne
            loopInt
    ret
```

```
.qlobl
           sumLong
sumLong:
           $0,
                  %rax
   movq
           $0,
   movl
                   %ecx
loopLong:
           %ecx, %esi
   cmpl
   jle
           endloop
   addq
          (%rdi,%rcx,8), %rax
   incl
           %ecx
           loopLong
   jmp
endloop:
   ret
```



Forgetting that the memory addresses of contiguous array cells differ by the size of the content of these cells instead of by 1 is a common mistake in assembly language programming

Summary

- Arrays
 - Elements packed into contiguous region of memory
 - Use index arithmetic to locate individual elements
 - 1 1D array: address of A[i] = A + i * L

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