Lecture 3: Parsing

CS4200 Compiler Construction

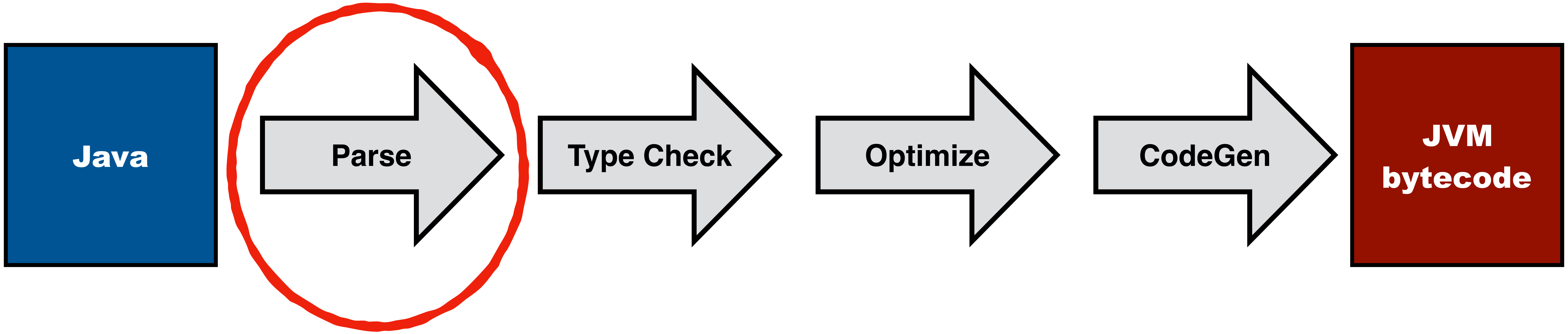
Eelco Visser

TU Delft

September 2019

TU Delft

This Lecture



Turning syntax definitions into parsers

This Lecture

Derivations

- Generating sentences and trees from context-free grammars

Grammar Transformations

- Eliminating ambiguity by transformation

Reductions

- From sentences to symbols

Parsing

- Shift/reduce parsing
- SLR parsing

Reading Material

These slides are the primary material for study of parsing in this course.

This provides a basic idea of parsing by studying derivations and one particular parsing algorithm.

It is by no means exhaustive.

There is a vast literature on parsing, and we encourage the interested student to read on.

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CS4200 Compiler Construction

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Classical compiler textbook

Chapter 4: Syntax Analysis

The material covered in this lecture corresponds to Sections 4.1, 4.2, 4.3, 4.5, 4.6

Pictures in these slides are copies from the book

Compilers: Principles, Techniques, and Tools, 2nd Edition

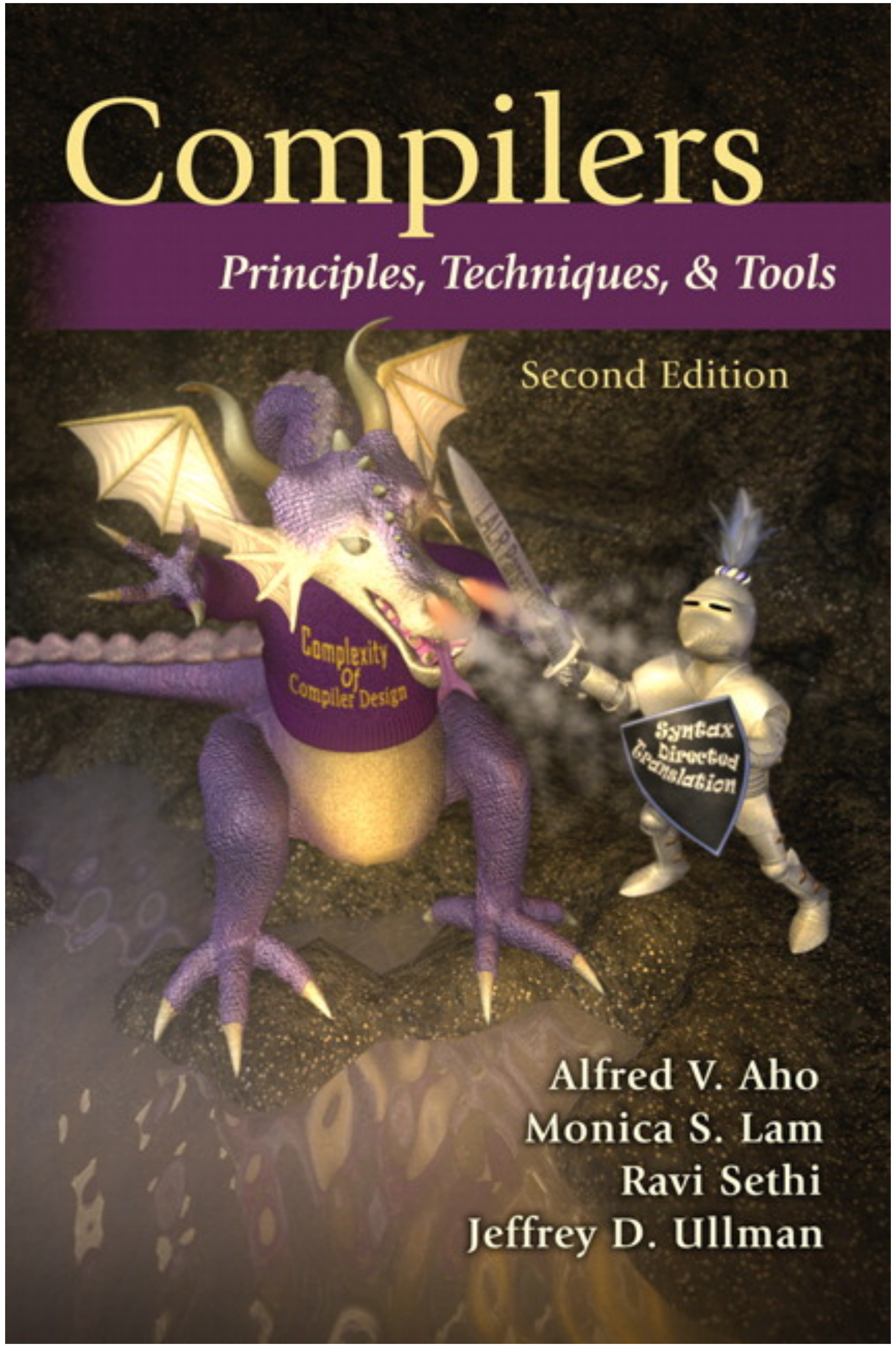
Alfred V. Aho, Columbia University

Monica S. Lam, Stanford University

Ravi Sethi, Avaya Labs

Jeffrey D. Ullman, Stanford University

2007 | Pearson



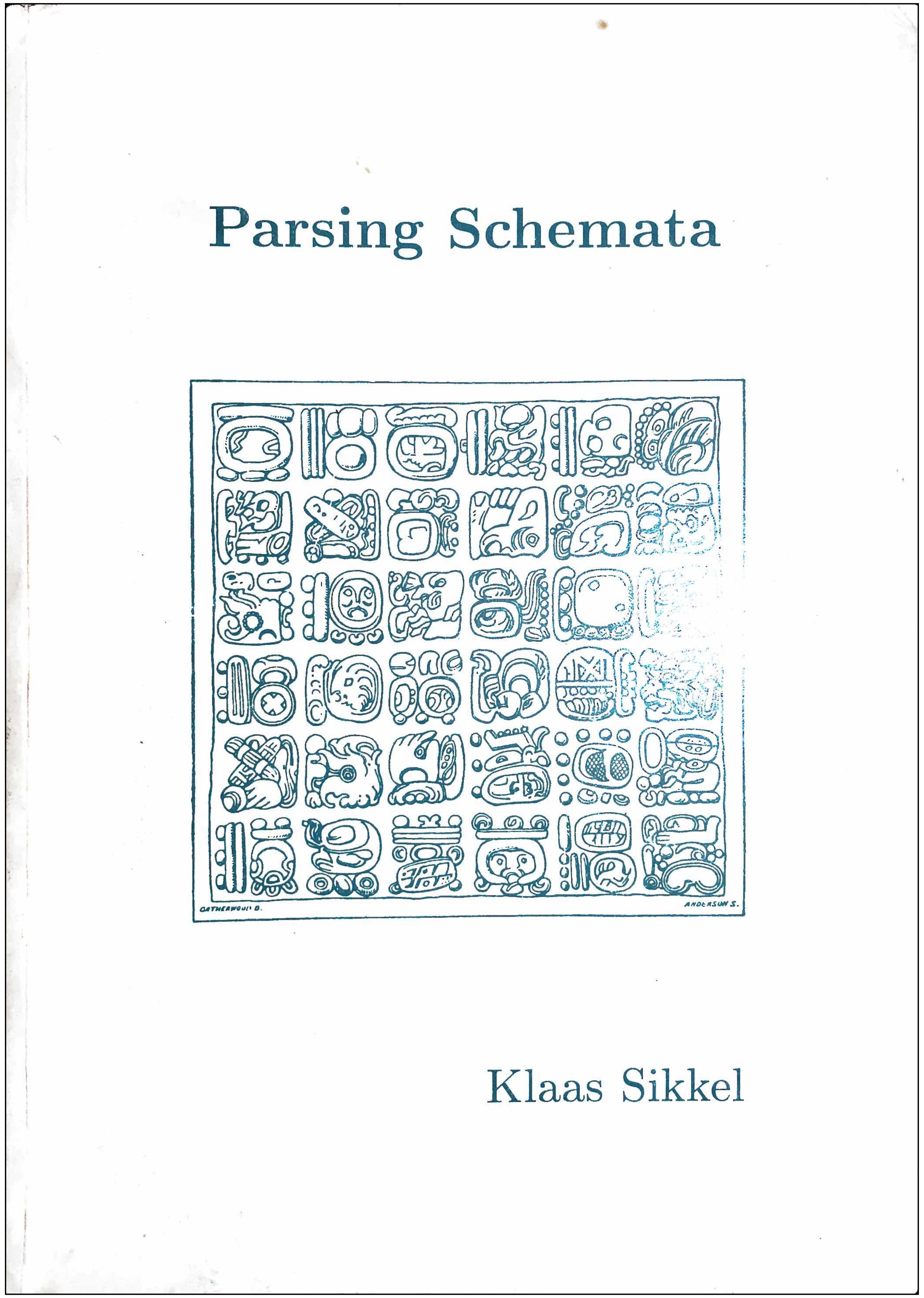
This PhD thesis presents a uniform framework for describing a wide range of parsing algorithms.

“Parsing schemata provide a general framework for specication, analysis and comparison of (sequential and/or parallel) parsing algorithms. A grammar specifies implicitly what the valid parses of a sentence are; a parsing algorithm specifies Elicitly how to compute these. Parsing schemata form a well-defined level of abstraction in between grammars and parsing algorithms. A parsing schema specifies the types of intermediate results that can be computed by a parser, and the rules that allow to Eand a given set of such results with new results. A parsing schema does not specify the data structures, control structures, and (in case of parallel processing) communication structures that are to be used by a parser.”

For the interested

Sikkel, N. (1993). *Parsing Schemata.*
PhD thesis. Enschede: Universiteit Twente

https://research.utwente.nl/en/publications/parsing-schemata



This paper applies parsing schemata to disambiguation filters for priority conflicts.

For the interested

https://ivi.fnwi.uva.nl/tcs/pub/reports/1995/P9507.ps.Z

A Case Study in Optimizing Parsing Schemata by Disambiguation Filters

Eelco Visser

Abstract

Disambiguation methods for context-free grammars enable concise specification of programming languages by ambiguous grammars. A disambiguation filter is a function that selects a subset from a set of parse trees—the possible parse trees for an ambiguous sentence. The framework of filters provides a declarative description of disambiguation methods independent of parsing. Although filters can be implemented straightforwardly as functions that prune the parse forest produced by some generalized parser, this can be too inefficient for practical applications.

In this paper the optimization of parsing schemata, a framework for high-level description of parsing algorithms, by disambiguation filters is considered in order to find efficient parsing algorithms for declaratively specified disambiguation methods. As a case study the optimization of the parsing schema of Earley's parsing algorithm by two filters is investigated. The main result is a technique for generation of efficient LR-like parsers for ambiguous grammars modulo priorities.

1 Introduction

The syntax of programming languages is conventionally described by context-free grammars. Although programming languages should be unambiguous, they are often described by ambiguous grammars because these allow a more natural formulation and yield better abstract syntax. For instance, the grammar

$$E \to E + E; E \to E * E; E \to a$$

gives a clearer description of arithmetic expressions than the grammar

$$E \to E + T; E \to T; T \to T * a; T \to a$$

To obtain an unambiguous specification of a language described by an ambiguous grammar it has to be disambiguated. For example, the grammar above can be disambiguated by associativity and priority rules that express that $E \to E * E$ has higher priority than $E \to E + E$ and that both productions are left associative.

Technical Report P9507, Programming Research Group, University of Amsterdam, July 1995.
Available as: ftp://ftp.fwi.uva.nl/pub/programming-research/reports/1995/P9507.ps.Z
This is an extended version of [Vis95].

This paper defines a declarative semantics for associativity and priority declarations for disambiguation.

The paper provides a safe semantics and extends it to deep priority conflicts.

The result of disambiguation is a contextual grammar, which generalises the disambiguation for the dangling-else grammar.

For the interested

The paper is still in production. Ask me for a copy of the draft.

A Direct Semantics for Declarative Disambiguation of Expression Grammars

LUIS EDUARDO S. AMORIM, Delft University of Technology
TIMOTHÉE HAUDEBOURG, Université Rennes, Inria, IRISA
MICHAEL J. STEINDORFER, Delft University of Technology
EELCO VISSER, Delft University of Technology

Associativity and priority rules provide a mechanism for disambiguation of context-free grammars that is based on concepts familiar to programmers from high-school mathematics. However, there is no standardized and declarative semantics for such rules, in particular for the disambiguation of the more complex expression grammars encountered in programming languages.

In this paper, we provide a *direct* semantics of disambiguation of *expression grammars* by means of associativity and priority declarations by defining the subtree exclusion patterns corresponding to each declaration. We introduce *deep priority conflicts* to solve ambiguities that cannot be captured by fixed-depth patterns such as ambiguities due to low priority prefix or postfix operators, dangling prefix, dangling suffix, longest match, and indirect recursion. We prove that the semantics is *safe* (preserves the language of the underlying context-free grammar) and *complete* (solves all ambiguities) for classes of expression grammars of increasing complexity that have only *harmless overlap*, provided a total set of disambiguation rules is provided.

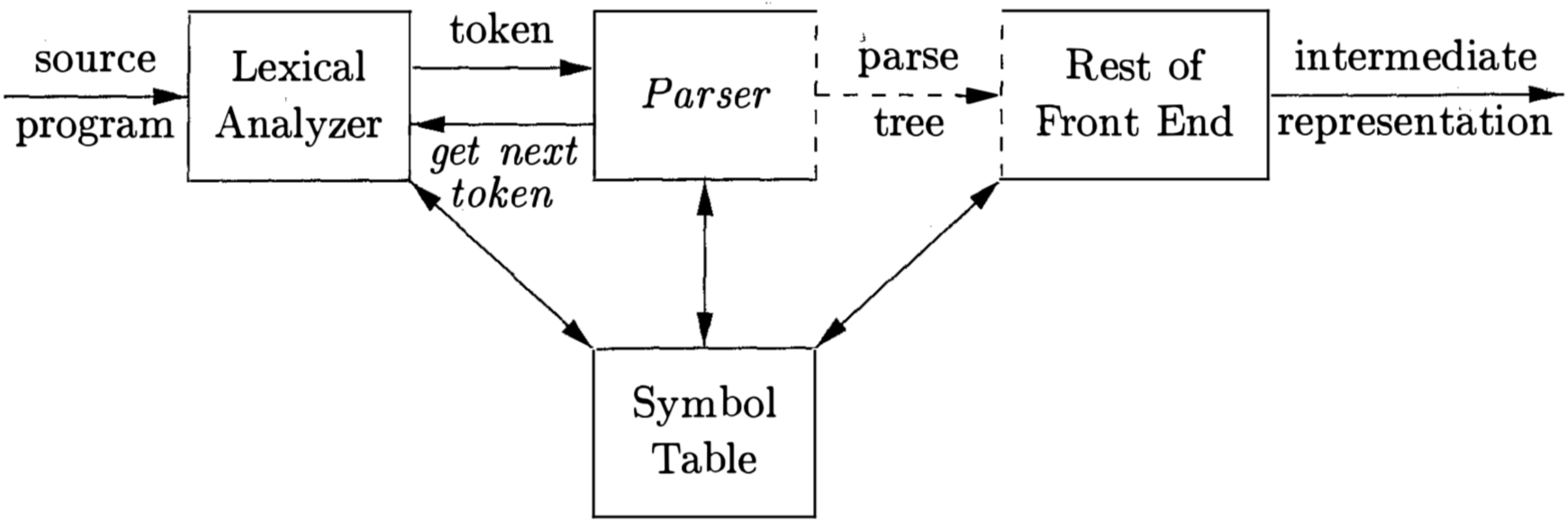
We define a canonical implementation that transforms a grammar with disambiguation rules into a disambiguated grammar. We also describe alternative implementations that interpret disambiguation rules during parser generation or during parsing. We have evaluated the approach by applying it to the grammars of five languages. Finally, we demonstrate that our approach is practical by measuring the performance of a parser that implements our disambiguation techniques, applying it to a corpus of real-world Java and OCaml programs.

CCS Concepts: •**Software and its engineering** → **Syntax; Parsers;**

ACM Reference format:
Luis Eduardo S. Amorim, Timothée Haudebourg, Michael J. Steindorfer, and Eelco Visser. 2019. A Direct Semantics for Declarative Disambiguation of Expression Grammars. *ACM Trans. Program. Lang. Syst.* 1, 1, Article 1 (February 2019), 80 pages.
DOI: 10.1145/nnnnnnn.nnnnnnn

Parser Architecture

Traditional Parser Architecture



Source: Compilers Principles, Techniques & Tools

Context-Free Grammars

Context-Free Grammars

Terminals

- Basic symbols from which strings are formed

Nonterminals

- Syntactic variables that denote sets of strings

Start Symbol

- Denotes the nonterminal that generates strings of the languages

Productions

- A = X … X
- Head/left side (A) is a nonterminal
- Body/right side (X … X) zero or more terminals and nonterminals

Example Context-Free Grammar

```
grammar
  start S
  non-terminals E T F
  terminals "+" "*" "(" ")" ID
  productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

Abbreviated Grammar

```
grammar
  start S
  non-terminals E T F
  terminals "+" "*" "(" ")" ID
  productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
grammar
  productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

Nonterminals, terminals can be derived from productions

First production defines start symbol

Notation

```
A, B, C: non-terminals
l: terminals
a, b, c: strings of non-terminals and terminals
    (alpha, beta, gamma in math)
w, v: strings of terminal symbols
```

Meta: Syntax of Grammars

```
context-free syntax // grammars

  Grammar.Grammar = <
    grammar
      <Start?>
      <Sorts?>
      <Terminals?>
      <Productions>
  >
```

```
context-free syntax

  Production.Prod = <
    <Symbol><Constructor?> = <Symbol*>
  >

  Symbol.NT = <<ID>>
  Symbol.T  = <<STRING>>
  Symbol.L  = <<LCID>>

  Constructor.Con = <.<ID>>
```

```
context-free syntax

  Start.Start = <
    start <ID>
  >

  Sorts.Sorts = <
    sorts <ID*>
  >

  Sorts.NonTerminals = <
    non-terminals <ID*>
  >

  Terminals.Terminals = <
    terminals <Symbol*>
  >

  Productions.Productions = <
    productions
      <Production*>
  >
```

Derivations: Generating Sentences from Symbols

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Derivations

```
grammar
  productions
    E = E "+" E
    E = E "*" E
    E = "-" E
    E = "(" E ")"
    E = ID
```

```
// derivation step: replace symbol by rhs of production
// E = E "+" E
// replace E by E "+" E
//
// derivation:
// repeatedly apply derivations
```

```
derivation
  E
  => "-" E
  => "-" "(" E ")"
  => "-" "(" ID ")"
```

```
derivation // derives in zero or more steps
  E =>* "-" "(" ID "+" ID ")"
```

Meta: Syntax of Derivations

```
context-free syntax // derivations

  Derivation.Derivation = <
    derivation
      <Symbol> <Step*>
  >

  Step.Step   = [=> [Symbol*]]
  Step.Steps  = [=>* [Symbol*]]
  Step.Steps1 = [=>+ [Symbol*]]
```

Left-Most Derivation

```
grammar
  productions
    E = E "+" E
    E = E "*" E
    E = "-" E
    E = "(" E ")"
    E = ID
```

```
derivation // left-most derivation
  E
  => "-" E
  => "-" "(" E ")"
  => "-" "(" E "+" E ")"
  => "-" "(" ID "+" E ")"
  => "-" "(" ID "+" ID ")"
```

Left-most derivation: Expand left-most non-terminal at each step

Right-Most Derivation

```
grammar
  productions
    E = E "+" E
    E = E "*" E
    E = "-" E
    E = "(" E ")"
    E = ID
```

```
derivation // left-most derivation
  E
  => "-" E
  => "-" "(" E ")"
  => "-" "(" E "+" E ")"
  => "-" "(" ID "+" E ")"
  => "-" "(" ID "+" ID ")"
```

```
derivation // right-most derivation
  E
  => "-" E
  => "-" "(" E ")"
  => "-" "(" E "+" E ")"
  => "-" "(" E "+" ID ")"
  => "-" "(" ID "+" ID ")"
```

Right-most derivation: Expand right-most non-terminal at each step

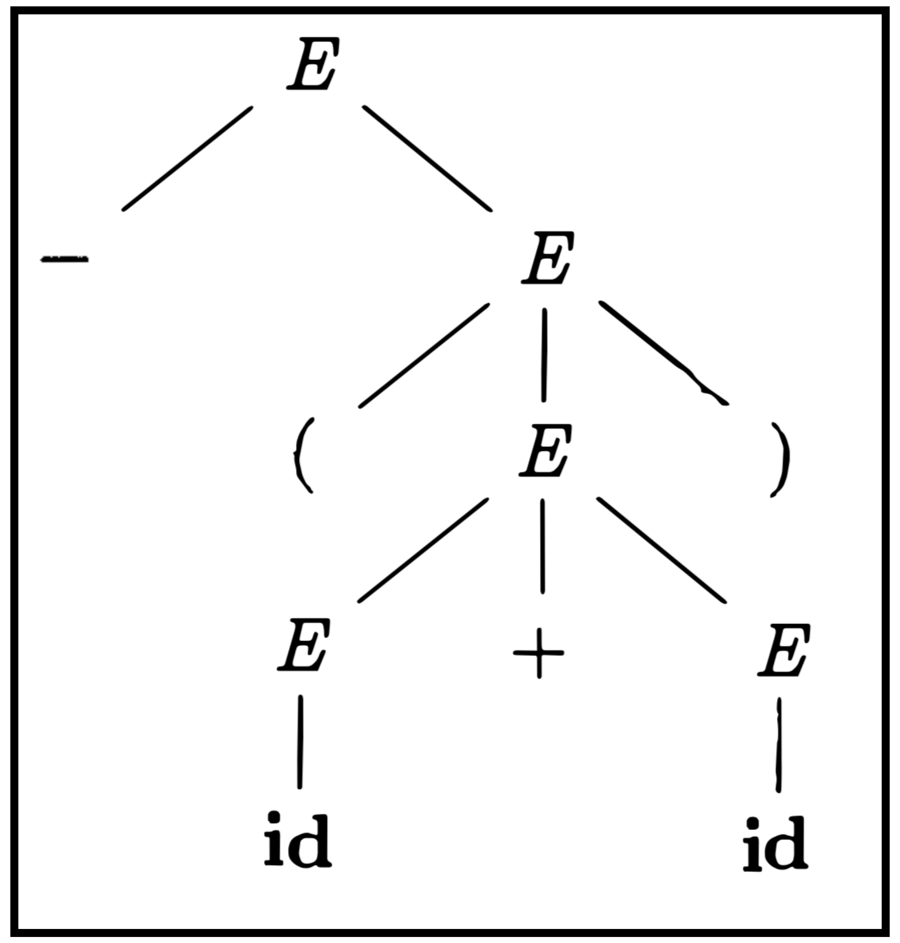
Meta: Tree Derivations

```
context-free syntax // tree derivations

  Derivation.TreeDerivation = <
    tree derivation
      <Symbol> <PStep*>
  >

  PStep.Step   = [=> [PT*]]
  PStep.Steps  = [=>* [PT*]]
  PStep.Steps1 = [=>+ [PT*]]

  PT.App = <<Symbol>[<PT*>]>
  PT.Str = <<STRING>>
  PT.Sym = <<Symbol>>
```

Left-Most Tree Derivation

```
grammar
  productions
    E.A = E "+" E
    E.T = E "*" E
    E.N = "-" E
    E.P = "(" E ")"
    E.V = ID
```

```
derivation // left-most derivation
  E
  => "-" E
  => "-" "(" E ")"
  => "-" "(" E "+" E ")"
  => "-" "(" ID "+" E ")"
  => "-" "(" ID "+" ID ")"
```




```
tree derivation // left-most
  E
  => E["-" E]
  => E["-" E["(" E ")"]]
  => E["-" E["(" E[E "+" E] ")"]]
  => E["-" E["(" E[E[ID] "+" E] ")"]]
  => E["-" E["(" E[E[ID] "+" E[ID]] ")"]]
```

Left-Most Tree Derivation


```
tree derivation // left-most
  E
  => E["-" E]
  => E["-" E["(" E ")"]]
  => E["-" E["(" E[E "+" E] ")"]]
  => E["-" E["(" E[E[ID] "+" E] ")"]]
  => E["-" E["(" E[E[ID] "+" E[ID]] ")"]]
```

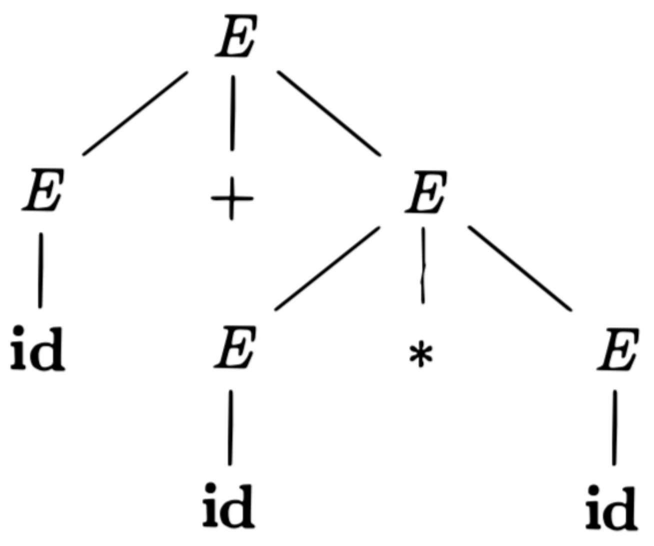
Ambiguity: Deriving Multiple Parse Trees

```
grammar
  productions
    E.A = E "+" E
    E.T = E "*" E
    E.N = "-" E
    E.P = "(" E ")"
    E.V = ID
```

```
derivation
  E =>* ID "+" ID "*" ID
```

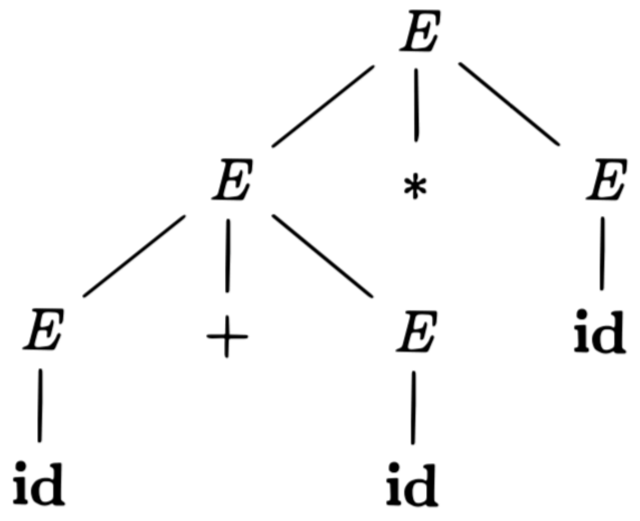
```
derivation
  E
  => E "+" E
  => ID "+" E
  => ID "+" E "*" E
  => ID "+" ID "*" E
  => ID "+" ID "*" ID
```

```
tree derivation
  E =>* E[E[ID] "+" E[E[ID] "*" E[ID]]]
```




```
derivation
  E
  => E "*" E
  => E "+" E "*" E
  => ID "+" E "*" E
  => ID "+" ID "*" E
  => ID "+" ID "*" ID
```

```
tree derivation
  E =>* E[E[E[ID] "+" E[ID]] "*" E[ID]]
```



Ambiguous grammar: produces >1 parse tree for a sentence

Meta: Term Derivations

```
context-free syntax // term derivations

  Derivation.TermDerivation = <
    term derivation
      <Symbol> <TStep*>
  >

  TStep.Step   = [=> [Term*]]
  TStep.Steps  = [=>* [Term*]]
  TStep.Steps1 = [=>+ [Term*]]

  Term.App = <<ID>(<{Term ","}*>)>
  Term.Str = <<STRING>>
  Term.Sym = <<Symbol>>
```

Ambiguity: Deriving Abstract Syntax Terms

```
grammar
  productions
    E.A = E "+" E
    E.T = E "*" E
    E.N = "-" E
    E.P = "(" E ")"
    E.V = ID
```

```
derivation
  E =>* ID "+" ID "*" ID
```

```
derivation
  E
  => E "+" E
  => ID "+" E
  => ID "+" E "*" E
  => ID "+" ID "*" E
  => ID "+" ID "*" ID
```

```
term derivation
  E
  => A(E, E)
  => A(V(ID), E)
  => A(V(ID), T(E, E))
  => A(V(ID), T(V(ID), E))
  => A(V(ID), T(V(ID), V(ID)))
```

```
derivation
  E
  => E "*" E
  => E "+" E "*" E
  => ID "+" E "*" E
  => ID "+" ID "*" E
  => ID "+" ID "*" ID
```

```
term derivation
  E
  => T(E, E)
  => T(A(E, E), E)
  => T(A(V(ID), E), E)
  => T(A(V(ID), V(ID)), E)
  => T(A(V(ID), V(ID)), V(ID))
```

Parse Trees Represent Derivations

```
List<String> YIELD(T : Tree) {
  T match {
    [A = Ts] => YIELDs(Ts);
    Str => [Str];
  };
}

List<String> YIELDs(Ts : List<Tree>) {
  Ts match {
    [] => "";
    [T | Ts] => YIELD(T) + YIELD(Ts);
  };
}
```

```
S =>* PT

iff

S =>* YIELD(PT)
```

Language Defined by a Grammar

$$\boxed{L(G) = \{\ w \mid S \Rightarrow^* w\ \}}$$

$$\boxed{T(G) = \{\ T \mid S \Rightarrow^* T\ \}}$$

$$\boxed{L(G) = YIELD(T(G))}$$

Grammar Transformations

Grammar Transformations

Why?

- Disambiguation
- For use by a particular parsing algorithm

Transformations

- Eliminating ambiguities
- Eliminating left recursion
- Left factoring

Properties

- Does transformation preserve the language (set of strings, trees)?
- Does transformation preserve the structure of trees?

Ambiguous Expression Grammar

```
grammar
  productions
    E.A = E "+" E
    E.T = E "*" E
    E.M = "-" E
    E.B = "(" E ")"
    E.V = ID
```

```
derivation
  E =>* ID "*" ID "+" ID
```

```
term derivation
  E
  => A(E, E)
  => A(T(E, E), E)
  => A(T(E, E), E)
  => A(T(V(ID), E), E)
  => A(T(V(ID), V(ID)), E)
  => A(T(V(ID), V(ID)), V(ID))
```

```
term derivation
  E
  => T(E, E)
  => T(E, E)
  => T(V(ID), E)
  => T(V(ID), A(E, E))
  => T(V(ID), A(V(ID), E))
  => T(V(ID), A(V(ID), V(ID)))
```

Associativity and Priority Filter Ambiguities

```
grammar
  productions
    E.A = E "+" E
    E.T = E "*" E
    E.M = "-" E
    E.B = "(" E ")"
    E.V = ID
```

```
grammar
  productions
    E.A = E "+" E {left}
    E.T = E "*" E {left}
    E.M = "-" E
    E.B = "(" E ")"
    E.V = ID
  priorities
    E.M > E. T > E.A
```

```
derivation
  E =>* ID "*" ID "+" ID
```

```
term derivation
  E
  => A(E, E)
  => A(T(E, E), E)
  => A(T(E, E), E)
  => A(T(V(ID), E), E)
  => A(T(V(ID), V(ID)), E)
  => A(T(V(ID), V(ID)), V(ID))
```

```
term derivation
  E
  => T(E, E)
  => T(E, E)
  => T(V(ID), E)
  => T(V(ID), A(E, E))
  => T(V(ID), A(V(ID), E))
  => T(V(ID), A(V(ID), V(ID)))
```

Define Associativity and Priority by Transformation

```
grammar
  productions
    E.A = E "+" E {left}
    E.T = E "*" E {left}
    E.M = "-" E
    E.B = "(" E ")"
    E.V = ID
  priorities
    E.M > E. T > E.A
```

```
grammar
  productions
    E.A = E "+" T
    E   = T
    T.T = T "*" F
    T   = F
    F.V = ID
    F.B = "(" E ")"
```

```
derivation
  E =>* ID "*" ID "+" ID
```

```
term derivation
  E
  => A(E, E)
  => A(T(E, E), E)
  => A(T(E, E), E)
  => A(T(V(ID), E), E)
  => A(T(V(ID), V(ID)), E)
  => A(T(V(ID), V(ID)), V(ID))
```

```
term derivation
  E
  => T(E, E)
  => T(E, E)
  => T(V(ID), E)
  => T(V(ID), A(E, E))
  => T(V(ID), A(V(ID), E))
  => T(V(ID), A(V(ID), V(ID)))
```

Define Associativity and Priority by Transformation

```
grammar
  productions
    E.A = E "+" E {left}
    E.T = E "*" E {left}
    E.M = "-" E
    E.B = "(" E ")"
    E.V = ID
  priorities
    E.M > E. T > E.A
```

```
grammar
  productions
    E.A = E "+" T
    E   = T
    T.T = T "*" F
    T   = F
    F.V = ID
    F.B = "(" E ")"
```

Define new non-terminal for each priority level:
E, T, F

Add ‘injection’ productions to include priority level n+1 in n:
E = T
T = F

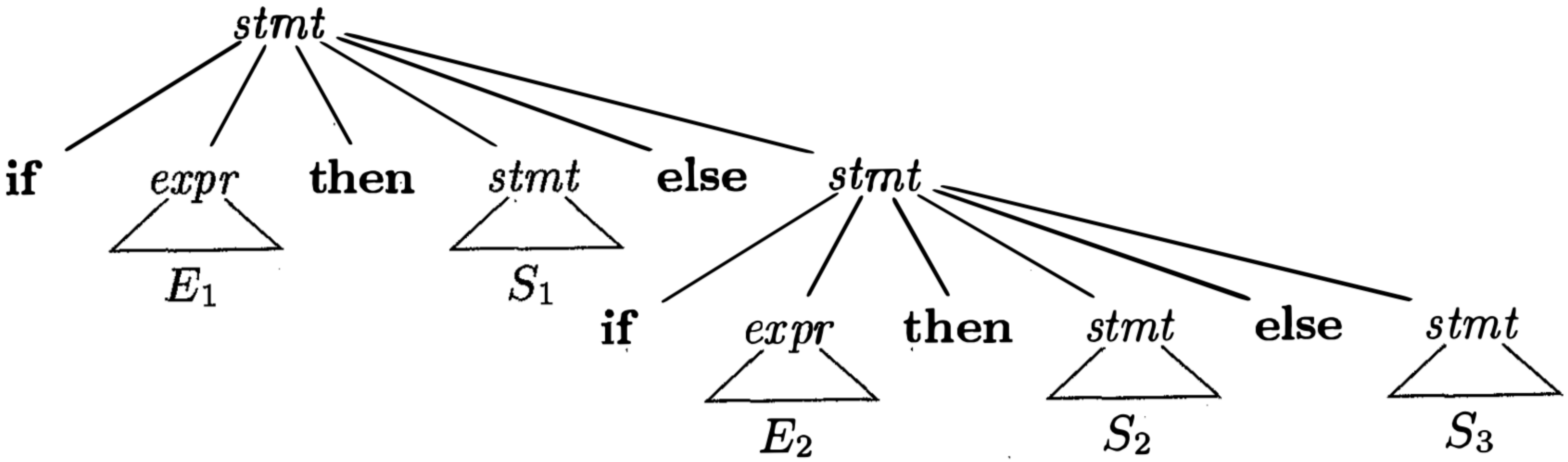
Change head of production to reflect priority level
T = T “*” F

Transform productions
Left: E = E “+” T
Right: E = T “+” E

Dangling Else Grammar

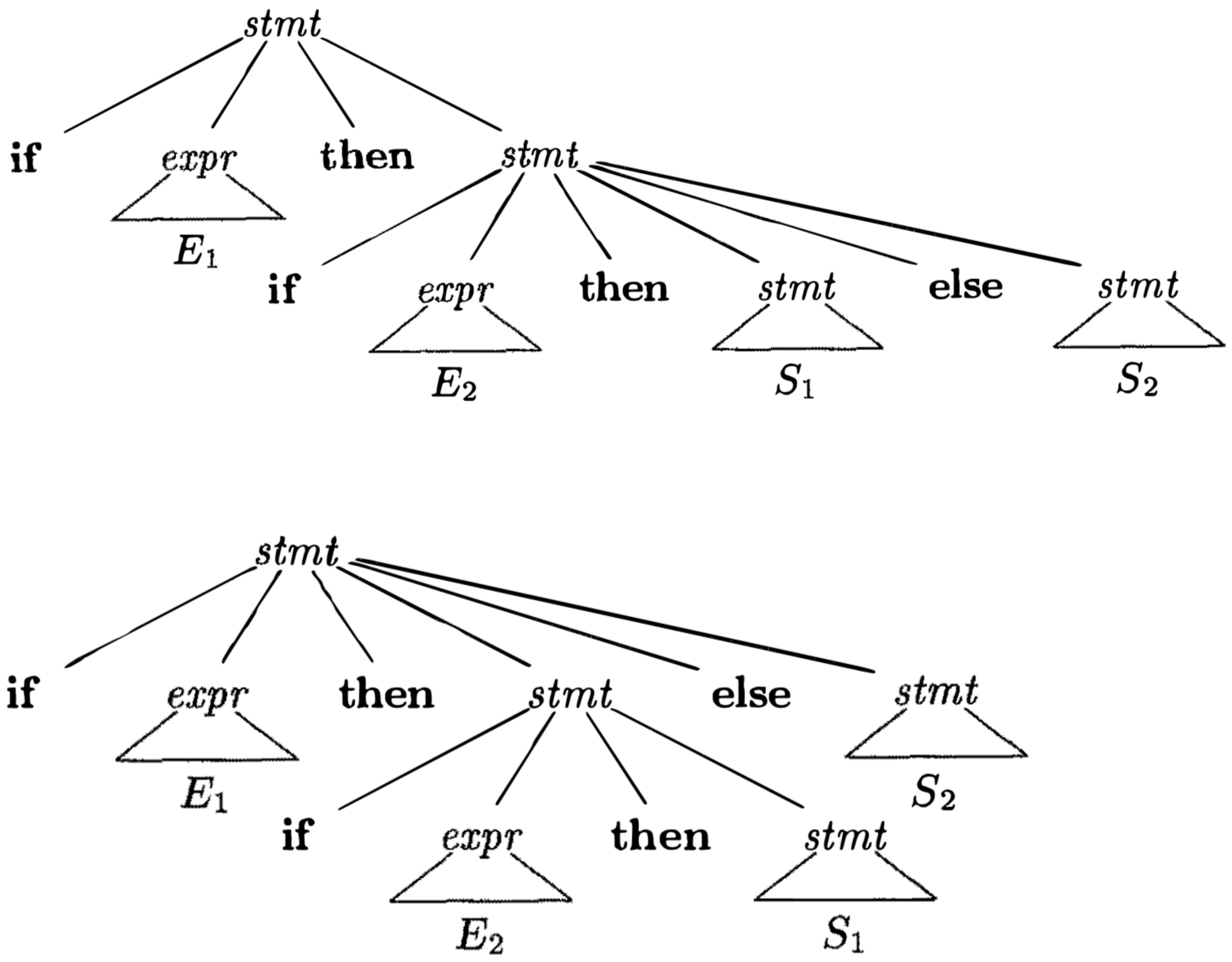
```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
derivation
  S =>* if E1 then S1 else if E2 then S2 else S3

term derivation
  S =>* IfE(E1, S1, IfE(E2, S2, S3))

term derivation
  S
  => IfE(E1, S1, S)
  => IfE(E1, S1, IfE(E2, S2, S3))
```



Dangling Else Grammar is Ambiguous

```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```




```
derivation
  S =>* if E1 then if E2 then S1 else S2
```

```
term derivation
  S
  => If(E1, S)
  => If(E1, IfE(E2, S1, S2))

derivation
  S
  => if E1 then S
  => if E1 then if E2 then S1 else S2
```

```
term derivation
  S
  => IfE(E1, S, S2)
  => IfE(E1, If(E2, S1), S2)

derivation
  S
  => if E1 then S else S2
  => if E1 then if E2 then S1 else S2
```

Eliminating Dangling Else Ambiguity

```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
grammar
  productions
    S.If    = if E then S
    S.IfE   = if E then SE else S
    S       = other
    SE.IfE = if E then SE else SE
    SE      = other
```

Generalization of this transformation: contextual grammars

Eliminating Left Recursion

```
grammar
  productions
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
grammar
  productions
    E  = T E'
    E' = "+" T E'
    E' =
    T  = F T'
    T' = "*" F T'
    T' =
    F  = "(" E ")"
    F  = ID
```

```
grammar
  productions
    A = A a
    A = b
```

```
// b followed by a list of as
```

```
grammar
  productions
    A  = b A'
    A' = a A'
    A' = // empty
```

Eliminating Left Recursion using Regular Expressions

```
grammar
  productions
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
grammar
  productions
    E  = T ("+" T)*
    T  = F ("*" F)*
    F  = "(" E ")"
    F  = ID
```

```
grammar
 productions
    A = A a
    A = b
```

```
grammar
  productions
    A = b a*
```

```
// b followed by a list of as
```

Left Factoring

```
grammar
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
grammar
  sorts S E
  productions
    S.If      = if E then S S'
    S'.Else   = else S
    S'.NoElse = // empty
    S         = other
```

```
grammar
  productions
    A = a b1
    A = a b2
    A = c
```

```
grammar
  productions
    A  = a A'
    A' = b1
    A' = b2
    A  = c
```

Properties of Grammar Transformations

Preservation

- Preserves set of sentences
- Preserves set of trees
- Preserves tree structure

Systematic

- Algorithmic
- Heuristic

Reducing Sentences to Symbols

Meta: Reductions

```
context-free syntax

  Reduction.Reduction = <
    reduction
      <Symbol*> <RStep*>
  >

  RStep.Step   = [<= [Symbol*]]
  RStep.Steps  = [<=* [Symbol*]]
  RStep.Steps1 = [<=+ [Symbol*]]
```

```
context-free syntax

  Reduction.TreeReduction = <
    tree reduction
      <PT*> <PRStep*>
  >

  PRStep.Step   = [<= [PT*]]
  PRStep.Steps  = [<=* [PT*]]
  PRStep.Steps1 = [<=+ [PT*]]
```

```
context-free syntax

  Reduction.TermReduction = <
    term reduction
      <Term*> <TRStep*>
  >

  TRStep.Step   = [<= [Term*]]
  TRStep.Steps  = [<=* [Term*]]
  TRStep.Steps1 = [<=+ [Term*]]
```

A Reduction is an Inverse Derivation

```
grammar
  sorts A
  productions
    A = b
```

```
reduction
  a b c <= a A c
```

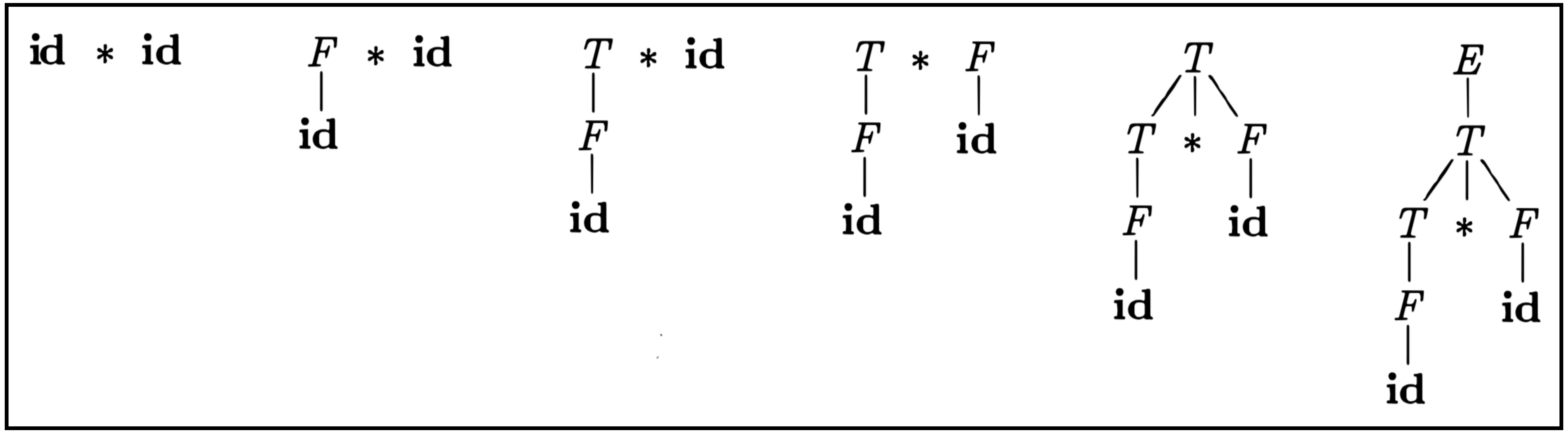
Reducing to Symbols

```
grammar
  sorts E T F ID
  productions
    E.P = E "+" T
    E.E = T
    T.M = T "*" F
    T.T = F
    F.B = "(" E ")"
    F.V = ID
```

```
reduction
  ID "*" ID
  <= F "*" ID
  <= T "*" ID
  <= T "*" F
  <= T
  <= E
```

Reducing to Parse Trees

```
grammar
  sorts E T F ID
  productions
    E.P = E "+" T
    E.E = T
    T.M = T "*" F
    T.T = F
    F.B = "(" E ")"
    F.V = ID
```

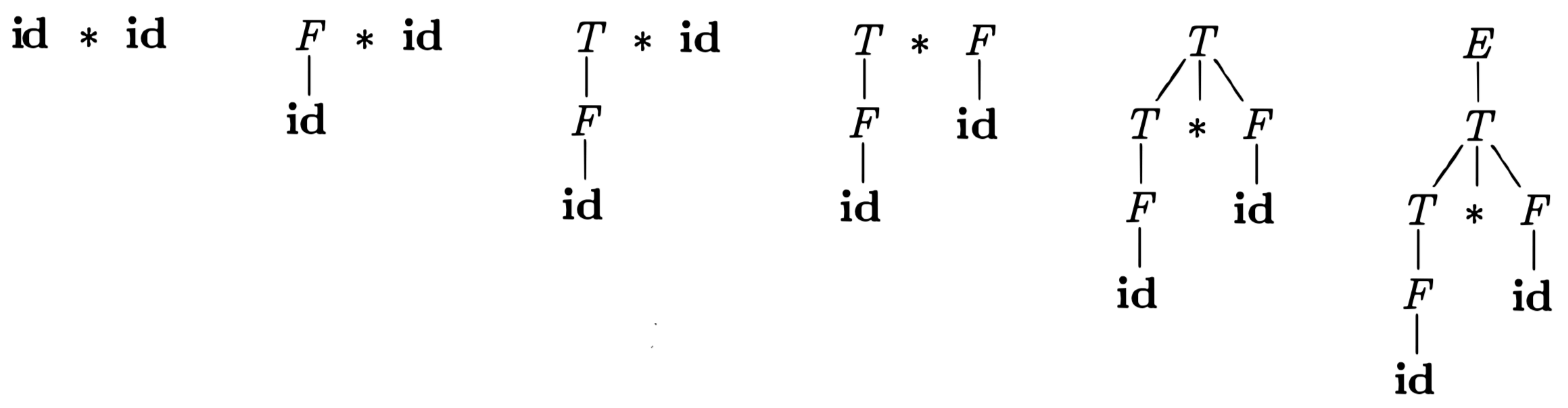



```
reduction
  ID "*" ID
  <= F "*" ID
  <= T "*" ID
  <= T "*" F
  <= T
  <= E
```

```
tree reduction
  ID "*" ID
  <= F[ID] "*" ID
  <= T[F[ID]] "*" ID
  <= T[F[ID]] "*" F[ID]
  <= T[T[F[ID]] "*" F[ID]]
  <= E[T[T[F[ID]] "*" F[ID]]]
```

Reducing to Abstract Syntax Terms

```
grammar
  sorts E T F ID
  productions
    E.P = E "+" T
    E.E = T
    T.M = T "*" F
    T.T = F
    F.B = "(" E ")"
    F.V = ID
```




```
reduction
  ID "*" ID
  <= F "*" ID
  <= T "*" ID
  <= T "*" F
  <= T
  <= E
```

```
tree reduction
  ID "*" ID
  <= F[ID] "*" ID
  <= T[F[ID]] "*" ID
  <= T[F[ID]] "*" F[ID]
  <= T[T[F[ID]] "*" F[ID]]
  <= E[T[T[F[ID]] "*" F[ID]]]
```

```
term reduction
  ID "*" ID
  <= V(ID) "*" ID
  <= T(V(ID)) "*" ID
  <= T(V(ID)) "*" V(ID)
  <= M(T(V(ID)), V(ID))
  <= E(M(T(V(ID)), V(ID)))
```

Handles

```
grammar
  productions
    A = b
```

```
derivation // right-most derivation
  S =>* a A w =>  a b w
```

```
reduction
  a b w <= a A w <=* S
```

```
// Handle

// sentential form: string of non-terminal and terminal symbols
// that can be derived from start symbol

// S =>* a

// right sentential form: a sentential form derived by a right-most derivation

// handle: the part of a right sentential form that if reduced
// would produce the previous right-sential form in a right-most derivation
```

Shift-Reduce Parsing

TU Delft

Shift-Reduce Parsing Machine

```
shift-reduce parse
  $ stack | input $ action
```

```
     $ a   | l w $ shift
  => $ a l |   w $
  // shift input symbol on the stack
```

```
    $ a b | w $ reduce by A = b
 => $ a A | w $
  // reduce n symbols of the stack to symbol A
```

```
$ S | $ accept
// reduced to start symbol; accept
```

```
 $ a | w $ error
  // no action possible in this state
```

Shift-Reduce Parsing

```
grammar
  productions
    E.P = E "+" T
    E.E = T
    T.M = T "*" F
    T.T = F
    F.B = "(" E ")"
    F.V = ID
```

```
shift-reduce parse
     $           | ID "*" ID $ shift
  => $ ID        |    "*" ID $ reduce by F = ID
  => $ F         |    "*" ID $ reduce by T = F
  => $ T         |    "*" ID $ shift
  => $ T "*"     |        ID $ shift
  => $ T "*" ID  |           $ reduce by F = ID
  => $ T "*" F   |           $ reduce by T = T "*" F
  => $ T         |           $ reduce by E = T
  => $ E         |           $ accept
```

Shift-reduce parsing constructs a right-most derivation

Shift-Reduce Conflicts

```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
shift-reduce parse
      $                    | if E then S else S $ shift
  =>  $ if                 |    E then S else S $ shift
  =>* $ if E then S        |             else S $ shift
  =>  $ if E then S else   |                  S $ shift
  =>  $ if E then S else S |                    $ reduce by S = if E then S else S
  =>  $ S                  |                    $ accept
```

```
shift-reduce parse
      $                    | if E then S else S $ shift
  =>  $ if                 |    E then S else S $ shift
  =>* $ if E then S        |             else S $ reduce by S = if E then S
  =>  $ S                  |             else S $ error
```

Shift-Reduce Conflicts

```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
shift-reduce parse
     $                                | if E then if E then S else S $ ...
  =>* $ if E then if E then S         |                      else S $ shift
  =>  $ if E then if E then S else    |                           S $ shift
  =>  $ if E then if E then S else S  |                             $ reduce by S = if E then S else S
  =>  $ if E then S                   |                             $ reduce by S = if E then S
  =>  $ S                             |                             $ accept
```

```
shift-reduce parse
     $                                | if E then if E then S else S $ ...
  =>* $ if E then if E then S         |                      else S $ reduce by S = if E then S
  =>  $ if E then S                   |                      else S $ shift
  =>  $ if E then S else              |                           S $ shift
  =>  $ if E then S else S            |                             $ reduce by S = if E then S else S
  =>  $ S
```

Shift-Reduce Conflicts (with Trees)

```
grammar
  sorts S E
  productions
    S.If  = if E then S
    S.IfE = if E then S else S
    S     = other
```

```
shift-reduce tree parse
     $                                  | if E then if E then S else S $ ...
  =>* $ if E then if E then S           |                      else S $ shift
  =>  $ if E then if E then S else      |                             S $ shift
  =>  $ if E then if E then S else S    |                               $ reduce by S = if E then S else S
  =>  $ if E then S[if E then S else S] |                               $ reduce by S = if E then S
  =>  $ S[if E then S[if E then S else S]] |                            $ accept
```

```
shift-reduce tree parse
     $                                  | if E then if E then S else S $ ...
  =>* $ if E then if E then S           |                      else S $ reduce by S = if E then S
  =>  $ if E then S[if E then S]        |                      else S $ shift
  =>  $ if E then S[if E then S] else   |                             S $ shift
  =>  $ if E then S[if E then S] else S |                               $ reduce by S = if E then S else S
  =>  $ S[if E then S[if E then S] else S] |                            $ accept
```

Simple LR Parsing

How can we make shift-reduce parsing deterministic?

```
grammar
  productions
    E.P = E "+" T
    E.E = T
    T.M = T "*" F
    T.T = F
    F.B = "(" E ")"
    F.V = ID
```

```
shift-reduce parse
     $            | ID "*" ID $ shift
  => $ ID         |    "*" ID $ reduce by F = ID
  => $ F          |    "*" ID $ reduce by T = F
  => $ T          |    "*" ID $ shift
  => $ T "*"      |        ID $ shift
  => $ T "*" ID   |           $ reduce by F = ID
  => $ T "*" F    |           $ reduce by T = T "*" F
  => $ T          |           $ reduce by E = T
  => $ E          |           $ accept
```

Is there a production in the grammar that matches the top of the stack?

How to choose between possible shift and reduce actions?

LR(k) Parsing

- L: left-to-right scanning
- R: constructing a right-most derivation
- k: the number of lookahead symbols

Motivation for LR

- LR grammars for many programming languages
- Most general non-backtracking shift-reduce parsing method
- Early detection of parse errors
- Class of grammars superset of predictive/LL methods

SLR: Simple LR

Items

`E = E "+" T`

`[E = . E "+" T]`

We expect to see an expression

`[E = E . "+" T]`

We have seen an expression
and may see a “+”

`[E = E "+" . T]`

`[E = E "+" T .]`

Production

Items

Item indicates how far we have progressed in parsing a production

Item Set

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
{
  [F = "(" . E ")"]
  [E = . E "+" T]
  [E = . T]
  [T = . T "*" F]
  [T = . F]
  [F = . "(" E ")"]
  [F = . ID]
}
```

Item set used to keep track where we are in a parse

Closure of an Item Set

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
{
  [F = "(" . E ")"]
}
```

```
{
  [F = "(" . E ")"]
  [E = . E "+" T]
  [E = . T]
  [T = . T "*" F]
  [T = . F]
  [F = . "(" E ")"]
  [F = . ID]
}
```

```
Set<Item> Closure(I) {
  J := I;
  for(each [A = a . B b] in J)
    for(each [B = c] in G)
      if(not [B = . c] in J)
        J := J + [B = .c];
  return J;
}
```

Goto

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
{
  [S = . E]
  [E = . E "+" T]
  [E = . T]
  [T = . T "*" F]
  [T = . F]
  [F = . "(" E ")"]
  [F = . ID]
}
```

```
Set<Item> Goto(I, X) {
  J := {};
  for(each [A = a . X b] in I)
    J := J + [A = a X . b];
  return Closure(J);
}
```

```
{
  [F = "(" . E ")"]
}
```

```
{
  [F = "(" . E ")"]
  [E = . E "+" T]
  [E = . T]
  [T = . T "*" F]
  [T = . F]
  [F = . "(" E ")"]
  [F = . ID]
}
```

Computing LR(0) Automaton

```
Set<Set<Item>> Items(G) {
  C := {Closure({[Sp = . S]})};
  W := C;
  while(W = {I | W'}) {
    W := W';
    for(each X in G) {
      J := Goto(I, X);
      if(not J = {} and not J in C)
        C := C + J;
        W := W + J;
      }
  }
  return C;
}
```

Computing LR(0) Automaton

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
state 0 {
    [S = . E]
  }
```

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
```

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
```

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
```

```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
```

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
  shift T to 2
```

```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
```

```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
```

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
  shift T to 2
  shift F to 3
```

```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
```

```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
```

```
state 3 {
    [T = F . ]
  }
  reduce T = F
```

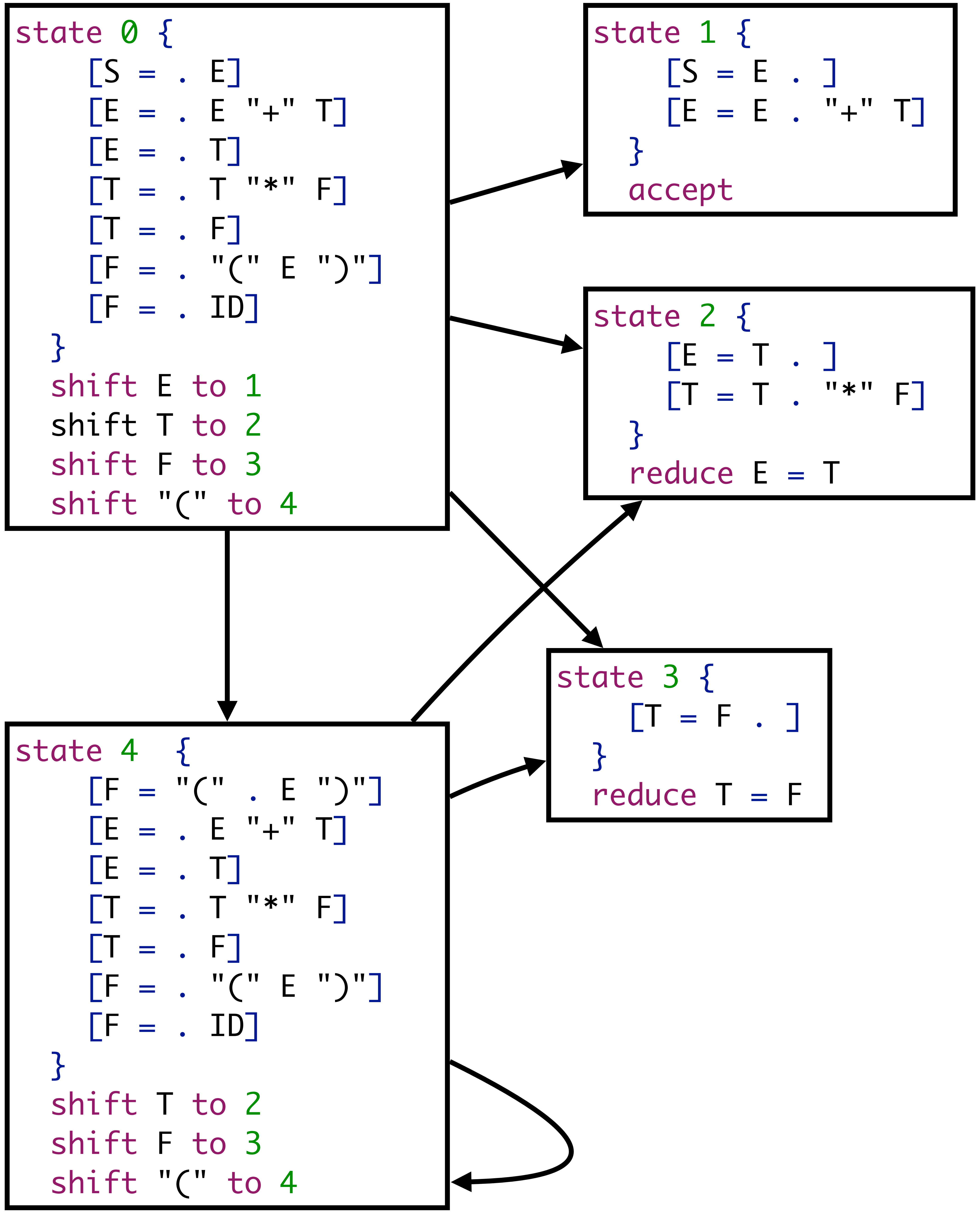
```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
  shift T to 2
  shift F to 3
  shift "(" to 4
```

```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
```

```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
```

```
state 3 {
    [T = F . ]
  }
  reduce T = F
```

```
state 4  {
    [F = "(" . E ")"]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
```


state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 2
shift F to 3
shift "(" to 4

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
  shift T to 2
  shift F to 3
  shift "(" to 4
  shift ID to 5
```

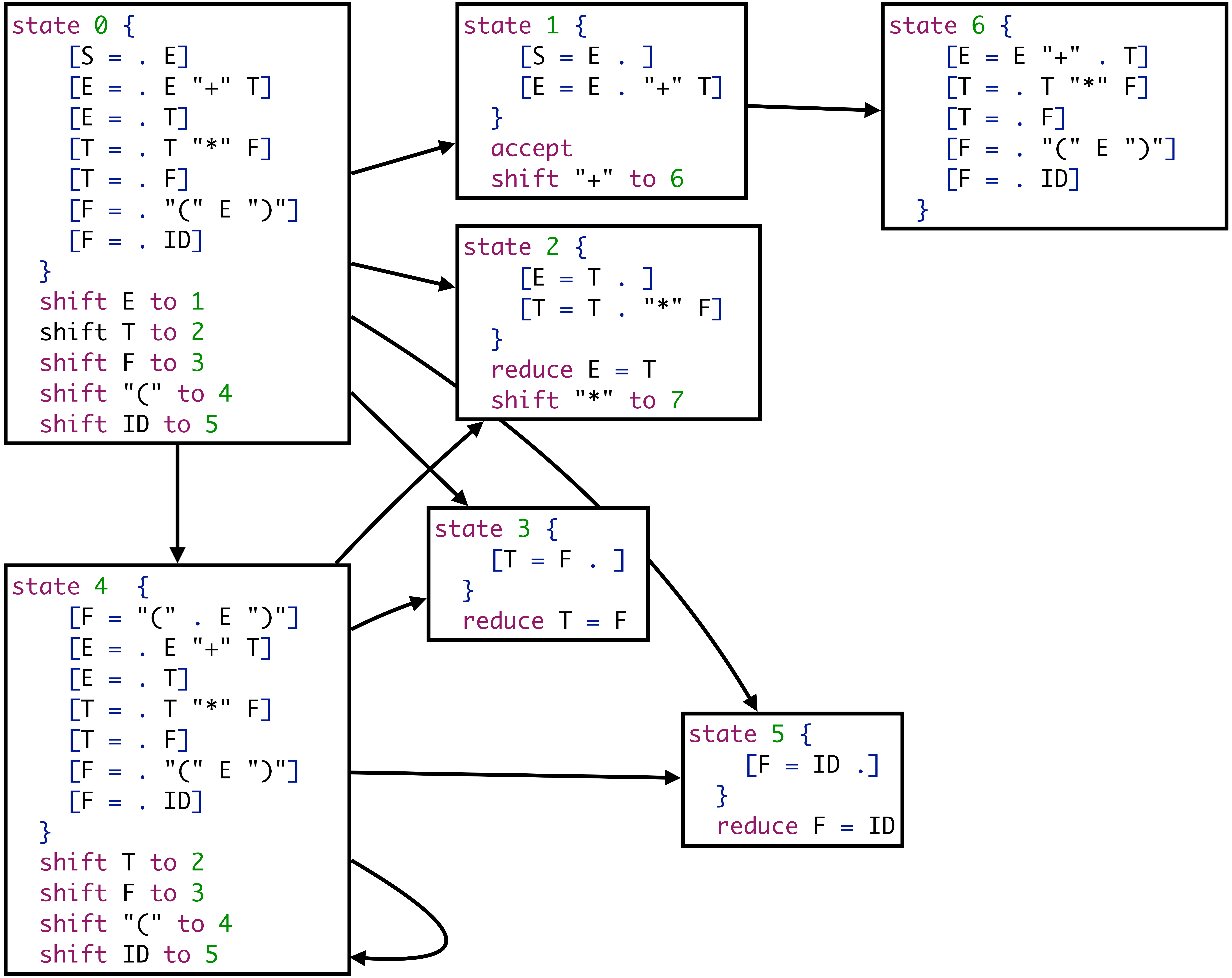
```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
  shift "+" to 6
```

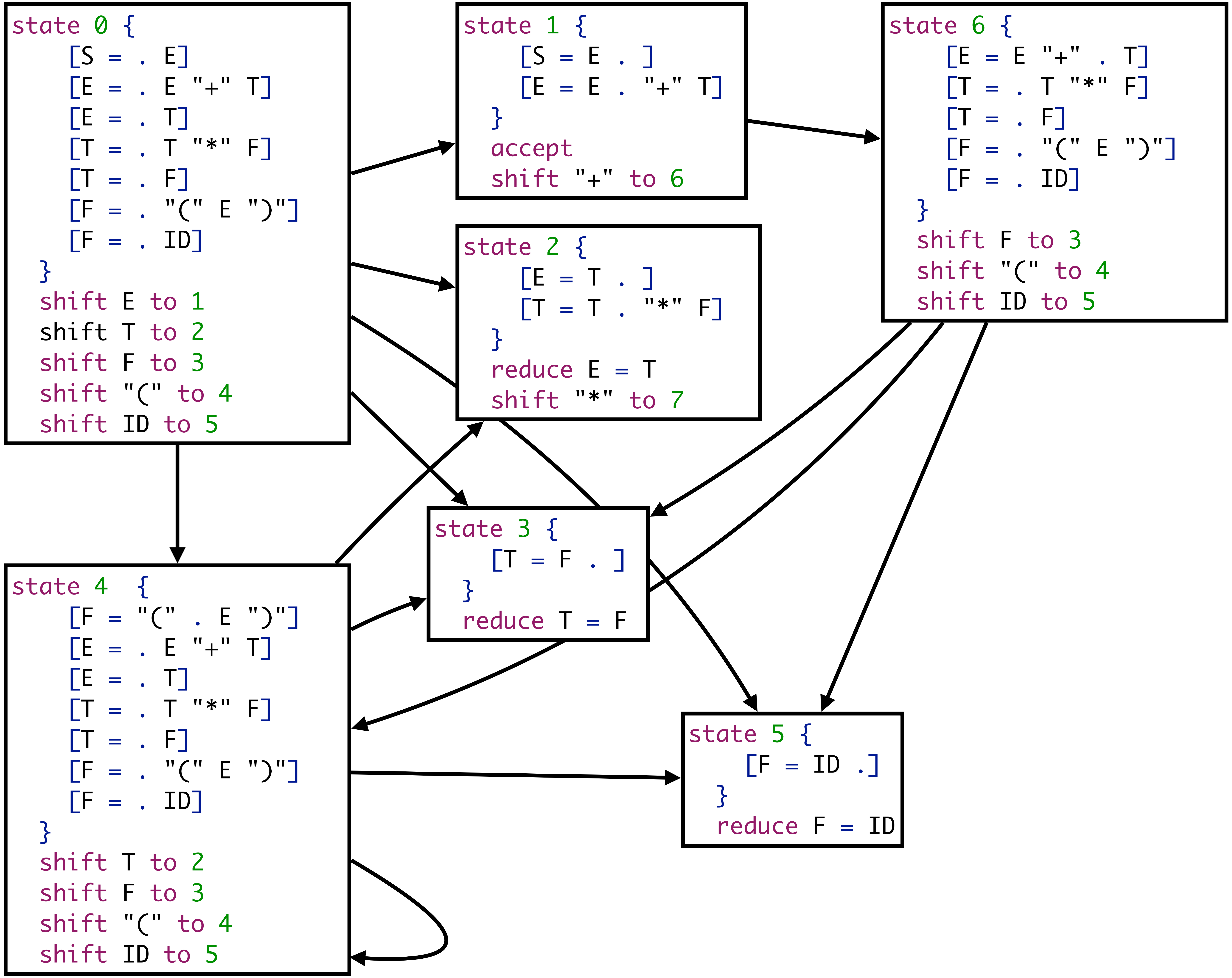
```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
  shift "*" to 7
```

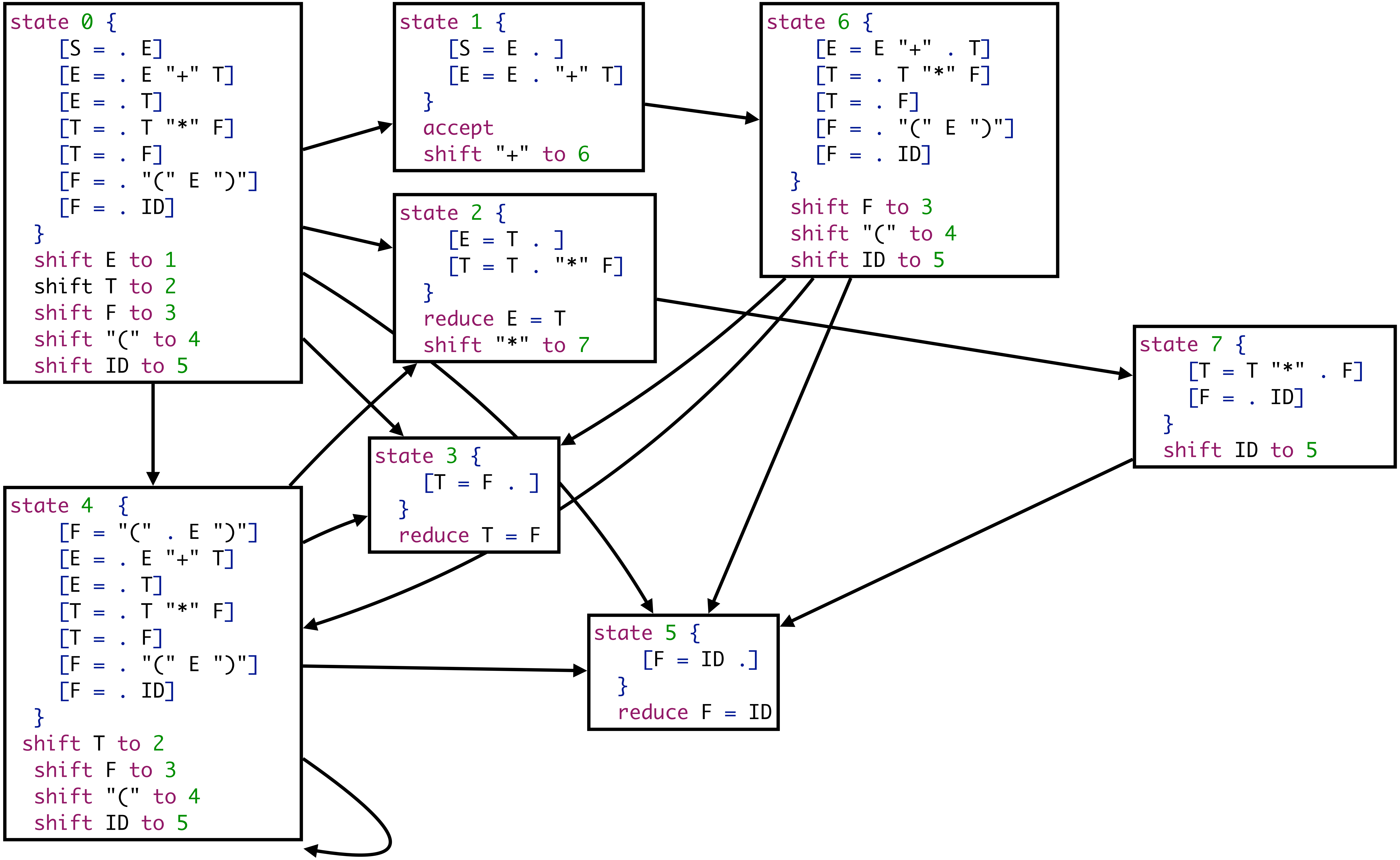
```
state 3 {
    [T = F . ]
  }
  reduce T = F
```

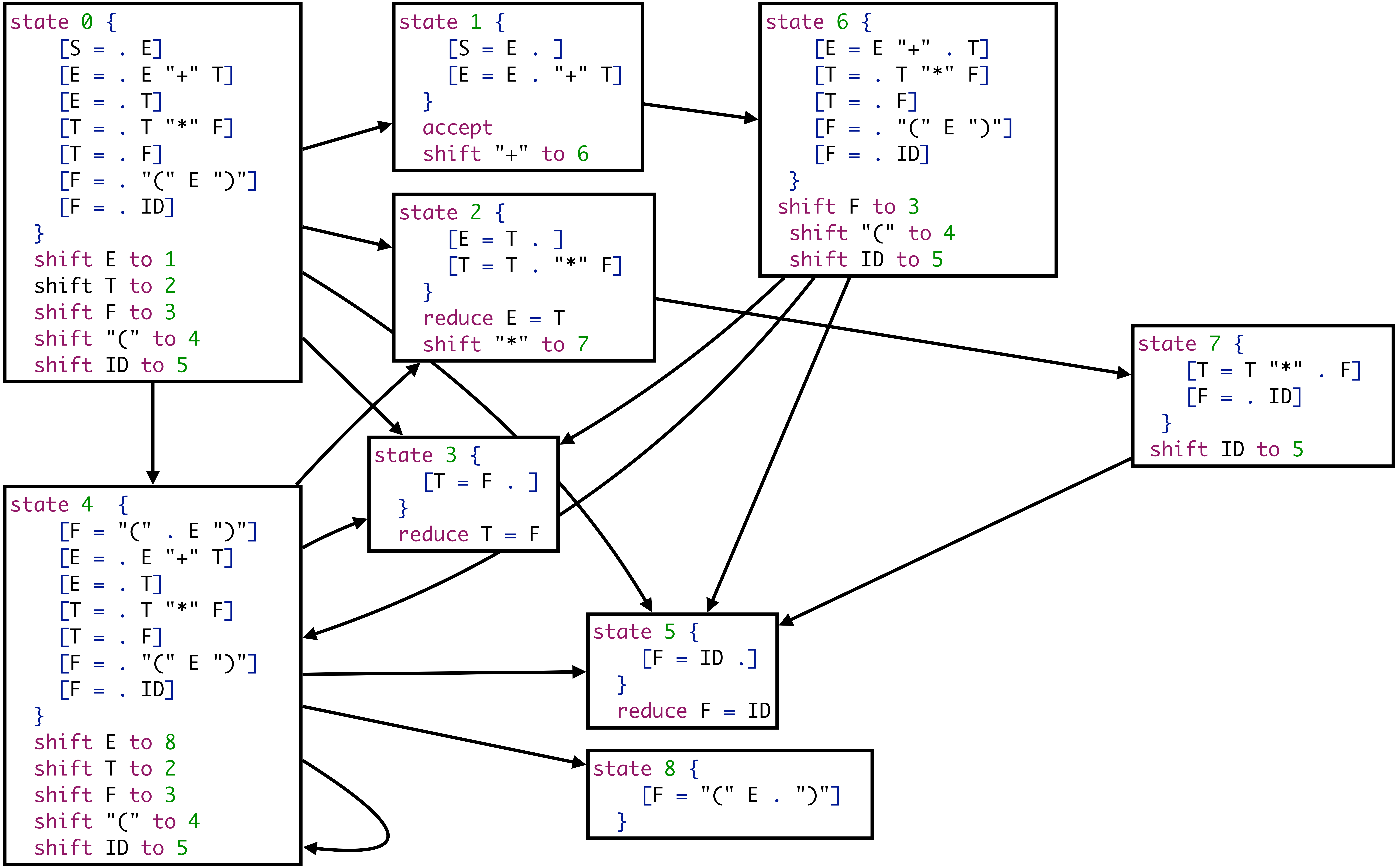
```
state 4  {
    [F = "(" . E ")"]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift T to 2
  shift F to 3
  shift "(" to 4
  shift ID to 5
```

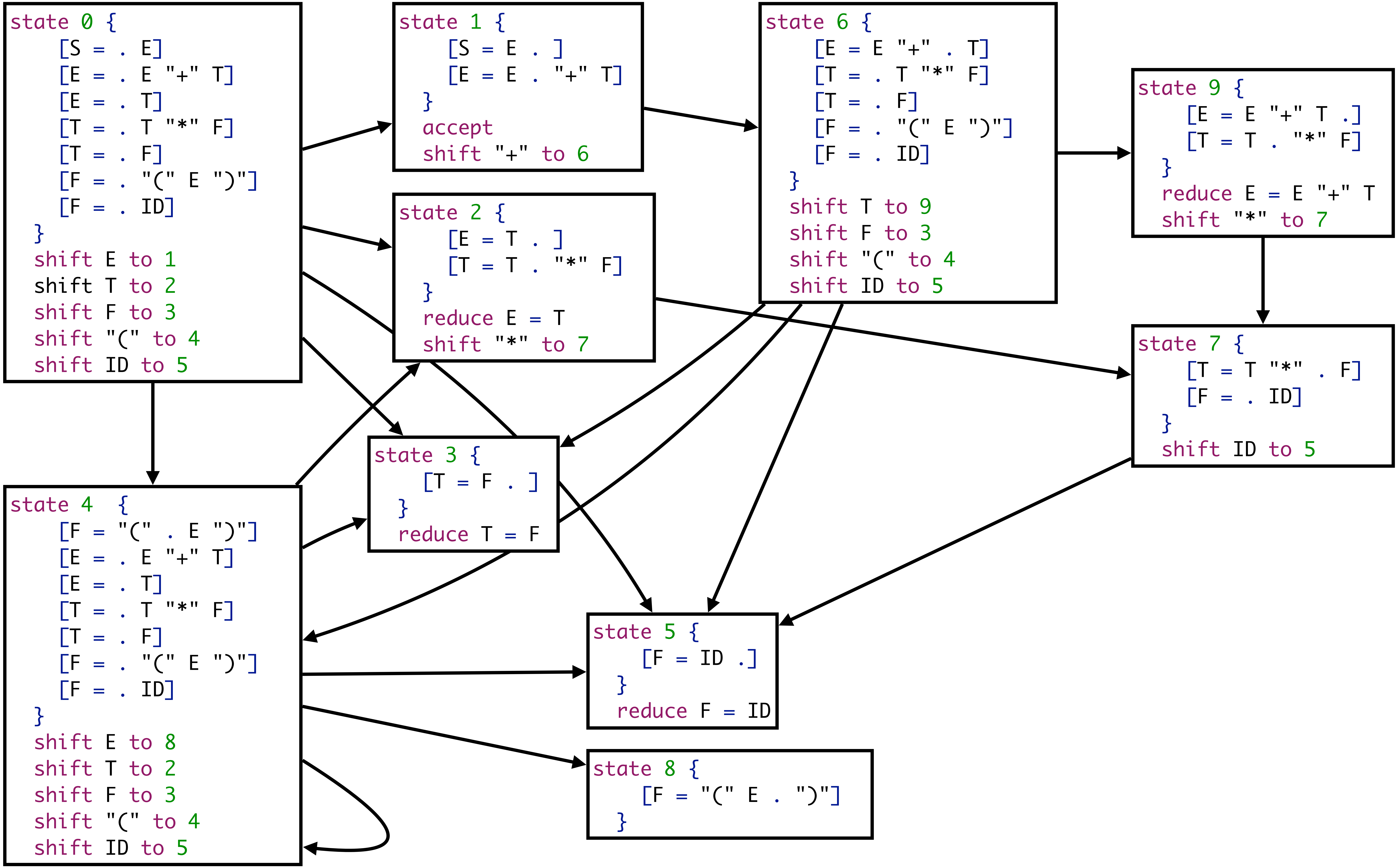
```
state 5 {
    [F = ID .]
  }
  reduce F = ID
```

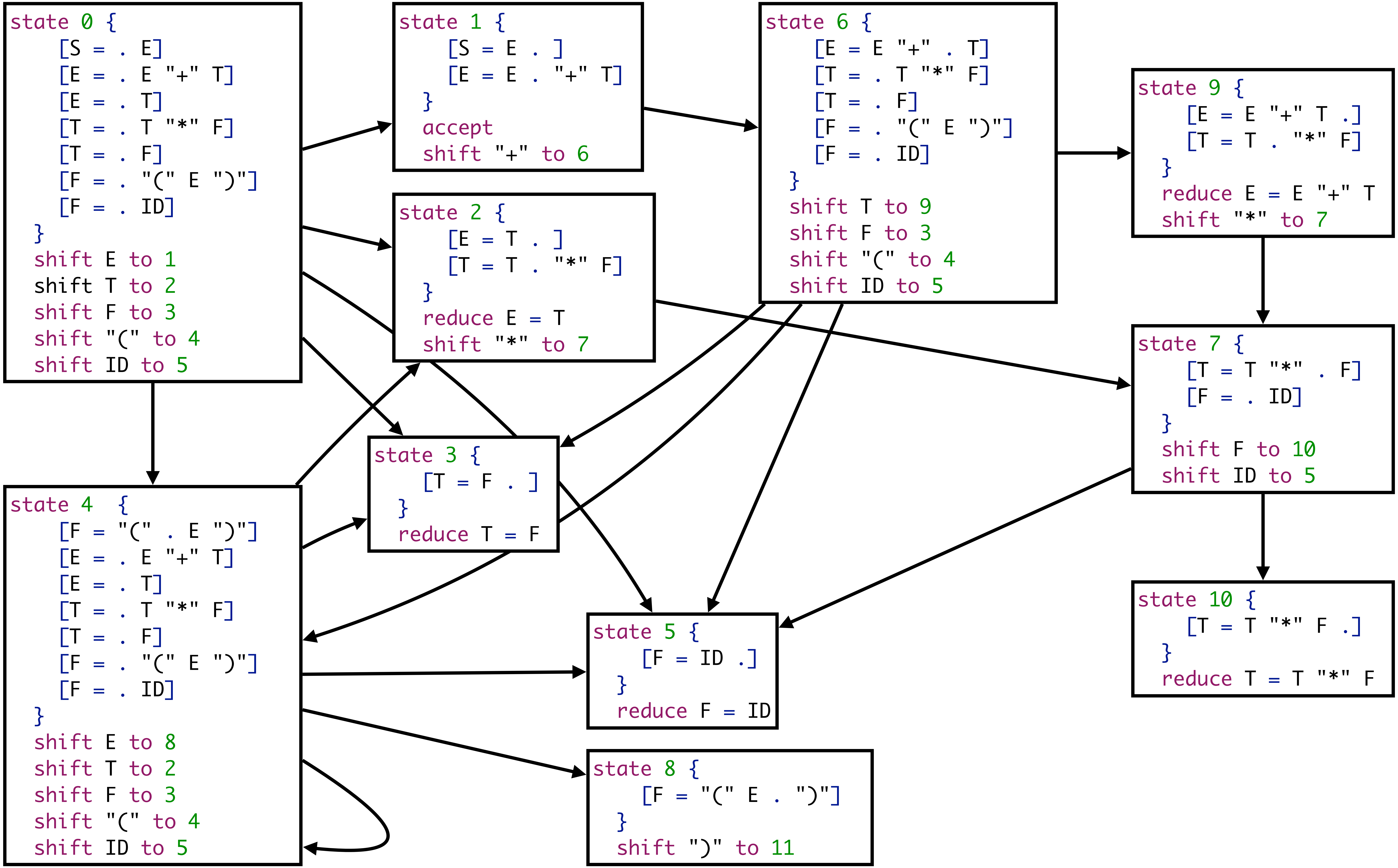

state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID

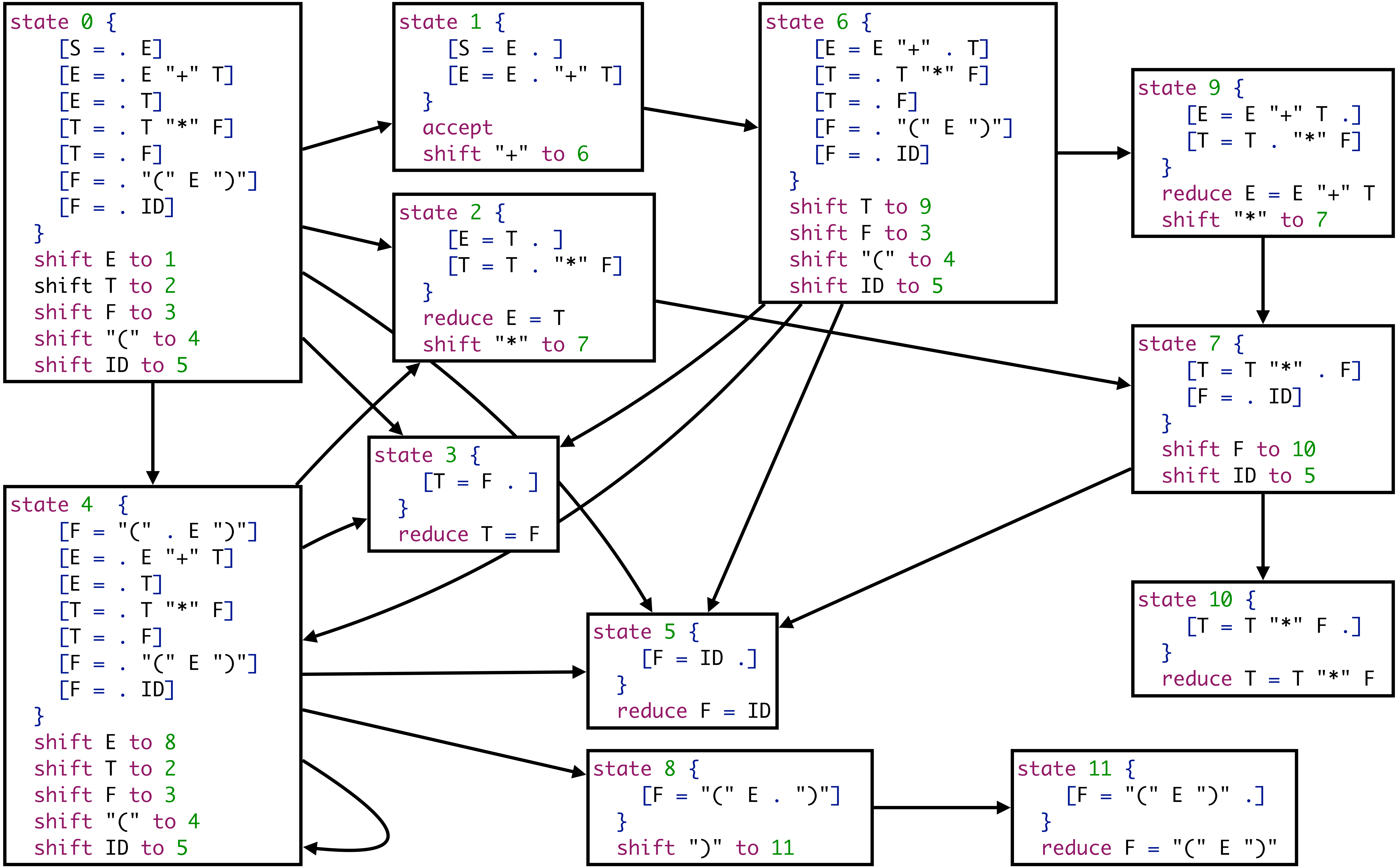

state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift F to 3
shift "(" to 4
shift ID to 5
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID


state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift F to 3
shift "(" to 4
shift ID to 5
state 7 {
[T = T "*" . F]
[F = . ID]
}
shift ID to 5
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID


state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift F to 3
shift "(" to 4
shift ID to 5
state 7 {
[T = T "*" . F]
[F = . ID]
}
shift ID to 5
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 8
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID
state 8 {
[F = "(" E . ")"]
}


state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 8
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 9
shift F to 3
shift "(" to 4
shift ID to 5
state 7 {
[T = T "*" . F]
[F = . ID]
}
shift ID to 5
state 8 {
[F = "(" E . ")"]
}
state 9 {
[E = E "+" T .]
[T = T . "*" F]
}
reduce E = E "+" T
shift "*" to 7


state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 8
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 9
shift F to 3
shift "(" to 4
shift ID to 5
state 7 {
[T = T "*" . F]
[F = . ID]
}
shift F to 10
shift ID to 5
state 8 {
[F = "(" E . ")"]
}
shift ")" to 11
state 9 {
[E = E "+" T .]
[T = T . "*" F]
}
reduce E = E "+" T
shift "*" to 7
state 10 {
[T = T "*" F .]
}
reduce T = T "*" F



state 0 {
[S = . E]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 1
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 1 {
[S = E .]
[E = E . "+" T]
}
accept
shift "+" to 6
state 2 {
[E = T .]
[T = T . "*" F]
}
reduce E = T
shift "*" to 7
state 6 {
[E = E "+" . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift T to 9
shift F to 3
shift "(" to 4
shift ID to 5
state 9 {
[E = E "+" T .]
[T = T . "*" F]
}
reduce E = E "+" T
shift "*" to 7
state 7 {
[T = T "*" . F]
[F = . ID]
}
shift F to 10
shift ID to 5
state 3 {
[T = F .]
}
reduce T = F
state 4 {
[F = "(" . E ")"]
[E = . E "+" T]
[E = . T]
[T = . T "*" F]
[T = . F]
[F = . "(" E ")"]
[F = . ID]
}
shift E to 8
shift T to 2
shift F to 3
shift "(" to 4
shift ID to 5
state 5 {
[F = ID .]
}
reduce F = ID
state 10 {
[T = T "*" F .]
}
reduce T = T "*" F
state 8 {
[F = "(" E . ")"]
}
shift ")" to 11
state 11 {
[F = "(" E ")" .]
}
reduce F = "(" E ")"

SLR Parse

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
   $            | 0          | ID "*" ID $ shift ID to 5
=> $ ID         | 0 5        |    "*" ID $ reduce F = ID
=> $            | 0          |  F "*" ID $ shift F to 3
=> $ F          | 0 3        |    "*" ID $ reduce T = F
=> $            | 0          |  T "*" ID $ shift T to 2
=> $ T          | 0 2        |    "*" ID $ shift "*" to 7
=> $ T "*"      | 0 2 7      |        ID $ shift ID to 5
=> $ T "*" ID   | 0 2 7 5    |           $ reduce F = ID
=> $ T "*"      | 0 2 7      |         F $ shift F to 10
=> $ T "*" F    | 0 2 7 10   |           $ reduce T = T "*" F
=> $            | 0          |         T $ shift T to 2
=> $ T          | 0 2        |           $ reduce E = T
=> $            | 0          |         E $ shift E to 1
=> $ E          | 0 1        |           $ accept
```

```
$            | 0        | ID "*" ID $ shift ID to 5
=> $ ID         | 0 5      |    "*" ID $ reduce F = ID
=> $            | 0        |  F "*" ID $ shift F to 3
=> $ F          | 0 3      |    "*" ID $ reduce T = F
=> $            | 0        |  T "*" ID $ shift T to 2
=> $ T          | 0 2      |    "*" ID $ shift "*" to 7
=> $ T "*"      | 0 2 7    |        ID $ shift ID to 5
=> $ T "*" ID   | 0 2 7 5  |           $ reduce F = ID
=> $ T "*"      | 0 2 7    |         F $ shift F to 10
=> $ T "*" F    | 0 2 7 10 |           $ reduce T = T "*" F
=> $            | 0        |         T $ shift T to 2
=> $ T          | 0 2      |           $ reduce E = T
=> $            | 0        |         E $ shift E to 1
=> $ E          | 0 1      |           $ accept
```

0

```
state 0 {
    [S = . E]
    [E = . E "+" T]
    [E = . T]
    [T = . T "*" F]
    [T = . F]
    [F = . "(" E ")"]
    [F = . ID]
  }
  shift E to 1
  shift T to 2
  shift F to 3
  shift "(" to 4
  shift ID to 5
```

1

```
state 5 {
    [F = ID .]
  }
  reduce F = ID
```

2

```
state 3 {
    [T = F . ]
  }
  reduce T = F
```

3

```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
  shift "*" to 7
```

4

```
state 7 {
    [T = T "*" . F]
    [F = . ID]
  }
  shift F to 10
  shift ID to 5
```

5

```
state 5 {
    [F = ID .]
  }
  reduce F = ID
```

6

```
state 10 {
    [T = T "*" F .]
  }
  reduce T = T "*" F
```

7

```
state 2 {
    [E = T . ]
    [T = T . "*" F]
  }
  reduce E = T
  shift "*" to 7
```

8

```
state 1 {
    [S = E . ]
    [E = E . "+" T]
  }
  accept
  shift "+" to 6
```

Parsing: ID "+" ID "*" ID .

Solving Shift/Reduce Conflicts

SLR Parse

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
   $            | 0          | ID "*" ID $ shift ID to 5
=> $ ID         | 0 5        |    "*" ID $ reduce F = ID
=> $            | 0          |  F "*" ID $ shift F to 3
=> $ F          | 0 3        |    "*" ID $ reduce T = F
=> $            | 0          |  T "*" ID $ shift T to 2
=> $ T          | 0 2        |    "*" ID $ shift "*" to 7
=> $ T "*"      | 0 2 7      |        ID $ shift ID to 5
=> $ T "*" ID   | 0 2 7 5    |           $ reduce F = ID
=> $ T "*"      | 0 2 7      |         F $ shift F to 10
=> $ T "*" F    | 0 2 7 10   |           $ reduce T = T "*" F
=> $            | 0          |         T $ shift T to 2
=> $ T          | 0 2        |           $ reduce E = T
=> $            | 0          |         E $ shift E to 1
=> $ E          | 0 1        |           $ accept
```

Why did we choose shift?

SLR Parse

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
   $          | 0        | ID "*" ID $ shift ID to 5
=> $ ID       | 0 5      |    "*" ID $ reduce F = ID
=> $          | 0        |  F "*" ID $ shift F to 3
=> $ F        | 0 3      |    "*" ID $ reduce T = F
=> $          | 0        |  T "*" ID $ shift T to 2
=> $ T        | 0 2      |    "*" ID $ reduce E = T
=> $ E        | 0        |    "*" ID $ shift E to 1
=> $ E        | 0 1      |    "*" ID $ error
```

Reduce action is also possible, but leads to error

How can we avoid that?

SLR Parse

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
   $            | 0          | ID "*" ID $ shift ID to 5
=> $ ID         | 0 5        |    "*" ID $ reduce F = ID
=> $            | 0          |  F "*" ID $ shift F to 3
=> $ F          | 0 3        |    "*" ID $ reduce T = F
=> $            | 0          |  T "*" ID $ shift T to 2
=> $ T          | 0 2        |    "*" ID $ reduce E = T
=> $ E          | 0 2        |    "*" ID $ shift E to 1
=> $ E          | 0 1        |    "*" ID $ error
```

E can not be followed by a "*"!

Rule: only reduce with [A = Bs] if next token can follow A

First and Follow

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
first sets
  S: {"(", ID}
  E: {"(", ID}
  T: {"(", ID}
  F: {"(", ID}
  ID: { ID }
  "+": {"+"}
  "*": {"*"}
  "(": {"("}
  ")": {")"}
```

```
follow sets
  S: {$}
  E: {$,"+",")"}
  T: {$,"+",")","*"}
  F: {$,"+",")","*"}
  ID: {$,"+",")","*"}
```

First: the tokens that a phrase for a non-terminal can start with

Follow: the tokens that can follow a non-terminal

First

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
first sets
  S: {"(", ID}
  E: {"(", ID}
  T: {"(", ID}
  F: {"(", ID}
  ID: { ID }
  "+": {"+"}
  "*": {"*"}
  "(": {"("}
  ")": {")"}
```

First: the tokens that a phrase for a symbol can start with

If A = Bs is a production, then First[A] includes First[Bs]

A terminal starts with itself

First (Attempt 1)

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
first sets
  S: {"(", ID}
  E: {"(", ID}
  T: {"(", ID}
  F: {"(", ID}
  ID: { ID }
  "+": {"+"}
  "*": {"*"}
  "(": {"("}
  ")": {")"}
```

```
var First : Map<Symbol, Set<Symbol>>

FIRST(G) {
  First := {};
  repeat {
    First' := First;
    for(each [A = As] in G) {
      First[A] := First[A] + Firsts(As);
    }
  } until First = First';
}
Set<Symbol> Firsts(As) {
  As match {
    [] => {};
    [A Bs] =>
      if(Terminal(A))
        {A}
      else
        First[A];
  };
}
```

First (Attempt 1)

```
grammar A
  start S
  non-terminals Decl Mod Args
  terminals ID "(" ")" "static" ""
  productions
    S    = Decl
    Decl = Mod ID "(" Args ")"
    Mod  = "static"
    Mod  =
    Args = Args ID
    Args =
```

```
first sets of A
  S: {"static"}
  Decl: {"static"}
  Mod: {"static"}
  Args: {}
  ID: { ID }
  "(": {"("}
  ")": {")"}
  "static": {"static"}
```

```
var First : Map<Symbol, Set<Symbol>>

FIRST(G) {
  First := {};
  repeat {
    First' := First;
    for(each [A = As] in G) {
      First[A] := First[A] + Firsts(As);
    }
  } until First = First';
}
Set<Symbol> Firsts(As) {
  As match {
    [] => {};
    [A Bs] =>
      if(Terminal(A))
        {A}
      else
        First[A];
  };
}
```

Is Non-Terminal Nullable?

```
grammar A
  start S
  non-terminals Decl Mod Args
  terminals ID "(" ")" "static" ""
  productions
    S    = Decl
    Decl = Mod ID "(" Args ")"
    Mod  = "static"
    Mod  =
    Args = Args ID
    Args =
```

```
var Nullable : Map<Symbol, Bool>

NULLABLE(G) {
  for(each A in G) {
    Nullable[A] := False;
  }
  repeat {
    Nullable' := Nullable;
    for(each [A = As] in G) {
      Nullable[A] := Nullable[A] || Nullables(As);
    }
  } until Nullable = Nullable';
}

Bool Nullables(As) {
  As match {
    [] => True;
    [B Bs] => not Terminal(B)
           or (Nullable[B] and Nullables(Bs));
  };
}
```

What are the First Terminals of a Non-Terminal?

```
grammar A
  start S
  non-terminals Decl Mod Args
  terminals ID "(" ")" "static" ""
  productions
    S    = Decl
    Decl = Mod ID "(" Args ")"
    Mod  = "static"
    Mod  =
    Args = Args ID
    Args =
```

```
first sets of A
  S: {"static",ID}
  Decl: {"static",ID}
  Mod: {"static"}
  Args: {ID}
  ID: { ID }
  "(": {"("}
  ")": {")"}
  "static": {"static"}
```

```
var First : Map<Symbol, Set<Symbol>>

FIRST(G) {
  First := {};
  repeat {
    First' := First;
    for(each [A = As] in G) {
      First[A] := First[A] + Firsts(As);
    }
  } until First = First';
}
Set<Symbol> Firsts(As) {
  As match {
    [] => {};
    [A Bs] =>
      if(Terminal(A))
        {A}
      else
        First[A] + if(Nullable[A]) Firsts(Bs) else {};
  };
}
```

What Terminal can Follow a Non-Terminal Start with?

```
grammar
 productions
    S = E
    E = E "+" T
    E = T
    T = T "*" F
    T = F
    F = "(" E ")"
    F = ID
```

```
follow sets of A
  S: {$}
  E: {$,"+",")"}
  T: {$,"+",")","*"}
  F: {$,"+",")","*"}
  ID: {$,"+",")","*"}
```

```
var Follow : Map<Symbol, Set<Symbol>>

FOLLOW(G) {
  Follow := {};
  repeat {
    Follow' := Follow;
    for(each [A = As B Cs] in G) {
      Follow[B] := Follow[B]
                 + Firsts(Cs)
                 + if(Nullables(Cs)) Follow[A]
                   else {};
    }
  } until Follow = Follow';
}
```

What Terminal can Follow a Non-Terminal Start with?

```
grammar A
  start S
  non-terminals Decl Mod Args
  terminals ID "(" ")" "static" ""
  productions
    S    = Decl
    Decl = Mod ID "(" Args ")"
    Mod  = "static"
    Mod  =
    Args = Args ID
    Args =
```

```
follow sets of A
  S: {$}
  Decl: {$}
  Mod: {ID}
  Args: {")" ID}
  ID: {"(",")"}
```

```
var Follow : Map<Symbol, Set<Symbol>>

FOLLOW(G) {
  Follow := {};
  repeat {
    Follow' := Follow;
    for(each [A = As B Cs] in G) {
      Follow[B] := Follow[B]
                 + Firsts(Cs)
                 + if(Nullables(Cs)) Follow[A]
                   else {};
    }
  } until Follow = Follow';
}
```

Parsing: Summary

Parsing

Context-free grammars

- Productions define how to generate sentences of language
- Derivation: generate sentence from (start) symbol
- Reduction: reduce sentence to (start) symbol

Parse tree

- Represents structure of derivation
- Abstracts from derivation order

Parser

- Algorithm to reconstruct derivation

First/Follow

- Selecting between actions in LR parse table

Other algorithms

- Top-down: LL(k) table
- Generalized parsing: Earley, Generalized-LR
- Scannerless parsing: characters as tokens

Disambiguation

- Semantics of declarative disambiguation
- Deep priority conflicts

LL Parsing

Use LL(1) grammar

- Not left recursive
- Left factored

Top-down back-track parsing

- Predict symbol
- If terminal: corresponds to next input symbol?
- Try productions for non-terminal in turn

Predictive parsing

- Predict symbol to parse
- Use lookahead to deterministically chose production for non-terminal

Variants

- Parser combinators, PEG, packrat, …

Next: Editor Services