

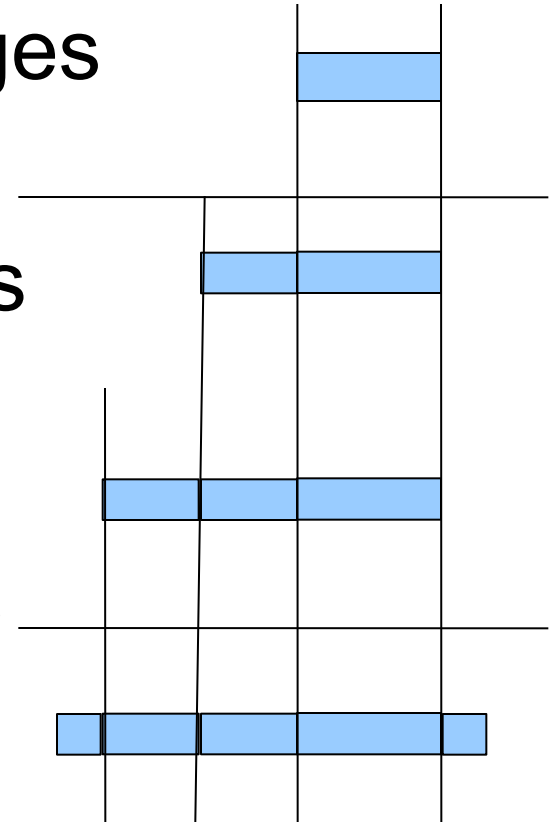
Computer Networks

Link Layer

Jianping Pan
Fall 2020

Review

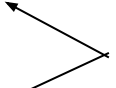
- Application layer: messages
 - HTTP, DNS
- Transport layer: segments
 - TCP, UDP
- Network layer: packets
 - IP, ICMP; RIP, OSPF, BGP
- Link layer: frames



Link layer services

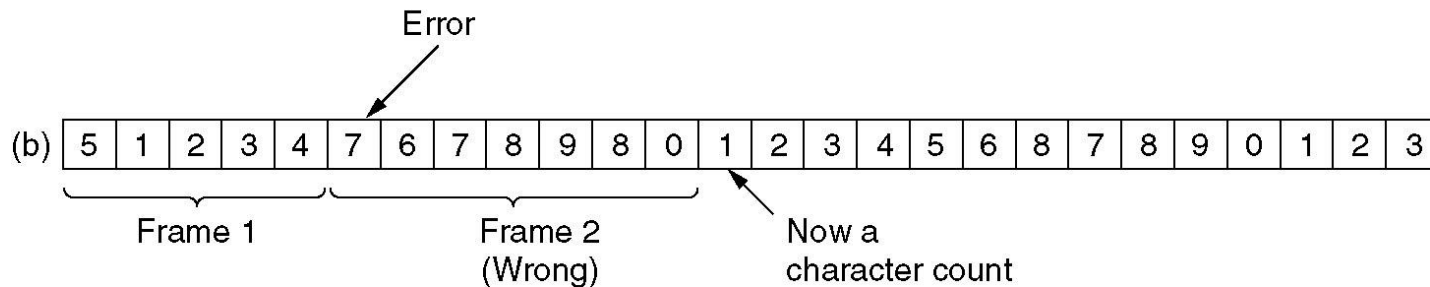
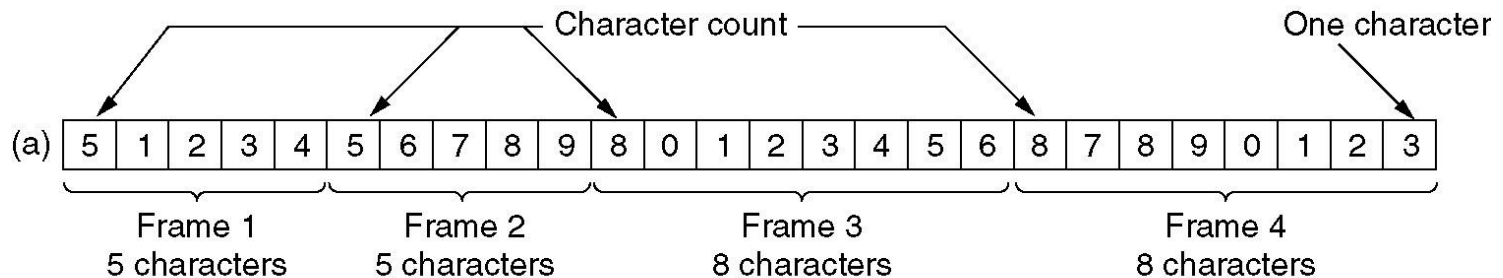
- Services provided by physical layer
 - bit delivery (recall: access networks)
 - hertz, baud, symbol-per-second, bit-per-second
- Services provided to network layer
 - frame control: framing
 - error control: how to deal with bit errors
 - flow control: fast sender vs slow receiver
 - medium access control (with shared medium)

Today's
topics



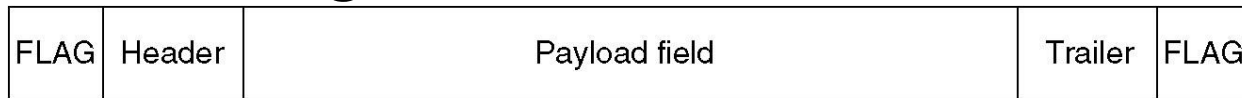
Frame control

- Character count
 - count error, and error propagation

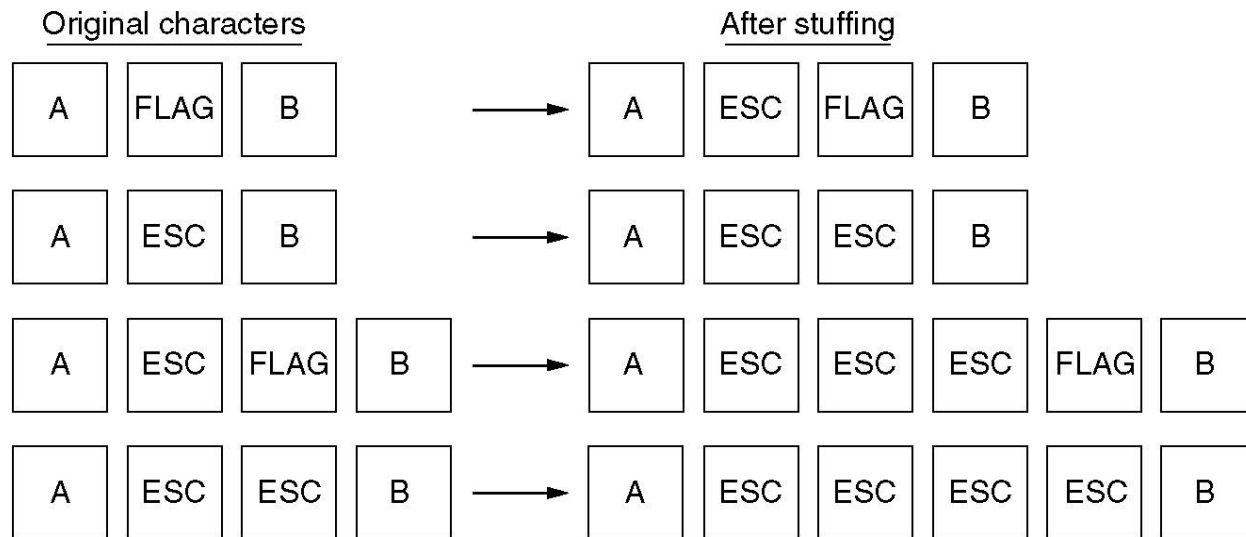


Byte-oriented framing

- Byte stuffing



(a)



(b)

CSc 361


Q: error propagation?

Bit-oriented framing

- Flag: 01111110
 - data transparency: bit stuffing
 - sender: insert a 0 after 5 1's
 - receiver: remove a 0 after 5 1's

(a) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

(b) 0 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 1 1 1 1 0 1 0 0 1 0



Stuffed bits

(c) 0 1 1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 0 0 1 0

Error control

- Hamming distance of codeword a and b

– number of *pairwisely* different bits

• number of bit flips needed to turn a to b

$$\begin{array}{r} 01010101 \\ \text{XOR) } 00100100 \\ \hline 01110001 \\ \downarrow \\ 4 \end{array}$$

- Hamming distance of codeword set $\{a_i\}$

– minimal distance btw a_i and a_j , where $i \neq j$

- e bit errors

– to detect: minimal Hamming distance $e+1$

– to correct: minimal Hamming distance $2e+1$

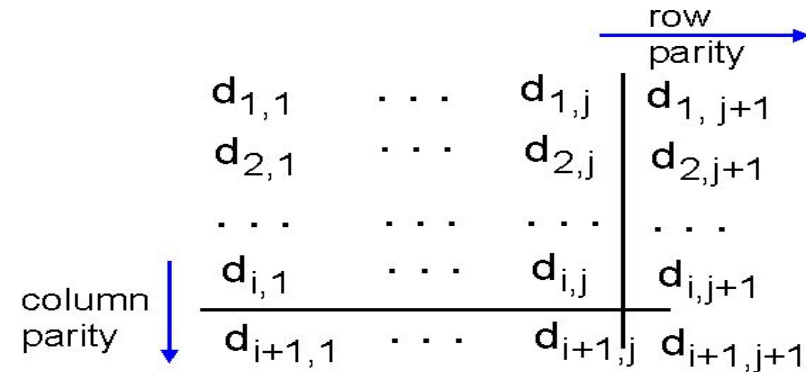
Parity check

- Parity bit

- even or odd parity
- i.e., the number of 1's
- e.g., 10101; check bit: 1 (even)
- Q: Hamming distance?
- detect 1-bit error

- 2-d parity

- correct 1-bit error



1	0	1	0	1	1
1	1	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

no errors

1	0	1	0	1	1
1	0	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

parity error

parity error

correctable single bit error

Hamming code

- Hamming code

- check bits

- at bit 1, 2, 4, 8

- data bits

- at bit 3, 5, 6, 7, 9, 10, 11

- e.g, 1001000

- correct 1-bit error

- Hamming code block

- correct up to block length

Char.

ASCII

Check bits

H

1001000

a

1100001

m

1101101

m

1101101

i

1101001

n

1101110

g

1100111

0100000

c

1100011

o

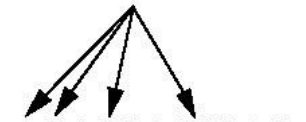
1101111

d

1100100

e

1100101



00110010000

10111001001

11101010101

11101010101

01101011001

01101010110

01111001111

10011000000

11111000011

10101011111

11111001100

00111000101

Order of bit transmission

Review: Internet checksum

- Checksum: widely used in upper layers
 - e.g., TCP checksum with pseudo header
 - optional UDP checksum with pseudo header
 - IP header checksum
- One's complement of one's complement sum
 - checksum generation
 - checksum verification
- When does checksum fail?

AA BB	AA BB
CC DD	CC DD
EE 00	EE 00
<u>00 00</u>	<u>9A 65</u>
+ 265 98	+ 2FF FD
65 9A	FF FF
<u>9A 65</u>	<u>00 00</u>

Cyclic Redundancy Check

- CRC: widely used in lower layers
 - e.g., IEEE 802.3 CRC-32-Ethernet
 - ITU-T X.25 CRC-16-CCITT
- Polynomial representation
 - message: $M(x)$; generator: $G(x)$ of order r
 - remainder: $R(x) = M(x) * 2^r \% G(x)$
 - CRC generation: $T(x) = M(x) * 2^r \text{ XOR } R(x)$
 - i.e., $T(x)$ is $G(x)$ divisible
 - error: $E(x)$ detected if not $G(x)$ divisible

CRC Example

Want:

$$D \cdot 2^r \text{ XOR } R = nG$$

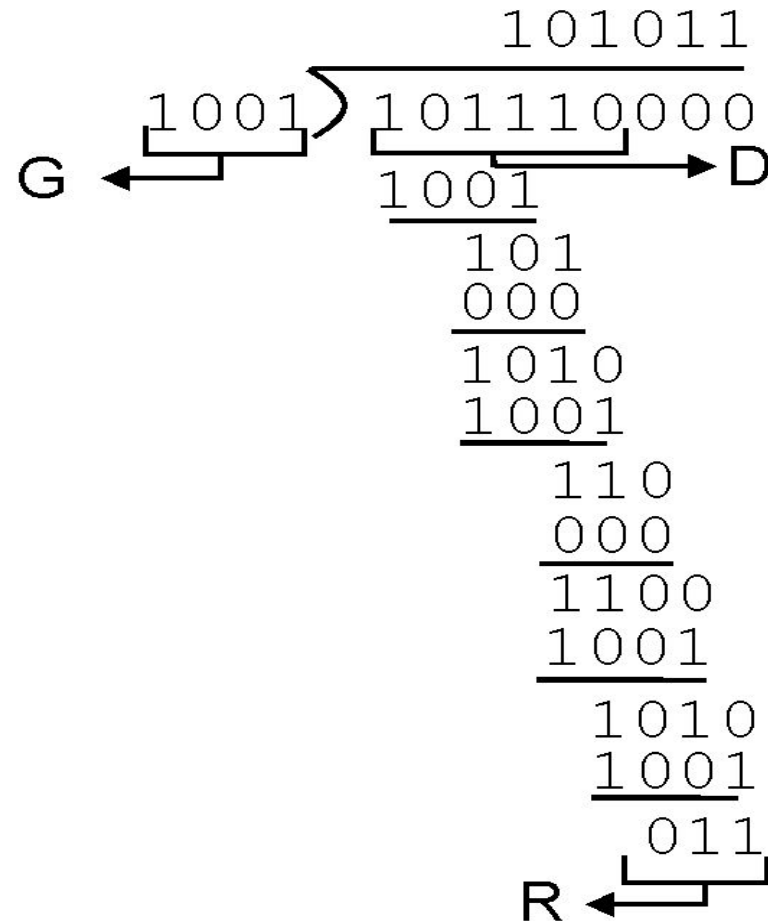
equivalently:

$$D \cdot 2^r = nG \text{ XOR } R$$

equivalently:

if we divide $D \cdot 2^r$ by G ,
want remainder R

$$R = \text{remainder}\left[\frac{D \cdot 2^r}{G}\right]$$



Error recovery

- Positive acknowledgment
 - cumulative acknowledgment
 - acknowledge packet x: acknowledge packets 1..x
 - when timeout, go-back-N
 - selective acknowledgment
 - acknowledge packet x: packet x is received OK
 - when timeout, selective repeat
- Negative acknowledgment
 - report: x is corrupted or *missing*

This lecture

- Link layer
 - framing
 - error control
 - error detecting, error correcting, error recovery
- Explore further
 - Information and Coding Theory
 - 1850s-1940s: check digit; 1940s-1960s: checksum
 - 1960s: Reed-Solomon; 1970s: LDPC codes
 - 1980s: Turbo codes; 1990s: Space-time code
 - 2000s: Polar code

Next lectures

- Flow control and LLC protocols
 - sliding window (1-bit, GBN, SR)
 - read K&R4: Computer Networking
 - Chapter 5 (except 5.7 and 5.8)
 - new editions: Chapter 6 (Link Layer)