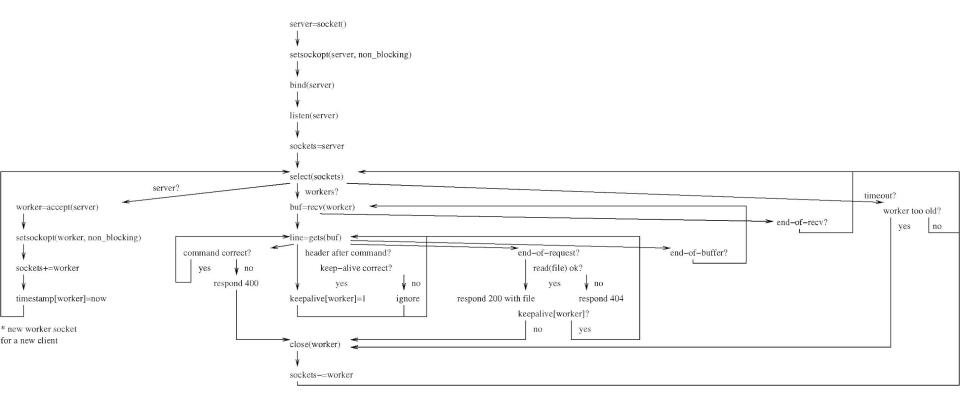


Our course reps

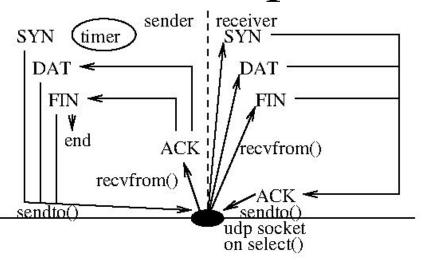
- Thanks to our diverse student volunteers!
 - B03: Chris Brandt (CS, EE, ITadmin, etc)
 - B04: Jennifer Cheng (CS+Stats)
 - B05: Emily Sluis (SEng, CSc110/1/5TA)
 - B06: Owen Thurston (CS, remote in Korea)
 - their uvic email address on connex
- AAA: Aggregate, Amplify and Anonymize
 - we will e-meet them this Thursday
 - we do welcome student feedback directly too

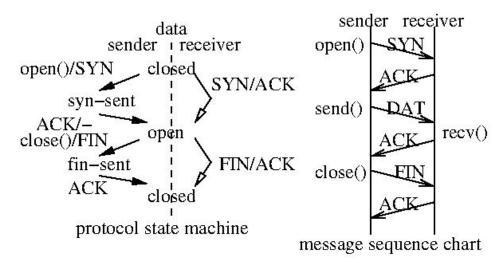
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One possible P1 flow chart



One possible P2 flow chart





sender pseudo code in open state

```
forever{
    on application write:
        packetize into packets
        send per receiver's window
        setup timer if not running
        update send_next

    on receiving ACK:
        cancel timer if covered
        setup timer if still unacked packets
        resend the oldest if enough dupacks
        send more if allowed by window

    on sender timeout:
        resend the oldest packet
        setup timer properly
```

receiver pseudo code in open state

```
forever{
    on receiving DAT:
    below acked?
    drop
    beyond acked+window?
    send RST; exit
    out of order?
    buffer or drop
    in order?
    buffer and update ackno
    enough in-order data?
    write to file
    update window size
    send ACK
}
```

Possible P3 interactions non-exhaustive nor complete

shortest-possible SYN|DAT|ACK|FIN SYN|DAT|ACK|FIN **ACK** SYN|DAT|ACK

ACK|FIN

ACK|FIN

ACK

ACK

SYN|DAT|ACK|FIN

ACK|FIN

RST

- so you need a protocol SYN|DAT|ACK
- state machine Be considerate
- on what you send
- Be accommodative SYN|DAT|ACK on what you receive

SYN|DAT|ACK|FIN **ACK**

SYN|DAT|ACK

SYN|DAT|ACK

ACK|FIN

ACK

ACK|FIN

FIN|ACK SYN TCP-like ACK DAT|ACK **ACK**

FIN|ACK

ACK

SYN

DAT

ACK

ACK

ACK

ACK

ACK

SYN|ACK

DAT|ACK

FIN|ACK

ACK SYN|ACK ACK DAT|ACK ACK FIN|ACK

^{*} see connex->forums->p3 for more discussion

Computer Networks

Wireless LAN

Jianping Pan Fall 2020

Review

- Media access control
 - Aloha
 - Slotted Aloha
 - CSMA
 - 1-persistent, p-persistent, non-persistent
 - CSMA/CD
 - IEEE 802.3 Ethernet
 - frame control, error detection

Today's topics

- WiFi: wireless fidelity
 - CSMA/CA
 - RTS/CTS
 - IEEE 802.11 family
 - frame control, error detection

Wireless LAN

- IEEE 802.11 family
 - 802.11: 2.4GHz, 2Mbps
 - 802.11a: 5GHz, 54Mbps, 30ft
 - 802.11b: 2.4GHz, 11Mbps, 100ft
 - 11 channels in North America
 - 3 non-overlapping channels: 1, 6, 11
 - 802.11g: 2.4GHz, 54Mbps, 100ft
 - OFDM: frequency division
 - 802.11n: new radio, 2.4GHz, 540Mbps



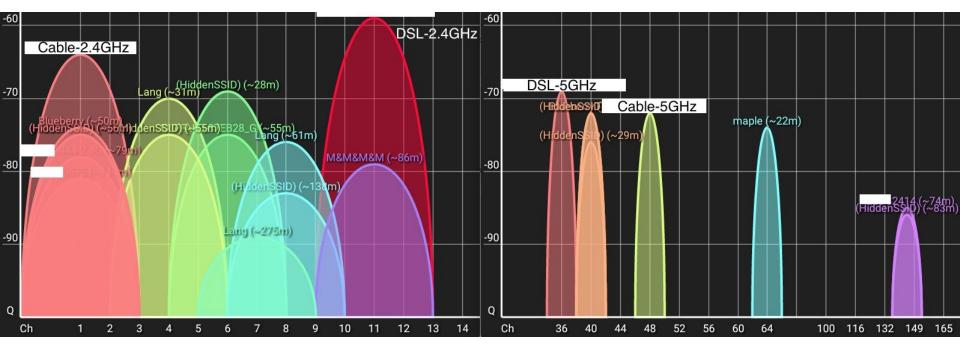


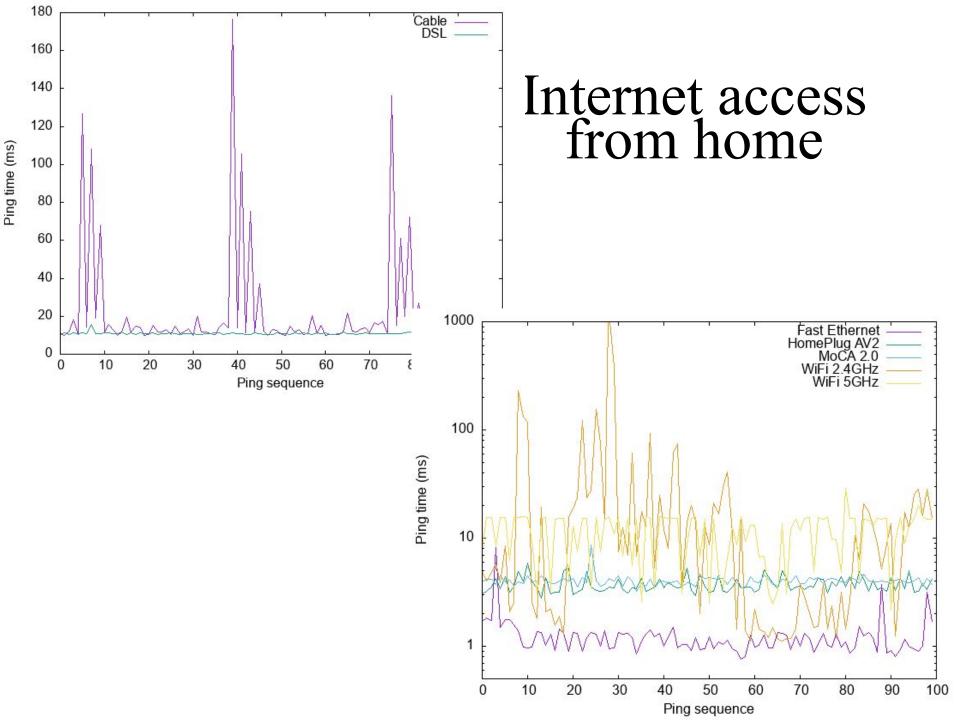
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802.11ac: gigabit WiFi?

A neighborhood scan

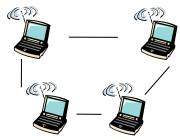
• 2.4GHz vs 5GHz

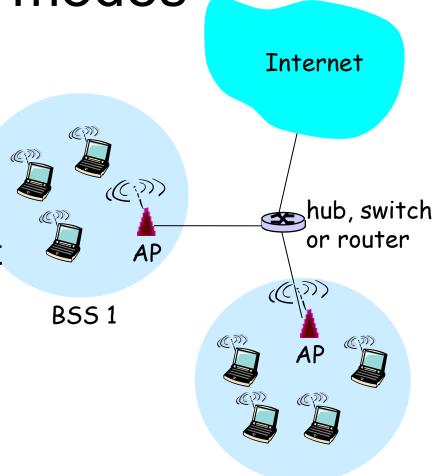




Operation modes

- Infrastructure mode
 - AP: access point
 - wireless station (sta)
 - BSS: basic service set
- Ad-hoc mode
 - no AP





BSS 2

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A hybrid network Coaxial cable Twisted pair DSL Cable Pass-thru Pass-thru Load-balancing Failover WAN2 WAN1 NAS Access Point Ethernet, MoCA, WDS

Operation procedures

- Association
 - channel scanning
 - beacon frame from AP
 - list and select AP to associate
- Authentication
 - network/user authentication
 - and possibly encryption
- Configuration
 - e.g., DHCP to configure network parameters

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Q: WEP (in)security

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Media access control

- Contention-free
 - PCF: point coordination function
 - e.g., AP
 - optional (not widely implemented)
- Contention-based
 - DCF: distributed coordination function
 - widely implemented
 - CSMA/CA

DCF

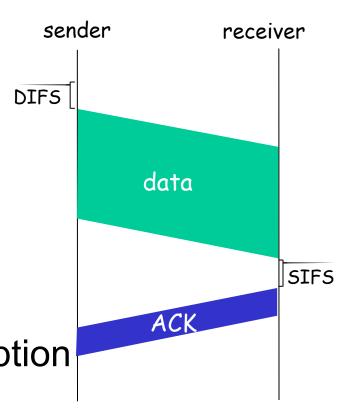
- Like Ethernet, uses CSMA:
 - random access
 - carrier sense: don't collide with ongoing transmission
- Unlike Ethernet:
 - no collision detection transmit all frames to completion
 - acknowledgment because without collision detection, you don't know if your transmission collided or not
- Why no collision detection?
 - difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
 - can't sense all collisions in any case: hidden terminal, fading
- Goal: avoid collisions: CSMA/C(ollision)A(voidance)

CSMA/CA

CSMA

CA: collision avoidance

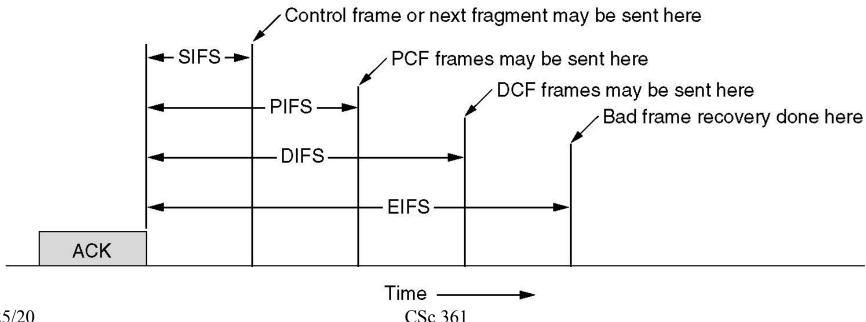
- if idle for DIFS, transmit
- if busy, random back-off
 - count down when idle
 - transmit when count to 0
- if no ack, collision or corruption
 - exponential backoff
 - CW: contention window CSc 361



Inter-frame spacing

SIFS: control frames or fragments

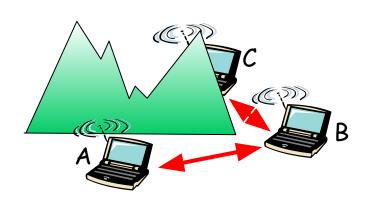
DIFS: DCF frames



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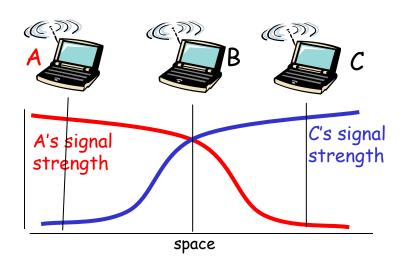
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Hidden terminal problems



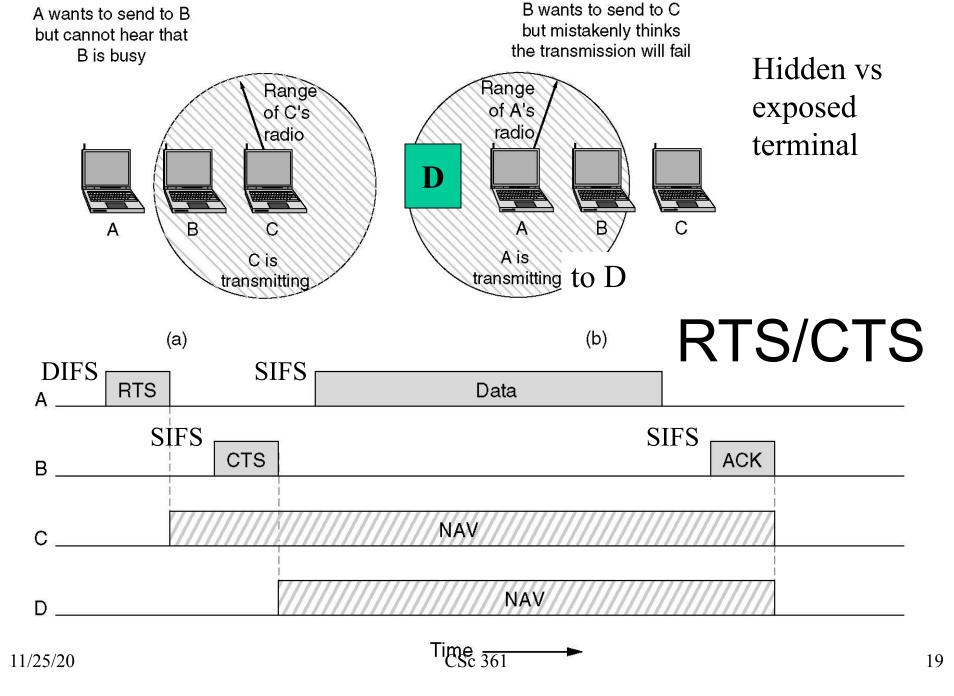
Hidden terminal problem

- B, A hear each other
- B, C hear each other
- A, C can not hear each other means A, C unaware of their interference at B



Signal fading:

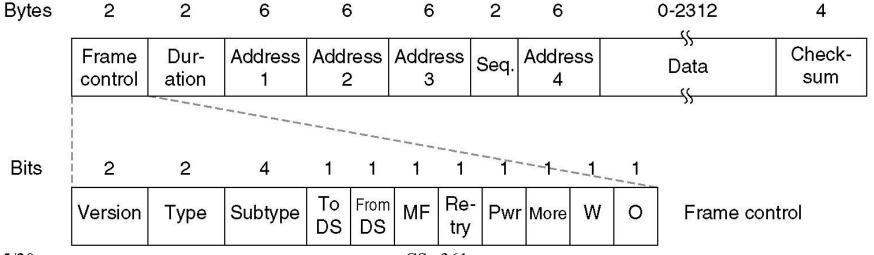
- □ B, A hear each other
- □ B, C hear each other
- A, C can not hear each other interfering at B



^{*} check whether RTS/CTS can improve your WiFi performance at home

802.11 frame

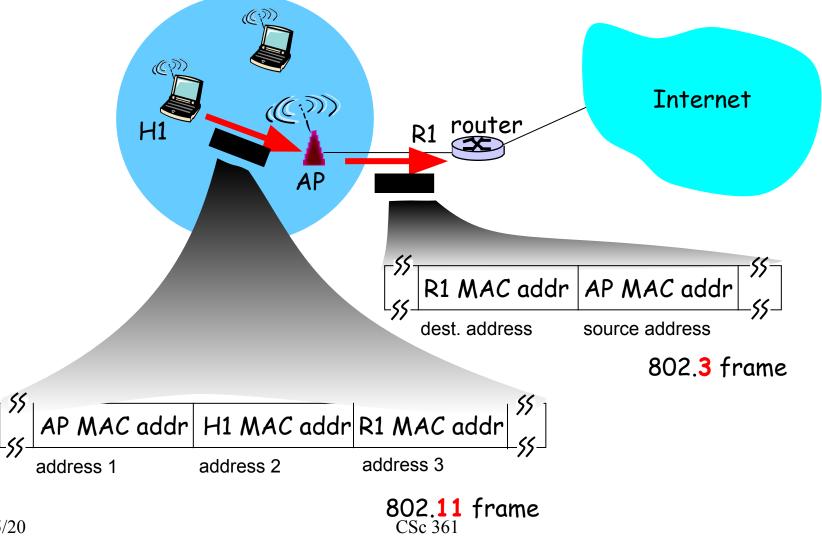
- Frame control
- Duration: NAV (network allocation vector)
- Addresses: dst, src, receiving, transmitting



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802.11 frame: addressing



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This lecture

- Wireless LAN
 - CSMA/CA
 - why not CSMA/CD
 - RTS/CTS
 - hidden vs exposed terminal
 - IEEE 802.11 family
 - 802.11a/b/g/n/ac/ax
- Explore further
 - CSC463: Wireless and Mobile Networks

Next few lectures

- This Friday (Nov 27)
 - M3 preparation
 - Guest lecture: UVicNet in Days of CoViD-19
 - by Ron Kozsan, UVic Director of Infrastructures
- Next Tuesday (Dec 1): Interworking
 - Put all things together
 - DNS (name -> IP), ARP (IP -> Ethernet), etc
- Next Wednesday (Dec 2): Term Review
- Next Friday (Dec 4): M3

Teaching from Home: Network Support Perspectives

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http://tinyurl.com/tfhnsp

1. Pedagogical challenges ▼

