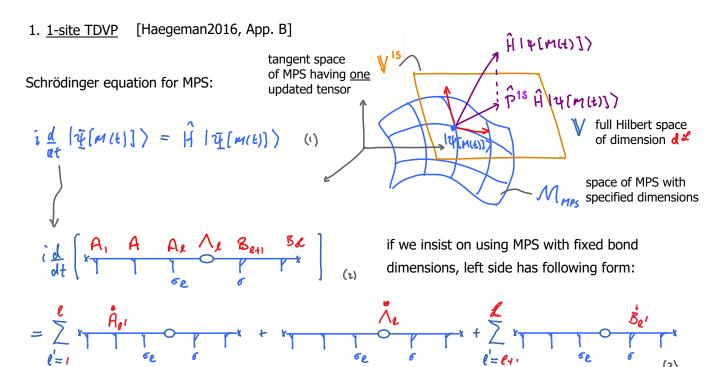
We consider time evolution using 'time-dependent variational principle' (TDVP)



Each term differs from $|\Psi(t)\rangle$ by precisely one site tensor or one bond tensor, so left side is a state in the tangent space, $|\Psi(t)\rangle$. But right side of (1) is <u>not</u>, since since $|\Psi(t)\rangle$ can have larger bond dimensions than $|\Psi(t)\rangle$.

So, project right side of (1) to
$$V^{15}$$
: $i \frac{d}{dt} \left| \overline{\psi} \left(m(t) \right) \right\rangle \approx P^{15} \left| \overline{\psi} \left(m(t) \right) \right\rangle$ tangent space approximation

Left and right sides of (4) are structurally consistent. To see this, consider bond ℓ

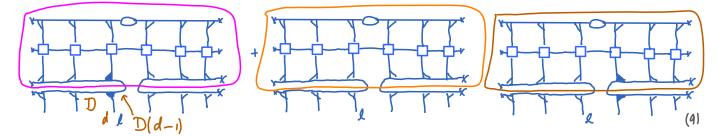
Left side of (4) contains:

$$\frac{d}{dt} \frac{A_{\ell} \wedge_{\ell} \beta_{\ell i_1}}{D |_{D} |_{D$$

Right side of (4) requires tangent space projector. Consider its form (TS-I.5.25):

$$P^{1S} = \frac{\ell'}{\tilde{\ell}} + \frac{1}{\tilde{\ell}} + \frac{1}{\tilde$$

The three terms with $\bar{\ell} = \ell$, $\ell' = \ell$, $\bar{\ell} = \ell + \ell$, applied to $\bar{H} \setminus \bar{\Psi}(\ell)$, yield



matching structure of (7). Thus, P^{ls} , applied to $H(\Psi(l))$, yields terms of precisely the right structure!

To integrate projected Schrödinger eq. (4), we write tangent space projector in the form (TS.5.26):

$$P^{IS} = \sum_{\ell=1}^{R} \frac{1}{\ell} \left(\frac{1}{\ell} \right) \left(\frac{1}{\ell} \right)$$

and write (4) as

 $i \stackrel{\mathcal{L}}{\geq} \frac{1}{|\mathcal{L}|} \stackrel{\mathcal{L}}{=} \frac{1}{|\mathcal{L}|} \stackrel{$

Right side is sum of terms, each specifying an update of one ψ_{ℓ}^{L} or ψ_{ℓ}^{L} on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$i \stackrel{Ce}{\leftarrow} := \stackrel{\wedge e}{\leftarrow} H_{\ell}^{1s}$$
 $i \stackrel{\wedge e}{\leftarrow} := - \stackrel{\wedge e}{\leftarrow} H_{\ell}^{b}$
 (12)

In site-canonical form, site ℓ involves two terms linear in C_{ℓ} : $i \in \ell$ $(t) = H_{\ell}^{15} C_{\ell} (t)$

Their contribution can be integrated exactly: replace $C_{\ell}(t)$ by $C_{\ell}(t+\tau) = e^{-it|_{\ell}^{15}\tau}$ $C_{\ell}(t)$ (14) forward time step

In bond-canonical form, site ℓ involves two terms linear in Λ_{ℓ} : $-i\Lambda_{\ell}(t) = \Pi_{\ell}^{b}\Lambda_{\ell}(t)$ (5)

Their contribution can be integrated exactly: replace $\bigwedge_{\ell} (t)$ by $\bigwedge_{\ell} (t-\tau) = e^{i + \frac{b}{\ell} \tau} \bigwedge_{\ell} (t)$ (6)

11-lecture-TDVP Page 2

In practice, $e^{-iH_{\ell}^{15}\tau}$ and $e^{iH_{\ell}^{b}\tau}$ are computed by using Krylov methods.

Build a Krylov space by applying \mathcal{H}^{13}_{ℓ} multiple times to \mathcal{C}_{ℓ} , set up the tridiagonal representation \mathcal{H}^{13}_{ℓ} in this basis, then compute the matrix exponential in this basis, and apply result to \mathcal{C}_{ℓ} . Likewise for \mathcal{H}^{13}_{ℓ} and \mathcal{M}_{ℓ} .

To successively update entire chains, alternate between site- and bond-canonical form, propagating forward or backward in time with $H_{\ell}^{l,s}$ or H_{ℓ}^{b} , respectively:

 $C_{\ell}(t) := \mathcal{T}$ $D_{\ell}(t)B_{\ell}(t) \dots B_{\ell}(t)$ 1. Forward sweep, for $l = 1, \dots, l = 1$, starting from (17)

$$C_{\varrho}(\xi) B_{\varrho+1}(\xi)$$

$$\downarrow^{t+\tau} \qquad A$$

$$\downarrow^{lS} \qquad C_{\varrho}(\xi+\tau) B_{\varrho+1}(\xi)$$

$$= A_{\varrho}(\xi+\tau) A_{\varrho}(\xi+\tau) B_{\varrho+1}(\xi)$$

$$= A_{\varrho}(\xi+\tau) A_{\varrho}(\xi+\tau) B_{\varrho+1}(\xi)$$

$$= A_{\varrho}(\xi+\tau) A_{\varrho}(\xi+\tau) A_{\varrho}(\xi+\tau) B_{\varrho+1}(\xi)$$

$$= A_{\varrho}(\xi+\tau) C_{\varrho+1}(\xi)$$

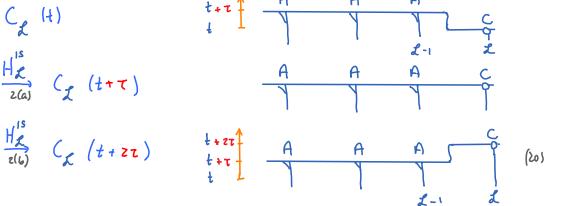
$$\downarrow^{t+\tau} A_{\varrho}(\xi+\tau) A_{\varrho$$

until we reach last site, and MPS described by

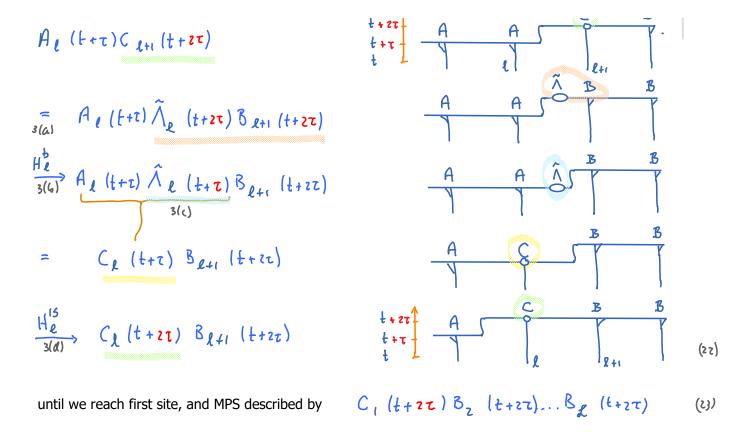
$$A_1(t+z) \dots A_{L_{2}}(t+z) C_{L}(t) \tag{19}$$

(18)

2. Turn around: (+) $\frac{H_{\mathcal{L}}^{13}}{2.6a}$ $(t+\tau)$



3. Backward sweep, for $\ell = \mathcal{L} - 1, \dots, 1$, starting from A_{ℓ} $(\ell + 2) \dots A_{\ell-1}$ $(\ell + 2) C_{\ell}$ $(\ell + 2) C_{\ell}$ $(\ell + 2) C_{\ell}$ $(\ell + 2) C_{\ell}$



The scheme described above involves 'one-site updates'. This has the (major!) drawback (as in one-site DMRG), that it is not possible to dynamically explore different symmetry sectors. To overcome this drawback, a 'two-site update' version of tangent space methods can be set up [Haegemann2016, App. C].

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!

A scheme for doing 1-site TDVP while nevertheless expanding bonds, called 'controlled bond expansion (CBE), was proposed in [Li2022] (see next lecture!).

The construction of tangent space V^{13} and its projector P^{13} can be generalized to n sites [Gleis2022a].

We focus on N = 2 (but general case is analogous). Define space of 2-site variations:

$$V^{2,1}$$
 $V^{2,5}$ = span of all states $V^{2,5}$ differing from $V^{2,5}$ on precisely 2 neighboring sites

$$= span \left\{ \left| \vec{\mathcal{V}} \right| \right\} = r + \frac{2 \text{ sites}}{\left(\left| \frac{1}{2} \right| + 1 \right)} \left\{ \left| \frac{1}{2} \right| \right\}$$
 (1)

formal definition: =
$$span \left\{ im \left(P_{\ell}^{2s} \right) \mid \ell \in [1, 2-1] \right\}$$

$$C_{image}$$
 (2)

Recall:

Recall:
$$\frac{|\text{ocal 2s projector:}}{|\text{local } 2\text{s projector:}} = \frac{1}{(\text{TS-I.4.9})}$$
(3)

Global 2s projector $\stackrel{?}{P}^{2s}$, such that $\stackrel{V}{V}^{2s} = i_{m} (\stackrel{?}{P}^{2s})$, can be found with a Gram-Schmidt scheme analogous to our construction of $\stackrel{?}{P}^{1s}$, see [Gleis2022a]:

compare (TS-I.5.22)
$$P^{2S} := \sum_{\ell=1}^{2} P_{\ell}^{2S} + P_{\ell}^{2S} + \sum_{\ell=1}^{2} P_{\ell}^{2S} \text{ for any } \ell \in [1, \mathcal{L}_{-1}]$$

$$P_{\ell}^{2S} - P_{\ell}^{1S} = P_{\ell, \ell+1}^{KD}$$

$$P_{\ell}^{2S} - P_{\ell}^{1S} = P_{\ell, \ell+1}^{KD}$$

$$p^{2S} = \sum_{\ell=1}^{\ell'-1} \frac{1}{\ell} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+2} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell}{\ell+1} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell} \\ \frac{\ell'-1}{\ell} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \frac{\ell'-1}{\ell'} \\ \frac{\ell'-1}{\ell'} \end{array} \right\} + \frac{\ell$$

All summands are mutually orthogonal, ensuring that $(P^{25})^2 = P^{25}$, and that $P^{25} = P_{\ell}^{25} = P_{\ell}^{25}$

Alternative expression: compare (TS.5.26)

$$P^{25} = \sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-2} P_{\ell+1}^{15} = \sum_{\ell=1}^{l-1} \frac{1}{|I|} \left(\sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-2} P_{\ell+1}^{15} \right) \left(\sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-1} P_{\ell+1}^{15} \right) \left(\sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-1} P_{\ell+1}^{15} - \sum_{\ell=1}^{l-1} P_{\ell+1}^{15} \right) \left(\sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-1} P_{\ell+1}^{15} - \sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-1}$$

This projector is used for 2-site TDVP (see TS-II.3)

Orthogonal n-site projectors

For any given MPS $|\psi(M)\rangle$, full Hilbert space of chain can be decomposed into mutually orthogonal subspaces:

$$V = V_1 \otimes \cdots \otimes V_k = \bigoplus_{N=0}^{k} V^{NL}$$
(8)

with
$$\bigvee^{\circ 1} := \bigvee^{\circ s} := \operatorname{Span} \{ | \Psi \rangle \}$$
 (9)

'irreducible'
$$\bigvee^{N\perp}$$
 is complement of $\bigvee^{(n-1)}$ in $\bigvee^{N} = \bigvee^{(n-1)} \oplus \bigvee^{N\perp}$ (6)

Solution = span of states differing from $| \psi \rangle$ on $| \psi \rangle$ contiguous sites, not expressible through subsets of $| \psi \rangle$ sites

Correspondingly, identity can be decomposed as:

$$1_{V} = 1_{A}^{OK} = \sum_{N=0}^{K} P^{NL}$$
completeness

orthogonality

(1)

where $P^{\perp N}$ is defined as the projector having $V^{N\perp}$ as image: $I_{N}(P^{N\perp}) \approx V^{N\perp}$ (12)

$$N \ge 1 := P^{NS} \left(1_V - P^{(n-1)S} \right) = P^{NS} - P^{(n-1)S}$$
 (14)

Consider n=1:

since
$$V^{(n-1)5}$$
 $\subset V^{n5} \Rightarrow im(P^{(n-1)5}) \subset im(P^{n5})$
 $\Rightarrow P^{n5} P^{(n-1)5} = P^{(n-1)5}$

(15)

(16)

$$= \sum_{\ell=1}^{\ell'} \frac{1}{\ell} \frac{1}{$$

- 4444

choose l'= L

$$= \sum_{\ell=1}^{l} P_{\ell,\ell+1}^{bk}$$
 projects onto all 1-site variations orthogonal to

$$= \sum_{\ell=1}^{\ell-1} \left(\begin{array}{c} x \\ \ell \end{array} \right) \left[\begin{array}{c} x \\ \ell \end{array} \right] \left[\begin{array}{c} x \\ \ell$$

$$- \underset{\ell}{\overset{\text{def}}{\longrightarrow}} \left(\begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right)$$

$$= \sum_{\ell=1}^{L-1} \frac{1}{\ell + \ell} = \sum_{\ell=1}^{L-1} P_{\ell,\ell+1}^{DD}$$

very important result! (26)

[Haegeman2016, Sec. V & App. C]

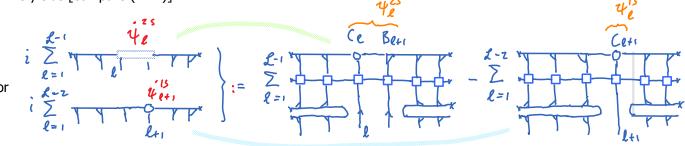
2-site tangent space methods are analogous to 1-site methods, but use a 2-site projector. There is a conceptual difference, though: the main reason for using 2-site schemes is that they allow sectors with new quantum numbers to be introduced if the action of H requires this. However, states with different ranges of quantum numbers live in different manifolds, hence this procedure 'cannot easily be captured in a smooth evolution described using a differential equation. However, like most numerical integration schemes, the aforementioned algorithm is intrinsically discrete by choosing a time step, and it poses no problem to formulate an analogous two-site algorithm'. [Haegeman2016, Sec. V]. In other words: the tangent space approach is conceptually not as clean for the 2-site as for the 1-site scheme.

Schrödinger equation, projected onto 2-site tangent space, now takes the form

$$i \frac{d}{at} | \Psi[M(t)] \rangle = \hat{\rho}^{2s} \hat{H} | \Psi[M(t)] \rangle$$

$$\hat{P}^{ZS} = \sum_{\ell=1}^{2-1} \frac{1}{\ell+1} \left| \frac{1}{\ell+1} - \sum_{\ell=2}^{2-1} \frac{1}{\ell+1} \right|$$

This yields [compare (1.11)]:



Right side is sum of terms, each specifying an update of one ψ_{ℓ}^{ts} or ψ_{ℓ}^{ts} on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$i \dot{\psi}_{\ell}^{2S} := \begin{array}{c} \psi_{\ell}^{2S} \\ \psi_{\ell}^{2S} \\ \psi_{\ell}^{2S} \end{array}, \quad \frac{i \dot{\psi}_{\ell+1}^{1S}}{\ell} := - \begin{array}{c} \psi_{\ell+1}^{1S} \\ \psi_{\ell+1}^{2S} \\ \psi_{\ell+1}^{2S} \end{array}$$

Right side is sum of terms, each linear in a factor appearing on the left. Can be integrated one site at a time:

In 2-site-canonical form, site
$$\ell$$
 involves two terms linear in Υ_{ℓ}^{rs} : $i \Psi_{\ell}^{rs}(t) = H_{\ell}^{rs} \Upsilon_{\ell}^{rs}(t)$ (1)

Their contribution can be integrated exactly: replace
$$\psi_{\ell}^{2s}(t)$$
 by $\psi_{\ell}^{2s}(t+\tau) = e^{-i H_{\ell}^{2s} \tau} \psi_{\ell}^{2s}(t)$ forward time step

In 1-site-canonical form, site
$$\ell$$
 involves two terms linear in $\Psi_{\ell+1}^{15}$: $i \Psi_{\ell+1}^{15}(t) = - H_{\ell+1}^{15} \Psi_{\ell+1}^{15}(t)$ (12)

Their contribution can be integrated exactly: replace
$$\psi_{\ell+1}^{(s)}(t)$$
 by $\psi_{\ell+1}^{(s)}(t-\tau) = e^{iH_{\ell+1}^{(s)}\tau} \psi_{\ell+1}^{(s)}(t)$ (3)

Their contribution can be integrated exactly: replace
$$\psi_{\ell+1}^{(s)}(t)$$
 by $\psi_{\ell+1}^{(s)}(t-\tau) = e^{iH_{\ell+1}^{(s)}\tau} \psi_{\ell+1}^{(s)}(t)$ backward(!) time step

To successively update entire chains, alternate between 2-site- and 1-site-canonical form, propagating forward or backward in time with H_{ℓ}^{2s} or H_{ℓ}^{2s} , respectively (analogously to 1-site scheme).

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!