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$LCS(x, y)$

1. $a \leftarrow \text{length}[x]$

$b \leftarrow \text{length}[y]$

for $x \leftarrow 0$ to a do

for $y \leftarrow 0$ to b do.

if ($x=0$ or $y=0$)

$\text{length}[LCS[i, j]] \leftarrow 0$

else if $x[i-1] = y[j-1]$

$\text{length}[LCS[i, j]] \leftarrow \text{length}[LCS[i-1, j-1]] + 1.$

else

$\text{length}[LCS[i, j]] \leftarrow \max(\text{length}[LCS[i-1, j]], \text{length}[LCS[i, j-1]])$

$\text{pointer} \leftarrow \text{length}[LCS[a, b]]$

$i \leftarrow a$

$j \leftarrow b.$

while $i > 0$ and $j > 0.$

do if $x[i-1] = y[j-1]$

$\text{seq}[\text{pointer} - 1] = x[i-1]$

$i \leftarrow i-1$

$j \leftarrow j-1$

pointer \leftarrow pointer - 1

else if $\text{lengthLSS}[i-1, j] > \text{lengthLCS}[i, j-1]$

$i \leftarrow i-1$

else $j \leftarrow j-1$

$z \leftarrow ''$

for $u=0$ to $\text{length}[seq]-3$ do

if $seq[u]=x$ and $seq[u+1]=y$ and $seq[u+2]=y$

$u \rightarrow u+2$

else

$z = z + seq[u]$

return z .

Space complexity = $O(n)$

Time complexity = $O(n \cdot n)$

2. We generate all the subsequences of a sequence, then we compare each subsequence with other sequences to check if they are common, and then find the subsequences of maximum length.

then we have a string of length n . subsequences of length $1, 2, 3, \dots, n$, so we can use permutation and combination.

$$\text{sub-sequences} = {}^nC_1 + {}^nC_2 + \dots + {}^nC_n.$$

So a string of length n has $2^n - 1$ different possible subsequences since we don't consider the subsequence with length 0.

to check if a subsequence is common to both the string or not.

then $O(n)$ time.

So total time needed to check all the sub-sequences $= n \cdot 2^n - 1$.

time complexity by brute force method is $O(n \cdot 2^n - 1)$