

① partition (A, p, q)

when $x \leftarrow A[p]$

$i \leftarrow p$

for $j = p+1$ to r

if $A[j] \leq x$

$i = i+1$

exchange $A[i]$ with $A[j]$

exchange $A[i]$ with $A[p]$

return i .

②. Inversions is that a sorted array has no inversions in it. since every element should be smaller than everything coming after it and larger than everything coming before it.

$$\begin{aligned} T(n) &= 1\% \cdot \Theta(n^2) + 99\% \cdot \Theta(n^2) \\ &= \Theta(n^2) \end{aligned}$$

③. best-case: i/sort.

$$O(n) + O(n \log n) + O(n-r) \\ = O(n)$$

worst-case.

$$O(n) + O(r^2) + O((n-r)^2) \\ = O(n^2)$$

average-case.

$$O(n) + O(n \log n) + O((n-r)^2) \\ = O(n^2)$$

(4) best case. when the data goes up, $O(n)$.

worst-case.

$$T(n) = \max_{1 \leq r \leq n} \{ T(r-1) + O(n-r)^2 + O(n) \}$$

$$T(n) = O(n^2) \leq Cn^2 \quad \text{when } C > 0.$$

$$\begin{aligned}
T(n) &= \max_{1 \leq r \leq n} \{ T(r-1) + a(n-r)^2 + an \} \\
&\leq \max_{1 \leq r \leq n} \{ C(r-1)^2 + a(n-r)^2 + an \} \\
&= \max \{ a(n-1)^2 + an, C(n-1)^2 + an \} \\
&= C(n-1)^2 + an \\
&= Cn^2 - 2cn + 1 + an \\
&= Cn^2
\end{aligned}$$

Worst case . $T(n) = \frac{1}{n} \sum_{r=1}^n \{ T(r-1) + O(n-r)^2 + O(n) \}$

$$T(n) = O(n^2) \leq Cn^2 \text{ for } C > a$$

$$\begin{aligned}
T_{\text{Avg}}(n) &= \frac{1}{n} \sum_{r=1}^n \{ T(r-1) + a(n-r)^2 + an \} \\
&= \frac{1}{n} \sum_{r=1}^n \{ C(r-1)^2 + a(n-r)^2 + an \} \\
&= \frac{1}{n} \left[\sum_{r=1}^n C(r-1)^2 + \sum_{r=1}^n a(n-r)^2 + Cn^2 \right] \\
&= \frac{1}{n} \left[(a+c) \int_0^n x^2 dx + an^2 \right] \\
&= \frac{1}{n} \left[(a+c) \frac{n^3}{3} + an^2 \right]
\end{aligned}$$

$$= \frac{1}{n} \left(\frac{an^3}{3} + \frac{bn^3}{3} + an^2 \right)$$

$$= \frac{an^2}{3} + \frac{bn^2}{3} + an^2$$

$$= O(n^2)$$