

Binary Search

Sunday, March 29, 2020 8:14 AM

Given a sorted list of elements/numbers, find whether a **key** occurs in the list or not. If the **key** is found, its index in the list, otherwise -1, is returned.

Algorithm:

1. Find the middle element of the **current list**.
2. If the middle element is equal to the **key**, return the key index.
3. If the middle element is greater than the **key**, make current list equal to the lower half of the **current list**. Repeat the Algorithm (from step 1.)
4. If the middle element is lesser than the **key**, make current list equal to the upper half of the **current list**. Repeat the Algorithm (from step 1.)
5. Otherwise, the **key** is not found.

Pseudo-Code

BinarySearch(Arr, key)

// Arr is sorted in the ascending order.

1. let start = 1
2. let end = Arr.length+1
3. while start != end
4. let mid = start + (end - start)/2
5. // assume integer division
6. if (Arr[mid] == key)
7. return mid
8. if (Arr[mid] < key)
9. start = mid + 1
10. else if (Arr[mid] > key)
11. end = mid
12. return -1

Arr: 1 3 4 6 8 12 34 56 67

Key: 67

No. Iteration	start	end	mid	Arr[mid]
1	1 6	10	5	8
2	6 9	10	8	56
3	9	10	9	67

Arr: 1 3 4 6 8 12 34 56 67

Key: 75

No. Iteration	start	end	mid	Arr[mid]
1	1 6	10	5	8
2	6 9	10	8	56
3	9 10	10	9	67
4	10	10		

Arr: 1 3 4 6 8 12 34 56 67

Key: 34

No. Iteration	start	end	mid	Arr[mid]
1	1 6	10	5	8
2	6	10 8	8	56
3	6	8	7	34

Worst Case Analysis:

Pseudo-Code

BinarySearch(Arr, key)

// Arr is sorted in the ascending order.

	cost	WCF
1. let start = 1	c1	1
2. let end = Arr.length+1	c2	1
3. while start != end	c3	log(n) (+ 1)
4. let mid = start + (end - start)/2	c4	log(n)
5. // assume integer division		
6. if (Arr[mid] == key)	c5	log(n)
7. return mid	c6	1 (0)
8. if (Arr[mid] < key)	c7	log(n)
9. start = mid + 1		
10. else if (Arr[mid] > key)		
11. end = mid		
12. return -1	c8	1

// Either line 7 or line 12 will be executed

Worst case arises when either the key does not occur in the list or first/last element of the list is the key.

$$\begin{aligned}T(n) &= (c3+c4+c5+c7)\log(n) + (c1+c2+c3+c6+c8) \\&= c10\log(n) + c11 \\&= \mathbf{O(\log(n))} \quad \mathbf{Logrithmic}\end{aligned}$$

1. for $l = 1$ to n
2. let $j = 1$
3. let $a = 1$
4. let $b = 1$
5. while $j \leq n$
6. $a = 10 + 5$
7. $b = a * 3$
8. $j = j * 2$

	cost	Freq
1. for l = 1 to n	c1	n+1
2. let j = 1	c2	n
3. let a = 1	c3	n
4. let b = 1	c4	n
5. while j <= n	c5	n(log(n)+1)
6. a = 10+5	c6	nlog(n)
7. b = a * 3	c7	nlog(n)
8. j = j * 2	c8	nlog(n)