CHAPTER 7

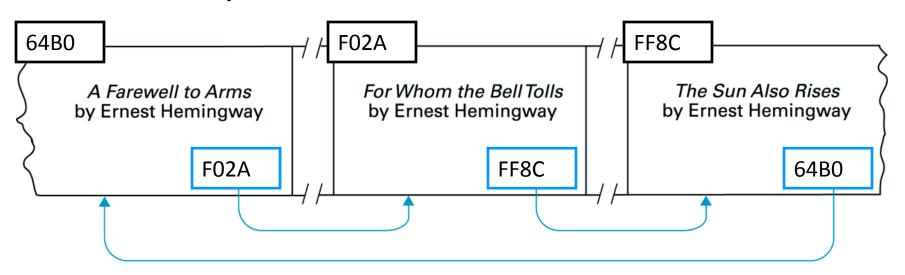
Data Structures

- Abstractions of the actual data organization in main memory
- Allow users to perceive data as 'logical units' (e.g.: arrangement in rows and columns)

7.1: Data Structure Basics: Pointers

• Pointers:

- pointer = location in memory that contains the address of another location in memory
- so: pointer *points* to data positioned elsewhere in memory



7.1: Static versus Dynamic Data Structures

• Static:

 shape & size of structure does not change over time

- example in C: int Table[2][9];

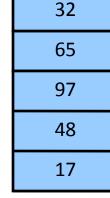
Table:

• Dynamic:

- shape & size may change

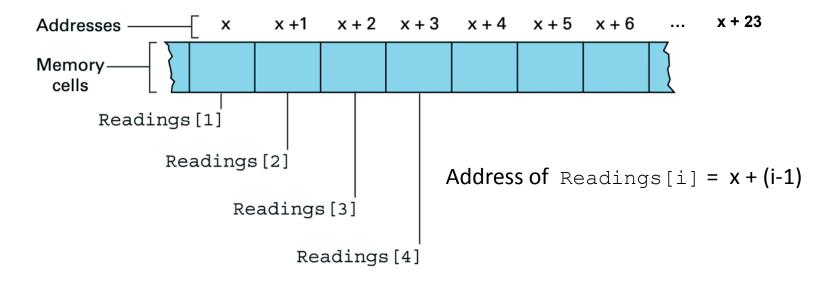
– example: Stack

Stack:



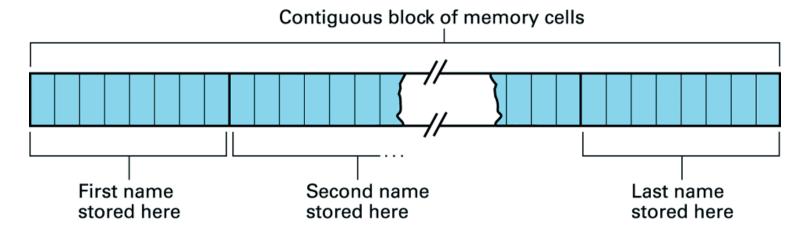
7.2: Arrays

- Example: to store 24 hourly temperature readings...
- ... a convenient storage structure is 1-D homogeneous array of 24 elements (e.g. in C: float Readings [24])
- In main memory:



7.3: Lists

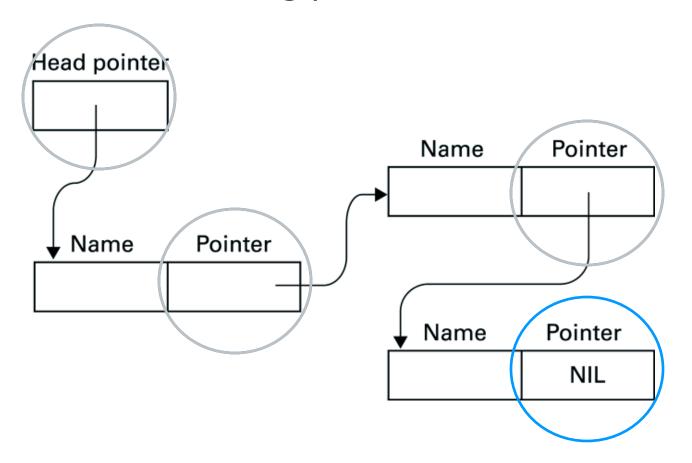
To store an ordered list of names we could use
 2-D homogeneous array (in C: char Names [10] [8])



- However:
 - addition & removal of names requires expensive data movements!

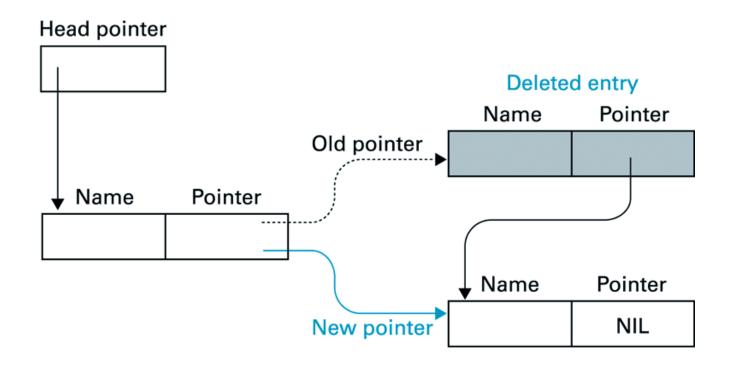
7.3: Linked Lists

 Data movements can be avoided by using a 'linked list', including pointers to list entries



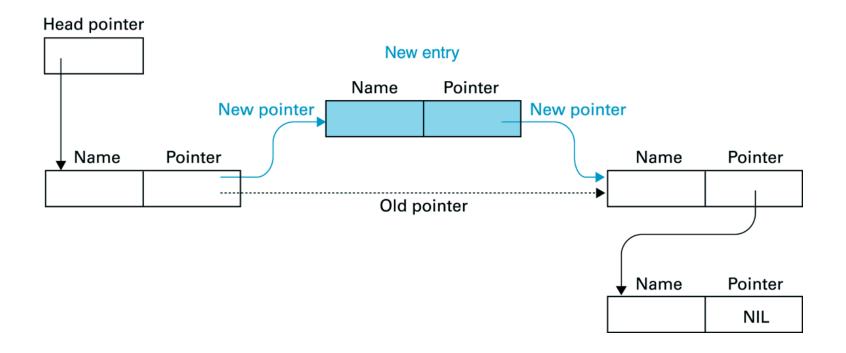
7.3: Deleting an Entry from a Linked List

 A list entry is removed by changing a single pointer:



7.3: Inserting an Entry into a Linked List

- A new entry is inserted by setting pointer of
 - (1) new entry to address of entry that is to follow
 - (2) preceding entry to address of new entry:



7.4: Stacks

- Disadvantage of contiguous array structures:
 - insertion / removal requires costly data movements
- Still okay if insertion / removal restricted to end of array => stack (with push & pop operations)
- Typical use:

 Procedure invoking order

 Main program

 Procedure A

 Procedure B

 Procedure C

 Call A

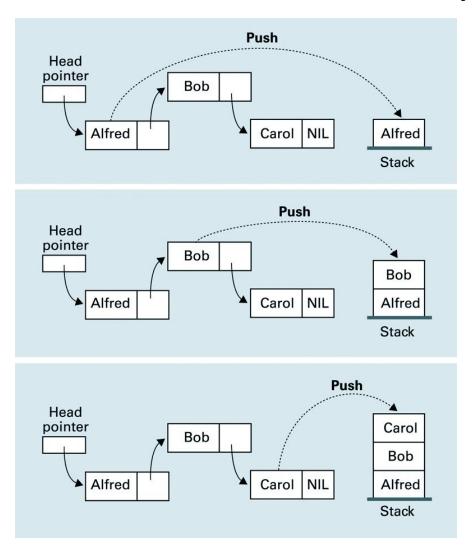
 Call B

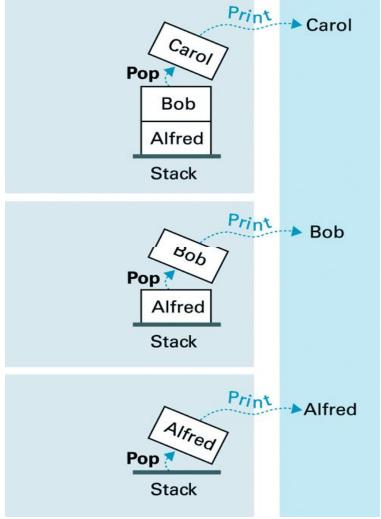
 Return from B

 Procedure C

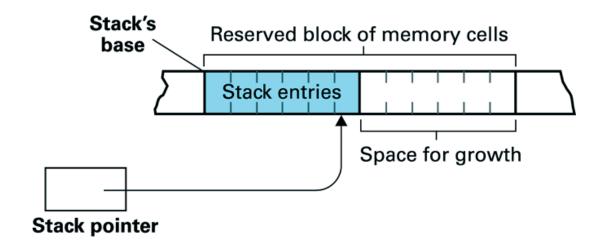
 Procedure returning order

7.4: Push / Pop (to print inverse linked list)





7.4: A Stack in Memory



- Here:
 - conceptual structure close to identical to actual structure in memory
- If maximum stack-size unknown:
 - pointers can be used (=> conceptual = actúal structure)

Chapter 7 - Data Structures: Conclusions

- Pointers:
 - basic aid in definition of dynamic data structures
- Often used data structures:
 - Arrays
 - Lists
 - Stacks

– ...

CHAPTER 8

File Structures

Reference: Computer Science an Overview

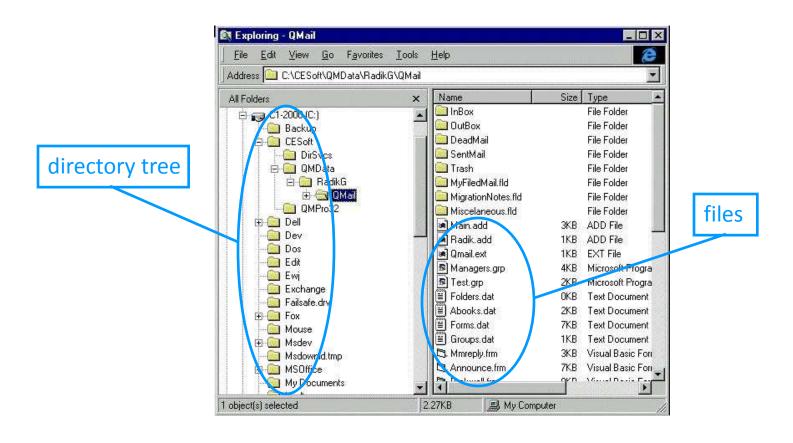
Author: J. Glenn Brook Shear

6th Edition

- Abstractions of the actual data organization on mass storage
- Again: differences between conceptual and actual data organization

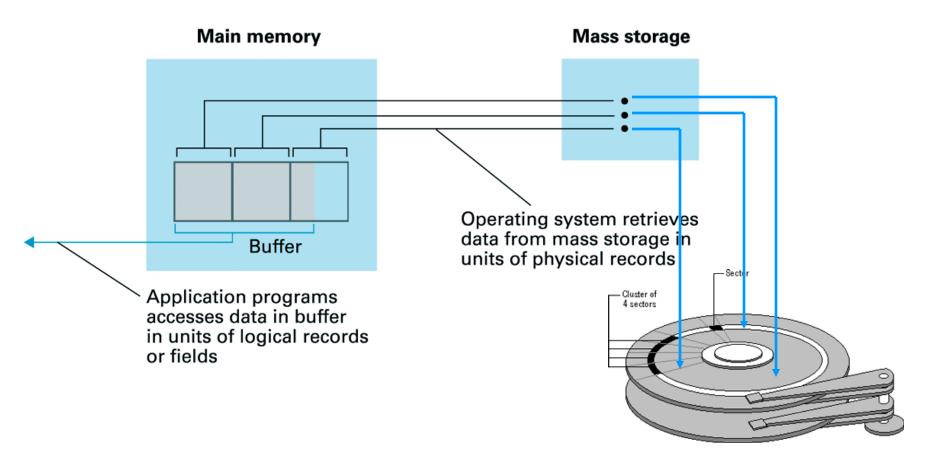
8.1: Files, Directories & the Operating System

- OS storage structure:
 - conceptual hierarchy of directories and files



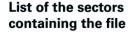
8.1: Files: Conceptual vs. Actual View

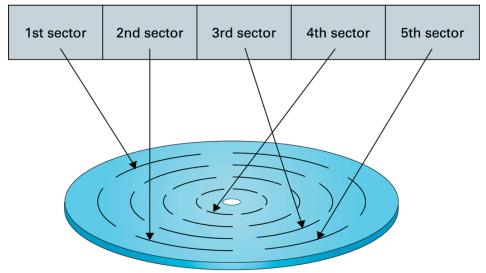
- View at OS-level is conceptual
 - actual storage may differ significantly!



8.2: Sequential Files

To 'remember' where data resides on disk, the
 OS maintains a list of sectors for each file



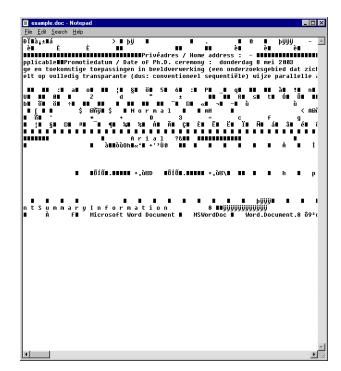


Disk sectors containing the file

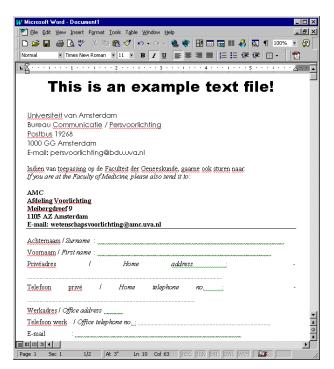
Result: sequential view of scattered set of data

8.2: Text Files

- Sequential file consisting of long string of encoded characters (e.g. ASCII-code)
 - But: character-string still interpreted by word processor!

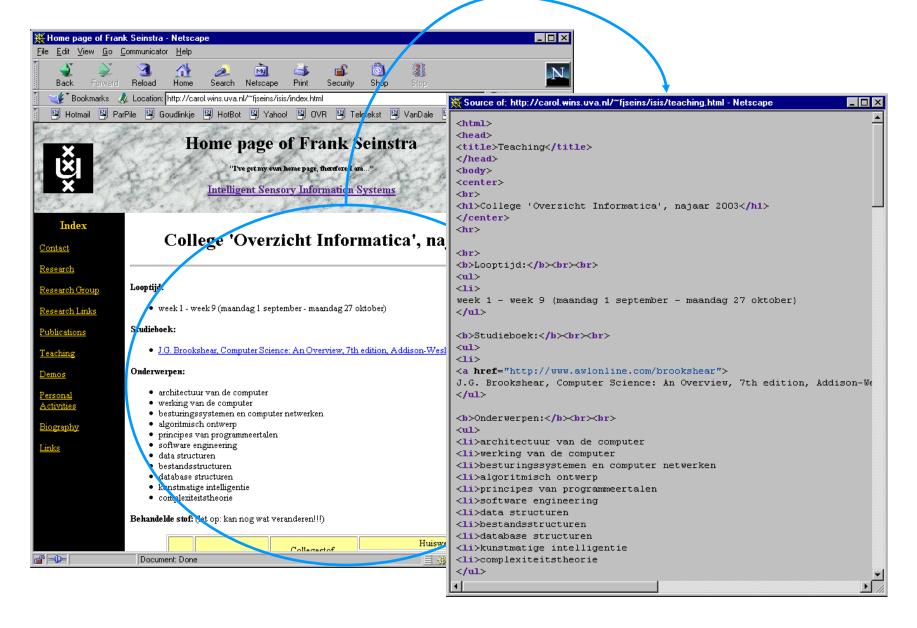


File in "Notepad"

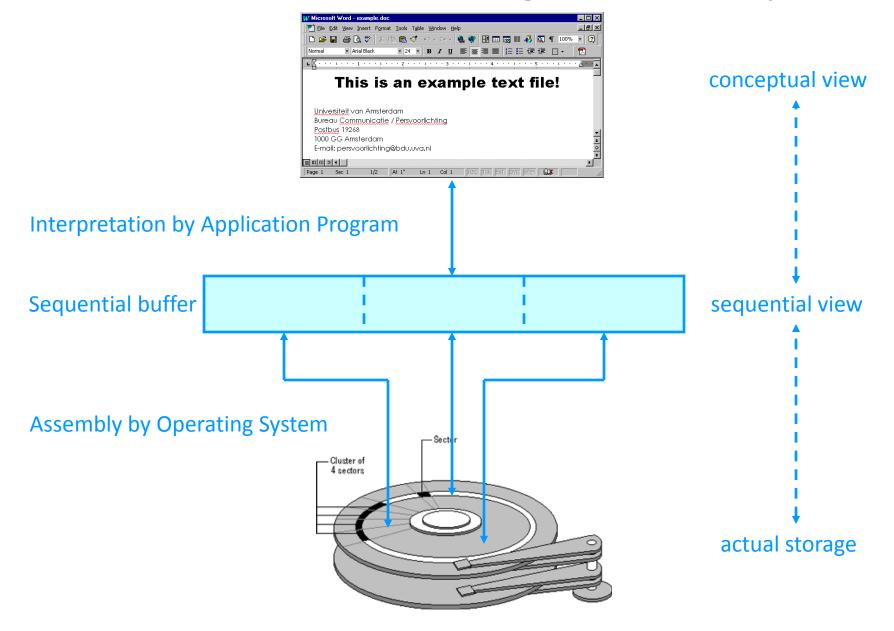


Same file in "MS Word"

8.2: Text files & Markup Languages (e.g. HTML)

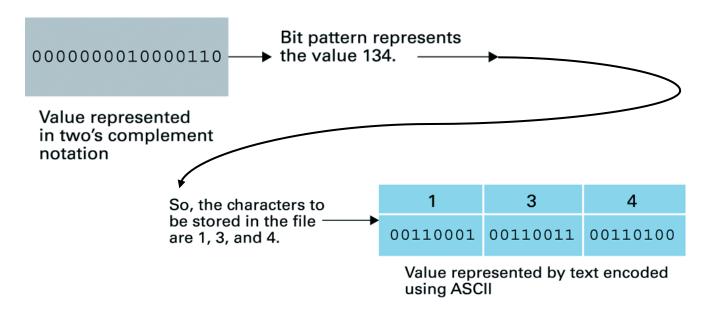


8.2: From actual storage to conceptual



8.2: Data Conversion

- When programming: note that data transfer to/from file may involve data conversion:
 - e.g., from two's complement notation to ASCII:



So: again it's about the interpretation of data

8.3: Quick File Access

- Disadvantage of sequential files:
 - no quick access to particular file data
- Two techniques to overcome this problem:
 - (1) *Indexing* or (2) *Hashing*
- Indexing:

Indexed File

/12N	12N67 John Smith		23-Jul-71	17,000.00	New York	
13C	08	Andrew White	27-Jun-70	24,500.00	Boston	
23G	19	Mary Jackson	5-Mar-39	41,000.00	San Francisco	
24X	17	Eleanor Tracy	17-Sep-63	9,635.00	Fort Lauderdale	
26X	28	Michael Flanagan	1-Nov-44	18,800.00	Washington	
32E	76	Glenn White	29-Feb-68	17,000.00	Detroit	
36Z0	05	Virginia Moore	27-Jun-70	32,000.00	San Francisco	
	: /	:	:	:	:	
	: /	:	:	:	:	
		i :	:	:	:	
	_					

loaded into main memory when opened

Index

12N67	location
13C08	location
23G19	location
24X17	location
26X28	location
32E76	location
36Z05	location
:	:
:	:
	/

keys

<u>Chapter 8 - Problem 10</u>

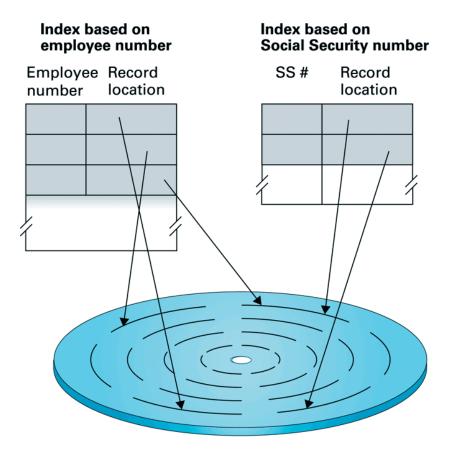
Why is a 'patient identification number' a better choice for a key field than the last name of each patient?

- If key unique:
 - additional sequential search never required

Patient's last name is not always unique

8.3: Inverted Files

Variation to (single) indexing: inverted file



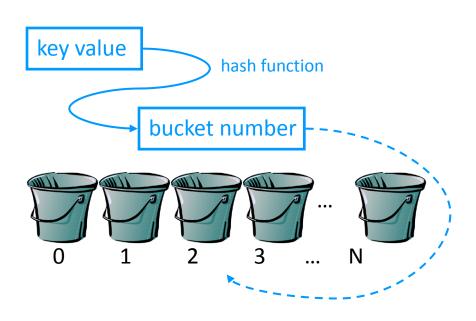
Records stored on disk

8.4: Hashing

- Disadvantage of indexing is... the index
 - requires extra space
- Solution: 'hashing'
 - finds position in file using a key value (as in indexing)...
 - ... simply by identifying location directly from the key

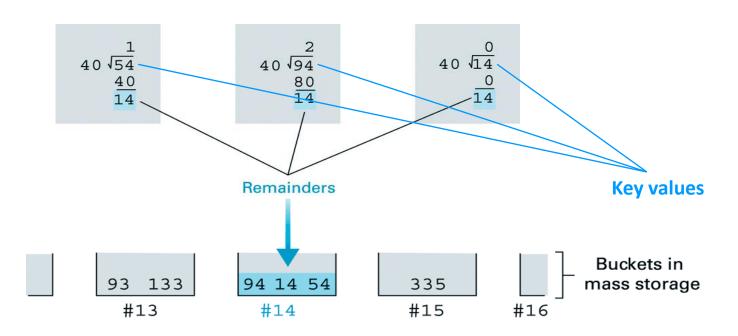
How?

 define set of 'buckets' & 'hash function' that converts keys to bucket numbers



8.4: Hash Function: Example

- If storage space divided into 40 buckets and hash function is division:
 - key values 14, 54, & 94 all map onto same bucket (collision)



Chapter 8 - File Structures: Conclusions

- File Structures:
 - abstractions of actual data organization on mass storage
- Changes of 'view':
 - actual storage -> sequential view by OS -> conceptual view presented to user
- Quick access to particular file data by
 - (1) indexing (many forms)
 - (2) hashing (requires no index, but requires bucket search!)

