1.

(a)
$$64k = 2^{16}$$

umber	of bits	for	block =	2 bits	s = 4 p	osition.		
umbor	of hita	f_{Or}	indox —	gizo c	f cocho	(block	form) -	8 hita

2 bit relational position

nu number of bits for index = size of cache (block form) = 8 bits. number of bits for tag = 6 bits.

- (b) 16 bits.
- (c) $2^{10}/4 = 2^8 = 256$ blocks.

6 bit tag

(d)
$$T_{access} = 100(h_1) + (1 - h_1)(1200) \le 120 \to 1080 \le 1100h_1 \to 1 \ge h_1 \ge 0.981818$$

8 bit index

2. size of the cache is 16 blocks.

requested address	tag	hit 1	hit 2	hit 3
0	0	X	✓	✓
21	1	X	✓	✓
23	1	\checkmark	\checkmark	\checkmark
35	2	X	\checkmark	\checkmark
76	4	X	\checkmark	✓
1	0	\checkmark	\checkmark	✓
66	4	\checkmark	\checkmark	\checkmark
80	5	X	\checkmark	✓
54	3	X	\checkmark	✓
36	2	✓	✓	✓
24	1	✓	✓	✓
23	1	✓	✓	✓
75	4	√	√	√

In the second and the third round all of the addresses are hitted because all of the requested addresses are less than 256.

Miss ratio =
$$6/42 = 1/7$$

3.

(a)

requested addresses	tag	hit	number of fulled	row
0	0	X	1	0
1	0	√	1	0
15	3	X	2	0
14	3	√	2	0
14	3	\checkmark	2	0
15	3	√	2	0
16	4	X	1	1
2	0	√	2	0
23	5	X	2	1
27	6	X	3	1
16	4	√	3	1
14	3	\checkmark	2	0
1	0	√	2	0
21	5	√	3	1
22	5	\checkmark	3	1
23	5	√	3	1
22	5	\checkmark	3	1
10	2	X	3	0
18	4	√	3	1
15	3	√	3	1
1	0	√	3	0
0	0	√	3	0
14	3	√	3	0
28	7	X	4	1
25	6	√	4	1

Hit ratio is equal to 0.72. Final state of cache:

0	1	2	3	null	null	null	null	8	9	10	11	12	13	14	15
16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31

(b,c) These two states are the same.

requested address	tag	hit
0	0	X
1	0	√
15	7	X
14	7 7 7	X ✓ ✓
14	7	√
15	7	√
16	8	X
2	1	X
23	11	X
27	13	X
16	8 7	√
14	1	✓ ✓
1	0	
21	10	X
22 23	11	X ✓ ✓ ✓ ✓
	11	√
22	11	
10	5	X 🗸
18	8	
15	7	√
1	0	√
0	0	√ √ √
14	7	√
28	14	X
25	12	X

The hit ratio is equal to 0.56

4. Time of not finding requested data in cache = 70ns. $T_{access} = 0.95*10 + 0.05*60 = 12.5$