EndTerm Project

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Data Integration and Synchronization Across Heterogeneous Systems

Hive, PostgreSQL, MongoDB

Introduction

In this project, we explore the concepts of data integration and synchronization across heterogeneous database systems. We work with three systems: Hive, PostgreSQL, and MongoDB. Each system stores the same dataset — a CSV file containing student course grades — in its native format. The schema includes five fields: Student ID, Course ID, Student Name, Mail, and Grade, with the composite key being the combination of Student ID and Course ID.

Our goal is to implement consistent CRUD functionality, focusing particularly on the GET (read) and SET (update) operations for the Grade field. Furthermore, we provide a mechanism to merge the state of one system with another using an operation log (oplog), thereby ensuring eventual consistency across all systems.

CRUD Operations and Supported Services

Hive

- Create: CREATE TABLE
- Read: SELECT Grade FROM table WHERE student_ID='...' AND course_ID='...'
- Update: INSERT OVERWRITE TABLE
- Delete: DROP TABLE

PostgreSQL

- Create: CREATE
- Read: SELECT Grade FROM table WHERE student_ID='...' AND course_ID='...'
- Update: UPDATE table SET Grade='...' WHERE student_ID='...' AND course_ID='...'
- Delete: DELETE

MongoDB

- Create: createCollection, insertOne, insertMany
- Read: findOne({ student_ID, course_ID }, { Grade })
- Update: updateOne({ student_ID, course_ID }, { \$set: { Grade } })
- Delete: deleteOne

System Architecture

The architecture consists of four Python modules:

- main.py
- hive_manager.py
- postgres_manager.py
- mongo_manager.py

Each manager module implements three primary methods: get(), set(), and merge().

If the (studentID, courseID) pair given from set(), get() operations from any system don't exist in the given dataset, a message will be displayed conveying the same.

Oplog Structure

Each system maintains an operation log (oplog) that records all read and write operations. The oplog is structured as follows:

```
new_database=# select * from oplogs;

timestamp | operation | student-ID | course-id | new_grade

2025-04-30 06:43:08.498 | SET | SID1033 | CSE016 | C
2025-04-30 06:43:29.349 | SET | SID1033 | CSE016 | B+
2025-04-30 06:44:21.796 | GET | SID1033 | CSE016 | X
(3 rows)
```

- Timestamp
- Operation Type: GET or SET
- Student ID
- Course ID
- Grade: Set value or X for GET

Only SET operations are considered during merge operations for synchronization. We also log the merge operations, but only if the operation requires a new value to overwrite an existing value in my Database. This is logged as a SET operation with the old timestamp from the merged database.

Merge Functionality

The merge() function in each system accepts another system as an argument. For example, hive.merge(sql) merges the state of the PostgreSQL database into Hive.

Implementation Details

- A key-value store is created with composite keys (student_ID, course_ID).
- The oplog of the calling system is traversed to populate the key-value store with the latest SET values based on timestamps.
- The structure of key is a tuple containing the studentID & courseID and value is a tuple containing timestamp, new grade & a flag which is equal to "local" to represent that the calling system has updated the log.

- The oplog of the other system is then processed, updating the key-value store with newer values. Here the flag is updated to "remote" to indicate that the other system has updated it.
- The final key-value pairs are then written into the calling system using set() operations, if they have the "remote" flag.

Merge Properties

The merge operation defined for each system satisfies the following mathematical properties, which are critical to ensuring consistency and convergence in distributed systems.

Commutativity

Merging two operation logs in different orders results in the same final state:

$$A \oplus B = B \oplus A$$

```
(myenv) iiitheliiith-NP-Laptop-15-d-d/Nxxx:-/MOSQL_PROJECTS python3 main.py
SQL: SET (SID1093, CSE016) > A
HIVE: SET (SID1093, CSE016) > B
Fetched 1 records from MongoDB.
DEBUG - filtered kv_store contents:
Key: (SID1033, CSE016) -> Value: (ts=2025-04-30 06:06:37.615000, grade=B, flag=renote)
trtgerred
Herged 1 records into HongoDB from HIVE.
HONGO: MERGE (HIVE)
Fetched 1 records from SQL.
Fetched 1 records from SQL.
Fetched 1 records from SQL.
Fetched 1 records from MongoDB.
```

Associativity

The grouping of merges does not affect the final state:

$$A \oplus (B \oplus C) = (A \oplus B) \oplus C$$

```
(nyenv) illibalilib-HP-Laptop-15-dw3xxx:-/HOSQL_PROJECT$ python3 main.py
HONGO: SET (SID1033, CSE016) >> B
SQL: SET (SID1033, CSE016) >> B
SQL: SET (SID1033, CSE016) >> B
SQL: SET (SID1033, CSE016) >> A
Fetched 1 records from MongoDB.

DEBUG - filtered_kv_store contents:
Key: (SID1033, CSE016) -> Value: (ts=2025-04-30 06:16:37.249000, grade=B, flag=renote)
trinnernd
Terminal records into MongoDB from HIVE.
Numour. mEMCE (HIVE)
Fetched 1 records from SQL.
Fetched 1 records from MongoDB.

DEBUG - filtered kv_store contents:
Key: (SID1033, CSE016) -> Value: (ts=2025-04-30 06:17:01.094000, grade=A, flag=renote)
trigerred
Merged 1 records into MongoDB from SQL.
MONGO: MERGE (SQL)
MONGO: SERGE (SQL)
MONGO: SERGE (SQL)
MONGO: SERGE (MONGO)
MONGO: SERGE (SQL) -> A
```

Idempotency

Merging a system with itself leaves it unchanged:

$$A \oplus A = A$$

```
(Myenv) iiiibaiiitb-HP-Laptop-15-dw3xxx:-/NOSQL_PROJECI$ python3 main.py
SQL: SET (SID1033, CSE016) -> A
Fetched I records from SQL.
Fetched I records from MongoOB.

Key: (SID1033, CSE016) -> Value: (ts=2025-04-30 06:21:41.701000, grade=A, flag=remote)
trigerred
Merged I records into MongoOB from SQL.
MONGO: HERCE (SQL)
MONGO: HERCE (SQL)
MONGO: TET (SID1033, CSE016) -> A
Fetched I records from MongoOB.

DEBUG - filtered_kv_store contents:
Key: (SID1033, CSE016) -> Value: (ts=2025-04-30 06:21:41.701000, grade=A, flag=remote)
MCGO: HERCE (SQL)
MONGO: MERCE (SQL)
```

Convergence

Once the views of all three databases became the same, they reached consistency.

```
(myenn) ittemetites implements and services and services pythona main.py
MONCO: SET (SID1033, CSE016) >> C
HIVE: SET (SID1033, CSE016) >> Value: (ts=2025-04-30 06:43:08.498000, grade=C, flag=renote) trigerred
Rerged 1 records into MongoDB From SQL.
MONCO: MERGE (SQL)
Fetched 1 records from HVe.
Fetched 1 records from HVe.
Fetched 1 records from HVe.
Fetched 1 records from MongoDB.
DEBUG - filtered_kv_store contents:
Key: (SID1033, CSE016) >> Value: (ts=2025-04-30 06:43:29.349000, grade=B+, flag=renote) trigerred
Merged 1 records from MongoDB from HIVE.
MONGO: MERGE (HIVE)
MONGO: MERGE (HIVE)
Fetched 1 records from MongoDB.
Fe
```

Challenges Faced

- Hive connectivity issues due to port misconfiguration.
- Designing a merge logic that respects commutative properties; resolved via a deterministic key-value update strategy using oplogs.
- Initially we explored an alternative design, which achieved eventual consistency, but didn't achieve SEC. (Instead of updating merge as a set operation in the oplog with the old timestamp, you use a new timestamp).

Code Explanation

main.py

This script acts as the central driver for executing operations across the three database systems. It performs the following key functions:

- Imports the three manager classes: HiveGradeManager, SQLGradeManager, and MongoDBGradeManager.
- Initializes each manager instance using a common CSV file (student_course_grades.csv) and the appropriate database connection strings.
- Maps each system identifier (e.g., HIVE, SQL, MONGO) to its corresponding manager instance in a dictionary (manager_map).
- Reads instructions from an input file named testcase_hive.in, where each line specifies an operation in the format SYSTEM.OPERATION.
- Parses and delegates SET, GET, and MERGE operations to the appropriate manager. The SET command updates a student's grade for a given course, GET retrieves the current grade, and MERGE synchronizes the state of the current system with another using its oplog.
- Introduces a one-second delay between operations to ensure unique timestamps in the oplogs for consistent merging.

This file abstracts the orchestration logic from the underlying database details, focusing purely on parsing and dispatching user-defined operations.

hive_manager.py

This file defines the HiveGradeManager class, which manages student grades in an Apache Hive database. Its main responsibilities include table initialization, data access (get/set), operation logging, and merging updates from external systems (SQL, MongoDB). The key functionalities are:

- Initialization: Connects to Hive and creates two tables in the new_database schema—grades for storing student grades and oplogs for recording operations.
- execute(): A generic helper method that runs SQL queries on Hive and returns results for SELECT queries.
- get(): Retrieves a student's grade for a specific course and logs the operation in oplogs.
- set(): Updates a student's grade in the grades table using INSERT OVERWRITE with conditional logic, and logs the operation.
- log2(): Similar to set(), but designed to apply externally sourced changes during a merge, using a provided timestamp for proper ordering.
- merge(): Imports the most recent SET operations from another system's oplog (either SQL or MongoDB), resolves conflicts based on timestamps, and applies newer external changes to the local Hive instance.
- Operation Logging: All GET and SET operations are recorded in the oplogs table with a millisecond-precision timestamp, enabling eventual consistency and cross-system merges.

Overall, this class ensures that Hive can act both as a primary data store and a replica target, maintaining consistency with other systems through a log-based synchronization strategy.

postgres_manager.py

This Python script defines a class SQLGradeManager that manages student grades using a PostgreSQL database, and supports data synchronization from Hive and MongoDB sources. It performs the following key functions:

- __init__: Initializes the database engine, metadata, and sets up tables using initialize_tables().
- initialize_tables: Defines and creates two tables:
 - grades holds student grades with a composite primary key (student-ID, course-id).
 - oplogs logs operations (GET, SET) with timestamps.

The method also loads initial data from a CSV file into the grades table.

- get(student_id, course_id): Retrieves a student's grade for a specific course and logs the read operation.
- set(student_id, course_id, new_grade): Updates the grade for a given student-course combination and logs the write operation.
- merge(source_system): Synchronizes updates from an external system (Hive or MongoDB) into the local SQL database. It performs conflict resolution by keeping the most recent grade (based on timestamps) and logs each applied remote change.
- _log_operation and _log_operation2: Internal methods to insert records into the oplogs table, either with the current or provided timestamp.

This design ensures consistency and auditability of grade updates across multiple systems, while resolving conflicts by preferring the most recent data.

mongo_manager.py

The mongo_manager.py script defines the MongoDBGradeManager class, which interfaces with a MongoDB database to manage student course grades. It performs the following key functions:

- Initialization: Connects to MongoDB, drops and recreates the grades and oplogs collections, sets up a compound index on student-ID and course-id, and loads initial data from a CSV file.
- Grade Retrieval (get): Fetches the grade for a specific student-course pair and logs the access in the oplogs collection.
- Grade Update (set and log2): Updates an existing grade and logs the update operation. The log2 method is a variant used during merge operations with an externally provided timestamp.
- Logging (_log_operation): Inserts a timestamped log entry into oplogs for each GET or SET operation.
- Merge Functionality: The merge method integrates grade updates from external sources (PostgreSQL or Hive) into MongoDB. It collects and compares the latest timestamped SET operations and applies only newer remote updates to MongoDB while logging them using log2.

This class ensures synchronization across databases and maintains a detailed log of all interactions.

Conclusion

This project demonstrates a robust approach to synchronizing data across heterogeneous systems using abstracted operations and operation logs. Through consistent and commutative merges, we can ensure eventual consistency and correctness across all platforms involved.

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