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→ if  $f$  is computed by  
a decision tree of size  $\ell$

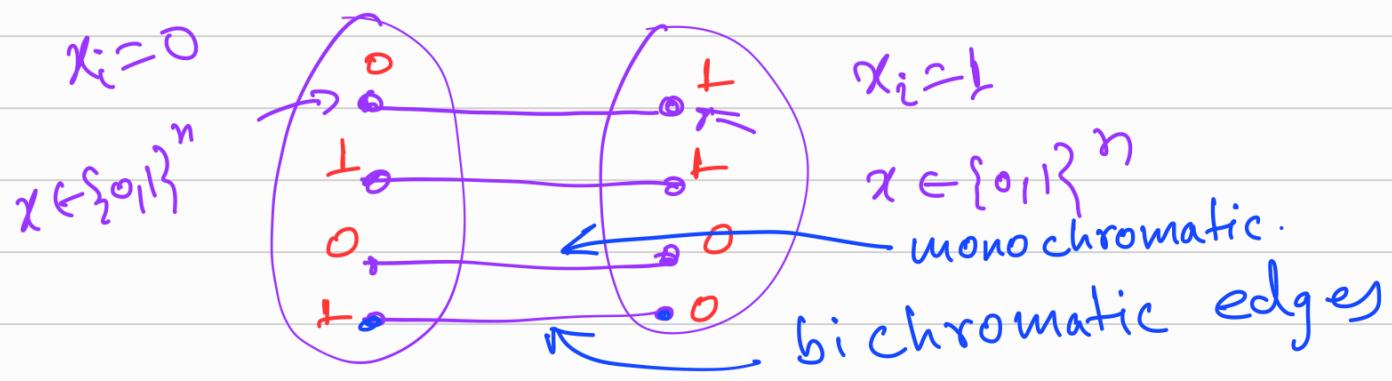
then  $\|f\|_1 := \sum_{S \subseteq [n]} |f(S)| \leq \ell$ .

→ Defn:- Influence of a Variable.

Let  $f: \{0,1\}^n \rightarrow \{0,1\}$

$\text{Inf}_i(f)$  :=  $\Pr_{x \sim \{0,1\}^n} [f(x^i) \neq f(x)]$

( $x^i$  is  $x$  with  $i$ -th bit -flipped)



$$\Pr_{x \sim \{0,1\}^n} [f(x') \neq f(x)]$$

$$= \left( \frac{1}{2^n} + \frac{1}{2^n} \right) = \frac{2 \cdot \text{\# bichromatic edges in } i\text{-th direction}}{2^n}$$

Can  $\text{Inf}_i(f) = 0$ ?

Then function  $f$  doesn't depend on variable  $x_i$ .

$$\text{Inf}_i(f) \neq 0 \Rightarrow \text{Inf}_i(f) \geq \frac{1}{2^{n-1}}$$

Defn:- Total Influence of a function  $f$

$$\text{Inf}(f) := \sum_{i=1}^n \text{Inf}_i(f)$$

$$\text{Average Sensitivity} = \text{Inf}(f).$$

$$:= \mathbb{E}_{\substack{x \sim \{0,1\}^n}} [s(f, x)]$$

$$\text{as}(f) = \mathbb{E}_x [s(f, x)]$$

define  $s_i(f, x) = \begin{cases} 0 & \text{if } f(x) = f(x^i) \\ 1 & \text{if } f(x) \neq f(x^i) \end{cases}$

$$= \frac{1}{2^n} \sum_{x \in \{0,1\}^n} s(f, x)$$

$$= \frac{1}{2^n} \sum_{x \in \{0,1\}^n} \sum_{i=1}^n s_i(f, x)$$

$$= \frac{1}{2^n} \sum_{i=1}^n \sum_{x \in \{0,1\}^n} s_i(f, x)$$

$$= \sum_{i=1}^n \left[ \frac{1}{2^n} \sum_{x \in \{0,1\}^n} s_i(f, x) \right]$$

$$= \sum_{i=1}^n \Pr_x [f(x) \neq f(x^i)]$$

$$= \sum_{i=1}^n \text{Inf}_i(f) = \text{Inf}(f).$$

$\Rightarrow$  if  $\text{Inf}(f) \neq 0$

$$\Rightarrow \text{Inf}(f) \geq \frac{n}{2^{n-1}}.$$

$$\Rightarrow \text{Inf}(f) \geq \frac{2 \cdot \# \text{ bichromatic edges}}{2^n} \geq \frac{2 \cdot n}{2^n}$$

$$\text{Inf}(\text{XOR/Parity}) = n.$$

$$0 \leq \text{Inf}(f) \leq n$$

Thm:- [KKL '88]

Kahn-Kalai-Linial.

$$f: \{\pm 1\}^n \rightarrow \{\pm 1\}$$

$$\text{Inf}_i(f) := \Pr_{\alpha} [f(\alpha) \neq f(\alpha^i)]$$

Then,

$$\boxed{\text{Inf}_i(f) = \sum_{S: i \in S} \hat{f}(S)^2}$$

Recall,  $f(\alpha) = \sum_{S \subseteq [n]} \hat{f}(S) \prod_{i \in S} \alpha_i$

$$\sum_{S \subseteq [n]} \hat{f}(S)^2 = 1 \quad [\text{Parseval's Thm}]$$

$$\sum_{S \subseteq [n]} \hat{f}(S)^2 = E[f^2]$$

$$f: \{\pm 1\}^n \rightarrow \mathbb{R}$$

Devise a function  $g_L : \{0, 1\}^{n-1} \rightarrow \{0, 1, -1\}$

$$|g_L(y)| = \begin{cases} 1 & \text{if } f(1, y) \neq f(0, y) \\ 0 & \text{o/w.} \end{cases}$$

$$g_L(y) := f(1, y) - f(0, y)$$

$$g_L : \{\pm 1\}^{n-1} \rightarrow \{\pm 1, 0\}$$

$$f : \{\pm 1\}^n \rightarrow \{\pm 1\}.$$

$$g_L(y) := \frac{f(1, y) - f(-1, y)}{2}$$

$$\text{Inf}_1[f] := \Pr_{x \sim \{\pm 1\}^n} [f(x) \neq f(x')]$$

$$= \Pr_{y \sim \{\pm 1\}^{n-1}} [g_L(y) \in \{\pm 1\}]$$

$$= \Pr_y [g_L^2(y) = 1]$$

$$= \mathbb{E}_y [g_1^2(y)]$$

$$\boxed{f(y_1, \dots, y_n) = \sum_{S \subseteq [n]} \hat{f}(S) \prod_{i \in S} y_i}$$

$$= \sum_{S: 1 \notin S} \hat{f}(S) \prod_{i \in S} y_i$$

$$+ \sum_{S: 1 \in S} \hat{f}(S) \prod_{i \in S} y_i$$

$$f(1, y_2, \dots, y_n) - f(-1, y_2, \dots, y_n)$$

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$$= \sum_{S: 1 \in S} \hat{f}(S) \prod_{i \in S \setminus \{1\}} y_i$$

$$g_1(y) = \sum_{S: 1 \in S} \hat{f}(S) \prod_{i \in S \setminus \{1\}} y_i$$

$$\text{Inf}_1(f) = \mathbb{E}_y [g(\omega)]^2$$

Parseval's equality

from  $\rightarrow = \sum \hat{f}(s)^2$

$s: 1 \in S$

$$f: \mathbb{S} \rightarrow \mathbb{R}$$

$$\mathbb{E}[f^2] = \sum_{s \subseteq [n]} \hat{f}(s)^2$$

$$\text{Inf}_i(f) = \sum_{s: i \in s} \hat{f}(s)^2$$

$$\text{Inf}(f) = \sum_{i=1}^n \text{Inf}_i(f)$$

$$= \sum_{i=1}^n \sum_{s: i \in s} \hat{f}(s)^2$$

$$\xrightarrow{\text{[KKL-Lemma]}} = \sum_{S \subseteq [n]} |S| \hat{f}(S)^2$$

Obs 1:  $\text{Inf}(f) \leq \deg(f)$

$$\begin{aligned} \text{Inf}(f) &= \sum_{S \subseteq [n]} |S| \hat{f}(S)^2 \\ &\leq \deg(f) \sum_{S \subseteq [n]} \hat{f}(S)^2 \\ &= \deg(f) \cdot 1 \end{aligned}$$

## Fourier Entropy-Influence

Conjecture. [Friedgut + Kalai 1998]

$$\sum_{S \subseteq [n]} \hat{f}(S)^2 = 1$$

$$H(f) = \sum_{S \subseteq [n]} \hat{f}(S)^2 \log \frac{1}{\hat{f}(S)^2}$$

$$Inf(f) = \sum_{S \subseteq [n]} |S| \hat{f}(S)^2$$

$\exists$  a universal constant  $C > 0$   
 $S \cdot f$ : for all  $f: \{\pm 1\}^n \rightarrow \{\pm 1\}$

$$H(f) \leq C \cdot Inf(f).$$

$$\sum \hat{f}(S)^2 \log \frac{1}{\hat{f}(S)^2} \leq C \cdot \sum |S| \hat{f}(S)^2$$

Q:- let  $f: \{0,1\}^n \rightarrow \{0,1\}$

be a Boolean fn over  
n vars and deg d.

What is an upper bound

on D in terms  
of d?

MUX :  $\{0,1\}^{n+2^n} \rightarrow \{0,1\}$

Address function

$\text{MUX}(x_1, \dots, x_n, y_0, \dots, y_n)$

$\equiv Y_{\text{bin}(x_1, \dots, x_n)}$

$$\deg \leq n+1$$

$$\#\text{Var} = n + 2^n := M$$

$$\deg = O(\log M)$$

Thm :- Let  $f: \{0,1\}^n \rightarrow \{0,1\}$

of deg d then

$$n \leq d \cdot 2^{d-1}$$

$$\Rightarrow d = \mathcal{O}(\log n).$$

when  $n$  is the no. of  
vars f depends on.

Thm:-  $\rightarrow$  multilinear.

let  $p: \{0,1\}^n \rightarrow \mathbb{R}$  of  
deg d.

Then  $\Pr_{x \sim \{0,1\}^n} [p(x) \neq 0] \geq \frac{1}{2^d}$

$f: \{0,1\}^n \rightarrow \{0,1\}$  of deg d.

$$\text{Inf}_i(f) = \sum_{S: i \in S} \hat{f}(S)^T x_i$$

$\text{Inf}_i(f)(x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n)$

$$\Pr_x [\text{Inf}_i(f) \neq 0] \geq \frac{1}{2^{d-1}}$$

$$\sum_{i=1}^n \text{Inf}_i(f) \leq d$$

W

$$\frac{n}{2^{d-1}}$$

$$\Rightarrow n \leq d \cdot 2^{d-1}.$$