

$$D(f) \leq C^0(f) \cdot C^1(f).$$

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Thm:- For all  $f: \{0,1\}^n \rightarrow \{0,1\}$

$$D(f) \leq C^0(f) \cdot C^1(f)$$

Proof:-

Claim:- Let  $(S, \alpha)$  and  $(T, \beta)$

be a 0-certificate and 1-certificate respectively.

Then  $S \cap T \neq \emptyset$ .

Furthermore,  $\exists i \in S \wedge T$

$$s.t. \alpha(i) \neq \beta(i).$$

Proof of Claim: Exercise!

Consider the following algo.

Step 0:-  $X = \{0, 1\}$

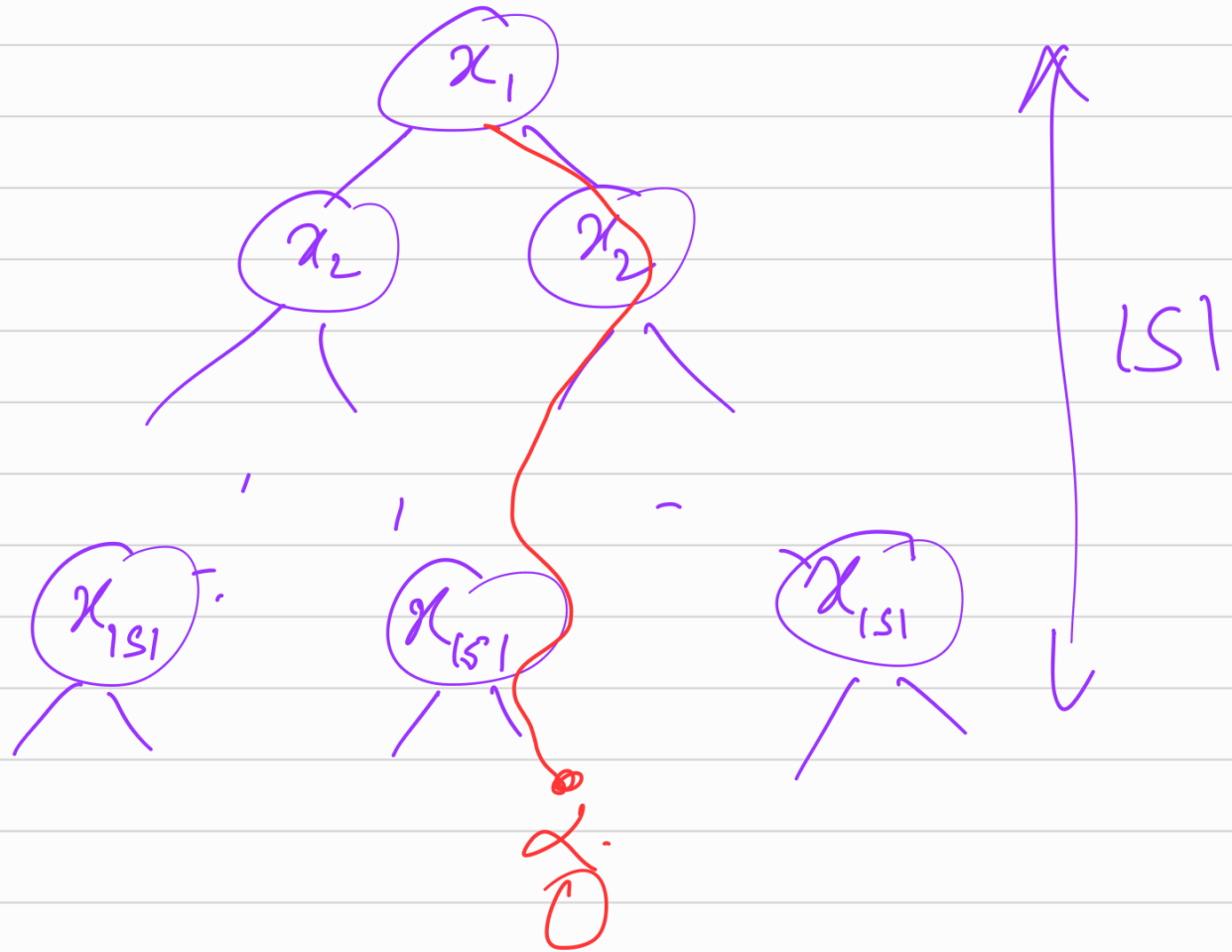
Repeat  $C'(f)$  times.

Step 1:- if  $f(x)$  is constant  
for all  $x \in X$ , then  
Output  $f(x)$ .

Otherwise, Pick a  $O$ -certificate.

$(\underline{S}, \alpha)$   
where  $S \subseteq [n]$  and  $\alpha: S \rightarrow \{0, 1\}$

Query the variables in  $S$ .



If you match with  $\alpha$  then  
Output 0 and halt.

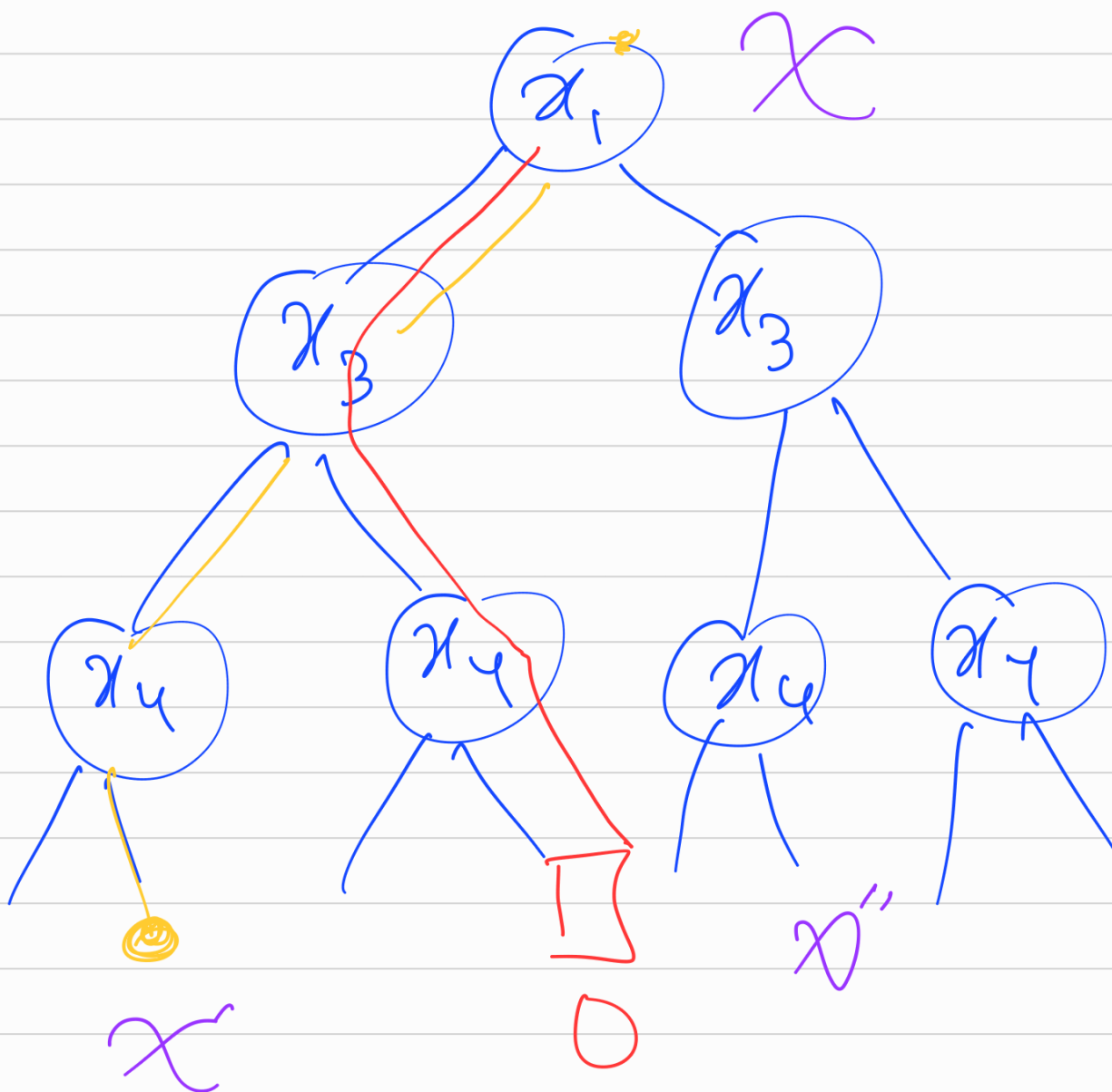
Else. Prune  $\mathcal{X}$  to  
be consistent with revealed  
answers.

Step 2: Pick a  $y \in \mathcal{X}$   
Output  $f(y)$  and halt.

Worst case

# queries made by  
the algo  $\leq C'(f) \cdot C^0(f)$

$$S = \{x_1, x_3, x_4\}$$



## Proof of Correctness:

Claim 1 - 0-certificate and 1-certificates

intersect in a contradictory way.

→ Claim 2 -

Let  $f_i$  be the function

after  $i^{\text{th}}$  iteration of

of Step 1.

Then,  $C'(f_i) \leq C'(f_{i-1}) - 1$

$$f_i = f(\dots, \dots, \dots)$$

with revealed  
bits set  
accordingly.

Proof:- Because every 0-certificates  
intersects every 1-certificates

After One iteration of  
Step 1 you know  
the value to one variable

from every 1-certificate.

Thm 1:-

$$D(f) \leq C^0(f) \cdot C^1(f)$$

$$\leq C(f)^2$$

P in decision tree

= NP in decision tree

$\wedge$  coNP in decision tree.