## Topics in Computing (CS5160): Problem Set 1

## Department of Computer Science and Engineering IIT Hyderabad

- Scan and upload your answer sheet on google classroom.
- Maintain academic honesty. If caught, you will get an F in the course.
- Please write "credit" or "audit" on your answer sheets depending on whether you are crediting or auditing the course.
- Due Date: 20 September (before 11:59pm).
- 1. Let  $f: \{0,1\}^n \to \{0,1\}$  be a Boolean function, and k be the largest natural number such that  $|f^{-1}(1)|$  is divisible by  $2^k$ . Show that  $\mathsf{D}(f) \geq n-k$ , where  $\mathsf{D}(f)$  is the decision tree complexity of f.
- 2. For every k, define  $f_k$  to be the following function taking inputs of length  $n=2^k$ :

$$f_k(x_1, \dots, x_n) = \begin{cases} f_{k-1}(x_1, \dots, x_{2^{k-1}}) \land f_{k-1}(x_{2^{k-1}+1}, \dots, x_{2^k}) & \text{if } k > 0 \text{ is even} \\ f_{k-1}(x_1, \dots, x_{2^{k-1}}) \lor f_{k-1}(x_{2^{k-1}+1}, \dots, x_{2^k}) & \text{if } k \text{ is odd} \\ x_i & \text{if } k = 0 \end{cases}$$

Show that  $D(f_k) = 2^k$ . (10 points)

3. Show that for a non-constant symmetric Boolean function  $f: \{0,1\}^n \to \{0,1\}$ ,

$$\mathsf{s}(f) \ge \left\lceil \frac{n+1}{2} \right\rceil,$$

where s(f) is the sensitivity of f. Also, give an example where this bound is tight. Recall, we say a function is symmetric if its value depends only on the number of 1s in the input. That is, it depends on the hamming weight of the input. (10 points)

4. Show that s(f) = bs(f) = C(f) for every monotone Boolean function f, where bs(f) and C(f) respectively denote the block sensitivity and certificate complexity of f. Recall we say a function is monotone if for any x, y such that  $x \leq y$ , that is,  $x_i \leq y_i$  for all  $i \in [n]$ ,  $f(x) \leq f(y)$ . (10 points)

- 5. Prove that when f is monotone,  $s(f) \le deg(f)$ . (10 points)
- 6. Define average sensitivity,  $\operatorname{as}(f)$ , of a Boolean function f to be the expected sensitivity of f on a random input:  $\operatorname{as}(f) = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} \operatorname{s}(f,x)$ . Let T be a decision tree and let  $\ell_x$  be the length of the unique path in T consistent with x. Define average depth of T to be  $\frac{1}{2^n} \sum_{x \in \{0,1\}^n} \ell_x$ . Show that the average depth of any decision tree for f is at least  $\operatorname{as}(f)$ . (10 points)