

11/09/23

Thm :- [Nisan-Szegedy]

For all $f: \{0,1\}^n \rightarrow \{0,1\}$,

$$\deg(f) = \Omega(\sqrt{\text{bs}(f)}).$$

Proof:-

Let $a \in \{0,1\}^n$ be an input s.t. $\text{bs}(f, a) = \text{bs}(f) = b$.

Furthermore, assume WLOG,

$$f(a) = 0.$$

Let B_1, \dots, B_b be the disjoint sensitive blocks at

a which witnesses the block sensitivity at a equals $b = \text{bs}(f)$.

Let $p(x_1, \dots, x_n)$ be the polynomial that represents f .

$$\deg(p) = \deg(f).$$

Obtain $q(y_1, \dots, y_b)$ from $p(x_1, \dots, x_n)$ by the following substitution.

(i) if $x_i \notin \bigcup_{j=1}^b B_j$ then

set $x_i = a_i$ in β .

(ii) if $x_i \in B_j$ and

$$a_i = 0$$

then set $x_i = y_j$ in β .

(iii) if $x_i \in B_j$ and

$$a_i = 1$$

then set $x_i = 1 - y_j$ in β .

After these substitutions

we have obtained a

polynomial $q_r(y_1, \dots, y_b)$.

Properties of q_r

(i) q_r is multilinear polynomial

of $\deg(q_r) \leq \deg(\phi)$

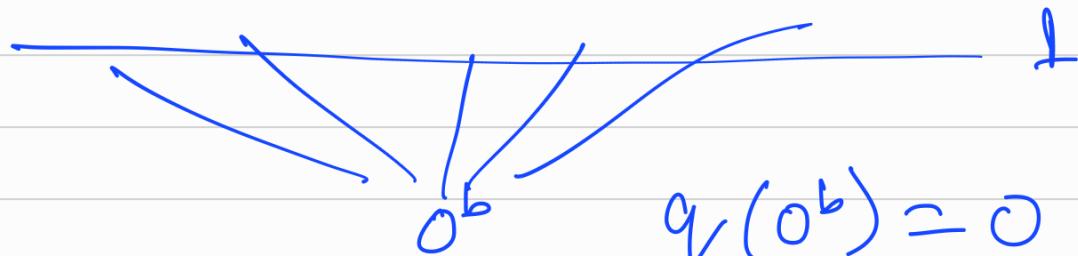
$\leq b$

$$q_r: \{0,1\}^b \rightarrow \{0,1\}$$

(ii) $q_r(0^b) = \phi(a) = f(a) = 0$

(iii) $q_r(e_i) = \phi(a^{B_i}) = f(a^{B_i}) = 1$

$\forall i \in \{1, 2, \dots, b\}$



Claim :- $\deg(q) \geq \sqrt{\frac{b}{2}}$

define

$$q_{\text{sym}}(y_1, \dots, y_b) = \underbrace{\sum_{\sigma \in S_b} q(\sigma(y))}_{b!}$$

where S_b is the set of permutations over

$$\{1, \dots, b\}$$

and $\sigma(y) = \sigma(y_1, \dots, y_b) = (y_{\sigma(1)}, y_{\sigma(2)}, \dots, y_{\sigma(b)})$

(Symmetrization of a polynomial)

(Minsky & Pabst '63).

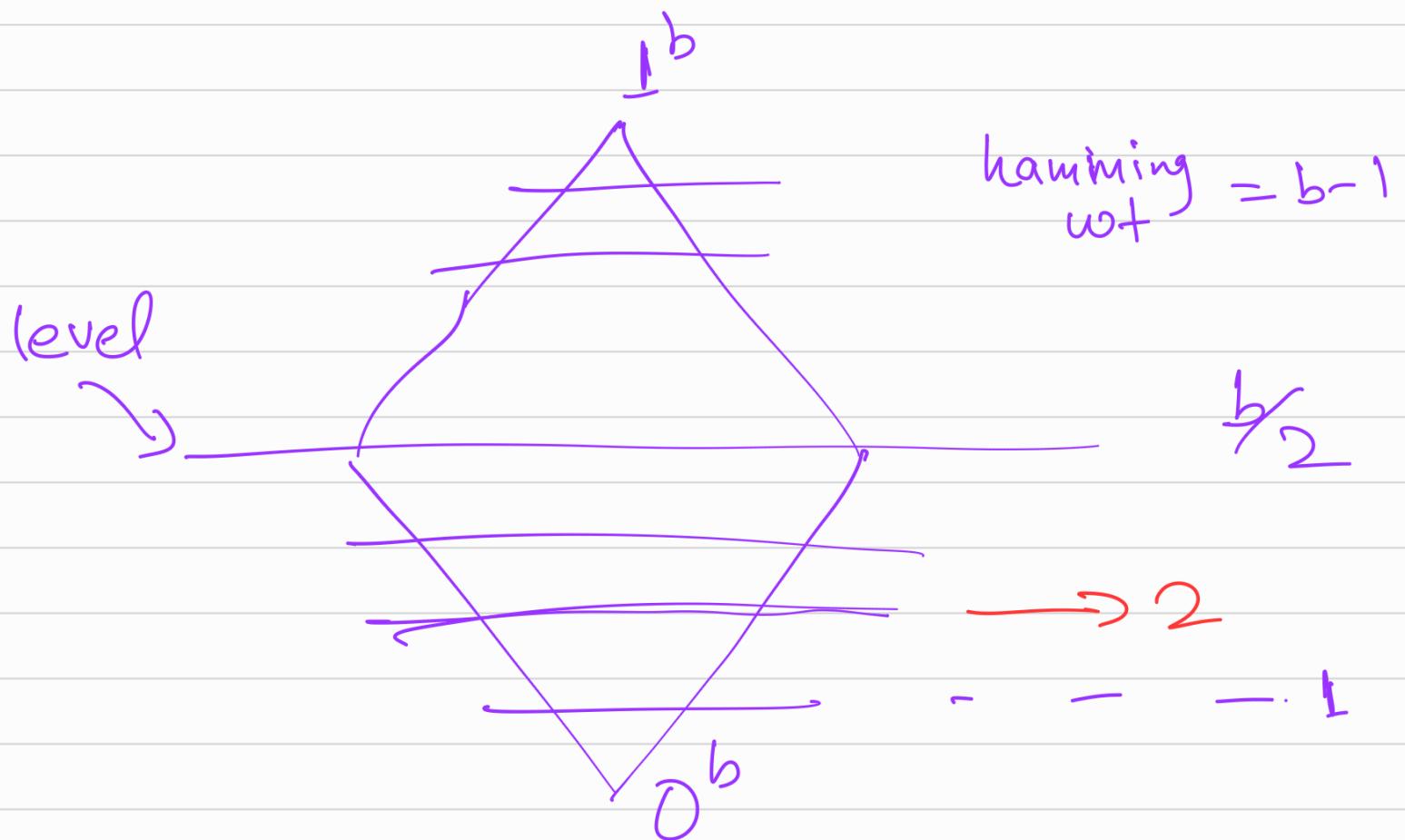
Properties of g_{sym}

① $g_{\text{sym}}(0^b) = g(0^b) = 0$

② $g_{\text{sym}}(e_i) = 1.$

③ $\deg(g_{\text{sym}}) \leq \deg(g).$

④ g_{sym} is a univariate polynomial.



$(10^{b-2}) \rightarrow \text{all inputs}$
 in level 2.

$$q_{\text{sym}}((10^{b-2})) = \frac{1}{b!} \sum_{\sigma \in S_b} q(\sigma(10^{b-2}))$$

$$= \frac{1}{b!} \sum_{\substack{y \in \{0,1\}^b \\ |y|=2}} x \cdot q(y)$$

$$\lambda := \left| \{ z \mid \sigma(110^{b-2}) = 0110^{b-3} \} \right|$$

$$\deg(g_{\text{sym}}) \leq \deg(g) \leq \deg(\phi)$$

Lemma :- Let $h : \mathbb{R} \rightarrow \mathbb{R}$

be a univariate polynomial

s.t. for every integer k ,

where $0 \leq k \leq n$,

$$b_1 \leq h(k) \leq b_2.$$

And $\exists a \in [0, n]$

$$\text{s.t. } |h'(z)| \geq c.$$

$$\text{Then } \deg(h) \geq \sqrt{\frac{c \cdot n}{c + b_2 - b_1}}$$

(Back to proof).

$$\Rightarrow \forall 0 \leq k \leq b$$

$$0 \leq q_{\text{sym}}(k) \leq 1$$

$$\Rightarrow q_{\text{sym}}(0) = 0$$

$$\Rightarrow q_{\text{sym}}(1) = 1$$

\Rightarrow By intermediate value theorem

\exists a point $z \in [0, 1]$

such that $q'_{\text{sym}}(z) \geq \frac{q_{\text{sym}}(1) - q_{\text{sym}}(0)}{1 - 0}$

$$= \frac{1 - 0}{1 - 0} = 1.$$

From the Lemma, we have

$$\deg(q_{\text{sym}}) \geq \sqrt{\frac{1 \cdot b}{1 + 1 - 0}}$$
$$= \sqrt{\frac{b}{2}}.$$



OPEN :- Show \exists a

function $f: \{0,1\}^n \rightarrow \{0,1\}$

s.t. $\deg(f) = O(\sqrt{bs(f)})$

Boolean function
on n variables $= 2^{2^n}$

Q :- How many Boolean
functions over n variables

have $\deg = n$?

$l \rightarrow o(l)$

$g(m) \in o(1)$

iff $\frac{g(m)}{2} \rightarrow 0$

as $m \rightarrow \infty$

B.f. with $\deg = n$

2^{2^n}

e.g. $\frac{1}{m}, \frac{1}{2^m}$

$\frac{1}{\sqrt{m}}, \frac{1}{\log m}$

Q :- How to Count

B.f. with $\deg = n$?

Consider $f: \{0,1\}^n \rightarrow \{0,1\}$

It is representable by

a polynomial $f(x_1, \dots, x_n)$.

$$f(x_1, \dots, x_n) = \sum_{S \subseteq [n]} c_S \prod_{i \in S} x_i$$

where c_S is the coefficient
of monomial $\prod_{i \in S} x_i$

Q'':- How to compute

$$c_S ?$$

$$\text{Thm : } c_S = \sum_{T \subseteq S} (-1)^{|S|-|T|} \cdot f(T)$$

where $f(T) = f(\alpha)$ where
 $\alpha_i = 1$ iff $i \in T$.

Define $X_1^{\text{even}} = \{x \mid f(x)=1 \text{ and } |x| \text{ is even}\}$

$$X_{\text{odd}} = \left\{ x \mid \begin{array}{l} f(x) = 1 \\ |x| \text{ is odd} \end{array} \right\}$$

Thm :- $\deg(f) = n$ iff

$$\left| \begin{array}{c} X_{\text{even}} \\ \hline \end{array} \right| \neq \left| \begin{array}{c} X_{\text{odd}} \\ \hline \end{array} \right|$$

Proof: $C_S = \sum_{T \subseteq S} (-1)^{|T| - |S|} \cdot f(T)$

$$\deg(f) = n \quad \text{iff} \quad c_{[n]} \neq 0$$

$$C_{[n]} = \sum_{T \subseteq [n]} (-1)^{n-|T|} \cdot f(T)$$

$$= (-1)^n \sum_{T \subseteq [n]} (-1)^{|T|} \cdot f(T)$$

$$= (-1)^n \sum_{T \subseteq [n]} (-1)^{|T|}$$

$$f(T) = 1$$

$$= (-1)^n \left(\left| X_1^{\text{even}} \right| - \left| X_1^{\text{odd}} \right| \right)$$

~~too~~

Count # B.f. with

$\deg < n$

\therefore for such functions

$$|X_{\text{even}}| = |X_{\text{odd}}|$$

$$\sum_{i=0}^{n-1} \binom{2^{n-1}}{i} \cdot \binom{2^{n-1}}{i}$$

$$= \frac{\binom{2^n}{2^{n-1}}}{\sqrt{2^n}}$$

$$\left[\binom{k}{k_2} \right] \approx \frac{2^k}{\sqrt{k}}$$