



# CS-3002: Information Security

## **Lecture # 9: Authentication and Access Control**

Department of Software Engineering  
FAST-NUCES



# Overview

- *Authentication*
- *Passwords*
- *Secure ID*
- *Google 2-step Authentication*
- *Access Control*



# Authentication



- **Authentication = binding of an identity to a subject**
- **Confirmation of identity by ...**
  - Knowledge factors = what the entity knows
  - Ownership factors = what the entity has
  - Human factors = what the entity is
  - Location factors = where the entity is

# Example

- **Login to a computer**
  - Authentication by *knowledge (password)*
- **Online debit cards**
  - Authentication by *ownership (card) and knowledge (PIN)*
- **Offline debit cards**
  - Authentication by *ownership (card) and human factor (signature)*



# Multi-Factor Authentication

- **Authentication using multiple factors**
  - Example: *Scene from the movie “Mission Impossible”*

Ethan Hunt needs to

1. use a stolen chip card (ownership factor)
2. forge a fingerprint (human factor)
3. enter the terminal room (location factor)
4. enter a password (knowledge factor)



# Passwords

- **Password** = information confirming the identity of an entity
  - Knowledge of a secret word, phrase or number



- **Often combination with (a)symmetric cryptography**
  - e.g. password is mapped to key of symmetric cipher
  - e.g. password protects private key of public-key algorithm
- ~~Passwords are just great.~~ Wait, it's not that easy



# Problems with Passwords

- **Password snooping**
  - Eavesdropping of passwords in network traffic
  - Retrieval of passwords from hosts (e.g. via malware)
- **Password guessing (online) or cracking (offline)**
  - Dictionary attacks = guessing using dictionary of words
  - Brute-force attacks = guessing using all possible strings
- **Human deficiencies**
  - Weak and often re-used passwords



# Passwords Storage

- **Passwords should never be stored in clear**
  - Application of cryptographic one-way functions
  - Only encoded (hashed) passwords are stored
    - Sony data breach revealed clear text password.
    - Why twitter auto-reset the passwords recently ?
- **Example: `$stored_pw = hash($password);`**
  - Simple to validate: `hash($input) == $stored_pw?`
  - Hard to deduce password from strong hash functions
- **Efficient cracking of stored passwords still possible**
  - Brute-force or dictionary attack using hashed strings

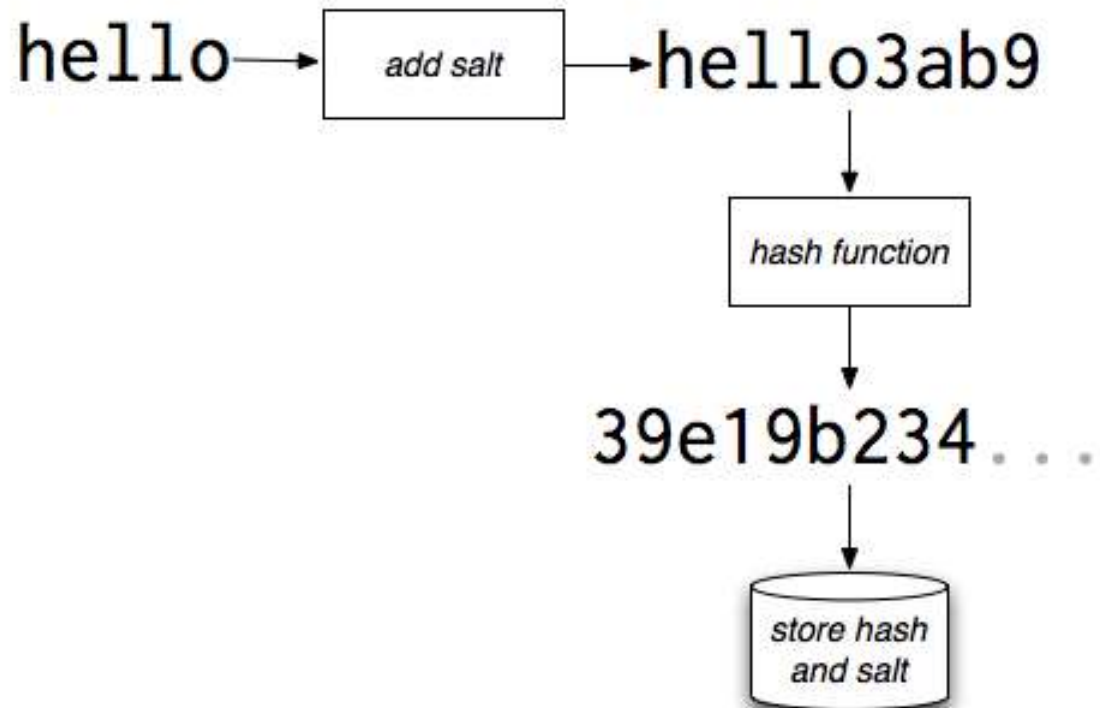




# Salted Passwords

- **Encoding of password with random string (salt)**
  - Example: `$stored_pw = hash($password+$salt);`
  - Salt value stored along with hashed password
- **Cracking of stored passwords more expensive**
  - Same password maps to different hash values
  - Without salt: cracking depends on # words
  - With salt: cracking depends on (# words  $\times$  # salts)
- **Security depends on quality of password, hash and salt**



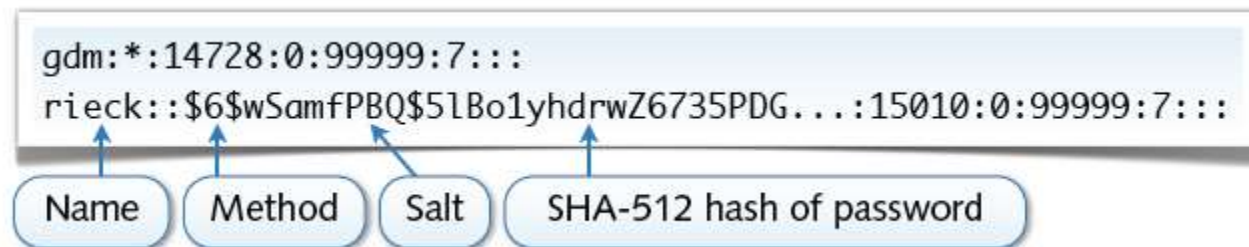


# Example: Unix Password

- **User credentials stored in two separate databases**
  - `/etc/passwd` Basic user information (publicly readable)



- `/etc/shadow` Salt and hashed passwords (protected)



# Good Password?

- **Testing for insecure passwords is *very easy***
  - A normal core i3 laptop can test 21 million MD5 hashes per hour
- **Passwords should be *very hard to guess***
  - No dictionary words, names, dates and patterns
  - Simple transformations (e.g. reversing) not sufficient
  - Minimum length and diversity of passwords
- **Study by Klein from 1989**
  - 21% of 13,797 passwords cracked within one week



# Selection of Passwords

- What about these? **Hmh40hcr.** and **DB:L,I4yF!**
- **Trick: first letters of memorable phrase**
  - “He made him an offer he can't refuse.” = **Hmh40hcr.**
  - “Darth Vader: Luke, I am your father!” = **DB:L,I4yF!**
- **Trick: interweave words of memorable phrase**
  - “My kingdom for a horse!” = **KiHor;NgSe**
- **Avoidance of too common phrases**
  - **2bon2b** found in 4 out of 30 million passwords



# One-time Passwords

- **Security of passwords “weakens” over time**
  - Password aging = enforced changing of passwords
  - One-time passwords = passwords used exactly once
- **Example: S/Key Algorithm**
  - User chooses initial key  $K_1$
  - Recursive hashing:  $H(K_1) = K_2, H(K_2) = K_3, \dots H(K_{n-1}) = K_n$
  - One-time passwords:  $P_1 = K_n, P_2 = K_{n-1}, \dots P_n = K_1$
  - Hard to deduce next password  $P_i$  from previous  $P_{i-1}$



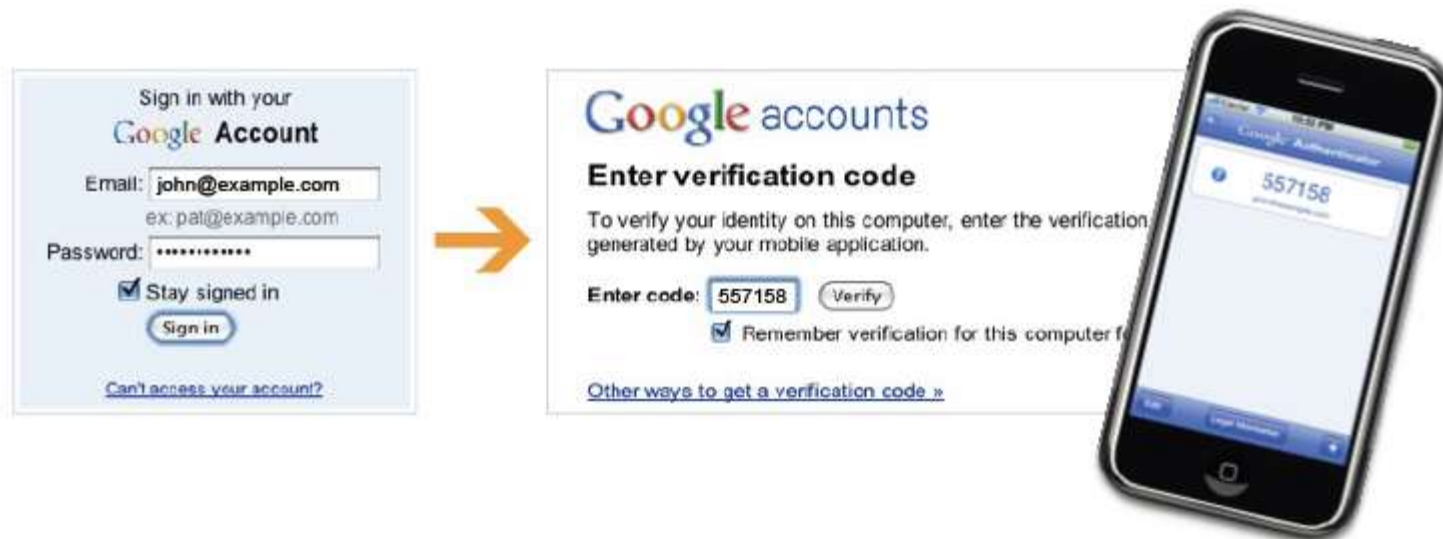
# Example: RSA SecurID

- **Security system using two-factor authentication**
  - Factors: knowledge (password) and ownership (device)
  - Device generates authentication code every 60 seconds
  - Authentication using *password and current code*
- **Code Generation**
  - Device initialized for each user with seed (random number)
  - Code computed from seed and current time (~one-time password)



# Example: Google 2-Step Verification

- **Security system by Google similar to SecureID**
  - Factors: knowledge (password) and ownership (phone)
  - Authentication code computed on mobile phone
  - Login at Google requires password and current code



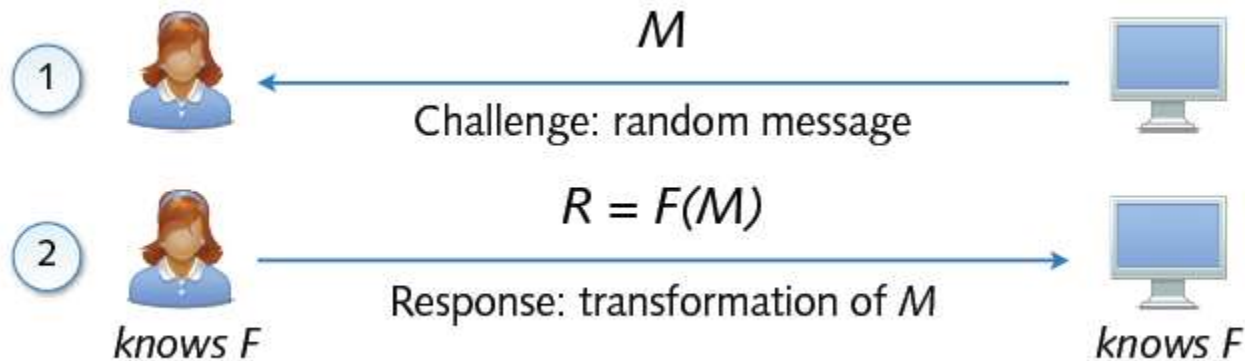
<https://blog.duosecurity.com/2013/02/bypassing-googles-two-factor-authentication/>





# Challenge-Response

- **Generic protocol scheme for authentication**
  - System and user share a secret function  $F$



- **Advantages over naive authentication methods**
  - Secret, e.g. password, is never transmitted in cleartext
  - Replay attacks against authentication not possible

# Challenge-Response (con't)

- **Secret function often parameterized by password**
  - $F = H(M + P)$  hash function  $H$  and password  $P$
  - $F = E_P(M)$  encryption function  $E$  and password  $P$
  - Hard to deduce  $P$  if  $F$  is cryptographically strong
- **Several methods related to challenge-response scheme**
  - One-time passwords
    - = challenge (index of password); response (password)
  - SecurID / Google 2-step
    - = challenge (current time); response (authentication code)



# Example: WPA2 (A Short Excursion)



# Wireless Networks

**Wireless network**  
e.g. WLAN (802.11)



- **Inherent security problems with wireless networks**
  - Communication over shared medium (air)
  - No physical access control and protection
  - Need for additional security measures (WEP, WPA, ...)



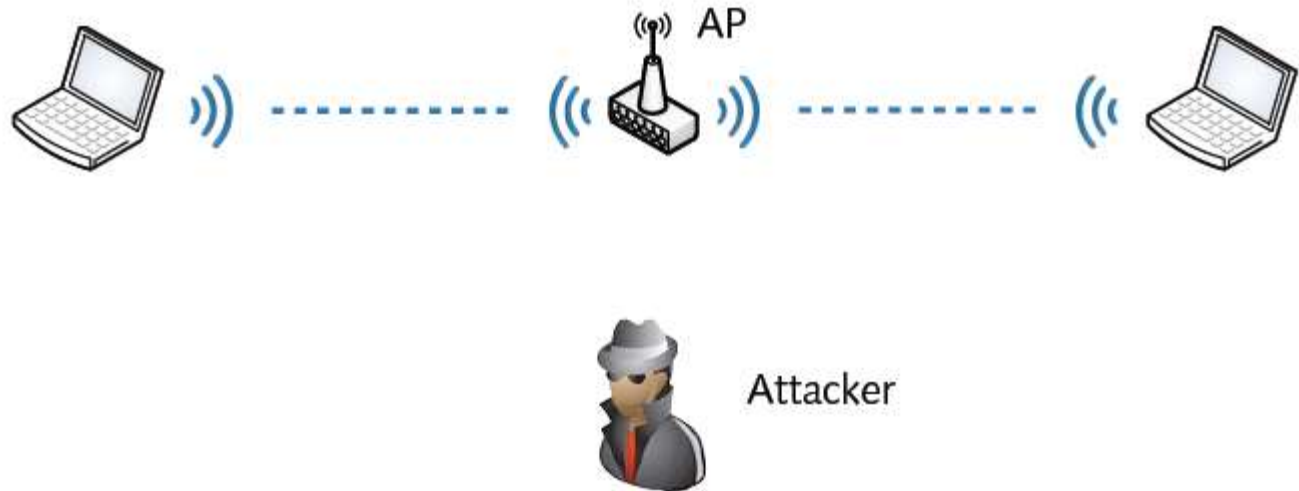
# A Closer Look at Attacks

- **Common attacks types**

- Masquerading and spoofing
- Eavesdropping of communication
- Tampering of messages

## Countermeasures

- ⚡ Authentication
- ⚡ Encryption
- ⚡ Integrity checks



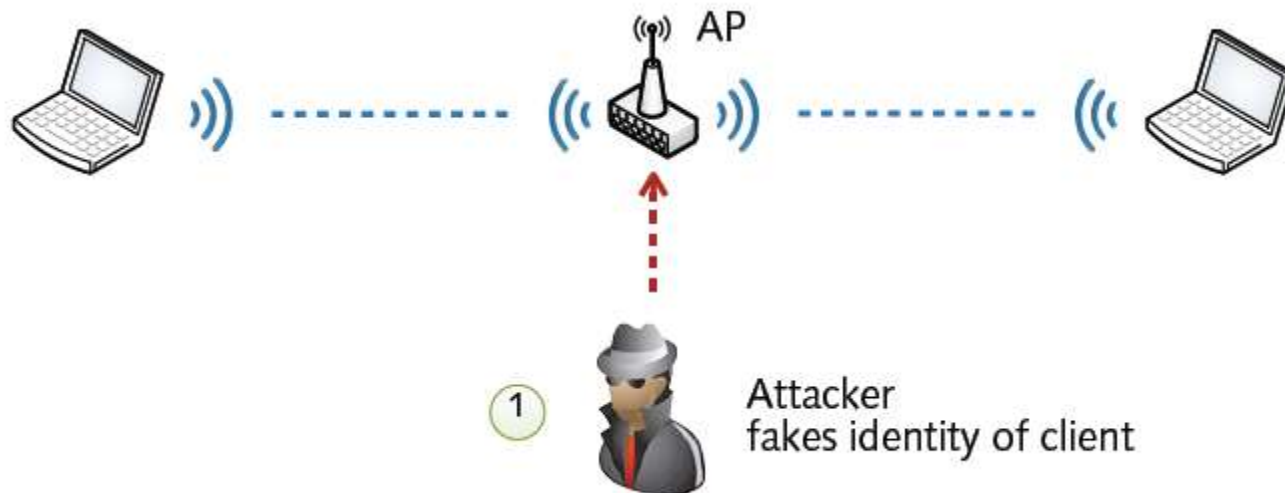
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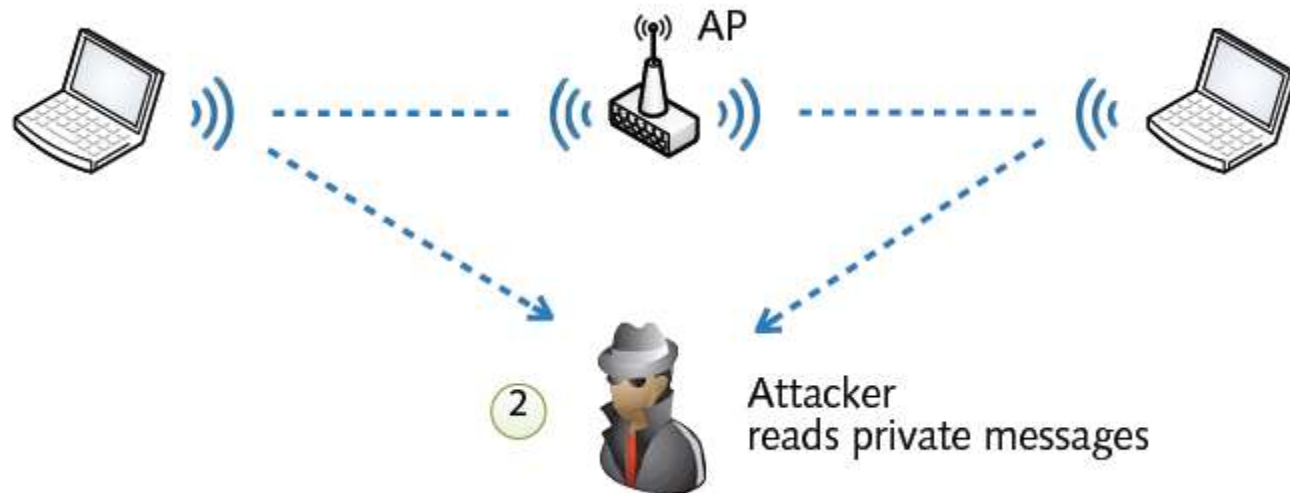
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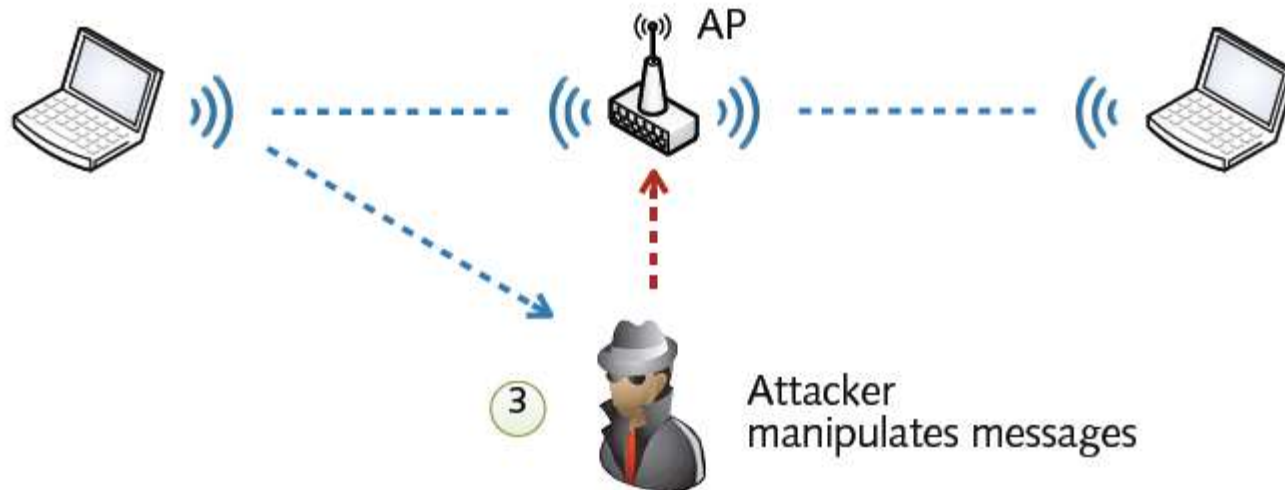
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# 802.11 and Security

	Authentication	Encryption	Integrity check
1997	<b>WEP</b> (Wired Equivalent Privacy) 802.11a		
	Shared keys medium	RC4 weak	CRC-32 weak
2003	<b>WPA</b> (Wi-Fi Protected Access) subset of 802.11i		
	Shared keys / 802.1x strong	TKIP medium	Michael medium
2004	<b>WPA2</b> (Wi-Fi Protected Access 2) 802.11i		
	Shared keys / 802.1x strong	AES-CCMP strong	AES-CCMP strong

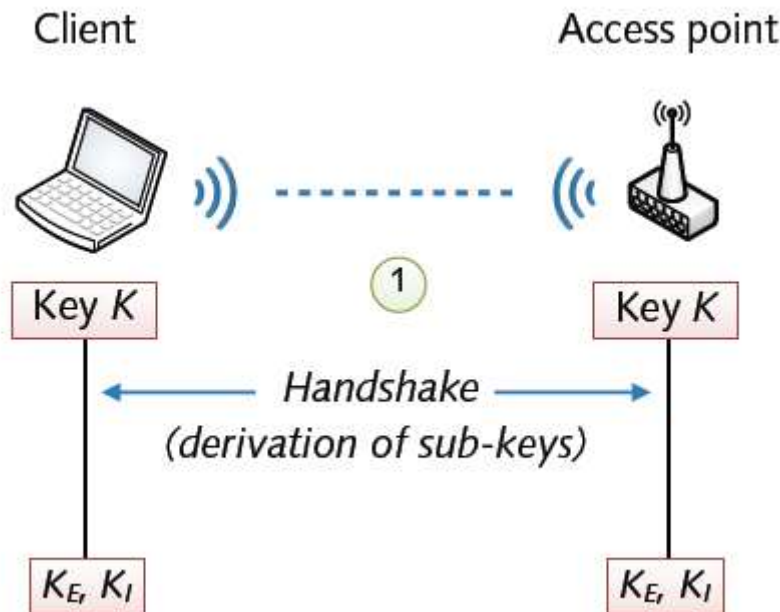
TKIP = Temporal Key Integrity Protocol

AES-CCMP = Counter Cipher Mode with Block Chaining Message Authentication Code Protocol



# WPA2 Authentication

- **Two different modes for authentication in WPA2**
  1. **Personal: Pre-shared keys (PSK) (aka “passwords”)**
  2. **Enterprise: 802.1x with Extensible Authentication Protocol**
- **Key Exchange:** The client and access point start with a **shared key (Key K)**.
- **Four-way Handshake:** This process generates unique sub-keys for encryption and integrity protection between the client and access point.
- **Sub-keys ( $K_e$ ,  $K_i$ ):** Derived from the original shared key, these sub-keys are used to encrypt and secure the communication between the devices.



# Chat gpt content for above slide

- The above slide explains the two different authentication modes in WPA2 and provides an overview of the handshake process involved in establishing a secure connection between a client (like a laptop) and an access point (e.g., a Wi-Fi router).
- **1.WPA2 Authentication Modes:**
- **Personal Mode:**
  - Uses **Pre-shared Key (PSK)**, which is commonly referred to as a "password."
  - This mode is typically used in home or small office networks.
  - The PSK is shared among all devices on the network, and each device must know this password to connect.
- **Enterprise Mode:**
  - Uses **802.1x with Extensible Authentication Protocol (EAP)**.
  - Commonly used in larger organizations, it provides stronger security than PSK.
  - In this mode, each user can have individual login credentials, which are validated through a central authentication server. This enables better control and tracking of users on the network.

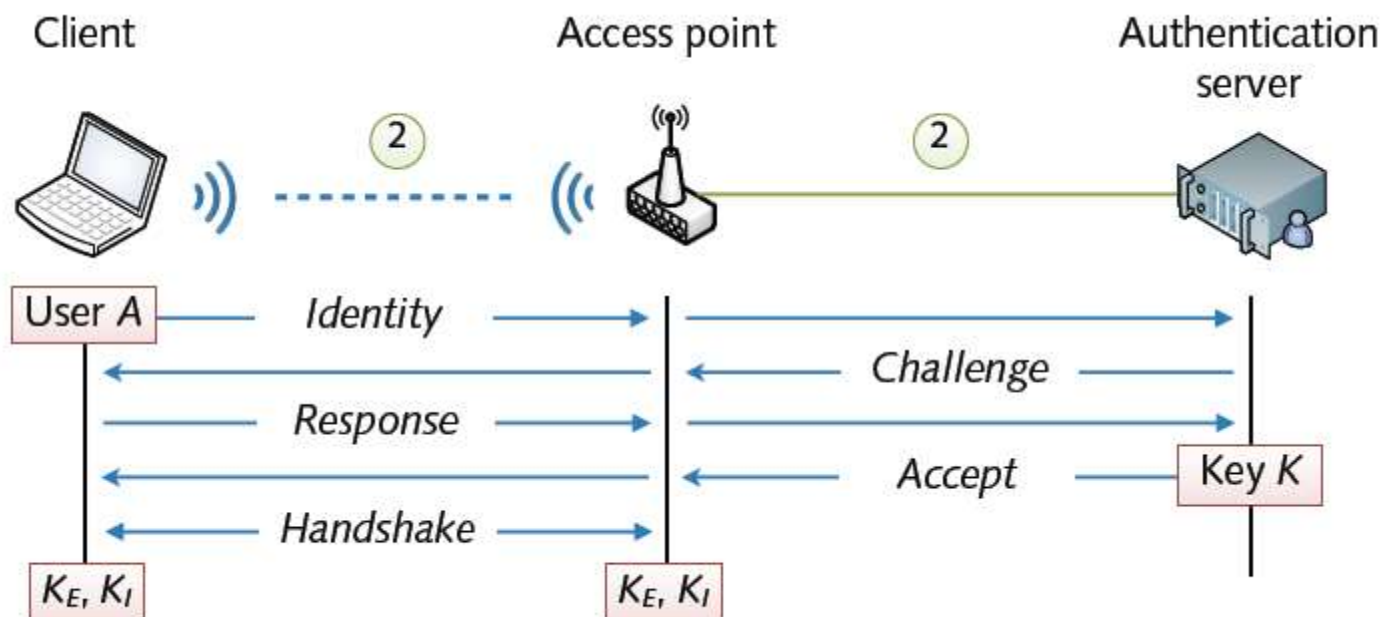
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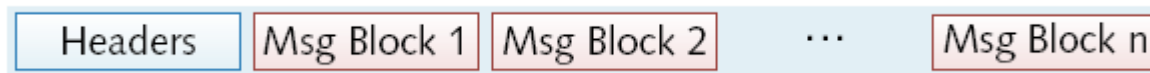
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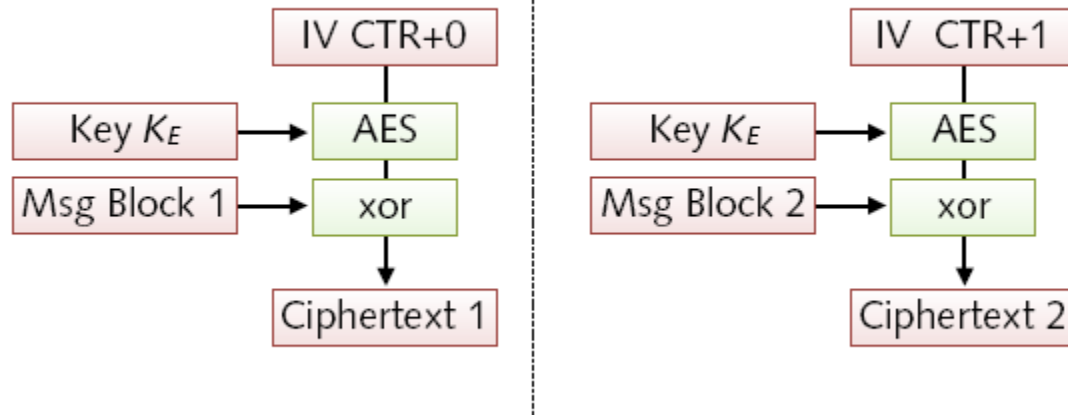


# WPA2 Encryption

- **Partitioning of each message in blocks**

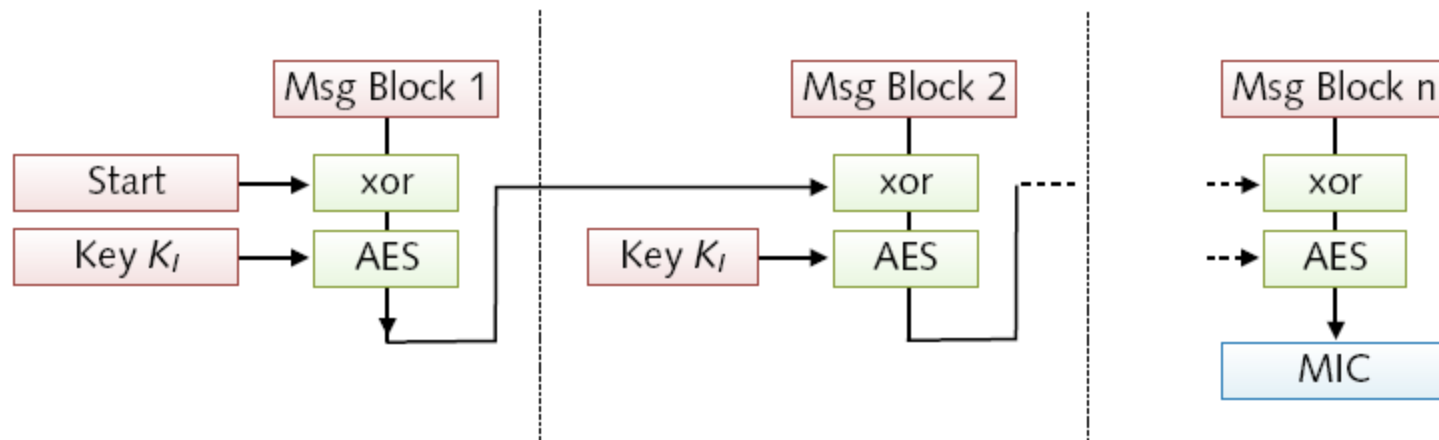


- **Encryption of each message block in counter mode**
  - Advanced Encryption Standard (AES) using key  $K_E$

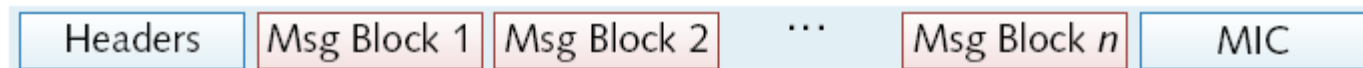


# WPA2 Integrity Check

- **Chaining of cipher blocks to a keyed hash value**
  - Message Integrity Code (MIC) using key  $K_I$



- **MIC appended to message prior to encryption**



# How secure is WPA2?

- **Attacks against WPA2**
  - (Almost) no attacks against cryptographic protocol
  - Best attack strategy so far: *brute-force attacks*
  - Target for potential attacks: Complexity of protocol
- **WPA2 security in practice**
  - Strength of passphrase in personal mode
  - Strength of authentication protocol in enterprise mode





# Access Control



# Access Control



- **Authorization and access control**
  - Control of what a subject is allowed to do
  - Management of permissions and capabilities
  - Often tight coupling with authentication
- **Examples**
  - Execution of programs, reading of files, ...

# Access Control Matrix

		Objects			
		File 1	File 2	Process 1	Process 2
Subjects	User 1	read		control, send	receive
	User 2	write	read	send	
	User 3	read, execute	read	control	control

Access control lists

Capabilities

- **Classic and simple representation for access control**
  - Mapping from subjects and objects to permissions



# Access Control Models

- **Access control non-trivial in practice**
  - Complex systems  $\rightsquigarrow$  complex access control models
- **Some characteristics of access control models**
  - Definition of objects and subjects  
E.g. subjects can be users, processes or hosts
  - Representation of permissions  
E.g. columns (access control lists), rows (capabilities)
  - Management of permissions  
E.g. discretionary, mandatory or role-based access control



# Representation: Access Control Lists

- **Access control lists (ACL)**
  - Attachment of permissions to objects (columns)
  - $\oplus$  Efficient and decentralize organization of permissions
  - $\ominus$  Listing of subject permissions very involved
- 
- ***Example: OpenBSD packet filter***
    - Deny access to the SSH service from any host
    - $\rightarrow$  *block in quick proto tcp from any to any port ssh*



# Representation: Capabilities

- **Capabilities**
- Attachment of permissions to subjects (rows)
- $\oplus$  Listing and control of subject permissions simple
- $\ominus$  Fine-grained permissions difficult to implement
- ***Example: Linux capabilities***
  - Restrict permissions to reboot system and load modules
  - $\rightarrow$  *lcap -z CAP\_SYS\_BOOT CAP\_SYS\_MODULE*



# Management of Permissions

- **Discretionary Access Control (DAC)**
  - Owner of an object controls access
  - Convenient but insecure if object changes owner
- **Mandatory Access Control (MAC)**
  - System globally enforces access control
  - Very secure but tedious to design and operate
- **Role-based Access Control (RBAC)**
  - System enforces access control using roles
  - In-between DAC and MAC models



# Example: UNIX Permissions

- **Discretionary access control of files**
  - Owner manages permissions of his files
- **Fixed-size access control lists: rwx rwx rwx**
  - Three subjects: user, group and other
  - Three permissions: read, write and execute

-rw-r-----	1	root	shadow	4321	17	Aug	00:23	/etc/shadow
-rw-r--r--	1	root	root	5086	16	Aug	00:20	/etc/passwd

Permissions	Owner	Group
-------------	-------	-------

- (a) Everybody can read the passwd file; root can write to it
- (b) Only root and the group shadow can read the shadow file





# Example: UNIX Permissions (con't)

- **Simple notation for management of permissions**
  - $\langle subjects \rangle + / - = \langle permissions \rangle$
  - Subjects: *u* (user), *g* (group), *o* (others), *a* (all)
  - Permissions: *r* (read), *w* (write), *x* (execute)
- **Examples**
  - Make file readable to everyone: *chmod a+r file*
  - Remove write permission from group: *chmod g-w file*
  - Make file readable by user only: *chmod u=r file*
- ***Alternative for UNIX gurus: octal encoding***



# Special Permissions

- **Some permissions with special semantics**
  - +x makes directories searchable
  - +t sticky bit (for directories deletion is restricted)
  - +s suid bit (change user id to file owner during execution)
- **A UNIX backdoor from the 1990s**

```
# cp /bin/sh /tmp/.backdoor  
# chown root:root /tmp/.backdoor  
# chmod u+s /tmp/.backdoor
```

- If it's bad, why do we need the suid bit?



# Acknowledgements

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