



Noakhali Science & Technology University

Noakhali-3814

Assignment On

Chapter – 12,13

Course Title : Information Security

Course Code : CSE 2205

Institute of Information Technology(IIT)

Submitted by:

Mir Mohammad Tahasin

Roll No : MUH2025007M

Year-2,Term-2

Institute of Information Technology

NSTU.

Submitted To:

Dr. Mohammad Nuruzzaman Bhuiyan

Assistant Professor

Institute of Information Technology

NSTU

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1. Write a routine (in pseudocode) to calculate the eighty constants in SHA-512 from the following table:

I	W_i
0..... 15	0,1,2,3,4,5,6,7,8,9,10,11,12, 13,14,15
16..... 31	0,4,8,12,1,5,9,13,2,6,10,14,3, 7,11,15

Question-01:

Solution:

(Routine in Pseudocode):

for i=0 to 79

 if i<=15

$w_i = i$

 else

$s_0 = (w_{i-15})$ right rotate 17

$s_1 = (w_{i-2})$ right rotate 19

$s_2 = (w_{i-7})$ right rotate 10

$w_i = s_0 \text{ xor } s_1 \text{ xor } s_2 \text{ xor } (w_{i-16})$

 end if

end for

Explanation:

SHA-512 is a hashing algorithm that performs a hashing function on some data given to it. Hashing functions take some data as input and produce an output (called hash digest) of fixed length for that input data. This output should, however, satisfy some conditions to be useful. So, SHA-512 does its work in a few stages. These stages go as follows:

1. Input formatting
2. Hash buffer initialization
3. Message Processing
4. Output

Here, the standard purposes a for circle to repeat through the numbers from 0 to 79. For every cycle, it first checks in the event that the ongoing worth of I is not exactly or equivalent to 15. Assuming it is, the consistent W_i is set to the worth of I . On the off chance that I is more prominent than 15, the standard purposes bitwise activities to compute W_i as the elite or (xor) of the aftereffects of turning the past upsides of W_i by specific sums, and afterward xor with W_{i-16} . The "right pivot" administrator turns the pieces of a number to the right by a predetermined number of positions, disposing of pieces that are moved off the end and bringing the pieces that were on the passed on finish to the right end.

Question-02:**Solution:****Signature:**

This process works as follows

1. The sender selects a random number r
2. The sender computes the first signature s_1 using $s_1 = \text{crimodp}$

3. The sender computes the second signature s_2 using the equation

$$s_2 = (M - dXs_1)Xr^{-1} \bmod (p-1)$$

Where P = large prime number

M = original message that needs to be signed

Here in the above problem :

Handwritten calculations on a piece of paper:

$$m = 320$$
$$s_1 = e_1^m = 2^{320} = 2083 \bmod 3119$$

we know,

$$\text{Second Signature, } s_2 = (m - dXs_1) \times r^{-1} = (320 - 127 \times 2083) \times 307^{-1}$$
$$= 2105 \bmod 3119$$

Alice send s_1, s_2 , and m to Bob. Bob uses the public key to calculate v_1 and v_2 .

$$v_1 = e_1^m = 2^{320} = 3006 \bmod 3119$$
$$v_2 = d^{s_1} \times s_2^{s_2} = 1702^{2083} \times 2083^{2105} = 3006 \bmod 3119$$

The verification part follows:

Verification:

This process works as follows.

1. The receiver performs the 1st part of verification called v_1 using the equation

$$V_1 = eM_1 \bmod P$$

2.The receiver performs the 2nd part of verification called as v2 using the equation

$$v2 = e s_1^2 s_2^{-1} \bmod p$$

Eg

$$V1 = e M_1 \bmod p = 1014 \bmod 19 = 16$$

And

$$V2 = e s_1^2 s_2^{-1} \bmod p = 43 * 34 \bmod p = 5184 \bmod 19 = 16$$

So, $v1 = v2$, the signature is valid.

