

## Lecture-3(a)

### Coin change:

1. Using DP:

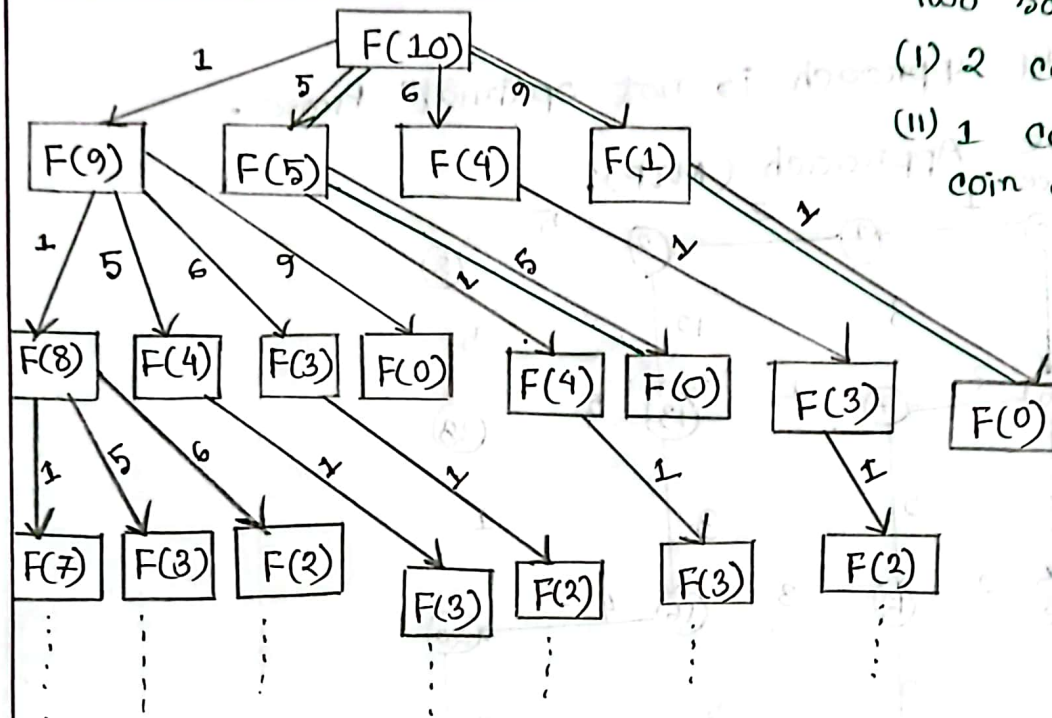
$C = \{1, 5, 6, 9\}$  Sum = 10

	0	1	2	3	4	5	6	7	8	9	10
1	0	1	2	3	4	5	6	7	8	9	10
5	0	1	2	3	4	min(5, 1+0) 1	min(6, 1+1) 2	min(7, 1+2) 3	min(8, 1+3) 4	min(9, 1+4) 5	min(10, 1+5) 2
6	0	1	2	3	4	1	min(2, 1+0) 1	min(3, 1+1) 2	min(4, 1+2) 3	min(5, 1+3) 4	min(2, 1+4) 2
9	0	1	2	3	4	1	1	2	3	min(4, 1+0) 1	min(2, 1+1) 2

$$5 - 5 = 0$$

Solution: 2 coin of "5"

### (1) Recursive Solution:



Two solve:

(1) 2 coin of "5"

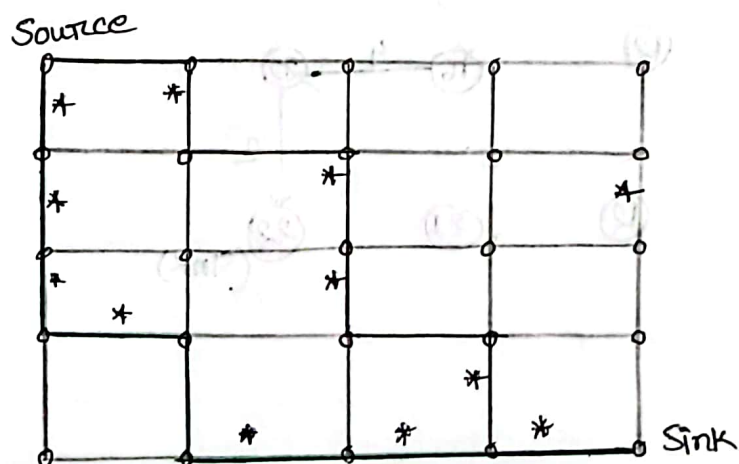
(11) 1 coin of "5" & 1 coin of "1"

### Manhattan Tourist Problem (MTP):

Seeking a path (from source to sink/destination) to travel (only eastward and southward) with most number of attractions (\*) in the Manhattan grid.

constraint: Only eastward & southward.

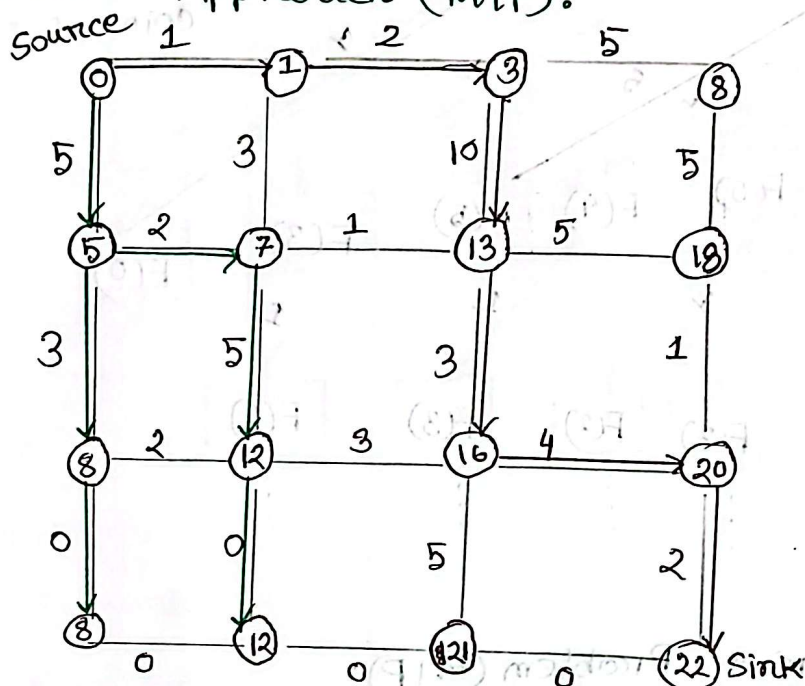
- The graph will always be directed acyclic graph.



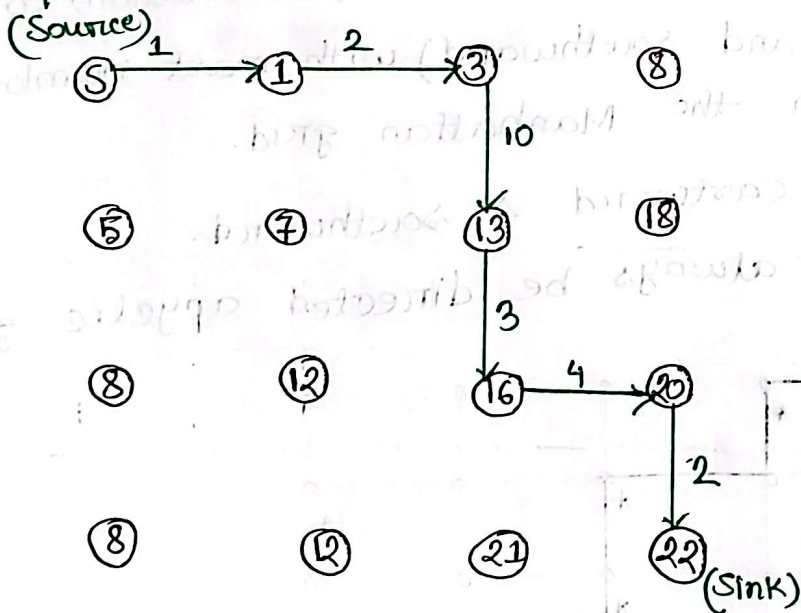
- Every intersection is a node and path is 'edge'.

~~14~~ "Greedy approach is not optimal" Prove.

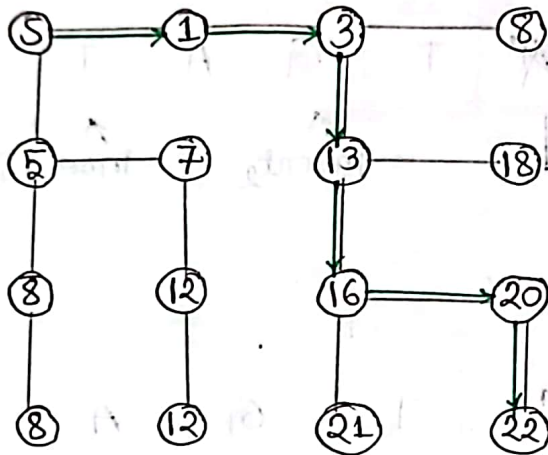
Recursive Approach (MTP):



Graph:



## Dynamic Programming (MTP):



## Alignment : 2 row representation:

V = A T C T G A T G  $n=8$

W = T G C A T A C  $m=7$

Let,

upper gap = deletion

lower gap = insertion

→ Deletion = 2

V	A	T	-	C	-	T	G	A	T	G
W	-	T	G	C	A	T	-	A	-	C

Mismatch = 1

Matching = 4

→ Insertion = 3



V=	A	T		C		T	G	A	T	G
W=	A	T	<del>X</del>	C	<del>X</del>	T	G	A	T	C

↑  
Insert<sub>1</sub>

↓  
Delete=2

↑  
Insert<sub>2</sub>

↑  
Insert<sub>3</sub>

W \ V	-	A	T	C	T	G	A	T	G
-	0	0	0	0	0	0	0	0	0
T	0	0	1	1	1	1	1	1	1
G	0	0	1	1	1	2	2	2	2
C	0	0	1	2	2	2	2	2	2
A	0	1	1	2	2	2	3	3	3
T	0	1	2	2	3	3	3	4	4
A	0	1	2	2	3	3	4	4	4
C	0	1	2	3	3	3	4	4	4

LCS = T C T A

↓  
Edit Distance

← ↑ ↑ ↑ ↑ ↑ ← ↑ ↑ ← ←  
V = A T - C - T G A - T G  
W = - T G C A T - A C - -

LCS does not show mis-match. That's why LCS is not preferable.

V = A T C T G A T C

W = T G C A T A C

W \ V	-	A	T	C	T	G	A	T	C
-	0	0	0	0	0	0	0	0	0
T	0	0	1	1	1	1	1	1	1
G	0	0	1	1	1	2	2	2	2
C	0	0	1	2	2	2	2	2	3
A	0	1	1	2	2	2	3	3	3
T	0	1	2	2	3	3	3	4	4
A	0	1	2	2	3	3	4	4	4
C	0	1	2	3	3	3	4	4	5

LCS = T C T A C

Edit Distance

← ↖ ↑ ↖ ↑ ↖ ← ↖ ← ↖  
 V = A T - C - T G A T C  
 W = - T G C A T - A - C

## Hamming Distance (Aligning Sequence without insertion and deletion):

P = N A F I S A

Q = N A B I L A

$$d_H = (0 + 0 + 1 + 0 + 1 + 0) = 2$$

- Sequence

- If two sequence's character match then '0' otherwise '1'.

## Edit Distance:

- Index by index always *ইনেক্স কন্ট্রোল*

- Count the insertion & deletion.

i<sup>th</sup> letter of V with

j<sup>th</sup> letter of W

V = A T A T A T A T

W = T A T A T A T A

$$d_H = (1 + 1 + 1 + 1 + 1 + 1 + 1 + 1) = 8$$

If shift W by 1

V = A T A T A T A T

W = - T A T A T A T A

$$d_H = (0 + 1 + 0 + 0 + 0 + 0 + 0 + 0 + 1) = 2$$

Formula of Edit distance:

$$\text{Edit distance}(v, w) = |v| + |w| - 2|LCS(v, w)|$$

→ distance of  $v$  &  $w$  will be count without gap.

~~Q~~ T G C A T A T → A T C C G A T in 5 steps

1. T G C A T A T → delete last T
2. T G C A T A → delete last A
3. T G C A T → Insert A at front
4. A T G C A T → substitute C for 3rd G
5. A T C C A T → Insert G before last A
6. A T C C G A T → Done



## Lecture-3(b)

Indel  $\rightarrow$  Insertion or deletion

Global alignment:

$$S_{i,j} = \max \begin{cases} S_{i-1,j-1} + 1 & \text{if } v_i = w_j \\ S_{i-1,j-1} - \mu & \text{if } v_i \neq w_j \\ S_{i-1,j} - \sigma & \text{Indel} \\ S_{i,j-1} - \sigma & \text{Indel} \end{cases}$$

Scoring Schema:

+1 : match premium

$-\mu$  : mismatch penalty

$-\sigma$  : indel penalty

#we scoring 1 for match & zero (0) for indel.

#match  $-\mu$  (#mismatch)  $-\sigma$  (#indel)

Scoring Matrix (Mismatch count):

W = C T A G C

V = C A G T C

Let, match = +2

mismatch = -1 =  $-\mu$

Indel = -3 =  $-\sigma$

V \ W	-	C	T	A	G	c
-	0	-3	-6	-9	-12	-15
C	-3	2	-6	-4	-9	-7
A	-6	-4	-1	1	-4	-7
G	-9	-7	-4	-2	0	-3
T	-12	-10	-7	-2	-3	-1
c	-15	-10	-8	-5	-3	2

$\nwarrow \quad \leftarrow \quad \nwarrow \quad \nwarrow \quad \uparrow \quad \nwarrow$   
 $W = \quad C \quad T \quad A \quad G \quad - \quad c$   
 $V = \quad c \quad - \quad A \quad G \quad T \quad c$

$$\text{Score} = 2 - 3 + 2 + 2 - 3 + 2$$

$$= 2 = \text{Final Score.}$$

### Local Alignment:

Here only consider diagonal arrows.

$$S_{i,j} = \max \begin{cases} S_{i-1,j-1} + 1 & ; \text{ if } v_i = w_j \\ S_{i-1,j-1} - \mu & ; \text{ if } v_i \neq w_j \\ S_{i-1,j} - \sigma \\ S_{i,j-1} - \sigma \end{cases} \text{ Indel}$$

Example:

V = A T C G

W = T C C

Let, match = +1

mismatch,  $\mu = +1$

indel,  $\sigma = +2$

$\begin{matrix} V \\ W \end{matrix}$	-	A	T	C	G
-	0	0	0	0	0
T	0	-1	1	-1	-1
C	0	-2	-2	2	-1
C	0	-1	-2	1	1

- এখানে কোনো box এ অর (-) value হলে আমরা ২য়০ বসাবো। এবং নাহলে max value হবে।

Find out the max and diagonal from it.

Multiple max value হলে অবশ্যই নিচের row হতে right এর column এ max choose করতে হবে।

$$V = \begin{matrix} A & T & C & G \\ & | & | & \\ W = & T & C & C \end{matrix}$$

### Gap Penalties:

1. Constant: (consider consecutive gap as 1 gap)

T A C C T A G

T - - - T A - - -  
 $\rightarrow$  3 consecutive gaps as 1.

Let,

match = +1

gap/indel = -1

$$\text{Score} = 1 - 1 + 1 + 1 - 1 \\ = 1$$

2. Linear (Least Realistic):

T A C C T A G

T - - - T A -

(Consider every gap individually).

Let,

match = +1

gap/indel = -1

$$\text{Score} = 1 - 1 - 1 - 1 + 1 + 1 - 1 \\ = -1$$



### 3. Affine:

- Gap opening: multiple gap পরপর হলে 1st gap ডিজেনা,
- Gap Extension: multiple gap এর Gap opening বাড়ানো যাকি ডিজেনা।

T A C C T A G

T T T T A T

← gap extension penalty ( $\sigma$ ) → Gap opening penalty ( $\rho$ )

Let,

match = +1,  $\rho = -2$ ,  $\sigma = -1$

$$\text{Score} = 1 - 2 - 1 - 1 + 1 + 1 - 2 \\ = -3$$

$$\text{Gap penalty} = \rho + \sigma(k-1)$$

$\Rightarrow k$  = number of consecutive gap

$$\text{if } k=3, \text{ gap penalty} = -2 + (-1)(3-1)$$

$$= -2 + (-1)(2)$$

$$= -2 - 2 = -4$$

$$\text{if } k=1, \text{ gap penalty} = -2 + (-1)(1-1)$$

$$= -2 + (-1)(0)$$

$$= -2 + 0$$

$$= -2$$