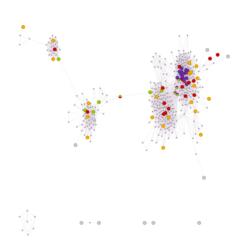
The Stochastic Block Model and Module Detection

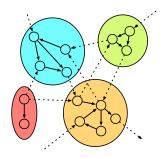
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January 16, 2014

Overview

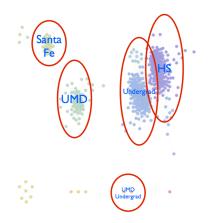
- Modules and Module Detection
- Vanilla Stochastic Block Model
- 3 Inference for the Stochastic Block Model
- 4 Extensions of the Stochastic Block Model



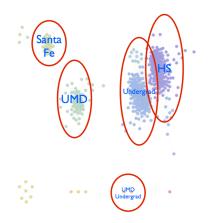


Basic Idea:

- A module or community is a collection of nodes defined by how its edges behave:
 - Edge Density: For social networks, we expect edge density to be greater within a community than without. (Assortative Community)
 - **Edge Weight:** For coexpression networks, we expect the correlations to be higher within a functional module than without.
 - Etc.



- The goal of module detection is to partition the network such that each module has a similar distribution over the nodes of within-community edge weights and without-community edge weights.
- The resulting partition will consist of K sets C_k , k = 1, ..., K, of nodes.
 - $\bigcup_{k=1}^{K} C_k = V = \{1, 2, 3, \dots, n\}$
 - $C_i \cap C_j = \emptyset, \quad i \neq j$
- Generalizations allow for coverings (partitions where we allow overlap), mixed membership, etc.



Questions:

- For a fixed K, how do we decide on a partition?
 - We need some sort of *goodness-of-fit* / *loss* function.
 - This tells us if our partition 'makes sense' / explains the data well.
- How do we choose K?
 - K is a tuning parameter, controls the flexibility of our partition.
 - Usual model checking (leave-one-out cross-validation, etc.) requires a bit of finessing due to the dependencies in the model.

Approaches:

- Modularity maximization
 - Choose a particular loss function based on a null model for the network (called the configuration model).
- Stochastic Block Models
 - Specify a probabilistic model for the network, and use standard techniques from statistical inference.
- Ad hoc / heuristic approaches
 - Try things out empirically and hope for the best.

Some notation:

- Let $G = (V, \mathbf{A})$ be a graph / network where:
 - $V = \{1, 2, ..., n\}$ indexes the nodes (vertices) in the network.
 - A is the (binary) adjacency matrix associated with G, such that

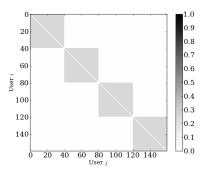
$$(\mathbf{A})_{uv} = a_{uv} = \left\{egin{array}{ll} 1 & : ext{ an edge exists between } u ext{ and } v \ 0 & : ext{ otherwise} \end{array}
ight.$$

Erdös-Rényi Random Graph:

- We imagine that our network is a realization of a random network where each edge a_{uv} is a realization of a Bernoulli random variable A_{uv} with bias p.
 - $A_{uv} \stackrel{\text{iid}}{\sim} \text{Bernoulli}(p)$.
- Thus, the number of edges connected to u is Binomial(n, p).
 - For each possible edge incident to u, of which there are n (including self-loops), we flip a coin.
 - Coin flips are 'governed' by the Binomial distribution.
- For small p and large n, we can approximate a Binomial(n, p) distribution with a Poisson(np) distribution.
 - Hence all the claims in Clauset's notes that the degree distribution of an Erdös-Rényi random graph is Poisson, since for real world networks n will be large and p will be small (assuming the network is sparse).

Stochastic Block Model:

 Basically the same model as an Erdös-Rényi random graph, but we allow p to vary between blocks of nodes. (Hence the name.)



Parameters of the Binary Stochastic Block Model:

- *K*, the number of communities / modules.
- $\mathbf{z} = (z_1, z_2, \dots, z_n)$, a vector giving the community membership for each node.
 - i.e. $z_u \in \{1, \dots, K\}, u = 1, \dots, n$.
- **M**, a $K \times K$ matrix where $m_{ij} = (\mathbf{M})_{ij}$ gives the probability of an edge between a node in community i and a node in community j.
 - i.e. For an edge a_{uv} , we use **z** to index into **M**, and read off $P(A_{uv}=1)=m_{z_uz_v}$.

See Clauset's notes for examples.

Using a Stochastic Block Model for Simulation:

- Given the parameters $\theta = (K, \mathbf{z}, \mathbf{M})$, we can easily simulate a graph from the model:
 - For each possible entry of **A**, e.g. between u and v, flip a coin with bias $m_{z_uz_v}$, and record an edge if it comes up 1.
- Similarly, we can write down the probability of generating any graph $G = (V, \mathbf{A})$, since each edge is independent:

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$$P(G = g; \theta) = \prod_{u,v} P(A_{uv} = a_{uv})$$

= $\prod_{u,v} m_{z_u z_v}^{a_{uv}} (1 - m_{z_u z_v})^{1 - a_{uv}}$

We'll need this for inference.

Forward vs. Inverse Probability:

- We know how to generate a network $G = (V, \mathbf{A})$ given $\theta = (K, \mathbf{z}, \mathbf{M})$.
- How do we do the opposite?
- Given $G = (V, \mathbf{A})$, how do we infer $\theta = (K, \mathbf{z}, \mathbf{M})$?
- As always: statistics.

Bayes v. Fisher:

- As with most problems in statistics, there are at least two ways:
 - The Frequentist Way: treat $\theta = (K, z, M)$ as fixed and G as random, and use maximum likelihood to get $\hat{\theta}$.
 - The Bayesian Way: treat θ as random and G as fixed, compute the posterior distribution of θ given G, and compute the posterior mean / median / mode for $\hat{\theta}$.
- It's interesting to think about what each of these interpretations mean in terms of a community / module.

Maximum Likelihood Estimation:

- There are two parts to Maximum Likelihood Estimation for the Stochastic Block Model:
 - Choose ẑ, a particular assignment of each of the nodes into the K communities.
 - This is the hard part. NP-complete in general cases.
 - Once we have chosen a $\hat{\mathbf{z}}$, estimate $\hat{\mathbf{M}}$.
 - This is easy, and we can write down the answer.
- Clauset sidesteps the first part. Approximation methods exist.
 - The Expectation-Maximization (EM) algorithm works here.

The Maximum Likelihood Estimator for M:

- See Clauset's notes for details.
- Let N_{ij} be the number of edges between community i and community j. Then

$$\hat{m}_{ij} = \frac{N_{ij}}{n_{ij}}$$

where n_{ij} is the number of possible edges between communities i and j.

• i.e. Compute the proportion of **potential** edges between *i* and *j* that show up in the network.

The Log Likelihood for the Stochastic Block Model:

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$$egin{aligned} P(G=g; heta) &= \prod_{u,v} P(A_{uv} = a_{uv}) \ &= \prod_{u,v} m_{z_u z_v}^{a_{uv}} (1 - m_{z_u z_v})^{1 - a_{uv}} \ &= \prod_{i,j \in \{1,\dots,K\}} m_{ij}^{N_{ij}} (1 - m_{ij})^{n_{ij} - N_{ij}} \end{aligned}$$

•

$$\log P(G; \hat{\theta}) = \sum_{i,j} N_{ij} \log \frac{N_{ij}}{n_{ij}} + (n_{ij} - N_{ij}) \log \left(\frac{n_{ij} - N_{ij}}{n_{ij}}\right).$$

 Allows us to measure the goodness-of-fit of the stochastic block model with respect to the observed network.

Extensions of the Stochastic Block Model

Problems with the Vanilla Stochastic Block Model:

- Degree distribution of any node will always be a mixture of Binomials (approximately a mixture of Poissons).
- To get long-tailed ('power law') degree distributions, the community memberships have to take a very particular form.
- Instead, allow each edge to have a 'propensity' of connecting to other nodes, γ_u , $u=1,\ldots,n$.
- This allows each node to have its own expected degree, which also depends on its community membership.
- This model is called the degree-corrected stochastic block model, and adds an additional parameter $\gamma = (\gamma_1, \dots, \gamma_n)$ to the model.

Extensions of the Stochastic Block Model

Networks with Weighted Edges:

- Instead of having the edges take on binary or non-negative integer values, we assume that edges are drawn from some continuous distribution, with the distribution differing between and within different communities.
- For example, we could assume we have a weighted adjacency matrix
 W where

$$W_{uv} \sim N(\mu_{z_u z_v}, \sigma_{z_u z_v}^2).$$

• We have to work a bit harder to infer $\hat{\mathbf{W}}$, but (very recent) methods do exist.