ACM CCS'17 Tutorial SGX Security and Privacy

Taesoo Kim (Georgia Tech), Zhiqiang Lin (UT Dallas) Chia-Che Tsai (UC Berkeley/Stony Brook)

> 11/02/2017 Dallas. Texas

Agenda

Why SGX

- Part 1: SGX Introduction and Applications (by Zhiqiang)
- Part 2: SGX Shielding Frameworks and Development Tools (by Chia-che)
- Part 3: Security Issues on SGX (by Taesoo)

References

SGX Introduction and Applications

Zhiqiang Lin

UT Dallas

Outline

- Why SGX
- SGX Introduction
 - Instructions and Data Structures
 - Software Development Model
 - Performance Overhead
- SGX Applications
 - Server Side Applications
 - Client Side Applications
 - Distributed Computing Applications
- Summary
- References

Outline

- Why SGX
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Summary

The Evolutions of Using Isolation for Malware Defense



Summary



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Virtualization



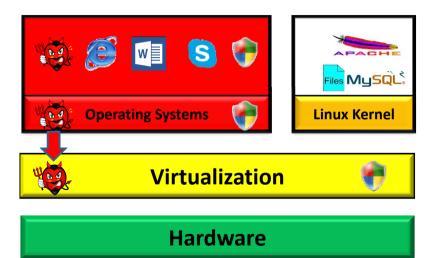


Virtualization

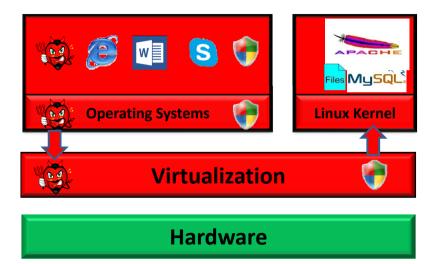


Summary

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Virtualization



Hardware



References







Virtualization









Virtualization



Summary

Hardware Assured Security

- Secure coprocessors [Yee94]
 - IBM 4758 [SW99]
- Aegis secure processor [SCG+03]
- Trusted Platform Module (TPM) [TPM03]
- Trust Zone [Alv04]
- AMD SVM [VD06]
- Intel Trusted Execution Technology (TXT) [FG13]

Hardware Assured Security

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Intel Software Guard eXtension (SGX) [MAB+13] Providing Hardware Assured Security w/ Minimized Attack Surface

Existing Computer Systems

Apps must trust

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- OS/VMM
- BIOS, SMM
- Trust relies on software

Summary

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Intel Software Guard eXtension (SGX) [MAB+13] Providing Hardware Assured Security w/ Minimized Attack Surface

Existing Computer Systems

- Apps must trust
 - OS/VMM
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Computer Systems w/ SGX

- Apps must trust
 - SGX hardware
- Trust excludes OS/VMM/BIOS/SMM

Intel Software Guard eXtension (SGX) [MAB+13] Providing Hardware Assured Security w/ Minimized Attack Surface

Existing Computer Systems

- Apps must trust
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Computer Systems w/ SGX

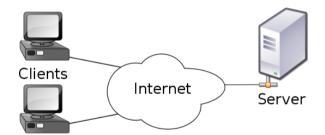
- Apps must trust
 - SGX hardware
- Trust excludes OS/VMM/BIOS/SMM

With SGX, for the first time, apps gain the ability to manage its own secret, without relying on the underlying systems software

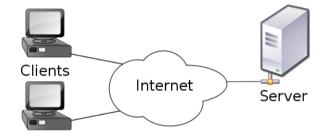
References

SGX provides a perfect platform for **Secure Remote Execution** [CD16]

SGX provides a perfect platform for **Secure Remote Execution** [CD16]



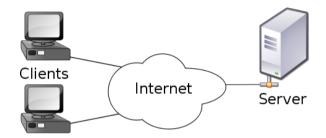
Why SGX





Client: End Users

- Cloud computing
 - (Secret-preserving) data analytics
 - Healthcare record processing



Client: End Users

- Cloud computing
 - (Secret-preserving) data analytics
 - Healthcare record processing

Server: Service Providers

- Computer game publishers
- Media streaming providers
- Software vendors (e.g., DRM)



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Why SGX



SGX	User Space	Kernel Space	Total
Version	enclu	encls	
SGX-v1	5	13	18
SGX-v2	8	16	24

SGX Instructions

Why SGX



SGX	User Space	Kernel Space	Total
Version	enclu	encls	
SGX-v1	5	13	18
SGX-v2	8	16	24

sysenter/sysexit Instructions

Why SGX

Privilege	Type	Instruction	Description	Version
Р	MEM	EADD	Add a page	v1
Р	MEM	EBLOCK	Block an EPC page	v1
Р	EXE	ECREATE	Create an enclave	v1
Р	DBG	EDBGRD	Read data by debugger	v1
Р	DBG	EDBGWR	Write data by debugger	v1
Р	MEM	EEXTEND	Extend EPC page measurement	v1
Р	EXE	EINIT	Initialize an enclave	v1
Р	MEM	ELDB	Load an EPC page as blocked	v1
Р	MEM	ELDU	Load an EPC page as unblocked	v1
Р	SEC	EPA	Add version array	v1
Р	MEM	EREMOVE	Remove a page from EPC	v1
Р	MEM	ETRACK	Activate EBLOCK checks	v1
Р	MEM	EWB	Write back/invalidate an EPC page	v1
Р	MEM	EAUG	Allocate page to an existing enclave	v2
Р	SEC	EMODPR	Restrict page permissions	v2
Р	EXE	EMODT	Change the type of an EPC page	v2
U	EXE	EENTER	Enter an enclave	v1
U	EXE	EEXIT	Exit an enclave	v1
U	SEC	EGETKEY	Create a cryptographic key	v1
U	SEC	EREPORT	Create a cryptographic report	v1
U	EXE	ERESUME	Re-enter an enclave	v1
U	MEM	EACCEPT	Accept changes to a page	v2
U	SEC	EMODPE	Enhance access rights	v2
U	MEM	EACCEPTCOPY	Copy page to a new location	v2

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enclu[EEXIT]:
 mov \$0x4,%rax
 enclu

SGX Data Structures

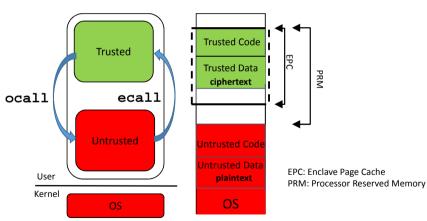
- SGX Enclave Control Structure (SECS)
- Thread Control Structure (TCS)
- State Save Area (SSA)
- Page Information (PAGEINFO)
- Security Information (SECINFO)
- Paging Crypto MetaData (PCMD)
- Enclave Signature Structure (SIGSTRUCT)
- EINT Token Structure (EINITTOKEN)
- Report (REPORT)
- Report Target Info (TARGETINFO)
- Key Request (KEYREQUEST)
- Version Array (VA)
- Enclave Page Cache Map (EPCM)



Summary

SGX Application Design, and Memory Layout

Application Design Memory Layout



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Two Ways of Having Software Protected by SGX

Approach-I: Developing SGX software from scratch through partitioning

- Partitioning software into trusted and untrusted component
- Using software development (e.g., SGX SDK [sgx]) tools to create an enclave module (shared object), which contains implementation for the trusted component
 - Defining enclave (executing the trusted component) interface
 - Using tools to generate stubs/proxies for ecalls and ocalls
 - Linking w/ SGX libraries
- Build and debug

Two Ways of Having Software Protected by SGX

Approach-I: Developing SGX software from scratch through partitioning

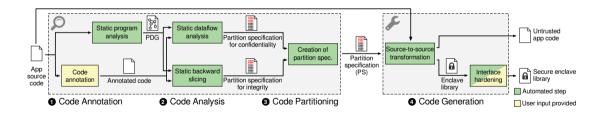
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Developing SGX software using a partitioning approach is what Intel intended



Teditious for the partitioning

Glamdring: Automatic App Partitioning for Intel SGX [LPM+17]



Source: https://www.usenix.org/system/files/conference/atc17/atc17-lind.pdf

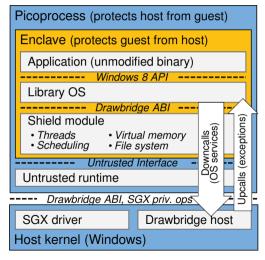
Two Ways of Having Software Protected by SGX

Approach-II: Enabling Legacy Software with SGX Protection w/o Parititioning

- SGX apps run at ring-3; system level code cannot be executed inside enclave.
- If SGX apps are executed with libOS (which is ring-3), then the problem gets solved
- No partitioning: the library OS approach
 - Haven [BPH14, BPH15]
 - @ Graphene-SGX [TPV17]

Summary

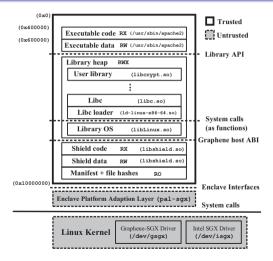
Haven [BPH14, BPH15]



References

Graphene-SGX [TPV17]

Why SGX



Two Ways of Having Software Protected by SGX

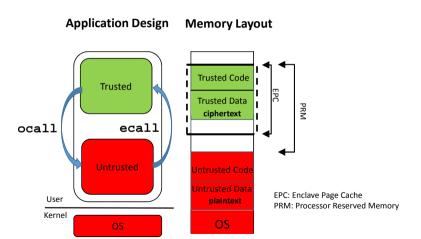
Approach	ТСВ	Effort	Library
Partitioning	Small	High	Intel SDK
No Partitoning	Large	Low	e.g., Graphene

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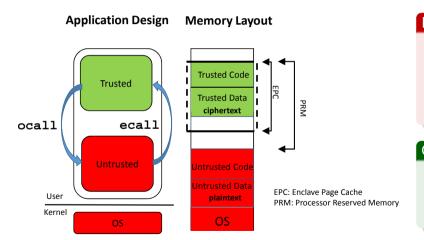


Performance Overhead



Performance Overhead

Why SGX



Runtime Overhead

- ecall
- 0 ocall
- enclave cache
- memory encryption

Other Overhead

- enclave creation
- enclave deletion
- SGX attestation

Instruction Level Overhead

Why SGX

Micro-benchmark	Description	Latency (cycles)
ecall (warm cache)	Calling an enclave w/o parameters, and immediately return	8,640
ecall (cold cache)	Same as above, but the entire cache is flushed	14,170
	to enclave	9,816
ecall (buffer)	Calling an enclave func, passing 2K buffer, from the enclave	11,172
	to and from enclave	10,827
ocall (warm cache)	Calling untrusted w/o parameters, and immediately return	8,314
ocall (cold cache)	Same as above, but the entire cache is flushed	14,160
	to untrusted	9,254
ocall (buffer)	Calling untrusted func, passing 2K buffer, from the untrusted	11,418
	to and from untrusted	9,801

Source: "Regaining Lost Cycles with HotCalls" [WBA17] http://www.ofirweisse.com/ISCA17_Ofir_Weisse.pdf

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Instruction Level Overhead

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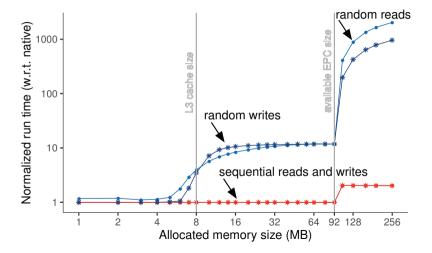
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Hyper-calls to the hypervisor take about 1,300 cycles (KVM hypervisor on x86) [DLL+16]. A round-trip time for an Exception-Less System gettsc is modest at 150 cycles [SS10].

Summary

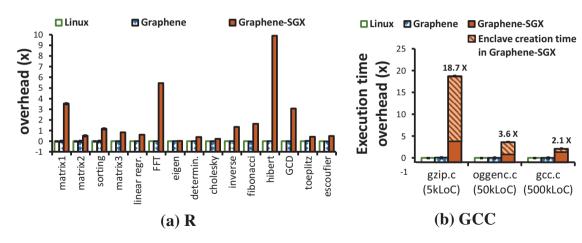
Memory Access Overhead Measured by SCONE [ATG+16]

Why SGX

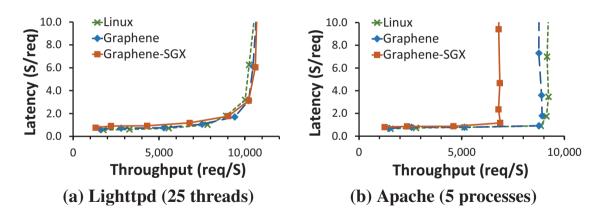


References

Desktop Applications: Tested by Graphene-SGX [TPV17]



Throughput of HTTP Servers: Tested by Graphene-SGX [TPV17]



Source: https://www.usenix.org/system/files/conference/atc17/atc17-tsai.pdf



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Cloud Applications

Why SGX



Source: https://www.rishabhsoft.com/wp-content/uploads/2014/10/Cloud_development_main.jpg



Paner Title

Server Side (e.g., Cloud) Applications

- rapel Title	venue	Application
Shielding Applications from an Untrusted Cloud with Haven	OSDI'14	App Hardening
VC3: trustworthy data analytics in the cloud using SGX	SP'15	Data Analytics
M2R: Enabling Stronger Privacy in MapReduce Computation	USENIX SEC'15	Data Analytics
Oblivious Multi-Party Machine Learning on Trusted Processors	USENIX SEC'16	Data Analytics
SCONE: Secure Linux Containers with Intel SGX	OSDI'16	App Hardening
SecureKeeper: Confidential ZooKeeper using Intel SGX	Middleware'16	Data Analytics
Attestation Transparency: Building secure Internet services for legacy clients	ASIACCS'16	Secure Service
S-NFV: Securing NFV states by using SGX	SDN-NFVSec'16	Cloud NFV
PANOPLY: Low-TCB Linux Applications with SGX Enclaves	NDSS'17	App Hardening
SGX-Log: Securing System Logs With SGX	ASIACCS'17	Securing Logging
Graphene-SGX: A Practical Library OS for Unmodified Applications on SGX	ATC'17	App Hardening
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SGX-BigMatrix: A Practical Encrypted Data Analytic Framework With SGX	CCS'17	Data Analytics
Enclave-Based Privacy-Preserving Alignment of Raw Genomic Information	SysTex'17	Genomic Computing

Venue Application

Venue

Application

Why SGX

Paper Title

raper ritte	venue	Application
Shielding Applications from an Untrusted Cloud with Haven	OSDI'14	App Hardening
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Venue

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Desktop Applications



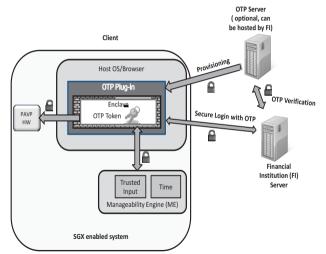
Client Side Applications

Why SGX

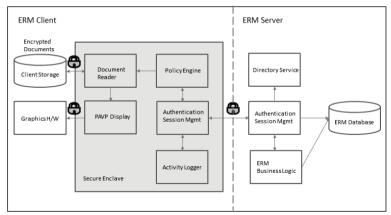
Paper/News Title	Venue	Application
Using Innovative Instructions to Create Trustworthy Software Solutions	Hasp'13	OTP, ERM, Video Conferencing
Password Manager with Intel SGX	Tutorial@16	Password Manager
Numecent to Show Off Pioneering Application Delivery Platform	IDF'16	DRM
A Case for Protecting Computer Games With SGX	SysTEX'16	Game Protection
TrustJS: Trusted Client-side Execution of JavaScript	EuroSec'17	Trusted Script Execution

Summary

Client Applications: One Time Pad (OTP)



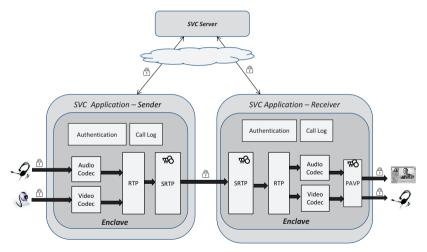
Client Applications: Enterprise Rights Management



Source: https://software.intel.com/sites/default/files/article/413936/hasp-2013-innovative-instructions-and-software-model-for-isolated-execution.pdf

Summary

Client Applications: Secure Video Chatting

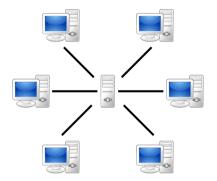


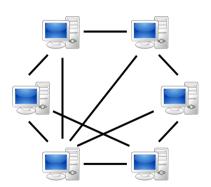
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Why SGX





Distributed Computing Applications

Why SGX

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Ryoan: a distributed sandbox for untrusted computation on secret data	OSDI'16	App Hardening
Proof of Luck: an Efficient Blockchain Consensus Protocol	SysTEX'16	Byzantine fault tolerance
Town Crier: An Authenticated Data Feed for Smart Contracts	CCS'16	Smart Contracts
Secure Content-Based Routing Using Intel SGX	Middleware'16	Content-Based Routing
Enhancing Security and Privacy of Tor's Ecosystem by Using TEE	NSDI'17	Tor nework
Hybrids on Steroids: SGX-Based High Performance BFT	EuroSys'17	Byzantine fault tolerance

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SGX Summary

A hardware assured TEE that secures applications w/ minimal attack surface

SGX Summary

- A hardware assured TEE that secures applications w/ minimal attack surface
- Two ways of developing SGX applications
 - Developing from scratch by partitioning applications
 - Executing the native (legacy) apps in an enclave w/o partitioning

Why SGX

- A hardware assured TEE that secures applications w/ minimal attack surface
- Two ways of developing SGX applications
 - Developing from scratch by partitioning applications
 - Executing the native (legacy) apps in an enclave w/o partitioning
- The major runtime overhead of SGX applications come from I/O, enclave encrypted memory access, and enclave cache misses.

SGX Summary

- A hardware assured TEE that secures applications w/ minimal attack surface
- Two ways of developing SGX applications
 - Developing from scratch by partitioning applications
 - 2 Executing the native (legacy) apps in an enclave w/o partitioning
- The major runtime overhead of SGX applications come from I/O, enclave encrypted memory access, and enclave cache misses.
- The killer applications of SGX is for securing remote execution, such as data analytics (e.g., MapReduce, Machine learning) in the cloud, and distributed computing (e.g., Tor, bitcoins).

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Why SGX

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