

Switching, Delays and Performance

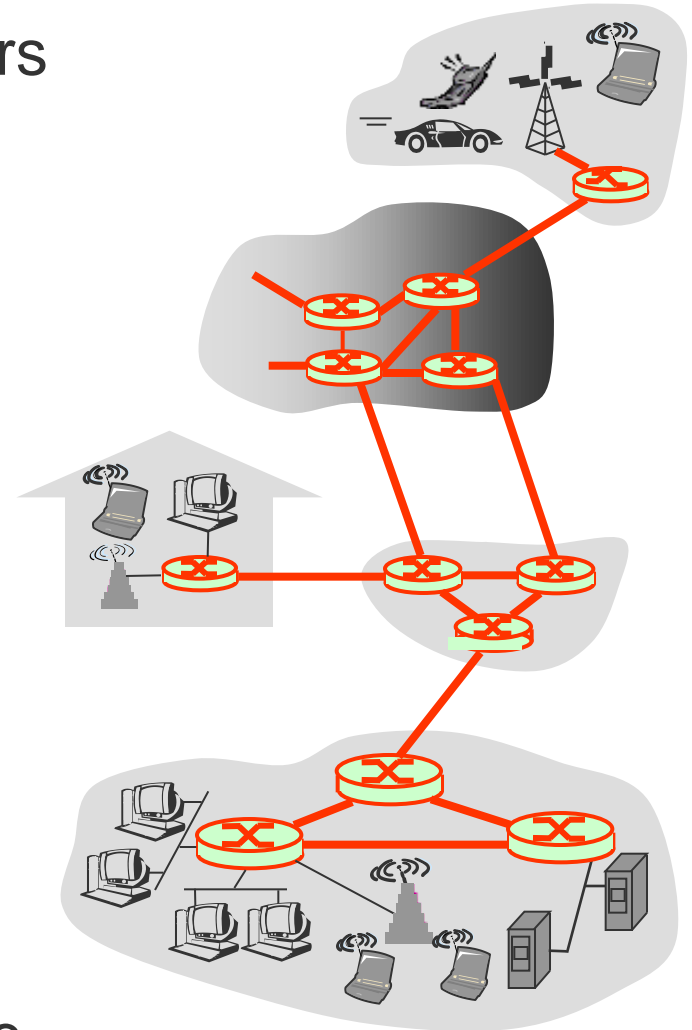
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Elements these slides come from Kurose and Ross, authors of "Computer Networking: A Top-down Approach", and are copyright Kurose and Ross

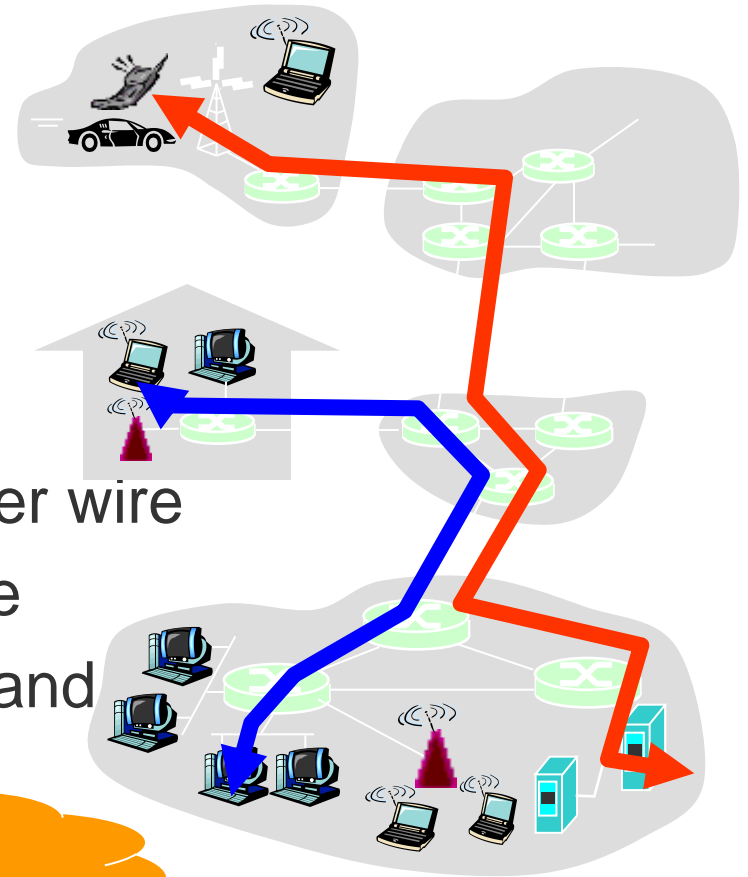
Switching in the Network Core

- Mesh of interconnected routers
- The fundamental question:
 - how is data transferred through net?
- Either:
 - circuit **switching**:
 - fixed path (channel)
 - dedicated resources
 - **packet-switching**, data sent:
 - in discrete “chunks”
 - when resource available



Switching: Traditional Circuit Switching

- End-end resources reserved for transfer
- Resources are dedicated:
 - no sharing
 - all link bandwidth
 - all switch capacity
- Acts as a circuit; e.g.
 - continuous piece of copper wire
- Get guaranteed performance
- Setup required to configure and establish path



Sensible approach?

Switching: Networking Circuit Switching

- Network resources shared by dividing into pieces
 - normally pieces are of a fixed size
 - number of pieces fixed when create network
- Bits **reserved** for (**allocated** to) end-to-end transfers
- Resource piece is *idle* if not used by owning transfer
- For a link, divide link bandwidth into pieces using:
 - frequency division multiplexing (FDM)
 - time division multiplexing (TDM)

Switching: FDM and TDM

WiFi

FDM

frequency
width = bandwidth



User gets continuous
low volume of data

TDM

frequency
width = bandwidth

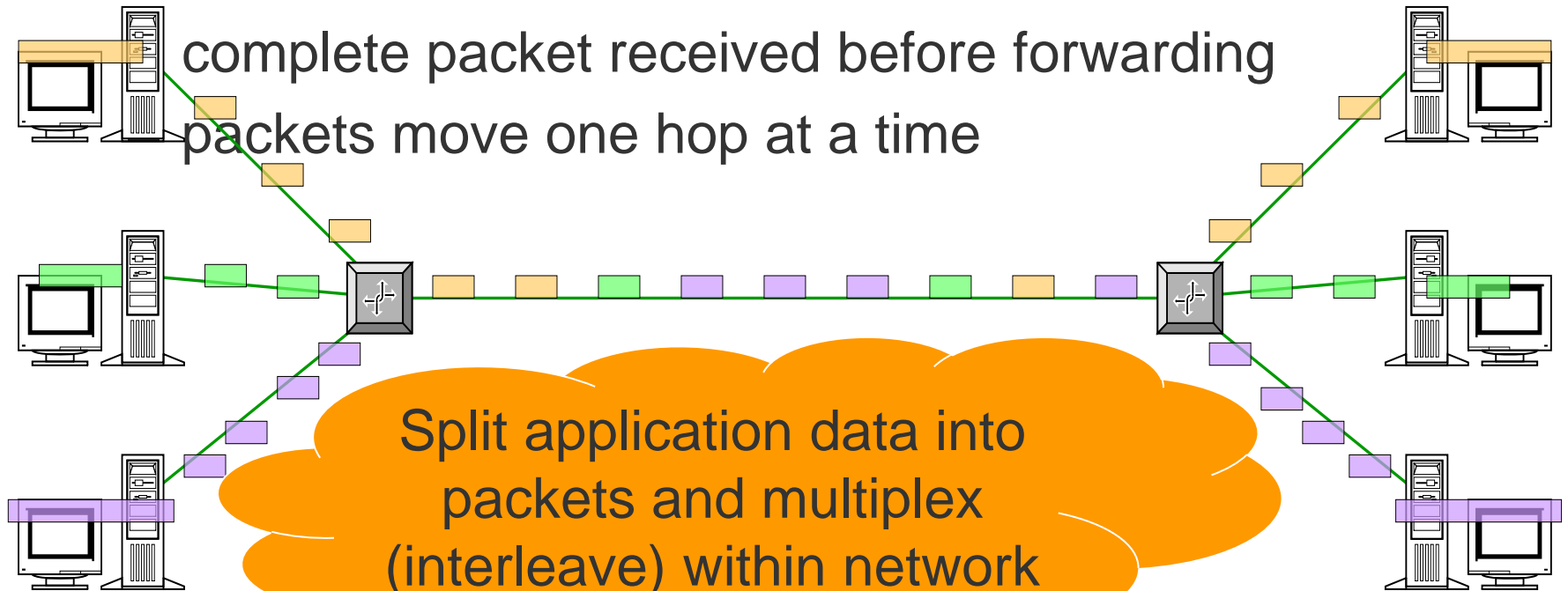


User periodically gets
burst of data

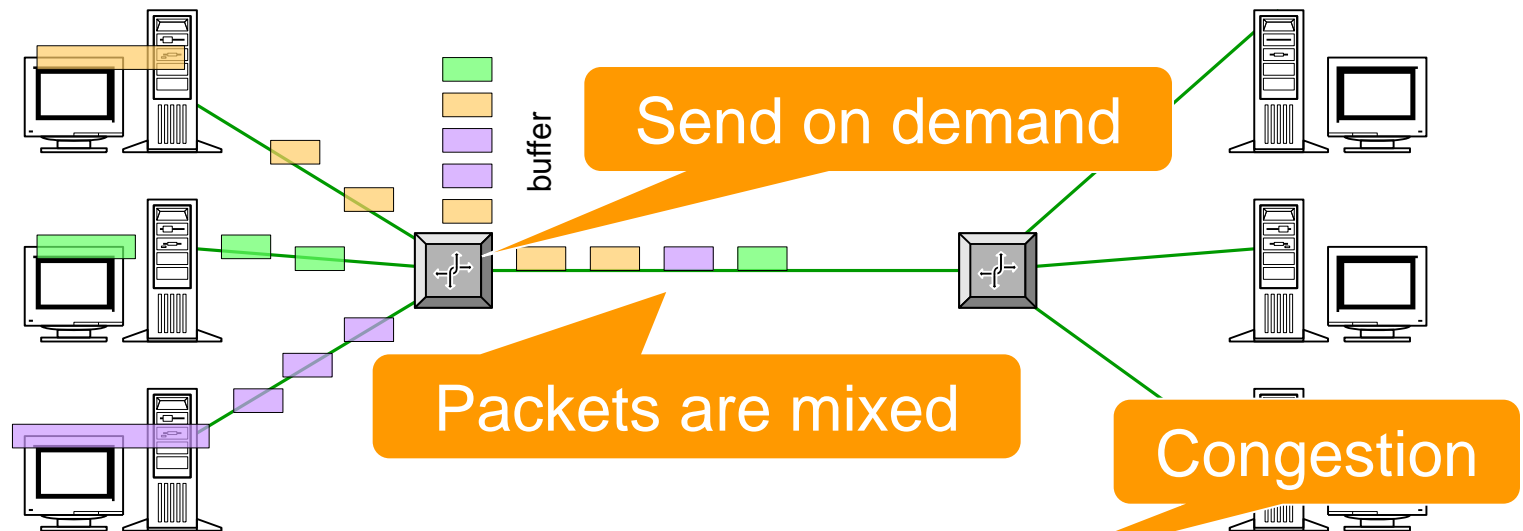


Switching: Packet Switching

- Each end-to-end data stream divided into **packets**
- Packets **share** network resources, use as needed
 - each packet uses full link bandwidth
- Store-and-forward:

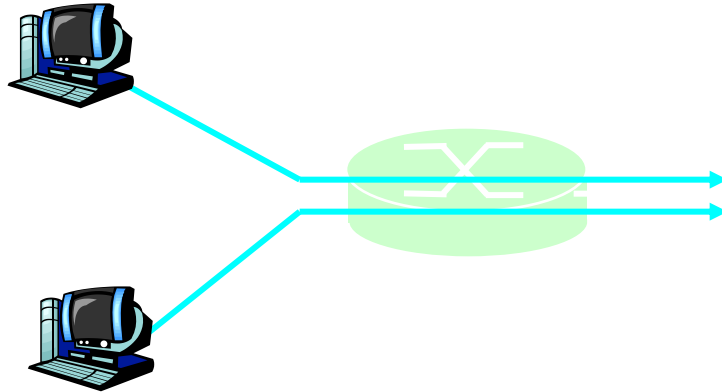


Switching: Statistical Multiplexing



- Packets can arrive faster than can send out (**buffer**)
- Overload normally transitory, will later clear buffer
- If buffer becomes full, **drop** (lose) packets
- Aim for low probability of packet dropping

Switching: Packet versus Circuit



Each user:

100 Kbps when “active”
active 10% of time

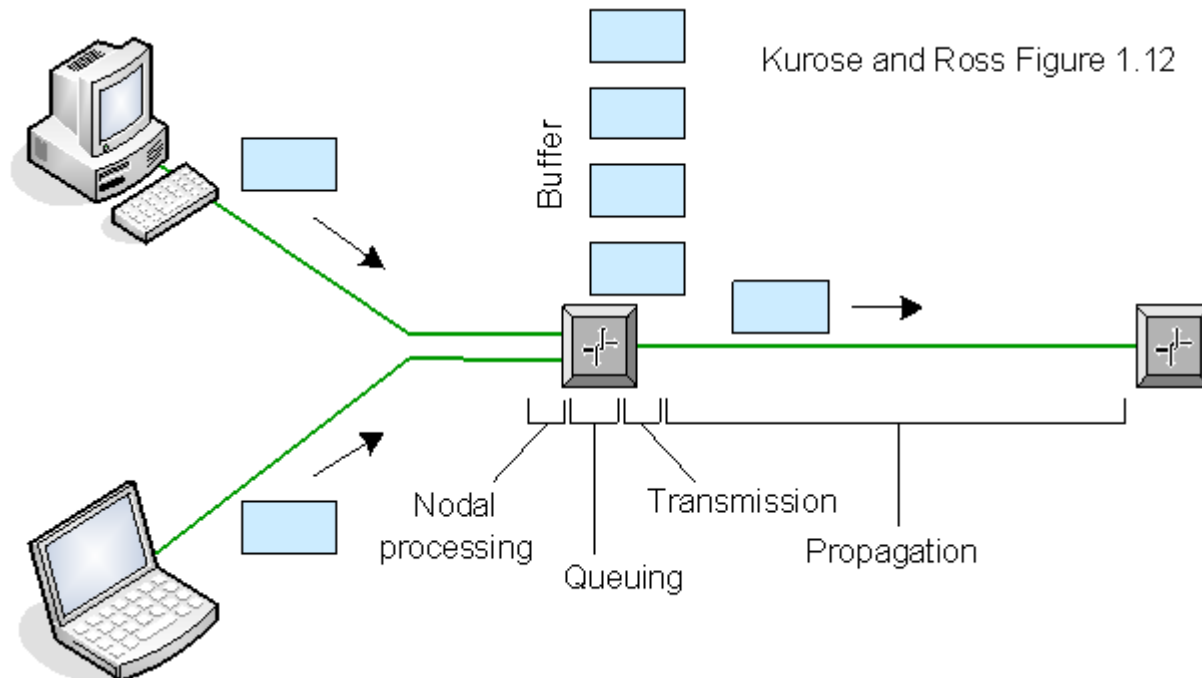
1 Mbps link

Figures are based
on assumptions

- Circuit-switching: $1\text{M}/100\text{K} = 10$ users
- Packet switching: 35 users
 - need statistics theory (why statistical multiplexing)
 - probability > 10 users active at same time $< .0004$
- Packet switching allows more users to use network

Delay, Loss and Throughput

- Data does not arrive instantaneously at destination
- Not all data sent arrives!
- Limit to amount of data can send per second



Delays: Transmission and Propagation

- Link has transmission rate: R (in bits per sec, bps)
 - depends on physical medium
 - often referred to as bandwidth of link
- For packet of L bits:
 - transmission delay (d_{trans}) = L/R
- Packets travel (propagate) with speed s (in m/s)
 - depends on physical medium
 - $X * c$ (c = speed of light, $\approx 3 \times 10^8$ m/s), $X < 1$
- For distance d (in meters):
 - propagation delay (d_{prop}) = d/s

Delays: Transmission vs. Propagation

≈ bit



≈ packet

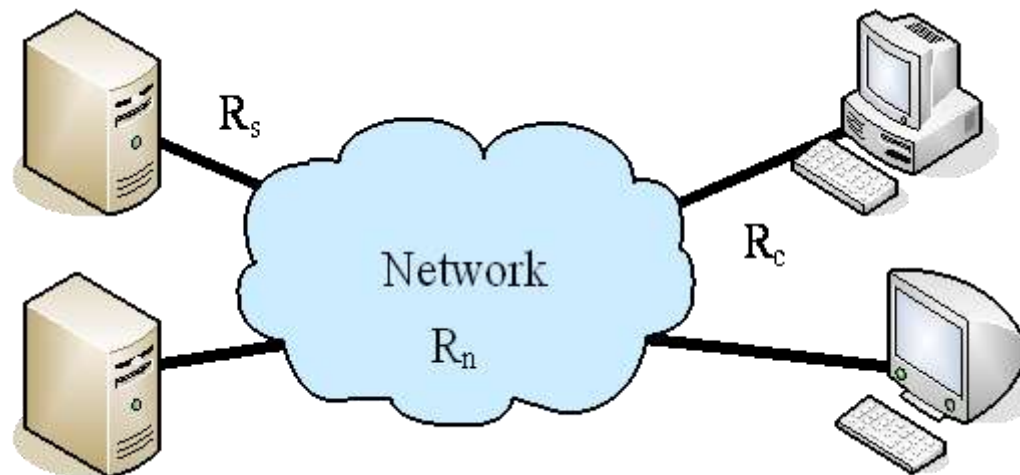
- Cars “propagate” at 100 km/hr
- Booth takes 12 sec to service car (transmission time)
- Q: How long for convoy to get to 2nd booth?
- “push” all cars through booth = $12 \times 10 = 120$ sec
- For last car to propagate from 1st to 2nd booth:
 - $100\text{km}/(100\text{km/hr}) = 1$ hr
- A: $120 \text{ sec} + 1\text{hr} = 62$ minutes

Delays: End-to-End Delay

- Delay between parts of application comes from:
 - individual node delays
 - processing delays in applications
 - waiting for access to shared medium
- Also use round-trip-time (RTT)
 - time source to destination and back
- If no queuing delays, dominant delay predictable
 - sending 1 byte on any link: propagation delay
 - sending 25Mb on 10Mbps link: transfer rate

Throughput in Computer Networks

- Rate at which data is transferred, GUI may report
- Measuring useful (application data) throughput
- For large transfers, determined by three parameters:
 - remote computer (link), network and local link

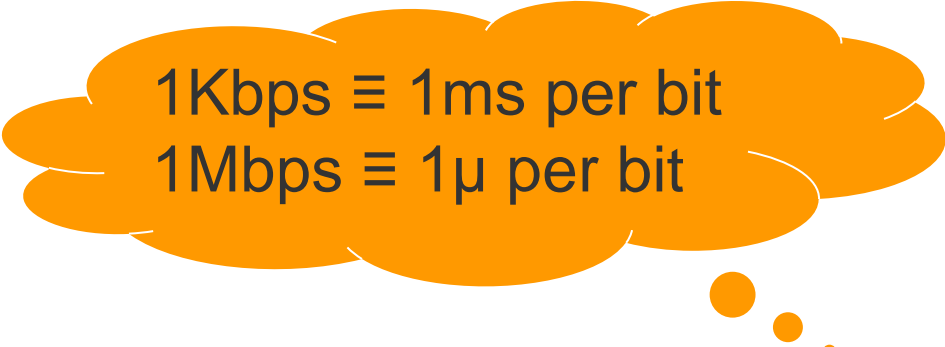


Units (Background)

- Link rates in bits per second (bps), with multipliers:

- Times often small quantities, with multipliers:

milli	m	10^{-3}
micro	μ	10^{-6}
nano	n	10^{-9}
pico	p	10^{-12}



1Kbps \equiv 1ms per bit
1Mbps \equiv 1 μ per bit

Summary

- Aim is cost effective connectivity for set of computers

Real result is a packet switched network

aim? use shared links and statistical multiplexing

- Communication services seen as process-to-process
- Delays and throughput define network performance

Course Structure

- **Lectures:** 3 per fortnight
 - Tuesdays week A, 9:00-10:00am
 - Thursdays, 3:00-4:00pm
- **Laboratories:** 1 per fortnight, needs preparation
 - week Bs
 - Moodle quiz, experiments in networking
- **Workshops:** 1 per fortnight, no preparation
 - Mondays (1.00pm) or Tuesdays (9.00am) week B
 - Moodle quiz, new material, examinable
- **Self assessment tests:** in your own time
 - Moodle quiz, recap lecture material

Laboratory: Overview

- Learn how to configure and use real networks
- Have complete control over network that configure
- But, cannot let you loose on the Internet!
 - physical private unconnected networks needs
 - 3+ machines per student
 - use virtual network of virtual machines:
 - whole virtual network runs on your lab machine
 - no link to the outside world



Laboratory: Core Tool (VNUML)

- Use virtual operating system to emulate machines
- Several virtual machines, so needs to be light weight
 - not VMware, Virtual Box, etc.
- Virtual Networking User Mode Linux (VNUML)
 - each VM runs as a user process
 - each VM has to boot: quite slow
 - you get a partially disabled console
 - username root, password xxxx
 - entirely command line driven
 - installed on School's Linux machines
 - not easy to install elsewhere

Based on
User Mode
Linux (UML)

Laboratory: Example Screen

Console
virtual
machine 1

Console
virtual
machine 2

Console
virtual
machine 3

```
[screen 0: ron]
View Terminal Tabs Help
mu-ns.mu +trace
Options: printcmd
999999 IN NS mu-ns.mu.
38 bytes from 172.16.3.1#53(172.16.3.1) in 39 ms
300 IN A 172.16.1.1
300 IN NS mu-ns.mu.
56 bytes from 172.16.1.1#53(mu-ns.mu) in 27 ms
ron:~#

[screen 0: mu-ns]
View Terminal Tabs Help
;; ADDITIONAL SECTION:
2081-ns.comp20081.scs.mu. 300 IN A 172.16.3.1
ec
1#53(172.16.3.1)
15:33:24 2008
2

[screen 0: server]
View Terminal Tabs Help
the exact distribution terms for each program are described in
the
individual files in /usr/share/doc/*/copyright.
Debian GNU/Linux comes with ABSOLUTELY NO WARRANTY, to the exte
nt
permitted by applicable law.
server:~#
```

Laboratory: Deliverables

- Deadline end of your laboratory session
 - extended deadline 11pm on day of laboratory
- Unlike programming laboratories:
 - there is no program to submit and demo
- Instead:
 - configuration files, commands used
 - output from commands, your analysis
- We'll supply a Moodle quiz:
 - fill in the results as you go along
 - submit at end of laboratory session
 - record time when submit performed

Laboratory: Testing To See If It Works

- Virtual machines minimally configured, available:
 - **ifconfig** shows configuration of network devices
 - **ping** sends “Are you there?” message, response:
 - Unknown Host – no address for given name
 - Network Unreachable – no path to machine
 - statistics – working
 - no output – not working
 - **arp** will show name/address of local machines
 - **netstat** shows statistics about the network
 - **traceroute** path through network (exercise 3...)
 - **dig** for information about names (exercise 4)

Reading

- . Kurose & K. Ross, Computer Networking: A Top-Down Approach, 6th Ed, Pearson, ISBN: 0-273-76896-4
- L. Peterson & B. S. Davie, Computer Networks, 5th Ed, Morgan Kaufmann, ISBN:0123850592.
- There are many good books on networking, if you do not like our selection then pick your own.