

COMP28411 Computer Networks

Nick Filer - Link Layer and Physical Layer

Tuesday	25/11	LL - 1 of 4	- Introduction
Thursday	27/11	LL - 2	- Switches, Routing and Medium Access
Thursday	4/12	LL - 3	- Framing, Ethernet, ATM Point-2-Point Protocol
Tuesday	9/12	PL - 4 of 4	- Physical Layer and Wireless
	2 and 4/12		- Error Detection Workshop

Some material from:

Kurose & Rose – Chapter 5 + Slides

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Lecture Summary

- Link Layer
 - Services
 - Where implemented
 - Packet Encapsulation – reminder
 - Flow Control
 - Link layer addressing
 - Mapping IP to/from MAC addresses.
 - Hubs and Switches

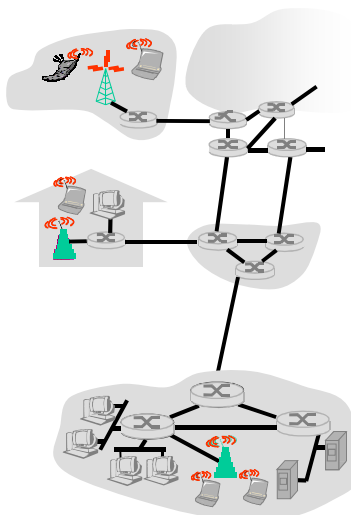
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Link Layer Overview

- **Terminology**
 - Hosts and Routers/Switches/Hubs are **NODES**.
 - Communications channel between nodes is a **LINK**.
 - Layer 2 packets are a **FRAME**
- **Responsibility**
 - Node to adjacent node transfer of layer 3 **datagram** over **link**.
 - Uses Layer 1 (Physical) methods to move **frames between directly interconnected hosts**.



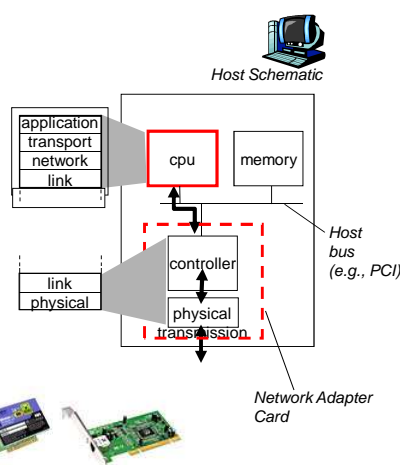
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Where is the link layer implemented?

- In every node.
- Link layer implemented in “adaptor” (*network interface card* NIC)
 - Ethernet card, PCMCIA/Express cards, 802.11 card
 - Implements link, physical layer
- Attaches into node’s system buses
- Combination of hardware, software, firmware.



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Adaptors Communicating

The diagram illustrates the communication between two nodes. On the left, a 'sending node' contains a 'controller' and a 'datagram'. The datagram is sent to the controller, which then encapsulates it into a 'frame'. This frame is transmitted over a link to the 'receiving node' on the right. The receiving node's controller receives the frame, extracts the 'datagram', and passes it to the upper layer.

- **Sending side:**
 - Encapsulates datagram in frame.
 - Adds error checking bits, flow control, etc.
- **Receiving side**
 - Looks for errors, flow control, etc.
 - Extracts datagram, passes to upper layer at receiving side.

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Virtual Links

The diagram shows a laptop with icons for Bluetooth, WiFi, and Ethernet. A person is shown sitting at a desk with a computer, representing a user or network administrator.

- **Alternatives:**
 - My computer talks Ethernet (Layer 2 + 1) .
 - My laptop talks either Ethernet, WiFi (Layer 2 + 1) or Bluetooth.
- **My Broadband (Last Mile Connection over telephone or cable) has:**
 - **IP input over Ethernet or WiFi.** Both IEEE 802.xx standards.
 - **ADSL2 physical layer with Asynchronous Transfer Mode (ATM)** at the data-link layer.
 - **VDSL (very high bit rate DSL) over fibre.** Often **PPPoE**.
 - On ADSL the ATM (layer 2) carries Point to Point Protocol over ATM (**PPPoA**) – also layer 2!.
 - Outside Europe many use PPPoE (Ethernet)
 - The PPP (layer 2) carries IP packets .
- We regard the PPPoA/E as providing a **virtual link**.
 - To higher layers it appears as a complete network.

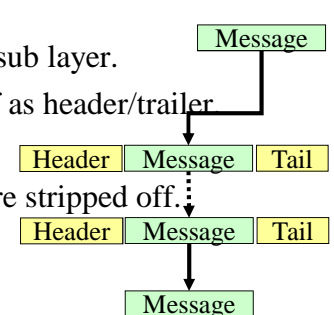
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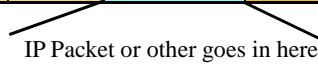
Encapsulation Review ?

- Protocol data + control is sent to host's sub layer for processing.
- Treated as a black box of bytes at next sub layer.
- Each layer adds control + data for itself as header/trailer
- Process is repeated at each layer.
- At destination, each layer's additions are stripped off.
 - So acts as multiplexor/demultiplexor.




Ethernet Frame Structure

Preamble	Start of frame delimiter	MAC Destination	MAC Source	Ethertype or Length	Data + Padding	CRC32	Interframe Gap
7 lots of 10101010	10101011	6 Octets	6 Octets	2 Octets	46-1500 Octets	4 Octets	12 Octets



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Flow Control

A2

- Is an optional service! Not always implemented at this layer.
- Purpose ?:
- Handshake:
 - Exchange stop/start messages.
 - Hardware: Request to Send/Clear To Send (RTS/CTS), Data Set Terminal Ready (DSR/DTR)
 - Software: X-ON/X-OFF
- Open-Flow:
 - By prior reservation – using Connection Admission Control (CAC)
 - To work must over resource – redundancy.
- Closed-Loop:
 - Some way of reporting resource availability and resource needs.
 - Ethernet PAUSE frame sent to special multicast address 01-80-C2-00-00-01 with 16 bit time request in 512 bit time quanta's.
 - Asynchronous Transfer Mode (ATM) has Available Bit Rate (ABR) guarantees minimum bit rate and reports congestion.

e.g. for real-time constant bit rate traffic.

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LAN/MAC Addresses



- 32-bit IP address:
 - *Network-layer* address.
 - Used to get datagram to destination IP subnet .
 - At least partially geographical and part ID.
- MAC (or LAN or physical or Ethernet) address:
 - Function: *Get frame from one interface to another physically-connected interface (on the same network).*
 - 48 bit (6 byte) MAC address (for most LANs)
 - Moving towards 64 bit in IPv6 network cards. Compatibility?
 - Burned in Network Interface Card (NIC) ROM, also sometimes software settable.
 - 3 Bytes Organization Identifier + 3 bytes NIC Identifier.
 - 2^{48} or 281,474,976,710,656 unique addresses, Expected to last until year 2100. Used in most IEEE 802 networks (Ethernet, WiFi), Bluetooth,
 - Does it matter if MAC addresses are re-used? _____

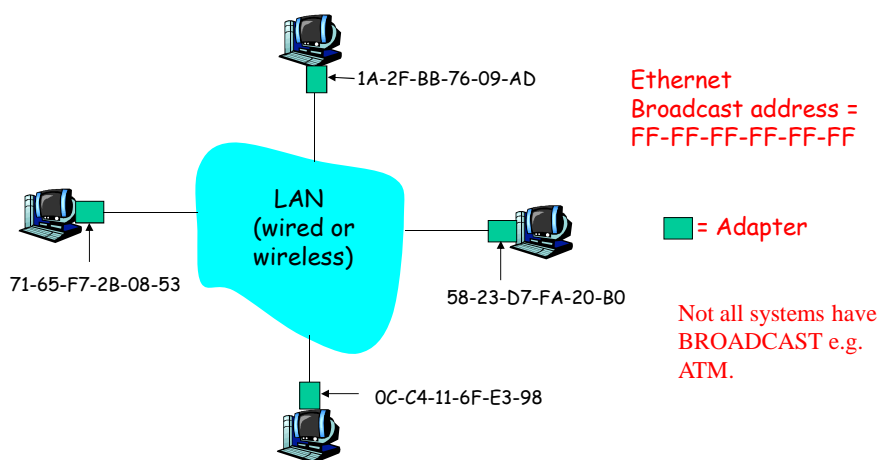
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LAN Addresses

Each adapter on LAN has unique LAN address.



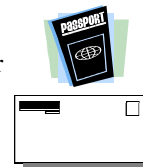
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LAN Address (more)

- MAC address allocation administered by IEEE
- Manufacturer buys portion of MAC address space (to assure uniqueness).
- Analogy:
 - (a) MAC address: like Passport Number or social security or National ID card.
 - (b) IP address: like postal address
- MAC flat address gives portability
 - Can move LAN card from one LAN to another
 - Switches auto-adapt to new card
- IP hierarchical address NOT portable
 - Address depends on IP subnet where node is attached (geography/location).
 - IP routing depends on attachment to network



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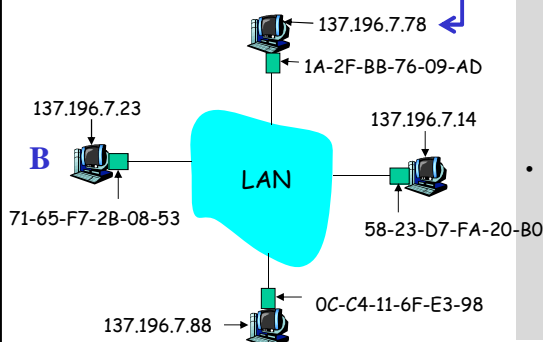
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ARP: Address Resolution Protocol

A1

Question: How to determine MAC address of B knowing B's IP address?



You should know this already!

- Each IP node (host, router) on LAN has an **ARP** table.
- ARP table: IP/MAC address mappings for some LAN nodes


```
< IP address;
  MAC address;
  TTL>
```
- TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min).

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Switch & Hubs



- Older networks were **Half-Duplex**, wires were shared by many machines. They used Hubs.
- Modern networks are **Full-Duplex** over **unshared point-to-point wires**. They use Switches.
 - Switches are **Link-layer** device: Smarter than hubs, take **active** role.
 - Store + forward frames.
 - Examine incoming frame's MAC address, **selectively** forward frame to one-or-more outgoing links when frame is to be forwarded on segment, uses collision avoidance techniques to access segment.
 - **Transparent**
 - Hosts are unaware of presence of switches.
 - **Plug-and-play, self-learning**
 - Switches do not need to be configured.
 - Remember that routers do need configuration!

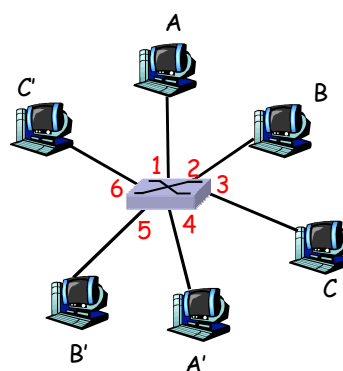
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Switch: Allows Multiple Simultaneous Transmissions

- Hosts have dedicated, direct connection to switch.
- Switches buffer (queue) packets.
- Ethernet protocol used on *each* incoming link, but no collisions; full duplex.
 - Each link is its own collision domain.
- **Switching**: A-to-A' and B-to-B' simultaneously, without collisions.
 - Not possible with dumb hub.



Switch with six interfaces
(1,2,3,4,5,6)

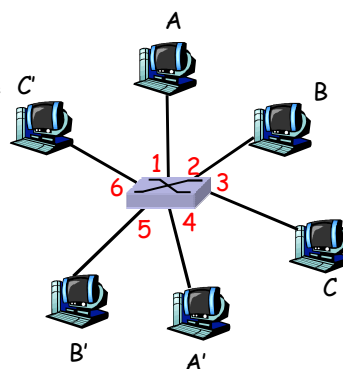
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Switch Table

- **Q:** How does switch know that A' is reachable via interface 4 and B' reachable via interface 5?



Switch with six interfaces
(1,2,3,4,5,6)

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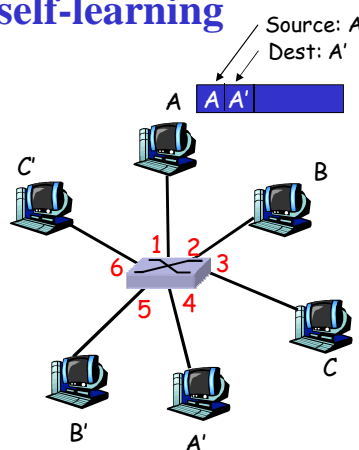
Switch: self-learning

- Switch **learns** which hosts can be reached through which interfaces
 - When a frame is received, the switch “learns” the location of the sender: Incoming LAN segment.
 - Records sender/location pair in switch table.
 - Example: Sending A → A' via switch.

Which network protocol probably initiates self-learning?

MAC addr	interface	TTL
A	1	60

Switch table
(initially empty)



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Switch: Frame Filtering/Forwarding

When frame is received:

1. Record link associated with sending host.
2. Index switch table using MAC dest address.
3. **IF** entry found for destination **THEN** {
 - IF** Destination is on segment from which frame arrived
 - THEN** Drop the frame
 - ELSE** Forward the frame on interface indicated
- ELSE** Flood

*forward on all but the interface
on which the frame arrived*

This is a simplification? **What
actually happens?** Try
Wikipedia Switch page.

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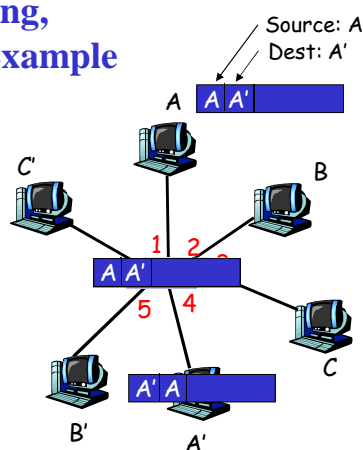
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Self-learning, Forwarding: Example

A6

- Example:
 - Sending A → A' via switch.
 - Then a frame A' → A.
- Frame destination unknown:
flood
- Destination A location
known:
Selective send



MAC addr	interface	TTL
A	1	60
A'	4	60

*Switch table
(initially empty)*

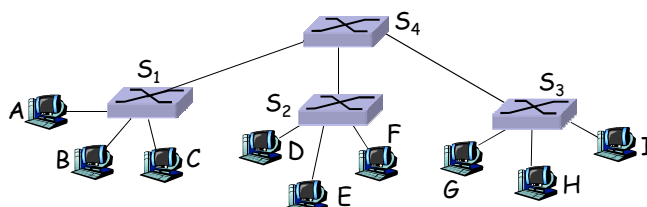
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Interconnecting Switches

- Switches can be connected together



- Q:** Sending from C to I →
How does S_1 know to forward frame destined to I via S_4 and S_3 ?
- A:** Self learning! (works exactly the same as in single-switch case!)

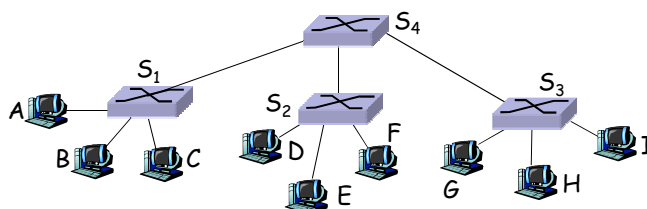
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Self-learning Multi-Switch Example

Suppose C sends frame to I, I responds to C – **Q:** How learn how to do it?




- Q:** Show switch tables and packet forwarding in S_1, S_2, S_3, S_4

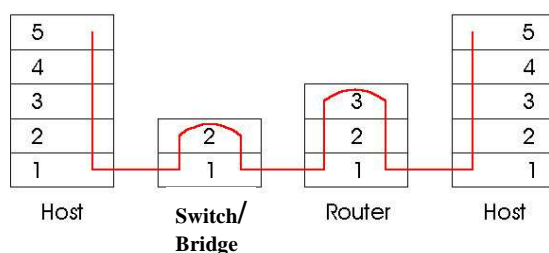
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Switches vs. Routers

- Both are store-and-forward devices
 - Routers: network layer devices (examine network layer headers)
 - Switches are link layer devices – sometimes called **bridges**.
- Routers maintain routing tables, implement routing algorithms
- Switches maintain switch tables, implement filtering, learning algorithms
- Bridges  connect separate switched networks or sub-nets together at layers 1 and 2.



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Lecture Summary

- Link Layer is concerned with node to node communications.
- It is normally implemented in hardware but may be soft coded
 - Services
 - Where implemented
 - Packet Encapsulation – reminder
 - Flow Control
 - Link layer addressing
 - Address Resolution
 - Hubs and Switches
- Next:
 - More on addresses and Sharing a network Link Layer Routing (between switches)
 - Multiple access – sharing a network.
 - Collisions and channel partitioning.
 - Random Access protocols

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Questions ?

- What information passes between:
 - The network layer and the link layer?
 - The link layer and the physical layer?
- If the transport layer can do error checking, why is error detection and correction important at the link layer?
- The network layer has the full address for the datagram, why doesn't the link layer use network addresses?
- What extra hardware is needed to implement a full-duplex link compared to a half-duplex link?
- If data is buffered at other layers, why does the link layer also do buffering?
 - Why does the network layer buffer?
 - Do you think the physical layer buffers? Why or why not?
- Why do you think IPv4 addresses were 32 bits whereas Ethernet MAC addresses were 48 bits?
 - Why are IPv6 IP addresses 128 bits whereas the IPv6 MAC address equivalent is 64 bits?
- What is "Inverse ARP" (InARP)? Why is it used?