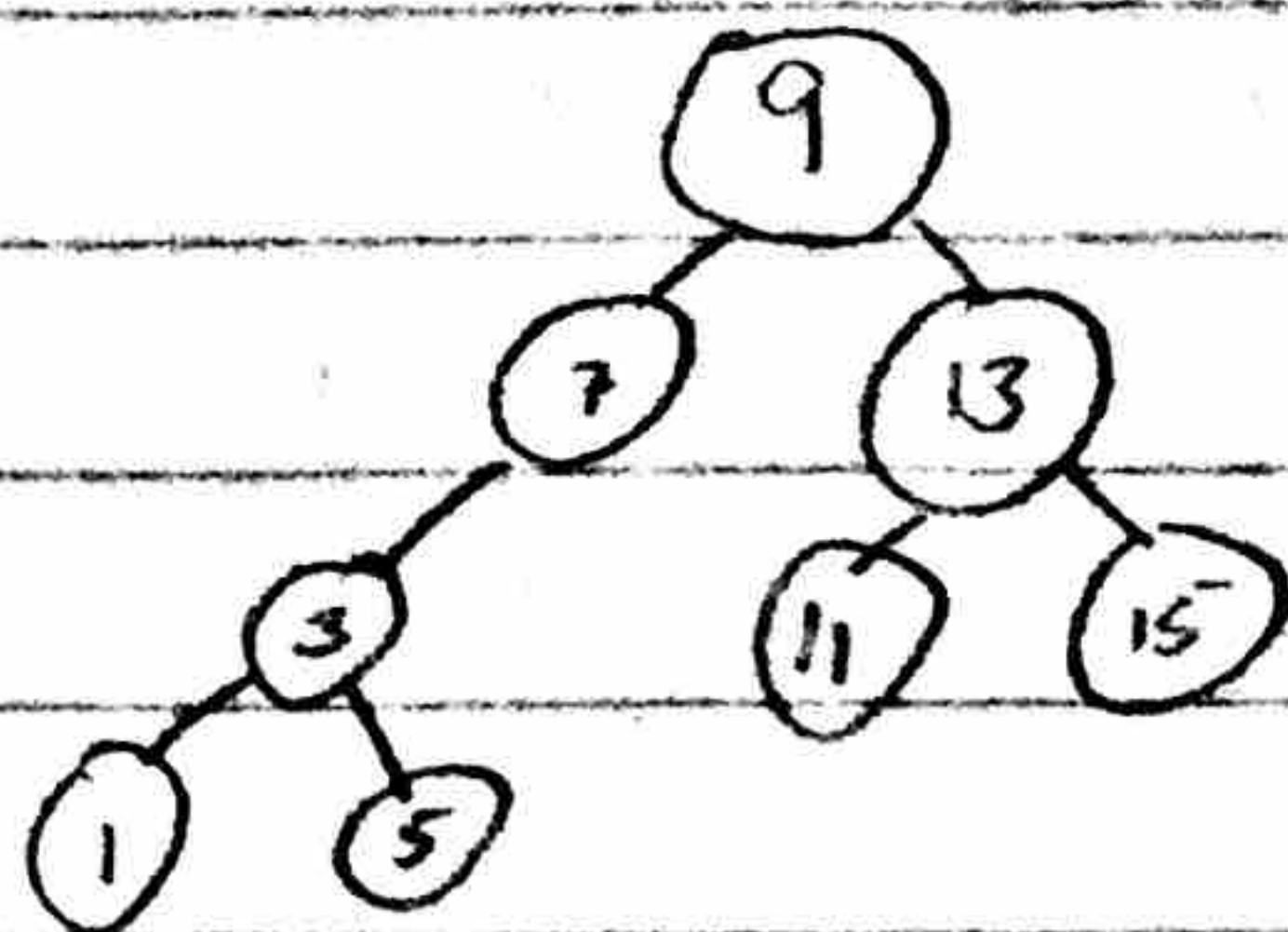
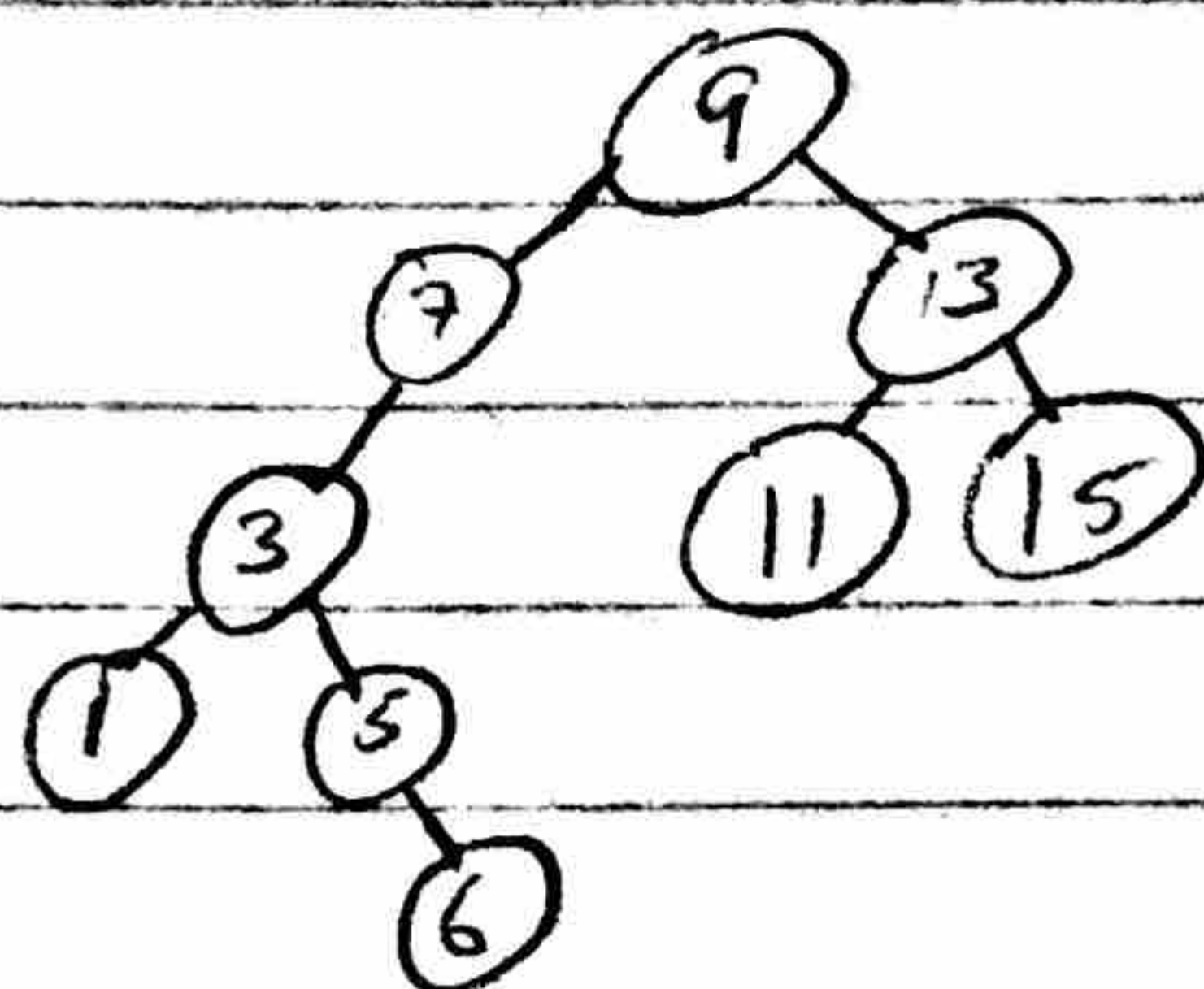


tg1632_hw8.pdf

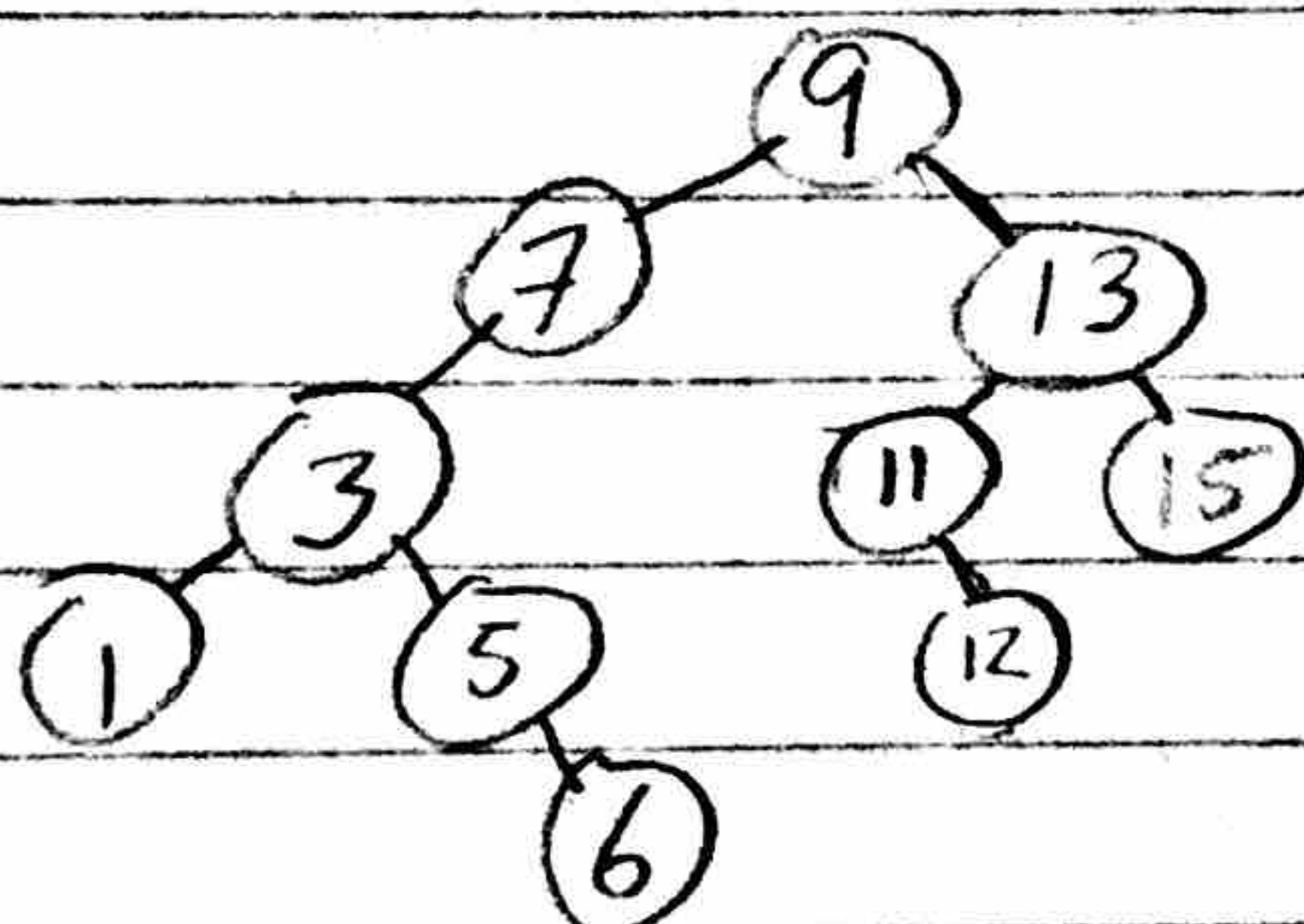
Question 1



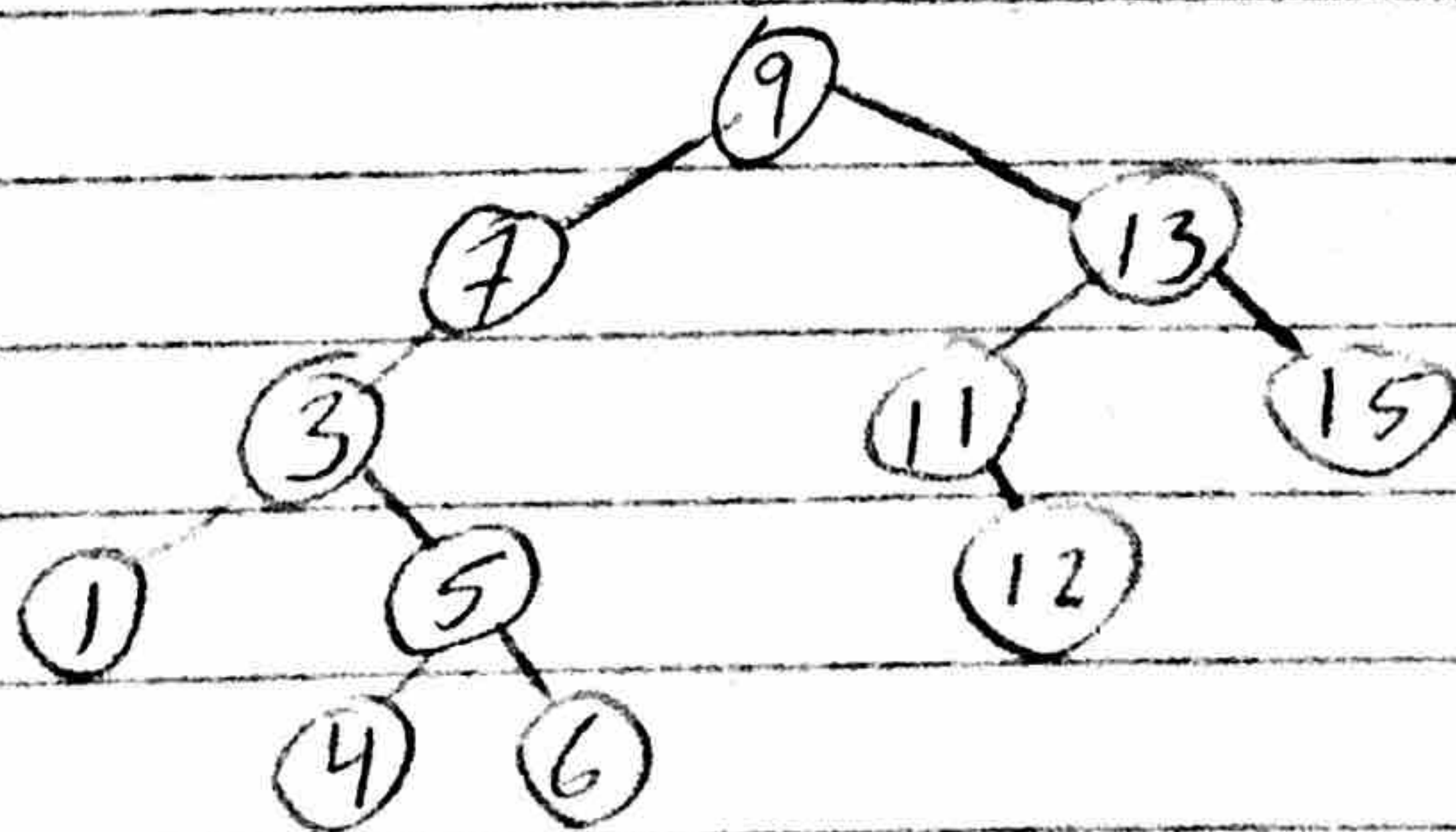
bst[6] = None



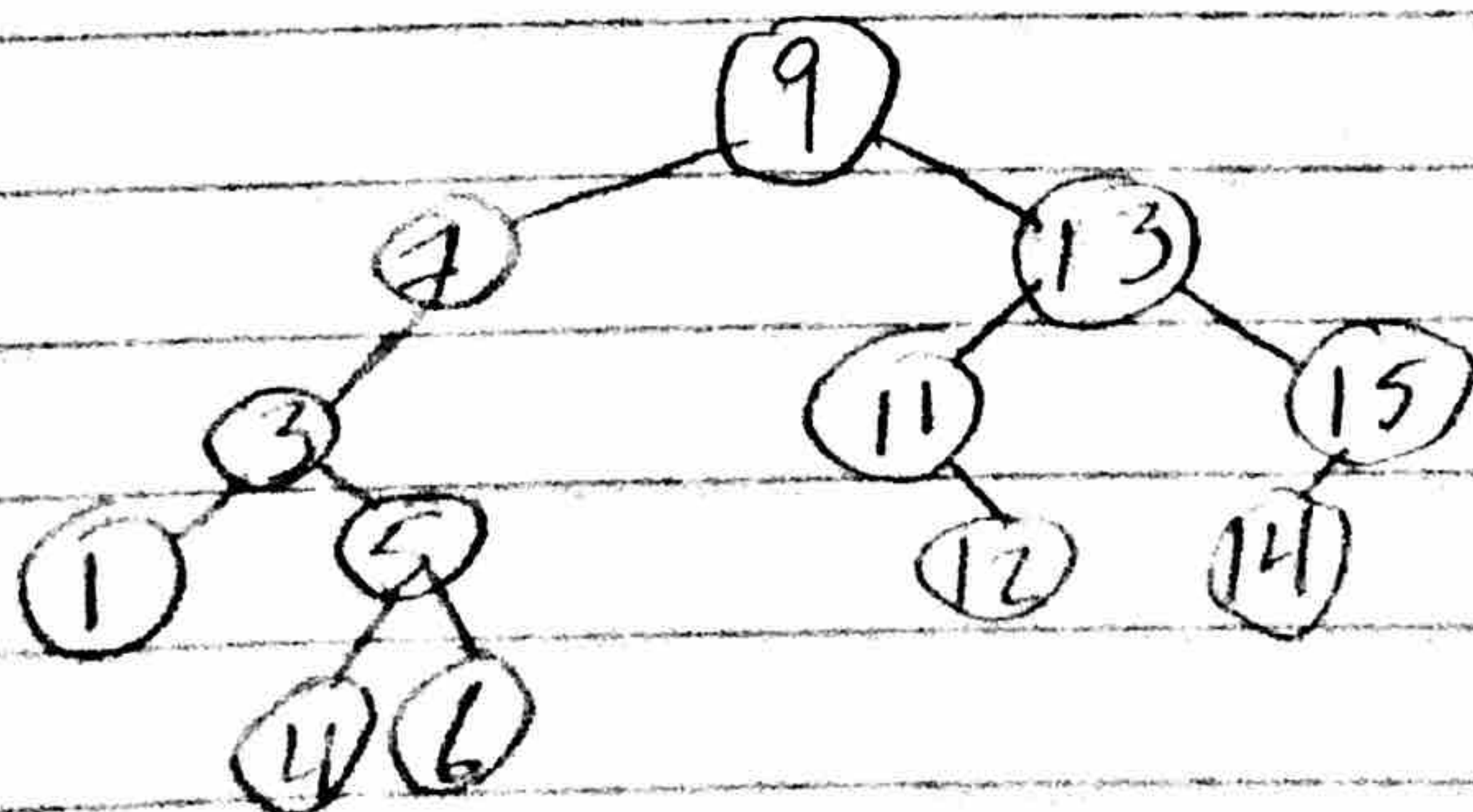
bst[12] = None



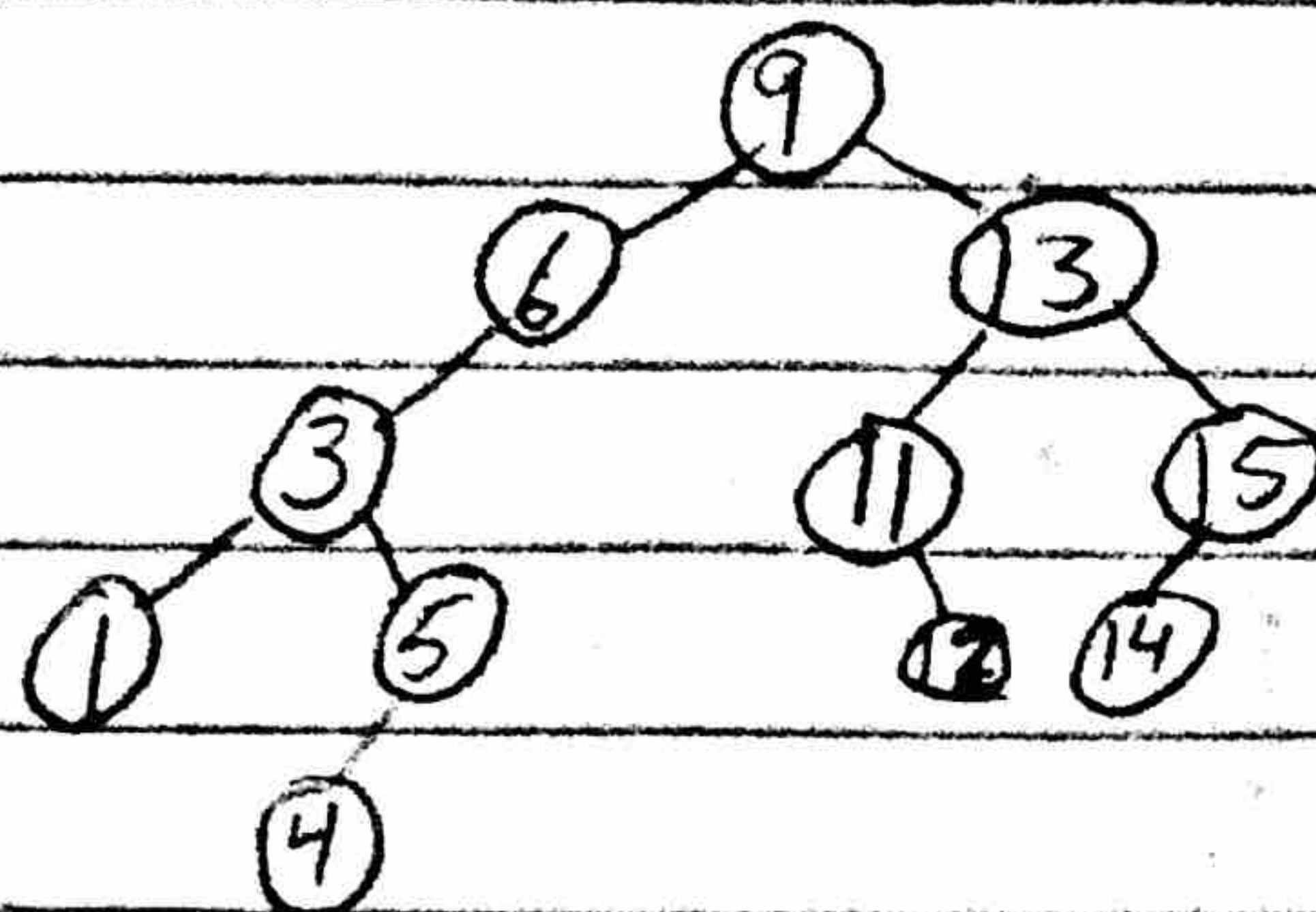
bst[4] = None



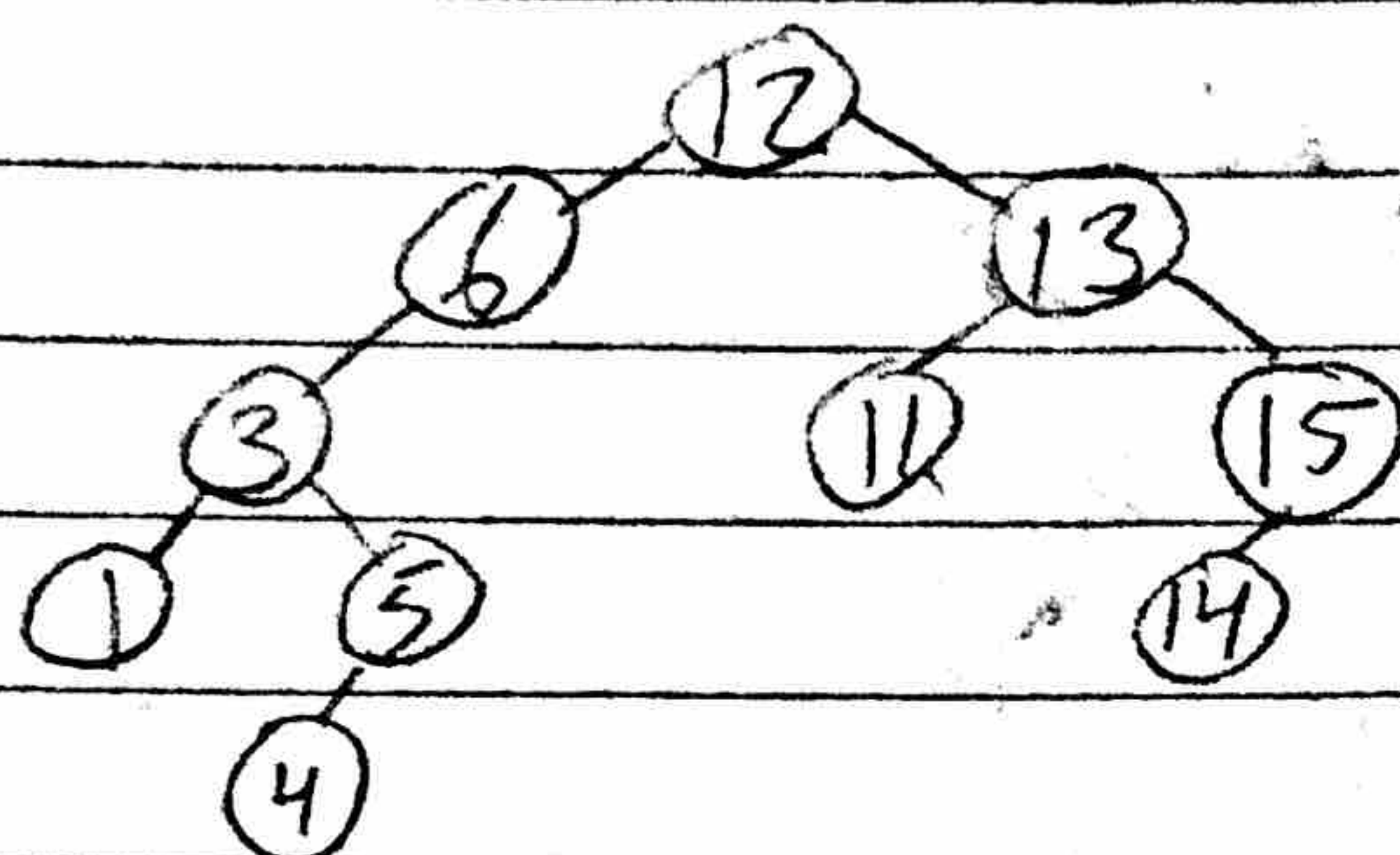
bst[14] = None



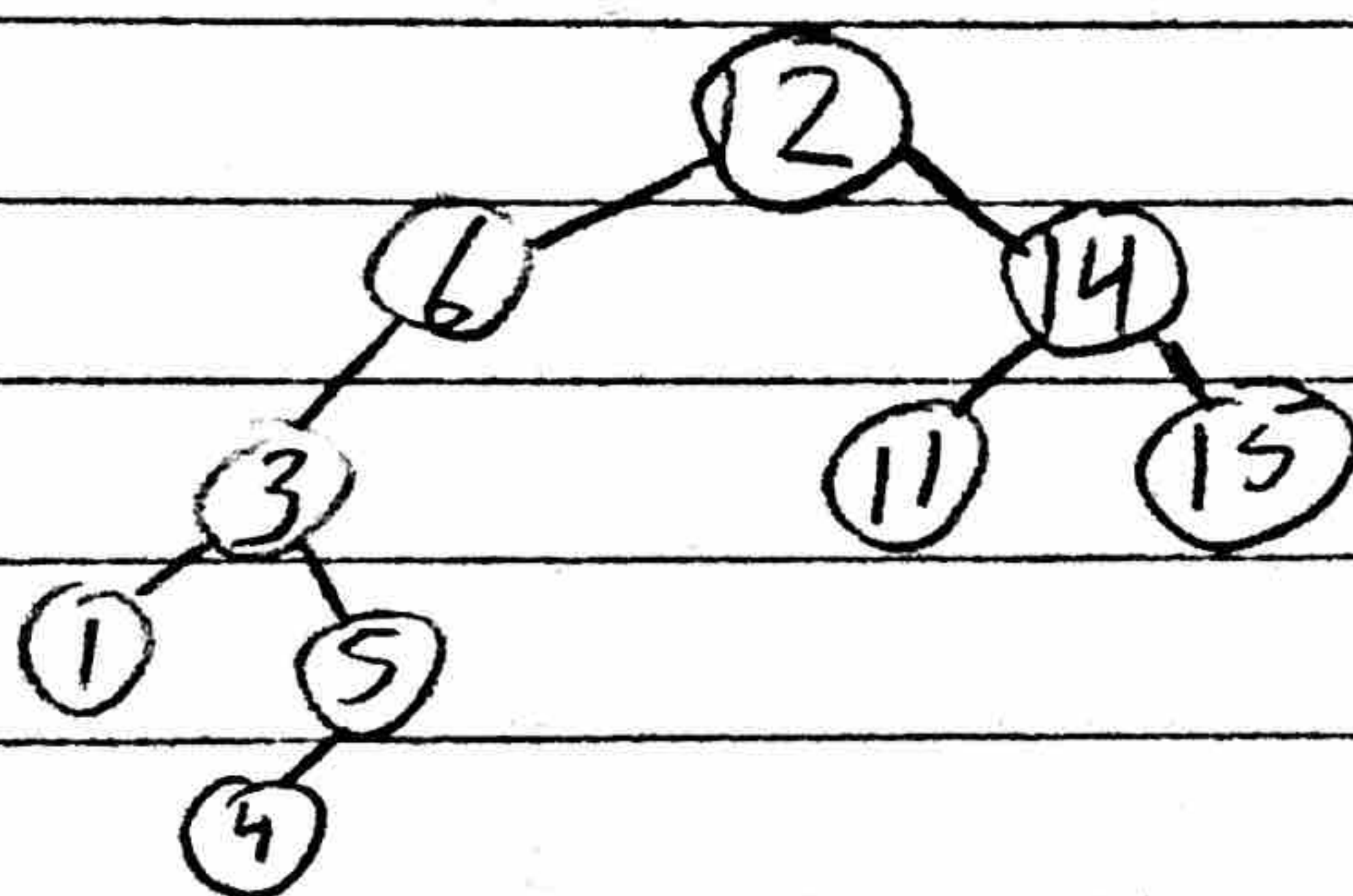
del bst[7]



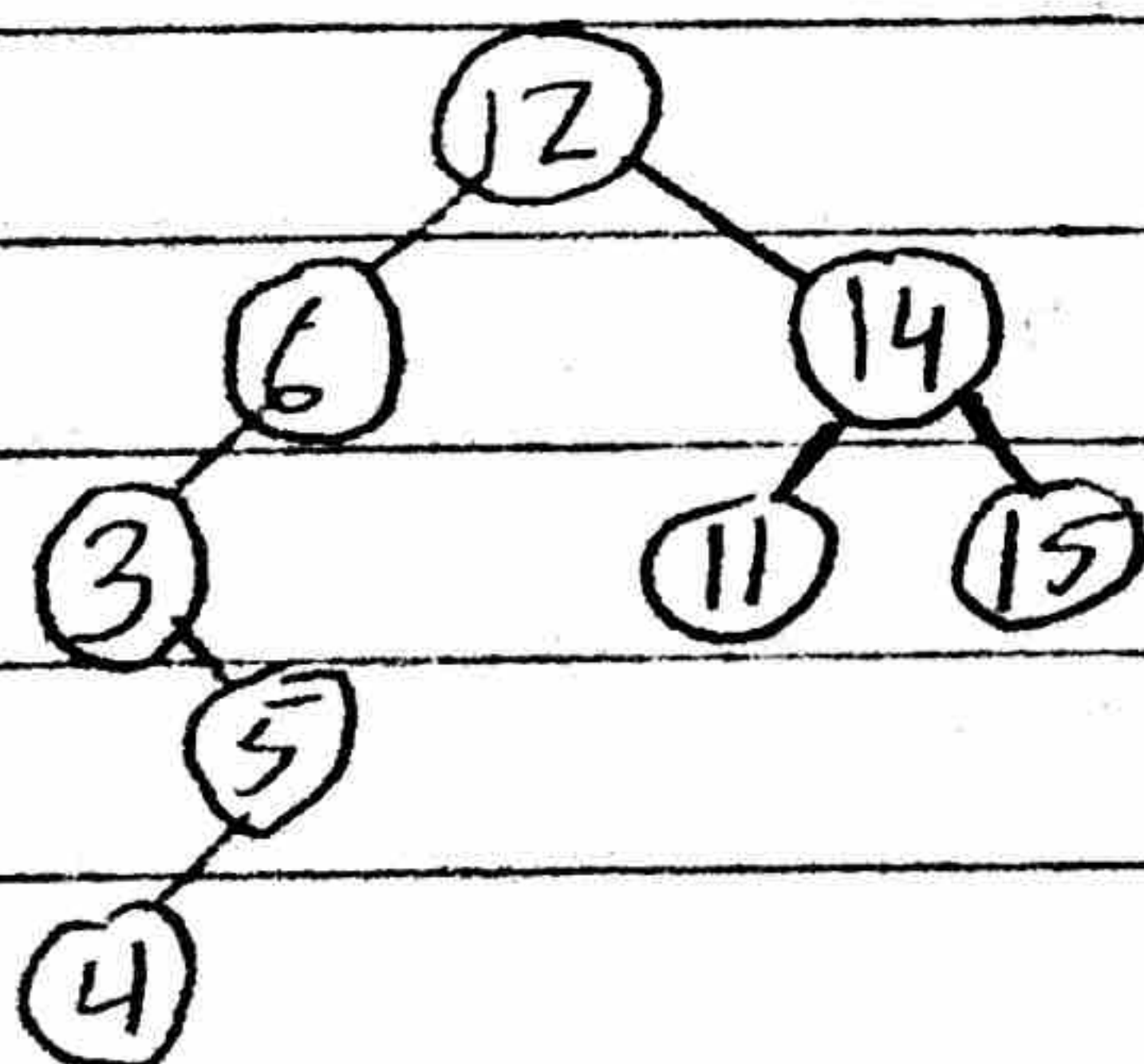
del bst[9]



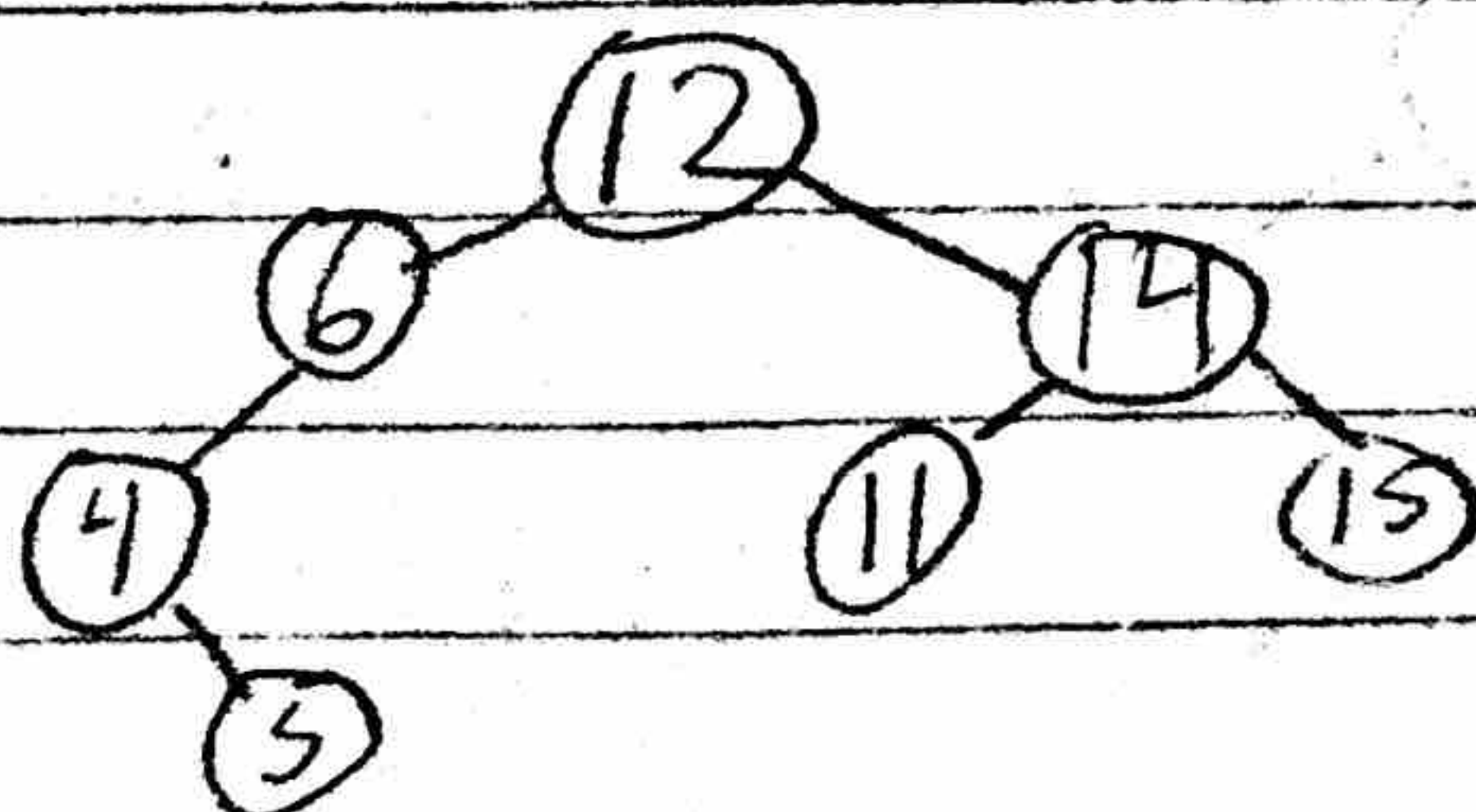
del bst[13]



del bst[1]



del bst[3]



Question 2(c)

- a) Since we are adding to the right every time we have to first go through 1 node, then 2, then 3, all the way up to n . This is:

$$1 + 2 + 3 + 4 + 5 + \dots + n$$

Simplifying you get:

$$\frac{\text{First} + \text{Last}}{2} (\text{no. of terms}) = \frac{1+n}{2} (n) \approx n^2$$

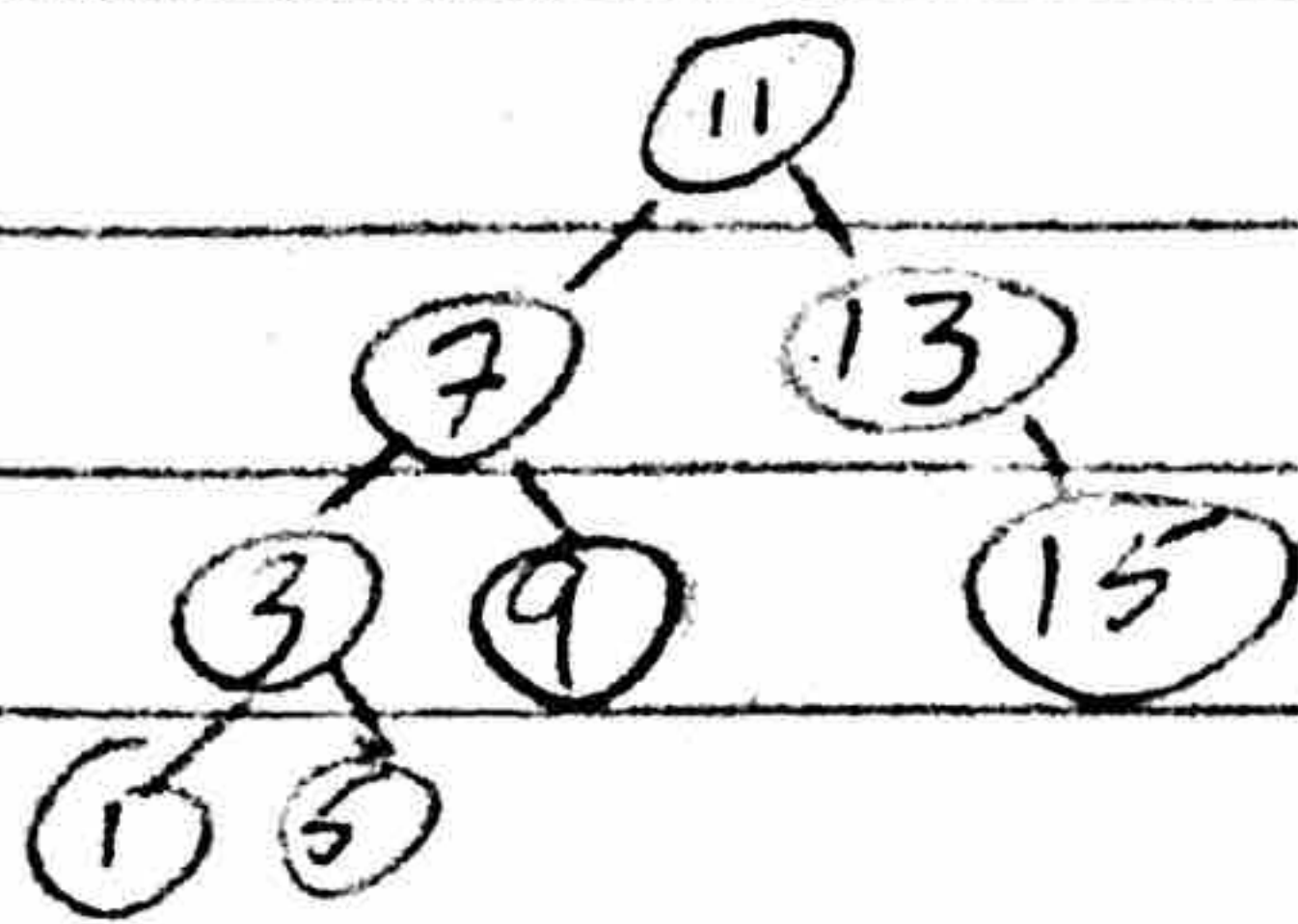
Hence $\boxed{O(n^2)}$

- b) Inserting in a complete binary search tree is $\log n$. You do this n times, because there are n items.

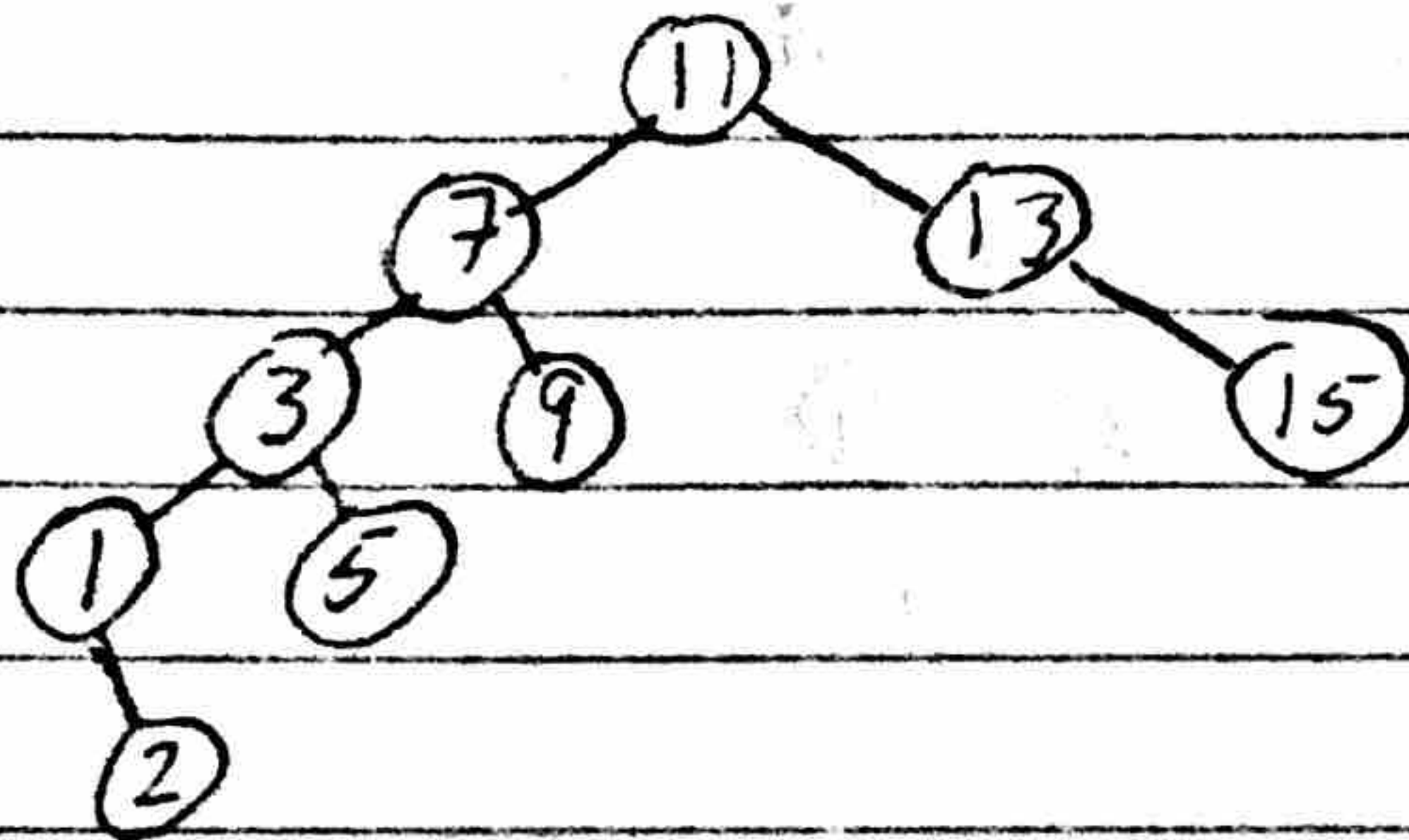
This means $\boxed{O(n \log n)}$

Part a)

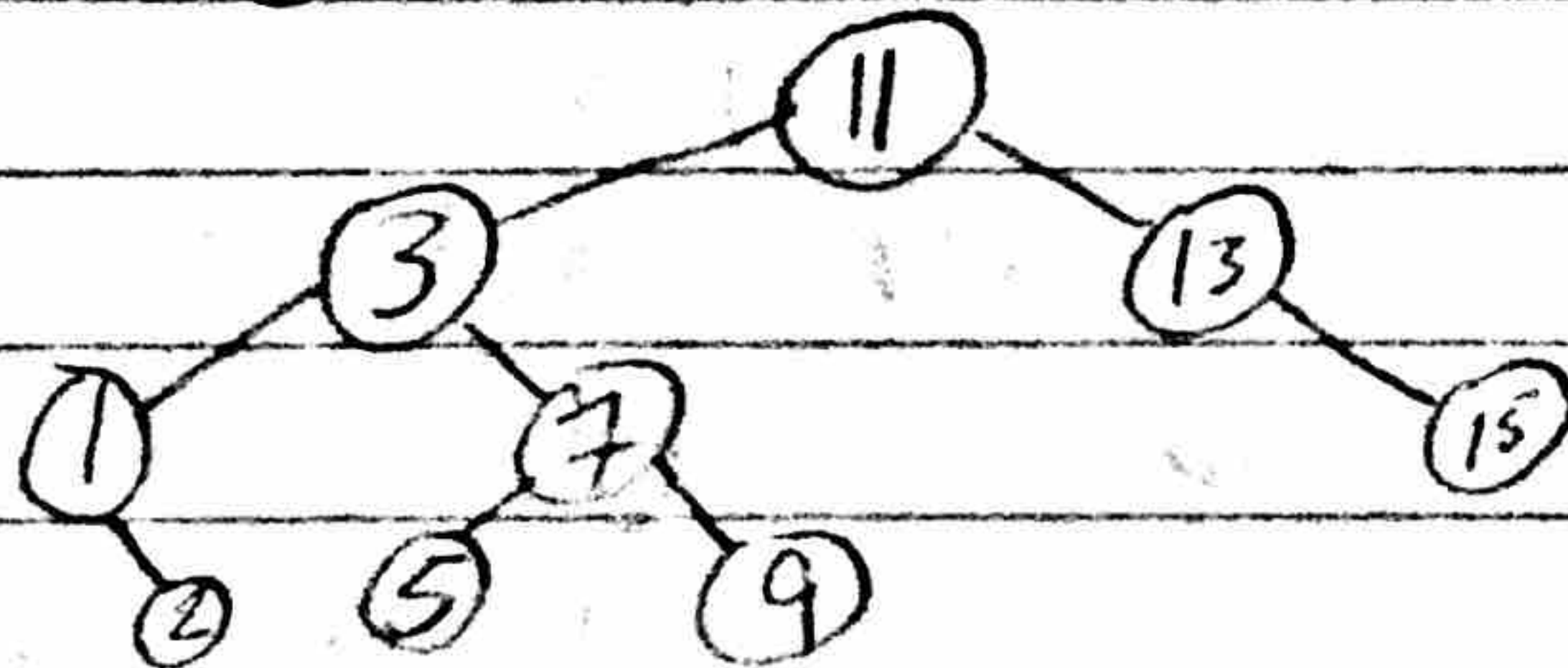
Question 6



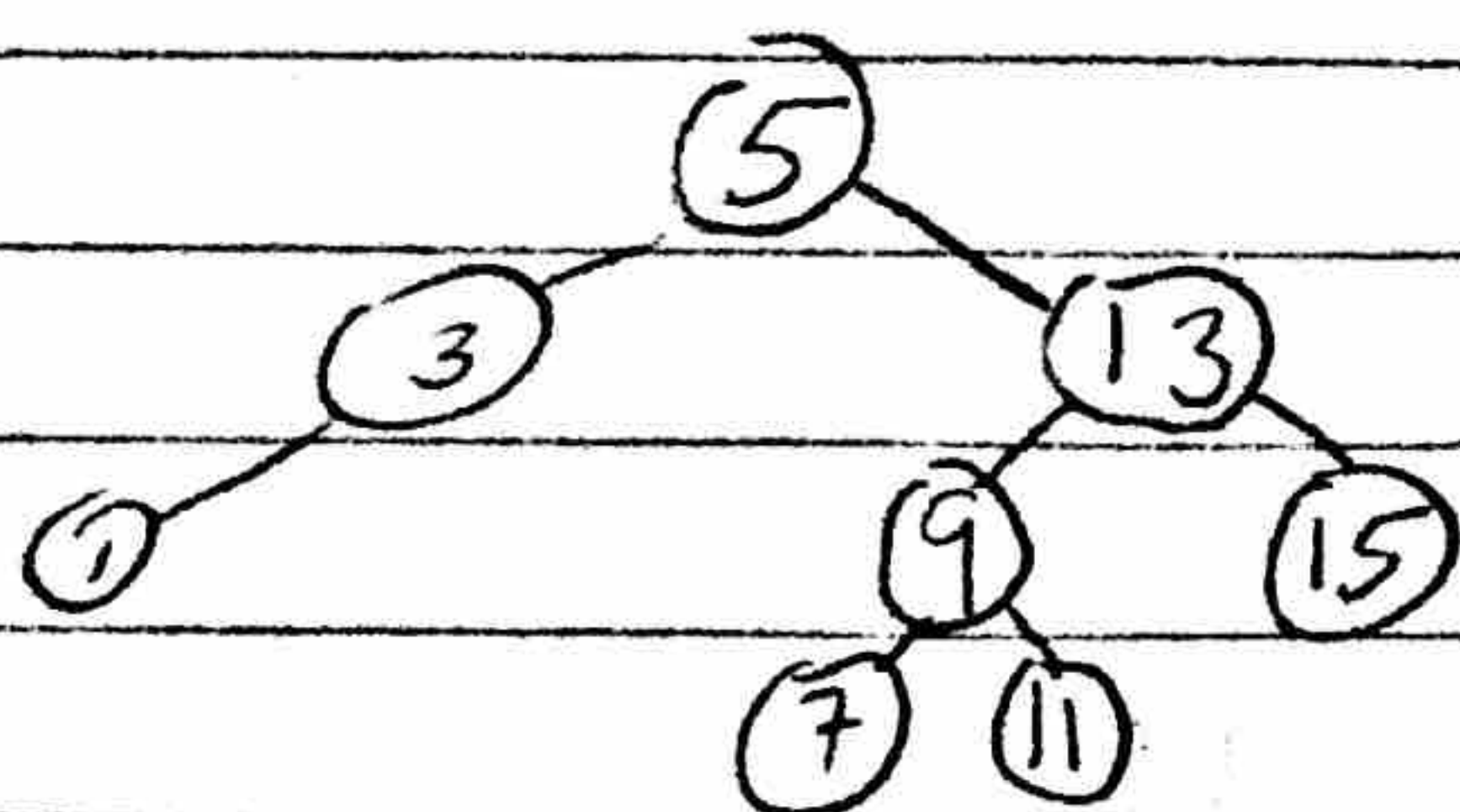
Insert 2:



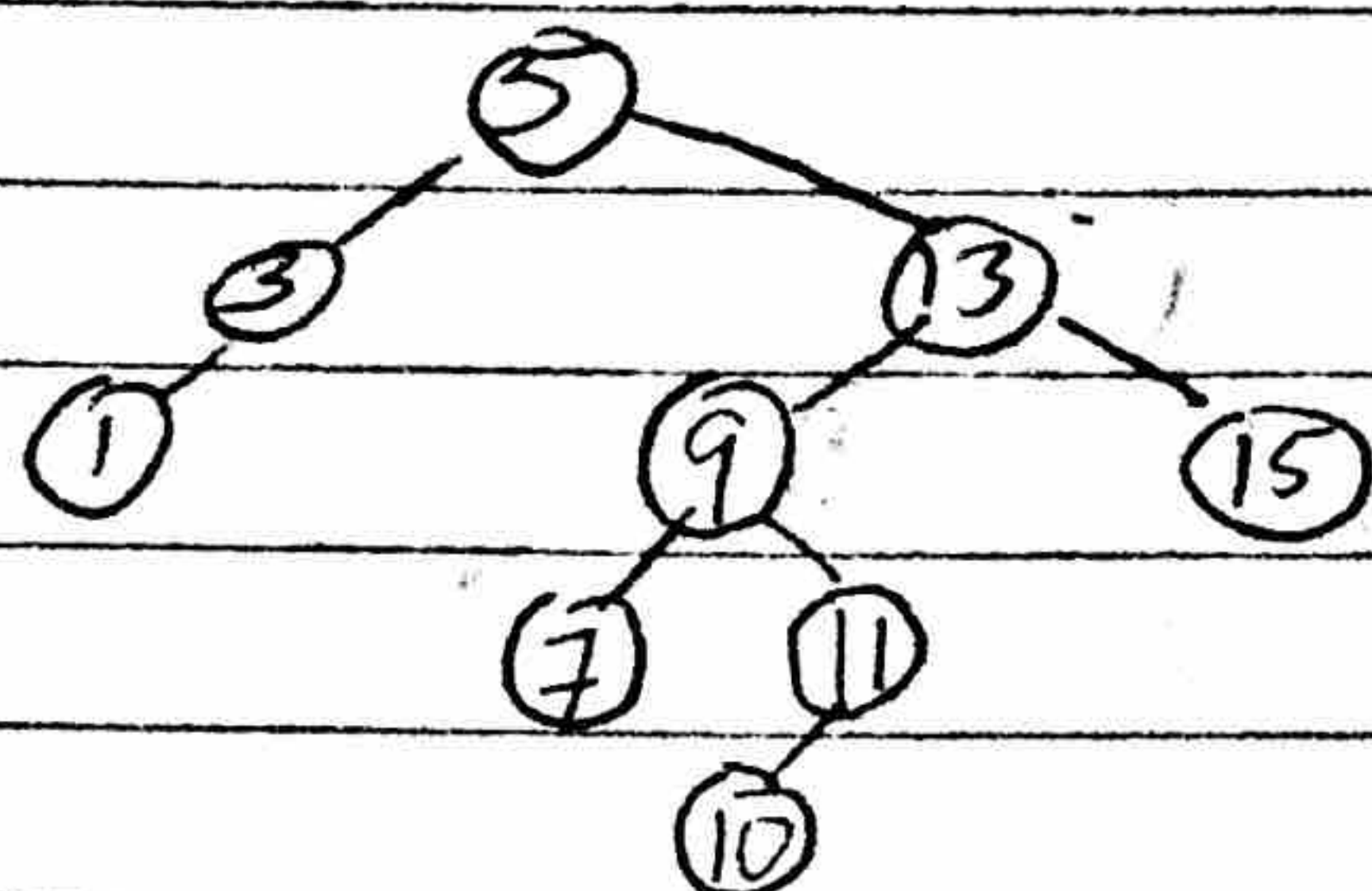
Right Rotate (7)



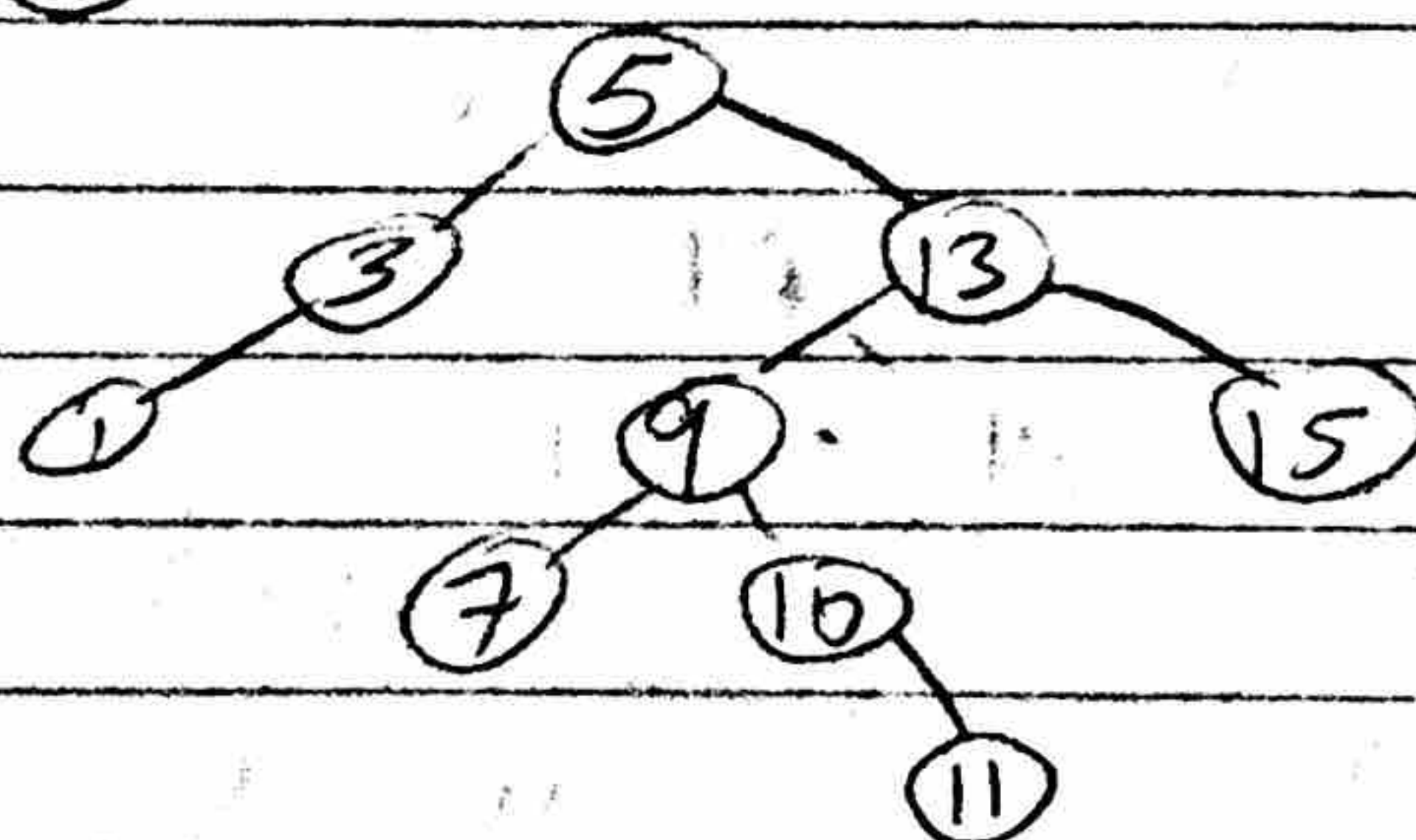
Part b)



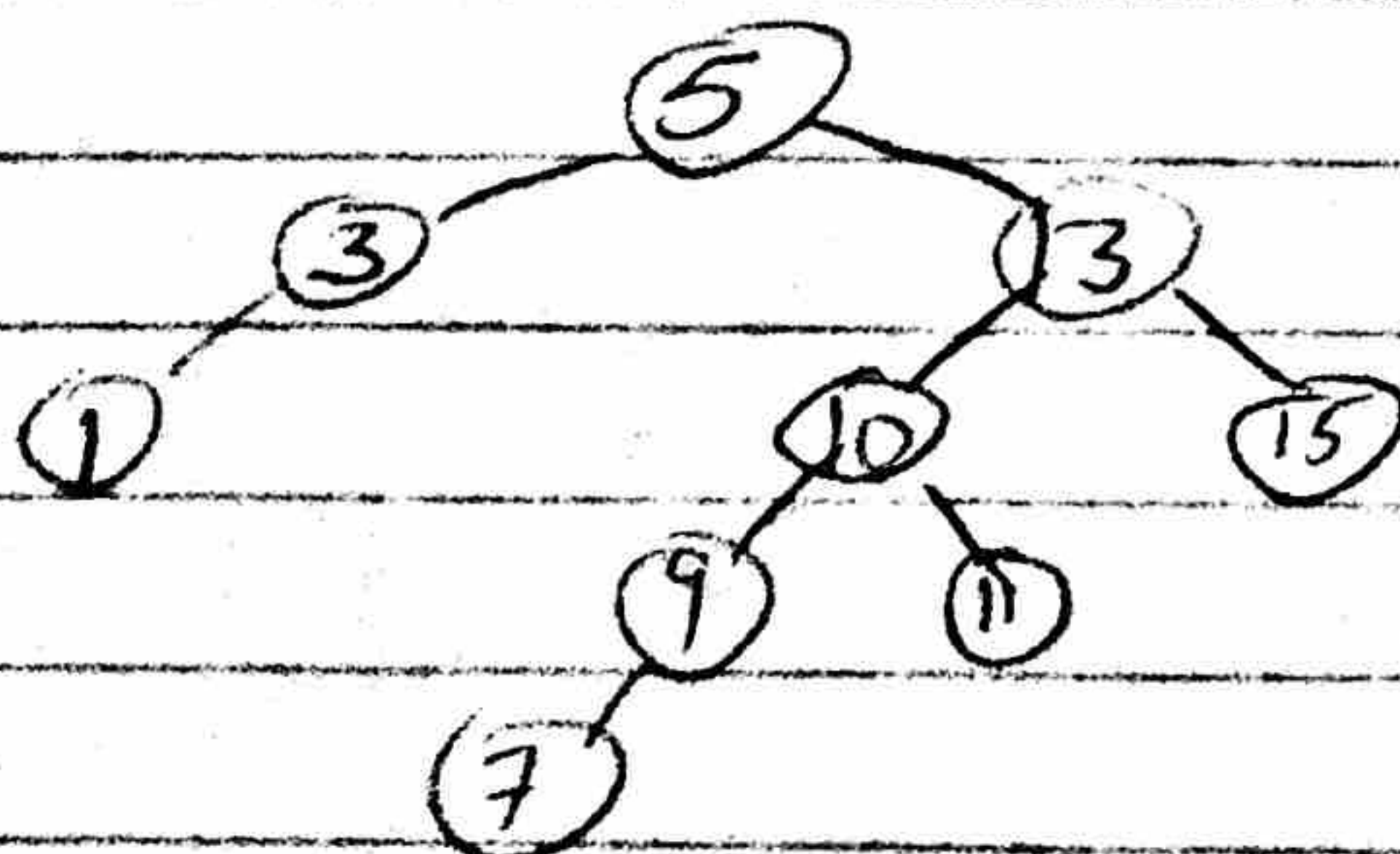
Insert 10:



Right Rotate 10



Left Rotate 9



Right Rotate 13

