

# **ECE220: Computer Systems and Programming**

## **Spring 2017 - Midterm Exam 2**

March 27, 2017

1. This is a closed book exam except for 1 sheet of hand-written notes
2. You may not use any personal electronic devices, such as cellphone and calculator
3. Absolutely no interaction between students is allowed
4. Illegible handwriting will be graded as incorrect

Select the Lecture Section You Will Attend to Pick Up Booklet

☐ Prof. Mitra (BL4, 9:30 am)

☐ Prof. Patel (BL2, 11:00 am)

☐ Prof. Chen (BL, 12:30 pm)

☐ Prof. Hu (BL3, 3:00 pm)

**Name:**\_\_\_\_\_

**NetID:**\_\_\_\_\_

**Room:**\_\_\_\_\_

Question 1 (20 points): \_\_\_\_\_

Question 2 (30 points): \_\_\_\_\_

Question 3 (20 points): \_\_\_\_\_

Question 4 (30 points): \_\_\_\_\_

**Total Score:** \_\_\_\_\_

## Problem 1 (20 points): Find Max Profit

For this problem, write a function `maxProfit` that takes as input an array of integers prices, and an integer `n` that indicates size of prices. The  $i$ -th element in the array is the price of a stock on the  $i$ -th day. The function should return the maximum profit achievable by buying one share of stock and then selling that share, possibly on the same day or a later day.

### Example 1:

Input: [2 7 3 1 2 8],  $n = 6$

Return value: 7

// Max profit = Buy at 1 on 4<sup>th</sup> day, Sell at 8 on 6<sup>th</sup> day

### Example 2:

Input: [5 4 3 2 1],  $n = 5$

Return value: 0

// Max profit = Buy at 1 on 5<sup>th</sup> day, Sell at 1 on 5<sup>th</sup> day

### Example 3:

Input: [4 3 2 3 6 4 2 1 3 4],  $n = 10$

Return value: 4

// Max profit = Buy at 2 on 3<sup>rd</sup> day, Sell at 6 on 5<sup>th</sup> day

Hint: you can use the following algorithm to solve this problem.

1. Assume buying the stock on the 1<sup>st</sup> day (first element) and then traverse the array to find max profit by selling on the same day or a later day.
2. Perform step 1 on each subsequent day (the rest of the elements in the array).

Rubric: There are two main ways of solving the problem: (i) with a single for loop and (ii) a double loop as suggested by the hint. For each of these solutions the rubric is as follows:

- Correct program 20 points
- Correct program with some logical mistakes lose (-2) each
  - Off by 1 in for loop or array indices
  - Taking min/max instead of max/min
  - Wrong polarity of subtraction
- Small syntax errors – 0.5 each
- Gibberish code (0 points)

Alternative methods for solving the problem could also get full marks.

```

int maxProfit(int *prices, int n)
{
/* YOUR CODE STARTS HERE */

//Implementation using the provided algorithm
    int i, j, temp, rc=0;
    for(i=0;i<size;i++){
        for(j=i+1;j<size;j++){
            temp = ptr[j]-ptr[i];
            if(temp>rc){
                rc = temp;
            }
        }
    }
    return rc;
}

// one for loop
int maxProfit(int *prices, int n){
    int i;
    int max_profit = 0;
    int min_price = INT_MAX;
    for(i = 0; i < n; ++i){
        if(min_price > prices[i])
            min_price = prices[i];
        if(max_profit < prices[i] - min_price)
            max_profit = prices[i] - min_price;
    }
    return max_profit;
}

```

## Problem 2 (30 points): C to LC-3

You will provide the data values that appear on the run-time stack during the execution of simple C program. Part of the stack frame for function main is shown in the memory table in Part A and Part B. R6 is the stack pointer and R5 is the frame pointer.

### *To ensure consistent answers:*

- a) after main's callee setup, R5 is 0xBCDB. Derive the value of R6 from this information
- b) place n above accum in the stack.

**Part A:** In the memory table, draw the stack at the point that JSR/JSRR begins transferring control to find\_ngon during the execution of the statement “z = find\_ngon(x, n);”. **Also indicate the values of R5 and R6 at this point of program execution.**

**Part B:** In the memory table, draw the stack right before “n = 0;” is executed; that is, after find\_ngon has set up the stack but before it has begun executing. **Also indicate the values of R5 and R6 at this point of program execution.**

**Part C:** Convert the find\_ngon function from C to an LC-3 subroutine with correct use of the run-time stack. Be aware of the requirements on code length (4 instructions each for callee setup and teardown) and execution (every local variable write is reflected in the stack frame).

```
int find_ngon(int perimeter, int side_len){
    int accum, n;
    n = 0;
    accum = perimeter;
    while (accum > 0) {
        accum -= side_len;
        n++;
    }
    return n;
}
```

```
int main() {
    int perim = 15;
    int side = 3;
    int z;

    z = find_ngon(perim, side);

    return 0;
}
```

**Part A (10 points): fill the stack with name (value is not needed)**

xBCC8	
xBCC9	
xBCCA	
xBCCB	
xBCCC	
xBCCD	
xBCCE	
xBCCF	
xBCD0	
xBCD1	
xBCD2	
xBCD3	
xBCD4	
xBCD5	
xBCD6	
xBCD7	perimeter (or x)
xBCD8	side_len (or n)
xBCD9	z
xBCDA	side
xBCDB	perim

Indicate the values of R5 and R6 at this point of program execution

R5 = xBCDB; R6 = xBCD7 (or xBCD6, but need to be consistent with part B)

**Rubric:** side (or side\_len), perim (or perimeter), R5, R6, each worth 2.5 pts  
(if doing caller save, it should be between z and first argument)

**Part B (10 points): fill the stack with name (value not needed)**

xBCC8	
xBCC9	
xBCCA	
xBCCB	
xBCCC	
xBCCD	
xBCCE	
xBCCF	
xBCD0	
xBCD1	
xBCD2	<b>n</b>
xBCD3	<b>accum</b>
xBCD4	<b>old frame pointer</b>
xBCD5	<b>return address</b>
xBCD6	<b>return value</b>
xBCD7	<b>perimeter (or x)</b>
xBCD8	<b>side_len (or n)</b>
xBCD9	<b>z</b>
xBCDA	<b>side</b>
xBCDB	<b>perim</b>

Indicate the values of R5 and R6 at this point of program execution

R5 = xBCD3; R6 = xBCD2 (or xBCD1, but need to be consistent with part A)

**Rubric:** side (or side\_len), perim (or perimeter), return value, return address, old frame pointer (or caller's frame pointer), accum, n each worth 1 pt; R5, R6 each worth 1.5 pts. (if doing callee save, it should be between old FP and accum)

**Part C (10 points):**

**;; FIND\_NGON Subroutine**

**; callee setup – push bookkeeping info and local variables**

**; This section shall be no longer than 4 instructions**

**Rubric: 2 points; note there are different possible correct ordering for storing R5, R6, R7**

ADD R6, R6, #-5 ; space bookkeeping info and local vars

STR R7, R6, #3 ; save return address

STR R5, R6, #2 ; save frame pointer

ADD R5, R6, #1 ; set frame pointer to first local var

**; function logic.**

**; Every variable write MUST be reflected in memory.**

**Rubric: 6 points; Loop with logic- 2 pts, update n in mem-1pts, update accum-1pts, load and initialize n, accum- 1 pt each**

LDR R1, R6, #5 ;var perimeter

LDR R2, R6, #6 ; var side\_len

ST R1, R5, #0 ; storing perim in accum so accum=perimeter

AND R3, R3, #0 ; n=0

ST R3, R5, #-1

NOT R2, R2

ADD R2, R2, #1 ; R2=-side\_len

LOOP

ADD R1, R1, #0

BRnz func\_end

ADD R1, R1, R2

ADD R3, R3, #1

STR R1, R5, #0

STR R3, R5, #1

BRnzp LOOP

**; callee tear-down – pop local variables, bookkeeping info**

**; This section shall be no longer than 4 instructions,**

**; INCLUDING any control flow transfers**

**Rubric: 2 points; the ordering here should be consistent with that in the first part above**  
**func\_end**

STR R3, R5, #3 ; push return value

ADD R6, R6, #4 ; pop local vars, return address frame pointer

LDR R7, R6, #-1 ; get return address

LDR R5, R6, #-2 ; get frame pointer

RET

### Problem 3 (20 points): Permutation using swap function

The goal of this problem is to print all possible permutations of a given string with no duplicate characters. For example, if the given string is “ABC”, your program should print out ABC, ACB, BCA, BAC, CAB and CBA.

In Lab 7 we did permutation using backtracking with a mask array of chosen elements. This time, you will use a different backtracking method which changes the characters in place in the string to generate permutations. The permutation function will use the familiar swap function, so let us start with that. The function declarations are provided; complete the functions.

**Part A (5 points):** Write the swap function to swap the character pointed by a and b.

```
void swap(char *a, char *b)
{
/* YOUR CODE STARTS HERE */
```

Rubric: correct answer 5 points

Small syntax error -0.5 each

This is the solution

```
{
    char temp;
    temp = *a;
    *a = *b;
    *b = temp;
}
```

```
/* YOUR CODE ENDS HERE */
}
```



**Part B (10 points):** The permutation function takes three parameters: `char* s` is the input string; `int left` and `right` are indices that define the substring that is being permuted. Here is the idea of the permutation function:

The base case occurs when the length of the substring is 1. In this case, there is nothing to permute and the input string is just printed out.

In the recursive case, for each character `c` in the range `[left, right]` the following three things are done: (i) first the `left` char is swapped with `c`, (ii) the permutation function is recursively called on the smaller substring defined by `[left+1, right]`, and finally, (iii) the swap in (i) is backtracked (i.e., undone).

Your implementation should not make copies of the input string. Our implementation used total 8 lines of C code.

```
void permutation(char *s, int left, int right)
{
    int i;
    if (left == right) /* 2 point Can also be left-
right==0 or left >= right*/
        printf("%s\n", s);
    else
    {
        for (i=left; i<=right; i++) /* 1 point(0.5 each) */
        {
            swap((s + left), (s + i)); /* 2 points Can also be
&s[left] and &s[i]*/
            permutation(s, left+1, right); /* 3 points */
            swap((s + left), (s + i)); /* 2 points Can also
be &s[left] and &s[i]*/
        }
    }
}
```

**Part C (5 points):** If the given string is “DEF”, what will be the output. Assume your implementation for the permutation is correct. (5 points)

```
int main()  
{  
    char str[] = "DEF";  
    permutation(str, 0, 2);  
    return 0;  
}
```

**Your Answer:**

1	DEF
2	DFE
3	EDF
4	EFD
5	FED
6	FDE

**Rubric:**

Order is correct: 5 pts

At least 3 in the right place: 2.5 pts

Else: 0 pts

## Problem4: Concepts (30 points)

**Part A (10 points):** print the addresses of an array

Danny created an array(arr1) that simply contains 5 integers. He wants to print out the addresses of arr1 by first assigning the addresses of arr1 to arr2 and then printing out using arr2 to see the arr1's addresses. However, the program does not work since he does not understand the concept of address, pointer, and array at all. As a good friend, you need to help him out by modifying his code. Note: (1) you can assume %p is the correct format to print address; (2) line 12 must not use arr1 after preprocessing.

line

```
0  #include <stdio.h>
1  int main() {
2      int arr1[5]={1,2,3,4,5};
3      int arr2[5];
4
5      int i;
6      for(i=0; i<5; i++){
7          arr2[i] = &(arr1+i);
8      }
9
10     for(i=0; i<5; i++){
11         /*assume %p is correct format to print address*/
12         printf("%p\n", &(arr2[i]) ); /*must not use arr1 after preprocessing*/
13     }
14
15 }
```

Please provide your answer below as shown in the example. (Must write the full syntax of the changed line)

**Example:**

**line 12 should be changed to**

**printf("%p\n", x);**

**Your answer:**

**Line 3 should be changed to int\* arr2[5]; (4 pts)**

**Line 7 should be changed to arr2[i] = arr+i; (3 pts)**

**Line 12 should be changed to printf("%d\n", arr2[i] ); (3 pts)**

**Part B (10 points):** Binary Search

An engineer is implementing a recursive binary search in C to use on an integer array sorted in **descending order** (largest to smallest). The code is as follows. Please help this engineer finish the function by filling in the correct parameters in line 5 and 6.

line

```
1    int binarySearch(int arr[], int start, int end, int item){
2        if (end >= start){
3            int mid = (end + start)/2;
4            if (arr[mid] == item) return mid;
5            if (arr[mid] > item) return binarySearch(arr, mid+1, end, item);
6            return binarySearch(arr, start, mid-1, item);
7        }
8        return -1;
9    }
```

Line 5: arr, mid+1 , end, item (5 pts) (0.5, 2, 2, 0.5)

Line 6: arr, start, mid-1 , item (5 pts) (0.5, 2, 2, 0.5)

**Part C (5 points):** How is a 2-D array stored in memory?

**Your Answer (no more than 30 words):**

**2-D array is stored as 1-D in row-major order in memory.**

---

---

---

---

**Part D (5 points):** What would happen if there's a bug in a recursive implementation, in which the code will never reach the base case?

**Your Answer (no more than 30 words):**

**The code will be stuck in an infinite loop, which will lead to segmentation fault**

**due to stack overflow.**

---

---

---

---

**Table E.2 The Standard ASCII Table**

ASCII			ASCII			ASCII			ASCII		
Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex	Character	Dec	Hex
nul	0	00	sp	32	20	@	64	40	`	96	60
soh	1	01	!	33	21	A	65	41	a	97	61
stx	2	02	"	34	22	B	66	42	b	98	62
etx	3	03	#	35	23	C	67	43	c	99	63
eot	4	04	\$	36	24	D	68	44	d	100	64
enq	5	05	%	37	25	E	69	45	e	101	65
ack	6	06	&	38	26	F	70	46	f	102	66
bel	7	07	'	39	27	G	71	47	g	103	67
bs	8	08	(	40	28	H	72	48	h	104	68
ht	9	09	)	41	29	I	73	49	i	105	69
lf	10	0A	*	42	2A	J	74	4A	j	106	6A
vt	11	0B	+	43	2B	K	75	4B	k	107	6B
ff	12	0C	,	44	2C	L	76	4C	l	108	6C
cr	13	0D	-	45	2D	M	77	4D	m	109	6D
so	14	0E	.	46	2E	N	78	4E	n	110	6E
si	15	0F	/	47	2F	O	79	4F	o	111	6F
dle	16	10	0	48	30	P	80	50	p	112	70
dc1	17	11	1	49	31	Q	81	51	q	113	71
dc2	18	12	2	50	32	R	82	52	r	114	72
dc3	19	13	3	51	33	S	83	53	s	115	73
dc4	20	14	4	52	34	T	84	54	t	116	74
nak	21	15	5	53	35	U	85	55	u	117	75
syn	22	16	6	54	36	V	86	56	v	118	76
etb	23	17	7	55	37	W	87	57	w	119	77
can	24	18	8	56	38	X	88	58	x	120	78
em	25	19	9	57	39	Y	89	59	y	121	79
sub	26	1A	:	58	3A	Z	90	5A	z	122	7A
esc	27	1B	;	59	3B	[	91	5B	{	123	7B
fs	28	1C	<	60	3C	\	92	5C		124	7C
gs	29	1D	=	61	3D	]	93	5D	}	125	7D
rs	30	1E	>	62	3E	^	94	5E	~	126	7E
us	31	1F	?	63	3F	_	95	5F	del	127	7F

NOTES: RTL corresponds to execution (after fetch!); JSRR not shown

ADD	0001	DR	SR1	0	00	SR2	ADD DR, SR1, SR2	LD	0010	DR	PCoffset9		LD DR, PCoffset9
DR ← SR1 + SR2, Setcc								DR ← M[PC + SEXT(PCoffset9)], Setcc					
ADD	0001	DR	SR1	1	imm5		ADD DR, SR1, imm5	LDI	1010	DR	PCoffset9		LDI DR, PCoffset9
DR ← SR1 + SEXT(imm5), Setcc								DR ← M[M[PC + SEXT(PCoffset9)]], Setcc					
AND	0101	DR	SR1	0	00	SR2	AND DR, SR1, SR2	LDR	0110	DR	BaseR	offset6	LDR DR, BaseR, offset6
DR ← SR1 AND SR2, Setcc								DR ← M[BaseR + SEXT(offset6)], Setcc					
AND	0101	DR	SR1	1	imm5		AND DR, SR1, imm5	LEA	1110	DR	PCoffset9		LEA DR, PCoffset9
DR ← SR1 AND SEXT(imm5), Setcc								DR ← PC + SEXT(PCoffset9), Setcc					
BR	0000	n	z	p	PCoffset9		BR{nzp} PCoffset9	NOT	1001	DR	SR	1111111	NOT DR, SR
((n AND N) OR (z AND Z) OR (p AND P)): PC ← PC + SEXT(PCoffset9)								DR ← NOT SR, Setcc					
JMP	1100	000	BaseR	000000			JMP BaseR	ST	0011	SR	PCoffset9		ST SR, PCoffset9
PC ← BaseR								M[PC + SEXT(PCoffset9)] ← SR					
JSR	0100	1	PCoffset11				JSR PCoffset11	STI	1011	SR	PCoffset9		STI SR, PCoffset9
R7 ← PC, PC ← PC + SEXT(PCoffset11)								M[M[PC + SEXT(PCoffset9)]] ← SR					
TRAP	1111	0000	trapvect8				TRAP trapvect8	STR	0111	SR	BaseR	offset6	STR SR, BaseR, offset6
R7 ← PC, PC ← M[ZEXT(trapvect8)]								M[BaseR + SEXT(offset6)] ← SR					

**End of ECE 220 Midterm Exam 2**