CSE 375: Assignment 3 Open-Hash

Design Framework:

The <u>sequential implementation</u> of Cuckoo Hash incorporates a 2D vector of pointers to bucket structs containing keys (which also represents the value for simplicity). The single-threaded application initializes the 2D vector with 2 hash tables of size key_max, and populates them with key_max / 2 random values. The program calls a driver function which executes contains/inserts/removes API functions respectively. If the number of iterations drastically exceeds the number of keys, then the likelihood of encountering a resize/rehash operation increases accordingly (currently the field for detecting a loop is set to 10).

The <u>concurrent implementation</u> of Cuckoo Hash incorporates a 2D vector of smart pointers (shared pointers) to a bucket struct of atomic keys. The multi-threaded application moves shared pointers around rather than copying atomics, and also implements a sort of lazy method for removing elements by resetting the bucket's key value. Using atomics, threads are able to load and store keys atomically using std::load and std::store attributes, and enforce a global memory order (sequential consistency) with std::memory_order_seq_cst.

The **transactional version** used 'synchronized {}' blocks, which is a synchronization mechanism that combines groups of statements in transactions that are atomic (either all statements occur, or nothing occurs). But it's important to note that synchronized blocks are NOT transactions. Transactional portions of the code (the parts previously wrapped in locks) wrapped in __transaction_atomic {} blocks are transactions in the C++ STM reference docs. I initially implemented these 'atomic transaction statements' with the __transaction_atomic keyword, but had too many 'unsafe function' errors that weren't being solved with the 'unsafe function' keyword, so I decided to switch to synchronized blocks.

In regards to rehash/resize operations, both the sequential and concurrent implementations implement the rehash() function as a single-threaded execution with single lock, locking the entire function and allowing only one thread to enter. Once inside the function, the thread locks all the buckets in the global hashtable_t struct (containing the individual hashtables) one at a time, and then performs the [1] **resize** and [2] **rehash** operations in order. The <u>sequential</u> and <u>concurrent</u> implementations differ in how they implement this function. Both create two separate temporary hash tables for storing values, resize the old hashtables, rehash the values in the temporary hashtable into the global hashtable, and then copy the values back into the appropriate index. But in the sequential implementation, .clear() is used to clear the hash tables, while in the concurrent implementation, a smarter technique is used: simply just adding new buckets to the appropriate newly sized locations.

Issues Encountered:

There were a number of issues encountered during testing. [1] Firstly, there were issues with .clear(), .resize(), and .shring_to_fit() operations in a multithreaded 2D vector. I suspect that some pointers were references to memory locations that were corrupted when clearing the vectors inside the rehash() function. This was solved by manually clearing and rehashing the hash tables by changing the values atomically to 0, and then pushing new buckets (rather than invoking the clear() and resize() functions). This doesn't seem to be a problem when defining 2 individual vectors, but seems to be an issue with a 2D vector definition. [2] There were deadlock and reference counting problems associated with my rehash() function in the concurrent implementation, again associated with the .clear() and .resize() calls. This was solved by removing those function calls, and redesigning my rehash() function by stripping down all the program components until I found the issue.

Performance Results:

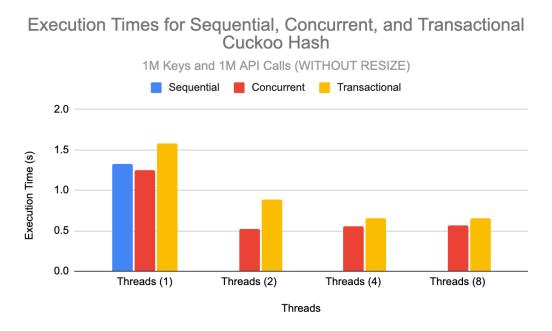


Figure 1: Sequential, Concurrent, and Transactional Implementation WTHOUT resize/rehash operations

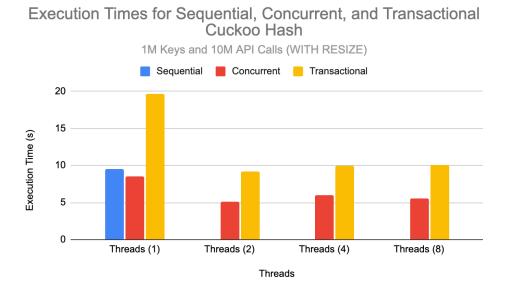


Figure 2: Sequential, Concurrent, and Transactional Implementation WITH resize/rehash operations

The results indicate that resizing and rehashing all the hashtable elements is an EXPENSIVE operation since it's single-threaded, resulting in a **10x** slowdown across the board from the sequential, concurrent, and transactional applications. The transactional implementation suffered the largest hit, since it's also using transaction-based synchronization primitives, which results in a larger slowdown in the overall program execution compared to both the sequential and concurrent implementations (as shown above).

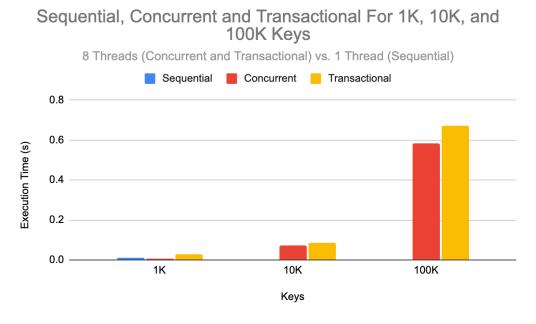


Figure 3: Sequential, Concurrent, and Transactional Implementation For Small Keys

Threads	Sequential	Concurrent	Transactional	Key Size	Data Structure
Threads (1)	1.32536	1.2523	1.58246	1M	W/O RESIZE
Threads (2)		0.525823	0.880877		
Threads (4)		0.555609	0.650723		
Threads (8)		0.565063	0.656737		
Threads	Sequential	Concurrent	Transactional	Key Size	Data Structure
Threads (1)	9.52855	8.5301	19.6046	1M	W/ RESIZE
Threads (2)		5.16483	9.13887		
Threads (4)		6.03343	9.96289		
Threads (8)		5.59087	10.052		
Keys	Sequential	Concurrent	Transactional		
1K	0.00861758	0.00601383	0.0287665		
10K		0.0725114	0.0841181		
100K		0.583123	0.673023		

Figure 4: Data Collection

In regards to the console output, executing the program prints the following to the console, for example. The distribution of key is as expected.

```
| Number of Tookup' operations: 4894 | Number of Tookup' appraised: 4894 | Number of Tookup' operations: 51 | Number of Tookup' operations: 52 | Operations operations: 52 | Operations
```

Figure 3: Sample Console Output

Note: The transactional version is compiled separately and includes its own main() function. Compile with 'g++ --std=c++17 -pthread -fgnu-tm transactional.cc -o transactional'