Contents						
1	All	Macros	1			
2	Data 2.1 2.2 2.3 2.4 2.5 2.6 2.7 2.8 2.9	MO with Update SparseTable (Rectangle Query)	1 1 1 2 2 2 3 3 3 4			
3	DP 3.1 3.2 3.3	Convex Hull Trick	4 4 4 5			
4	$4.1 \\ 4.2$	Point	5 5 6 7 9			
5	Grad 5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	LCA, ETT, VT SCC Euler Tour on Edge LCA In O(1) HLD Dinic Max Flow Min Cost Max Flow Bridge Tree	10 10 10 10 10 11 11 12 12 13			
6	Mat 6.1 6.2 6.3 6.4 6.5 6.6	Linear Sieve Pollard Rho Chinese Remainder Theorem Mobius Function	13 13 13 13 13 13			
7	7.1 7.2 7.3 7.4 7.5 7.6 7.7 7.8	Aho Corasick Double hash KMP Manacher's String Match FFT Suffix Array Trie	15 15 15 15 16 16 16 17			
8	Extr 8.1 8.2 8.3	Make File	17 17 17 17			

9.1 9.2 9.3 9.4 9.5	Stirling Numbers First Kind . Stirling Numbers Second Kind Other Combinatorial Identities	19 19 19 19 19 19

All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :20000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,
    ssse3,sse4,popcnt,abm,mmx,avx,tune=
    native")
#include <ext/pb_ds/assoc_container.hpp</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns
        iterator to the kth largest
        element counting from 0
   //order_of_key(val) --> returns the
        number of items in a set that
        are strictly smaller than our
template <typename DT>
using ordered_set = tree <DT, null_type</pre>
    , less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef LOCAL
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\',' << x <<</pre>
     '\'';}
void show(const string &x) {cerr << '\"</pre>
    ' << x << '\"';}
void show(bool x) {cerr << (x ? "true"</pre>
    : "false");}
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f';</pre>
    show(x.first); cerr << ", "; show(x | LT lazy[4 * N];
    .second); cerr << '}'; }
template<typename T>
void show(T x) {int f = 0; cerr << "{";</pre>
     for (auto &i: x) cerr << (f++ ? ",</pre>
     " : ""), show(i); cerr << "}";}
void debug_out(string s) {
   cerr << '\n';</pre>
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(),
         '), s.end());
   cerr << "
                  "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl:
   if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " <<</pre>
    __LINE__ << endl; debug_out(#x, x);
     cerr << endl;</pre>
#else
#define dbg(x...)
```

```
#endif
const int RANDOM = chrono::
    high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) ^ x;
unsigned hash_combine(unsigned a,
    unsigned b) { return a * 31 + b; }
struct chash {
 int operator()(int x) const { return
     hash_f(x); }
typedef gp_hash_table<int, int, chash>
    gp;
gp table;
mt19937_64 rng(atoi(argv[1]));
long long random(long long 1, long long
    r) {
   uniform_int_distribution<long long>
       dist(1, r);
   return dist(rng);
2
    Data Structure
```

2.1 Segment Tree

```
const int N = 1000006;
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N];
int L, R;
void pull(int s, int e, int node) {
 val[node] = val[node << 1] + val[node</pre>
       << 1 | 1];
void apply(const LT &U, int s, int e,
    int node) {
 val[node] += (e - s + 1) * U;
 lazy[node] += U;
void reset(int node) { lazy[node] =
    None; }
DT merge(const DT &a, const DT &b) {
   return a + b; }
DT get(int s, int e, int node) { return |}
     val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
 apply(lazy[node], s, s + e >> 1, node
       << 1):
 apply(lazy[node], s + e + 2 \gg 1, e,
     node << 1 | 1);
 reset(node);
void build(int s, int e, vector<DT> &v,
     int node = 1) {
 int m = s + e >> 1;
```

```
if (s == e) {
    val[node] = v[s];
   return;
  build(s, m, v, node * 2);
  build(m + 1, e, v, node * 2 + 1);
  pull(s, e, node);
void update(int S, int E, LT uval, int
    s = L, int e = R, int node = 1) {
  if (S > E) return;
  if (S == s and E == e) \{
    apply(uval, s, e, node);
    return;
  push(s, e, node);
  int m = s + e >> 1;
  update(S, min(m, E), uval, s, m, node
       * 2);
  update(max(S, m + 1), E, uval, m + 1,
       e, node *2 + 1);
  pull(s, e, node);
DT query(int S, int E, int s = L, int e
     = R, int node = 1) {
  if (S > E) return I;
  if (s == S and e == E) return get(s,
      e, node);
  push(s, e, node);
  int m = s + e >> 1;
  DT L = query(S, min(m, E), s, m, node
       * 2);
  DT R = query(max(S, m + 1), E, m + 1,
       e, node * 2 + 1);
  return merge(L, R);
void init(int _L, int _R, vector<DT> &v
  L = L, R = R;
  build(L, R, v);
```

2.2Persistent Segment Tree

```
struct Node {
   int 1 = 0, r = 0, val = 0;
tr[20 * N];
int ptr = 0;
int build(int st, int en) {
   int u = ++ptr;
   if (st == en) return u;
   int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   1 = build(st, mid);
   r = build(mid + 1, en);
   val = tr[l].val + tr[r].val;
   return u;
int update(int pre, int st, int en, int
     idx, int v) {
   int u = ++ptr;
   tr[u] = tr[pre];
   if (st == en) {
       tr[u].val += v;
       return u;
   int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   if (idx <= mid) {</pre>
```

```
r = tr[pre].r;
       1 = update(tr[pre].1, st, mid,
           idx, v);
   } else {
       1 = tr[pre].1;
       r = update(tr[pre].r, mid + 1,
           en, idx, v);
   tr[u].val = tr[l].val + tr[r].val;
   return u;
// finding the kth elelment in a range
int query(int left, int right, int st,
    int en, int k) {
   if (st == en) return st;
   int cnt = tr[tr[right].1].val - tr[
        tr[left].l].val;
   int mid = (st + en) / 2;
   if (cnt >= k) return query(tr[left].
        1, tr[right].1, st, mid, k);
   else return query(tr[left].r, tr[
       right].r, mid + 1, en, k - cnt)
int V[N], root[N], a[N];
int main() {
   map<int, int> mp; int n, q;
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >>
       a[i], mp[a[i]];
   int c = 0;
   for (auto x : mp) mp[x.first] = ++c,
        V[c] = x.first;
   root[0] = build(1, n);
   for (int i = 1; i <= n; i++) {</pre>
       root[i] = update(root[i - 1], 1,
            n, mp[a[i]], 1);
   while (q--) {
       int 1, r, k; cin >> 1 >> r >> k;
            1++, k++;
       cout << V[query(root[l - 1],</pre>
           root[r], 1, n, k)] << '\n';
   }
   return 0;
```

2.3 Lazy Persistent Segment Tree

```
struct node {
   int val;
   int lazy;
   int left, right;
} nodes[MAXN];
int sz = 0;
int a[MAXN];
int build(int s, int e) {
   int curr = sz;
   sz++;
   if(s == e) {
       nodes[curr].left = nodes[curr].
           right = -1;
       nodes[curr].val = a[s];
       nodes[curr].lazy = 0;
       return curr;
   int m = (s + e) / 2;
   nodes[curr].left = build(s, m);
   nodes[curr].right = build(m + 1, e);
```

```
nodes[curr].val = nodes[nodes[curr].
        left].val + nodes[nodes[curr].
        right].val;
   return curr;
int update(int prev, int s, int e, int
    1, int r, int val) {
    if(s > r or e < 1) return prev;</pre>
   int curr = sz;
   sz++;
   if(s \ge 1 \text{ and } e \le r) {
       nodes[curr].left = nodes[prev].
            left;
       nodes[curr].right = nodes[prev].
            right;
       nodes[curr].val = nodes[prev].
            val + val*(e-s+1);
       nodes[curr].lazy = nodes[prev].
            lazy + val;
       return curr;
   int m = (s + e) / 2;
   nodes[curr].left = update(nodes[prev
        ].left, s, m, l, r, val);
   nodes[curr].right = update(nodes[
        prev].right, m+1, e, l, r, val)
   nodes[curr].lazy = nodes[prev].lazy;
   nodes[curr].val = nodes[nodes[curr].
        left].val + nodes[nodes[curr].
        right].val + nodes[curr].lazy*(
        e-s+1);
   return curr;
11 query(int nd, int s, int e, int 1,
    int r, 11 c) {
    if(s > r or e < 1) return 0;</pre>
    if(s >= 1 and e <= r) return c*(e-s</pre>
        +1) + nodes[nd].val;
   int m = (s+e)/2;
   c += nodes[nd].lazy;
   return query(nodes[nd].left, s, m, 1
        , r, c) + query(nodes[nd].right
        , m+1, e, l, r, c);
```

2.4 Implicit Segment Tree

```
struct node {
   ll val;
   ll lazy;
   int left, right;
ነ:
node nodes[4*1000006];
int sz;
void reset(int nd) {
   nodes[nd].left = nodes[nd].right =
   nodes[nd].val = nodes[nd].lazy = 0;
int getLeft(int nd) {
   if(nodes[nd].left == -1) {
       nodes[nd].left = sz++;
       reset(nodes[nd].left);
   return nodes[nd].left;
int getRight(int nd) {
    if(nodes[nd].right == -1) {
       nodes[nd].right = sz++;
```

```
reset(nodes[nd].right);
   }
   return nodes[nd].right;
void update(int nd, int s, int e, int l
    , int r, ll val) {
   if(e < 1 or s > r) return;
   if(s \ge 1 and e \le r) {
       nodes[nd].val += val*(e-s + 1);
       nodes[nd].lazy += val;
       return;
   }
   int m = (s + e) / 2;
   update(getLeft(nd), s, m, l ,r, val)
   update(getRight(nd), m+1, e, 1, r,
   nodes[nd].val = nodes[nodes[nd].left
        ].val + nodes[nodes[nd].right].
        val + nodes[nd].lazy * (e-s+1);
ll query(int nd, int s, int e, int l,
    int r, ll c) {
   if(s > r or e < 1) return 0;</pre>
   if(s \ge 1 \text{ and } e \le r) {
       return nodes[nd].val + c*(e-s+1)
   int m = (s+e)/2;
   c += nodes[nd].lazy;
   return query(getLeft(nd), s, m, 1, r
        , c) + query(getRight(nd), m+1,
         e, 1, r, c);
```

DSU With Rollbacks 2.5

```
struct Rollback_DSU {
 int n;
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n,
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par
       [node])
     ; // no path compression
   return node;
 void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)])
        swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back())
     par[op.back().second] = op.back().
         second;
     sz[op.back().first] -= sz[op.back
         ().second];
 }
```

```
2.6 BIT-2D
```

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j \leftarrow -j & j)
       bit[x][j] += val;
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j & -j) ans
       += bit[x][j];
 return ans;
}
int get(int x1, int y1, int x2, int y2)
 return get(x2, y2) - get(x1 - 1, y2)
      - get(x2, y1 - 1) + get(x1 - 1,
      y1 - 1);
```

2.7 Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];
vector<LL> merge(vector<LL> v1, vector<
    LL> v2) {
 LL i = 0, j = 0;
 vector<LL> ret;
  while (i < v1.size() || j < v2.size()</pre>
      ) {
   if (i == v1.size()) {
     ret.push_back(v2[j]);
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
     i++;
   } else {
     if (v1[i] < v2[j]) {</pre>
       ret.push_back(v1[i]);
       i++;
     } else {
       ret.push_back(v2[j]);
       j++;
   }
 return ret;
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 }
 LL leftNode = 2 * node, rightNode = 2
       * node + 1:
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
```

```
Tree[node] = merge(Tree[leftNode],
      Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1,
    LL r, LL k) {
 if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin
        (), Tree[node].end(), k) -
          Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, 1, r,
        query(rightNode, mid + 1, ed, 1
            , r, k);
```

2.8 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs =
int arr[N], freq[2 * N], cnt[2 * N], id
    [N], ans[N];
struct query {
 int 1, r, t, L, R;
 query(int l = 1, int r = 0, int t =
      1, int id = -1)
     : 1(1), r(r), t(t), L(1 / sz), R(r)
           / sz) {}
  bool operator<(const query &rhs)</pre>
      const {
   return (L < rhs.L) or (L == rhs.L
       and R < rhs.R) or
          (L == rhs.L and R == rhs.R
              and t < rhs.t);</pre>
 }
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
 cnt[freq[arr[idx]]]++;
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[ int A[MAXN][MAXN];</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
```

```
arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0);
int main() {
  int n, q;
  cin >> n >> q;
  int counter = 0;
  map<int, int> M;
  for (int i = 1; i <= n; i++) {</pre>
    cin >> arr[i];
    if (!M[arr[i]]) M[arr[i]] = ++
        counter:
    arr[i] = M[arr[i]];
  iota(id, id + N, 0);
  while (q--) {
    int tp, x, y;
    cin >> tp >> x >> y;
    if (tp == 1)
     Q[++qi] = query(x, y, ui);
    else {
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = \{x, y, arr[x]\};
     arr[x] = y;
  }
  t = ui;
  cnt[0] = 3 * n;
  sort(id + 1, id + qi + 1, [&](int x,
      int y) { return Q[x] < Q[y]; });
  for (int i = 1; i <= qi; i++) {</pre>
    int x = id[i];
    while (Q[x].t > t) apply(++t);
    while (Q[x].t < t) undo(t--);
    while (Q[x].1 < 1) add(--1);
    while (Q[x].r > r) add(++r);
    while (Q[x].1 > 1) remove(1++);
    while (Q[x].r < r) remove(r--);
   ans[x] = mex();
 for (int i = 1; i <= qi; i++) cout <<</pre>
       ans[i] << '\n';
```

2.9 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
   for (int q = 1; (1 << q) <= N; q++)
```

```
int add = 1 << (q - 1);
     for (int j = 1; j + add <= N; j++)</pre>
       M[i][j][0][q] = max(M[i][j][0][q]
             - 1], M[i][j + add][0][q -
            1]);
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q
          ++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p]
              - 1][q], M[i + add][j][p -
               1] [q]);
     }
   }
 }
// returns max of all A[i][j], where x1
    =i<=x2 and y1<=j<=y2
int Query(int x1, int y1, int x2, int
    y2) {
  int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 \ll kX;
 int addY = 1 << kY;</pre>
 int ret1 = max(M[x1][y1][kX][kY], M[
      x1][y2 - addY + 1][kX][kY]);
  int ret2 = max(M[x2 - addX + 1][y1][
      kX][kY],
                M[x2 - addX + 1][y2 -
                    addY + 1][kX][kY]);
 return max(ret1, ret2);
```

2.10 Sparse Table

```
// O-based indexing, query finds in
    range [first, last]
#define lg(x) (31 - __builtin_clz(x))
const int N = 1e5 + 1;
const int K = lg(N);
int a[N], tr[N][K + 1];
namespace sparse_table {
 int f(int p1, int p2) { return min(p1
      , p2); }
 void build() {
   for(int i = 0; i < n; i++)</pre>
     tr[i][0] = a[i];
   for(int j = 1; j <= K; j++) {</pre>
     for(int i = 0; i + (1<<j) <= n; i</pre>
          ++)
       tr[i][j] = f(tr[i][j - 1], tr[i
            + (1<<(j - 1))][j - 1]);
   }
 }
 int query(int 1, int r) {
   int d = lg(r - l + 1);
   return f(table[1][d], table[r - (1<<</pre>
        d) + 1][d]);
 }
}
```

DP

3.1 Convex Hull Trick

```
struct line {
    11 m, c;
    line() {}
    line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
    vector<line> lines;
    int ptr = 0;
    convex_hull_trick() {}
    bool bad(line a, line b, line c) {
         return 1.0 * (c.c - a.c) * (a.m - b.
                    m) < 1.0 * (b.c - a.c) * (a.m - a.c) * (a.c) * (
                      c.m);
    }
    void add(line L) {
         int sz = lines.size();
         while (sz \ge 2 \&\& bad(lines[sz - 2],
                        lines[sz - 1], L)) {
              lines.pop_back();
         lines.pb(L);
    11 get(int idx, int x) { return (111
                * lines[idx].m * x + lines[idx].c
                ); }
    11 query(int x) {
         if (lines.empty()) return 0;
         if (ptr >= lines.size()) ptr = lines
                     .size() - 1;
          while (ptr < lines.size() - 1 && get
                     (ptr, x) > get(ptr + 1, x)) ptr
                     ++;
         return get(ptr, x);
   }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
    fastio;
    int t;
    cin >> t;
    while (t--) {
         int n, a, b, c;
         cin >> n >> a >> b >> c;
         for (int i = 1; i <= n; i++) cin >>
                     sum[i];
         for (int i = 1; i <= n; i++) dp[i] =</pre>
                       0, sum[i] += sum[i - 1];
         convex_hull_trick cht;
          cht.add(line(0, 0));
          for (int pos = 1; pos <= n; pos++) {</pre>
              dp[pos] = cht.query(sum[pos]) - 1
                         11 * a * sqr(sum[pos]) - c;
              cht.add(line(211 * a * sum[pos],
                          dp[pos] - a * sqr(sum[pos])));
         ll ans = (-111 * dp[n]);
         ans += (111 * sum[n] * b);
         cout << ans << "\n";
              Dynamic CHT
3.2
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 IS_QUERY = -(1LL << 62);</pre>
struct line {
  11 m. b:
  mutable function <const line*()> succ
  bool operator < (const line &rhs)</pre>
      const {
    if (rhs.b != IS_QUERY) return m <</pre>
        rhs.m;
    const line *s = succ();
    if (!s) return 0;
    11 x = rhs.m;
    return b - s -> b < (s -> m - m) * x
 }
};
struct HullDynamic : public multiset <</pre>
    line> {
  bool bad (iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y -> m == z -> m && y -> b
          \langle = z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x
        -> m && y -> b <= x -> b;
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z
        -> m - y -> m) >= 1.0 * (y -> b)
         -z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 }
  void insert_line (ll m, ll b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) ==
         end() ? 0 : &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(
        y))) erase(next(y));
    while (y != begin() && bad(prev(y)))
          erase(prev(y));
  11 eval (ll x) {
    auto 1 = *lower_bound((line) {x,
        IS_QUERY});
    return 1.m * x + 1.b;
 }
int main() {
  HullDynamic hull;
  hull.insert_line(1, 1);
  hull.insert_line(-1, 1);
  cout << hull.eval(69) << endl;</pre>
  cout << hull.eval(420) << endl;</pre>
  return 0;
```

3.3 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c)
LL calc(line L, LL x) { return 1LL * L.
    m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft =
      NULL, node *rt = NULL)
     : L(line(m, c)), lft(lft), rt(rt)
};
struct LiChao {
 node *root;
 LiChao() { root = new node(); }
 void update(node *now, int L, int R,
      line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(
        lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
           node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo;
     if (now->lft == NULL) now->lft =
         new node():
     update(now->lft, L, mid, hi);
 }
 LL query(node *now, int L, int R, LL
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x <= mid)</pre>
     return max(calc(now->L, x), query(
         now->lft, L, mid, x));
   else
     return max(calc(now->L, x), query(
          now->rt, mid + 1, R, x));
 }
};
```

Geometry

```
4.1 Point
typedef double Tf;
typedef double Ti; /// use long long
    for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ?</pre>
    0 : (x < 0 ? -1 : 1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y)
```

```
Point operator+(const Point& u) const
                  { return Point(x + u.x, y + u.y) | Point normal(Point a) {
                ; }
     Point operator-(const Point& u) const
                  { return Point(x - u.x, y - u.y) |}
     Point operator*(const LL u) const {
                return Point(x * u, y * u); }
     Point operator*(const Tf u) const {
                return Point(x * u, y * u); }
     Point operator/(const Tf u) const {
                return Point(x / u, y / u); }
     bool operator==(const Point& u) const
          return dcmp(x - u.x) == 0 \&\& dcmp(y)
                     -u.y) == 0;
      bool operator!=(const Point& u) const
                   { return !(*this == u); }
     bool operator<(const Point& u) const</pre>
          return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                    -u.x) == 0 && dcmp(y - u.y) <
                    0);
 Ti dot(Point a, Point b) { return a.x *
              b.x + a.y * b.y; }
 Ti cross(Point a, Point b) { return a.x |};
              * b.y - a.y * b.x; }
 Tf length(Point a) { return sqrt(dot(a,
              a)): }
 Ti sqLength(Point a) { return dot(a, a)
            ; }
 Tf distance(Point a, Point b) { return
           length(a - b); }
 Tf angle(Point u) { return atan2(u.y, u
            .x); }
  // returns angle between oa, ob in (-PI
            , PI]
 Tf angleBetween(Point a, Point b) {
     Tf ans = angle(b) - angle(a);
     return ans <= -PI ? ans + 2 * PI : (
                ans > PI ? ans - 2 * PI : ans);
  // Rotate a ccw by rad radians, Tf Ti
 Point rotate(Point a, Tf rad) {
     return Point(a.x * cos(rad) - a.y *
                sin(rad),
                                   a.x * sin(rad) + a.y * cos
                                              (rad));
 // rotate a ccw by angle th with cos(th
           ) = co \&\& sin(th) = si, tf ti same
 Point rotatePrecise(Point a, Tf co, Tf
      return Point(a.x * co - a.y * si, a.y
                  * co + a.x * si);
 Point rotate90(Point a) { return Point
            (-a.y, a.x); }
 // scales vector a by s such that
           length of a becomes s, Tf Ti same
Point scale(Point a, Tf s) { return a /
```

length(a) * s; }

```
// returns an unit vector perpendicular
     to vector a, Tf Ti same
 Tf 1 = length(a);
 return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if
    on ab && -1 if right of ab
int orient(Point a, Point b, Point c) {
     return dcmp(cross(b - a, c - a));
/// Use as sort(v.begin(), v.end(),
    polarComp(0, dir))
/// Polar comparator around O starting
    at direction dir
struct polarComp {
 Point O, dir;
 polarComp(Point 0 = Point(0, 0),
      Point dir = Point(1, 0): O(0),
      dir(dir) {}
 bool half(Point p) {
   return dcmp(cross(dir, p)) < 0 ||</pre>
          (dcmp(cross(dir, p)) == 0 \&\&
              dcmp(dot(dir, p)) > 0);
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) <</pre>
        make_tuple(half(q), cross(p, q)
struct Segment {
 Point a, b;
 Segment(Point aa, Point bb) : a(aa),
     b(bb) {}
typedef Segment Line;
struct Circle {
 Point o;
 Tf r;
 Circle(Point o = Point(0, 0), Tf r =
      0) : o(o), r(r) {}
  // returns true if point p is in ||
      on the circle
 bool contains(Point p) { return dcmp(
      sqLength(p - o) - r * r) \le 0;
  // returns a point on the circle rad
     radians away from +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value
   return Point(o.x + cos(rad) * r, o.y
        + sin(rad) * r);
 // area of a circular sector with
      central angle rad
 Tf area(Tf rad = PI + PI) { return
      rad * r * r / 2; }
  // area of the circular sector cut by
       a chord with central angle alpha
 Tf sector(Tf alpha) { return r * r *
      0.5 * (alpha - sin(alpha)); }
```

4.2 Linear

```
// **** LINE LINE INTERSECTION START
// returns true if point p is on
    segment s
```

```
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p))
      == 0 && dcmp(dot(s.a - p, s.b - p
// returns true if segment p && q touch |}
     or intersect
bool segmentsIntersect(Segment p,
    Segment q) {
 if (onSegment(p.a, q) || onSegment(p.
      b, q)) return true;
 if (onSegment(q.a, p) || onSegment(q.
      b, p)) return true;
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 &&
      dcmp(c3) * dcmp(c4) < 0;
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.
      a)) == 0;
// lines are represented as a ray from
    a point: (point, vector)
// returns false if two lines (p, v) &&
     (q, w) are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference, Tf Ti
     Same
bool lineLineIntersection(Point p,
    Point v, Point q, Point w, Point& o
 if (dcmp(cross(v, w)) == 0) return
      false:
 Point u = p - q;
 o = p + v * (cross(w, u) / cross(v, w)
      ));
 return true;
// returns false if two lines p && q
    are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference
bool lineLineIntersection(Line p, Line
    q, Point& o) {
 return lineLineIntersection(p.a, p.b
      - p.a, q.a, q.b - q.a, o);
// returns the distance from point a to |}
     line 1
// **** LINE LINE INTERSECTION FINISH
    ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(1.b - 1.a, p - 1.a)
      / length(1.b - 1.a));
// returns the shortest distance from
    point a to segment s
Tf distancePointSegment(Point p,
    Segment s) {
 if (s.a == s.b) return length(p - s.a
      );
 Point v1 = s.b - s.a, v2 = p - s.a,
      v3 = p - s.b;
 if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
```

```
else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
   return abs(cross(v1, v2) / length(v1
        )):
// returns the shortest distance from
    segment p to segment q
Tf distanceSegmentSegment(Segment p,
    Segment q) {
  if (segmentsIntersect(p, q)) return
      0;
 Tf ans = distancePointSegment(p.a, q)
 ans = min(ans, distancePointSegment(p
      .b, q));
 ans = min(ans, distancePointSegment(q
      .a, p));
  ans = min(ans, distancePointSegment(q
      .b, p));
 return ans;
// returns the projection of point p on
     line 1, Tf Ti Same
Point projectPointLine(Point p, Line 1)
 Point v = 1.b - 1.a;
 return 1.a + v * ((Tf)dot(v, p - 1.a)
       / dot(v, v));
4.3 Circular
// Extremely inaccurate for finding
    near touches
// compute intersection of line 1 with
    circle c
// The intersections are given in order
     of the ray (l.a, l.b), Tf Ti same
vector<Point> circleLineIntersection(
    Circle c, Line 1) {
 vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B *
       B - A * C;
  if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D
      + EPS)) / A);
  if (D > EPS) ret.push_back(1.a + b *
      (-B + sqrt(D)) / A);
 return ret;
// signed area of intersection of
    circle(c.o, c.r) &&
 // triangle(c.o, s.a, s.b) [cross(a-o,
    b-o)/21
Tf circleTriangleIntersectionArea(
    Circle c, Segment s) {
 using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
 if (dcmp(distancePointSegment(c.o, s)
       - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b -
         c.o) * (c.r * c.r) / 2.0;
  // triangle
  if (dcmp(OA - c.r) <= 0 && dcmp(OB -
      c.r) \ll 0
```

```
return cross(c.o - s.b, s.a - s.b) /
         2.0:
 // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect =
      circleLineIntersection(c, s);
 return circleTriangleIntersectionArea
      (c, Segment(s.a, Sect[0])) +
        circleTriangleIntersectionArea(
            c, Segment(Sect[0], Sect
            [1])) +
        circleTriangleIntersectionArea(
            c, Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c
    .r) && simple polyson(p[])
Tf circlePolyIntersectionArea(Circle c,
     Polygon p) {
 Tf res = 0;
 int n = p.size();
 for (int i = 0; i < n; ++i)</pre>
        circleTriangleIntersectionArea(
        c, Segment(p[i], p[(i + 1) % n
        ]));
 return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
    ----> -2
// interior tangents (d = R - r)
    ----> -1
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r)
    ) ----> 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
int circleCirclePosition(Circle c1,
    Circle c2) {
 Tf d = length(c1.o - c2.o);
 int in = dcmp(d - abs(c1.r - c2.r)),
      ex = dcmp(d - (c1.r + c2.r));
 return in < 0 ? -2 : in == 0 ? -1 :
      ex == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points
    between two circles c1 && c2, Tf Ti
vector<Point> circleCircleIntersection(
    Circle c1, Circle c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o);
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return
       ret;
 if (dcmp(abs(c1.r - c2.r) - d) > 0)
      return ret;
 Point v = c2.o - c1.o;
 Tf co = (c1.r * c1.r + sqLength(v) -
      c2.r * c2.r) / (2 * c1.r * length
      (v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co,
       -si), c1.r) + c1.o;
 Point p2 = scale(rotatePrecise(v, co,
       si), c1.r) + c1.o;
```

```
int circleCircleTangencyPoints(Circle
                                                                                     for (int i = 0; i < n; i++) A[i] = A[</pre>
 ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
                                              c1, Circle c2, vector<Point> &a,
                                                                                         i + 1] - A[i];
                                                                                     for (int i = 0; i < m; i++) B[i] = B[</pre>
                                                                      vector<
 return ret;
                                                                                          i + 1] - B[i];
                                                                           Point>
// intersection area between two
                                                                           &b) {
                                           a.clear(), b.clear();
                                                                                     Polygon C(n + m + 1);
    circles c1, c2
Tf circleCircleIntersectionArea(Circle
                                           int cnt = 0;
                                                                                     C[0] = A.back() + B.back();
    c1, Circle c2) {
                                           if (dcmp(c1.r - c2.r) < 0) {
                                                                                     merge(A.begin(), A.end() - 1, B.begin
 Point AB = c2.o - c1.o;
                                             swap(c1, c2);
                                                                                          (), B.end() - 1, C.begin() + 1,
 Tf d = length(AB);
                                             swap(a, b);
                                                                                           polarComp(Point(0, 0), Point(0,
 if (d >= c1.r + c2.r) return 0;
                                                                                               -1))):
 if (d + c1.r <= c2.r) return PI * c1.</pre>
                                           Tf d2 = sqLength(c1.o - c2.o);
                                                                                     for (int i = 1; i < C.size(); i++) C[</pre>
                                                                                         i] = C[i] + C[i - 1];
      r * c1.r;
                                           Tf rdif = c1.r - c2.r, rsum = c1.r +
 if (d + c2.r \le c1.r) return PI * c2.
                                                c2.r;
                                                                                     C.pop_back();
                                           if (dcmp(d2 - rdif * rdif) < 0)</pre>
                                                                                     return C;
      r * c2.r:
                                                return 0;
 Tf alpha1 = acos((c1.r * c1.r + d * d
                                           if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
                                                                                   // finds the rectangle with minimum
       -c2.r*c2.r) / (2.0*c1.r*d
                                                ) == 0) return -1;
                                                                                        area enclosing a convex polygon and
      ));
                                                                                   // the rectangle with minimum perimeter
 Tf alpha2 = acos((c2.r * c2.r + d * d
                                           Tf base = angle(c2.o - c1.o);
                                                                                         enclosing a convex polygon
       -c1.r*c1.r) / (2.0*c2.r*d)
                                           if (dcmp(d2 - rdif * rdif) == 0) {
                                                                                   // Tf Ti Same
                                                                                   pair<Tf, Tf>
      ));
                                             a.push_back(c1.point(base));
 return c1.sector(2 * alpha1) + c2.
                                             b.push_back(c2.point(base));
                                                                                        \verb|rotatingCalipersBoundingBox(const|\\
      sector(2 * alpha2);
                                             cnt++;
                                                                                       Polygon &p) {
                                             return cnt;
                                                                                     using Linear::distancePointLine;
// returns tangents from a point p to
                                                                                     int n = p.size();
                                                                                     int 1 = 1, r = 1, j = 1;
    circle c, Tf Ti same
                                                                                     Tf area = 1e100;
vector<Point> pointCircleTangents(Point
                                           Tf ang = acos((c1.r - c2.r) / sqrt(d2)
     p, Circle c) {
                                                                                     Tf perimeter = 1e100;
 vector<Point> ret;
                                           a.push_back(c1.point(base + ang));
                                                                                     for (int i = 0; i < n; i++) {</pre>
                                                                                       Point v = (p[(i + 1) \% n] - p[i]) /
 Point u = c.o - p;
                                           b.push_back(c2.point(base + ang));
                                                                                           length(p[(i + 1) % n] - p[i]);
 Tf d = length(u);
                                           cnt++;
 if (d < c.r)
                                           a.push_back(c1.point(base - ang));
                                                                                       while (dcmp(dot(v, p[r % n] - p[i])
                                           b.push_back(c2.point(base - ang));
                                                                                           - dot(v, p[(r + 1) % n] - p[i])
 else if (dcmp(d - c.r) == 0) {
                                                                                           ) < 0)
                                           cnt++;
   ret = {rotate(u, PI / 2)};
                                                                                         r++;
 } else {
                                           if (dcmp(d2 - rsum * rsum) == 0) {
                                                                                       while (j < r \mid | dcmp(cross(v, p[j %
   Tf ang = asin(c.r / d);
                                             a.push_back(c1.point(base));
                                                                                           n] - p[i]) -
                                                                                                           cross(v, p[(j +
   ret = {rotate(u, -ang), rotate(u,
                                             b.push_back(c2.point(PI + base));
        ang)};
                                             cnt++;
                                                                                                               1) % n] - p[
                                           } else if (dcmp(d2 - rsum * rsum) >
                                                                                                               i])) < 0)
 return ret;
                                                0) {
                                                                                         j++;
}
                                             Tf ang = acos((c1.r + c2.r) / sqrt(
                                                                                       while (1 < j ||
                                                                                             dcmp(dot(v, p[1 % n] - p[i])
// returns the points on tangents that
                                                 d2)):
    touches the circle, Tf Ti Same
                                                                                                  - dot(v, p[(1 + 1) % n] -
                                             a.push_back(c1.point(base + ang));
vector<Point> pointCircleTangencyPoints
                                             b.push_back(c2.point(PI + base + ang
                                                                                                  p[i])) > 0)
    (Point p, Circle c) {
                                                                                         1++;
 Point u = p - c.o;
                                                                                       Tf w = dot(v, p[r \% n] - p[i]) - dot
                                             cnt++;
 Tf d = length(u);
                                             a.push_back(c1.point(base - ang));
                                                                                           (v, p[1 % n] - p[i]);
 if (d < c.r)
                                             b.push_back(c2.point(PI + base - ang
                                                                                       Tf h = distancePointLine(p[j % n],
   return {};
                                                                                           Line(p[i], p[(i + 1) % n]));
 else if (dcmp(d - c.r) == 0)
                                             cnt++;
                                                                                       area = min(area, w * h);
                                           }
   return {c.o + u};
                                                                                       perimeter = min(perimeter, 2 * w + 2)
 else {
                                           return cnt;
                                                                                            * h);
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
                                                                                     return make_pair(area, perimeter);
                                               Convex
                                         4.4
   return {c.o + rotate(u, -ang), c.o +
        rotate(u, ang));
                                                                                   // returns the left side of polygon u
                                          /// minkowski sum of two polygons in O(
                                                                                        after cutting it by ray a->b
                                                                                   Polygon cutPolygon(Polygon u, Point a,
                                         Polygon minkowskiSum(Polygon A, Polygon
// for two circles c1 && c2, returns
                                                                                        Point b) {
                                               B) {
    two list of points a && b
                                                                                     using Linear::lineLineIntersection;
                                           int n = A.size(), m = B.size();
// such that a[i] is on c1 && b[i] is
                                           rotate(A.begin(), min_element(A.begin
                                                                                     using Linear::onSegment;
    c2 && for every i
                                                (), A.end()), A.end());
// Line(a[i], b[i]) is a tangent to
                                           rotate(B.begin(), min_element(B.begin
                                                                                     Polygon ret;
                                                                                     int n = u.size();
    both circles
                                                (), B.end()), B.end());
// CAUTION: a[i] = b[i] in case they
                                                                                     for (int i = 0; i < n; i++) {</pre>
    touch \mid -1 for c1 = c2
                                                                                       Point c = u[i], d = u[(i + 1) \% n];
                                           A.push_back(A[0]);
                                           B.push_back(B[0]);
```

```
if (dcmp(cross(b - a, c - a)) >= 0)
       ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0)
     Point t;
     lineLineIntersection(a, b - a, c,
          d - c, t);
     if (onSegment(t, Segment(c, d)))
         ret.push_back(t);
   }
 }
 return ret;
\ensuremath{//} returns true if point p is in or on
    triangle abc
bool pointInTriangle(Point a, Point b,
    Point c, Point p) {
 return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >=
      0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no
    three collinear points
// returns inside = -1, on = 0, outside
int pointInConvexPolygon(const Polygon
    &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p -
       pt[0])) > 0)
     lo = mid;
   else
     hi = mid;
 }
 bool in = pointInTriangle(pt[0], pt[
      lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p
       - pt[lo - 1])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p -
      pt[lo])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) %
       n], p - pt[(hi + 1) % n])) == 0)
   return 0;
 return -1;
}
// Extreme Point for a direction is the
     farthest point in that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly,
    Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1)
       % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 +
            n) \% n], u)) >= 0) {
     return c;
```

```
bool a_up = dcmp(dot(poly[(a + 1) %
       n] - poly[a], u)) >= 0;
   bool c_up = dcmp(dot(poly[(c + 1) %
       n] - poly[c], u)) >= 0;
   bool a_above_c = dcmp(dot(poly[a] -
       poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 if (dcmp(dot(poly[a] - poly[(a + 1) % // Calculate [ACW, CW] tangent pair
       n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n | constexpr int CW = -1, ACW = 1;
         ) % n], u)) > 0)
   return a;
 return b % n;
// For a convex polygon p and a line 1,
     returns a list of segments
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i
   ], p[(i + 1) modulo |p|])
// #1 If a segment is collinear with
    the line, only that is returned
// #2 Else if 1 goes through i'th point
    , the i'th segment is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(
    const Polygon &p, Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v))
 int rt = extremePoint(p, rotate90(v)
      * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(l.a, l.b, p[(idx - 1 + n)
        % n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
```

```
for (int i = 0; i < 2; ++i) {</pre>
    int lo = i ? rt : lf;
    int hi = i ? lf : rt;
    int olo = i ? ort : olf;
    while (true) {
      int gap = (hi - lo + n) \% n;
      if (gap < 2) break;</pre>
     int mid = (lo + gap / 2) % n;
     int omid = orient(l.a, l.b, p[mid
          ]);
     if (!omid) {
       lo = mid;
       break;
      if (omid == olo)
       lo = mid;
      else
       hi = mid;
   ret.push_back(lo);
  return ret;
    from an external point
bool isGood(Point u, Point v, Point Q,
    int dir) {
  return orient(Q, u, v) != -dir;
Point better (Point u, Point v, Point Q,
     int dir) {
  return orient(Q, u, v) == dir ? u : v
Point pointPolyTangent(const Polygon &
    pt, Point Q, int dir, int lo, int
    hi) {
  while (hi - lo > 1) {
    int mid = (lo + hi) / 2;
    bool pvs = isGood(pt[mid], pt[mid -
        1], Q, dir);
    bool nxt = isGood(pt[mid], pt[mid +
        1], Q, dir);
    if (pvs && nxt) return pt[mid];
    if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q,
           dir, mid + 1, hi);
     Point p2 = pointPolyTangent(pt, Q,
           dir, lo, mid - 1);
     return better(p1, p2, Q, dir);
    if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) ==
          dir)
       hi = mid - 1;
      else if (better(pt[lo], pt[hi], Q,
           dir) == pt[lo])
       hi = mid - 1;
      else
       lo = mid + 1;
    if (!nxt) {
      if (orient(Q, pt[mid], pt[lo]) ==
```

```
lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++)
      ret = better(ret, pt[i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(
    const Polygon &pt, Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt,
      Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q
      , CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &
    poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0
        && (ret.empty() || b != ret.
       back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p
     of n vertices
Tf signedPolygonArea(const Polygon &p)
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() -</pre>
   ret += cross(p[i] - p[0], p[i + 1] -
        p[0]);
 return ret / 2;
}
// given a polygon p of n vertices,
    generates the convex hull in in CCW
// Tested on https://acm.timus.ru/
    problem.aspx?space=1&num=1185
// Caution: when all points are
    colinear AND removeRedundant ==
    false
// output will be contain duplicate
    points (from upper hull) at back
Polygon convexHull(Polygon p, bool
    removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
```

```
p.erase(unique(p.begin(), p.end()), p
      .end());
  int n = p.size();
  Polygon ch(n + n);
  int m = 0; // preparing lower hull
  for (int i = 0; i < n; i++) {</pre>
    while (m > 1 &&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 1]))
               <= check)
     m--;
    ch[m++] = p[i];
  int k = m; // preparing upper hull
  for (int i = n - 2; i >= 0; i--) {
    while (m > k \&\&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 2]))
               <= check)
    ch[m++] = p[i];
  if (n > 1) m--;
  ch.resize(m);
  return ch;
// returns inside = -1, on = 0, outside
int pointInPolygon(const Polygon &p,
    Point o) {
  using Linear::onSegment;
  int wn = 0, n = p.size();
  for (int i = 0; i < n; i++) {</pre>
    int j = (i + 1) \% n;
    if (onSegment(o, Segment(p[i], p[j])
        ) || o == p[i]) return 0;
    int k = dcmp(cross(p[j] - p[i], o -
        p[i]));
    int d1 = dcmp(p[i].y - o.y);
    int d2 = dcmp(p[j].y - o.y);
    if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn
    if (k < 0 && d2 <= 0 && d1 > 0) wn
  }
 return wn ? -1 : 1;
// Given a simple polygon p, and a line
     1, returns (x, y)
// x = longest segment of 1 in p, y =
    total length of 1 in p.
pair<Tf, Tf> linePolygonIntersection(
    Line 1, const Polygon &p) {
  using Linear::lineLineIntersection;
  int n = p.size();
  vector<pair<Tf, int>> ev;
  for (int i = 0; i < n; ++i) {</pre>
    Point a = p[i], b = p[(i + 1) \% n],
        z = p[(i - 1 + n) \% n];
    int ora = orient(1.a, 1.b, a), orb =
         orient(1.a, 1.b, b),
       orz = orient(l.a, l.b, z);
    if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(
            d, 0);
       // else // Point Touch
```

```
} else if (orz)
     ev.emplace_back(d, orz);
   else if (orb)
     ev.emplace_back(d, orb);
 } else if (ora == -orb) {
   Point ins;
   lineLineIntersection(1, Line(a, b)
        , ins);
   ev.emplace_back(dot(ins - 1.a, 1.b
         - 1.a), 0);
}
sort(ev.begin(), ev.end());
Tf ans = 0, len = 0, last = 0, tot =
    0:
bool active = false;
int sign = 0;
for (auto &qq : ev) {
 int tp = qq.second;
 Tf d = qq.first; /// current Segment
       is (last, d)
                  /// On Border
 if (sign) {
   len += d - last;
   tot += d - last;
   ans = max(ans, len);
   if (tp != sign) active = !active;
   sign = 0;
 } else {
   if (active) { /// Strictly Inside
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
   if (tp == 0)
     active = !active;
   else
     sign = tp;
 last = d;
 if (!active) len = 0;
ans /= length(l.b - l.a);
tot /= length(l.b - l.a);
return {ans, tot};
```

4.6 Point Redifined

```
namespace geo {
   #define x real()
   #define y imag()
   #define setx(i) real(i)
   #define sety(i) imag(i)
   typedef complex<double> point;
   typedef point vec2;
   double dist(point a, point b) {
       return abs(b-a); }
   double dot(vec2 a, vec2 b) { return
        (conj(a) * b).x; }
   double cross(vec2 a, vec2 b) {
       return (conj(a) * b).y; }
   vec2 rotate90(vec2 a) { return point
        (-a.y, a.x); }
   double area(vec2 a, vec2 b) { return
        abs(cross(a,b)); }
   // double area(point a, point b,
       point c) { return area( getVec2
        (a,b), getVec2(a,c)); }
```

5 Graph

5.1 LCA, ETT, VT

```
#define lg(n) (31 - __builtin_clz(n))
const int N = 1e5 + 1;
const int K = lg(N) + 1;
vector<int> adj[N];
int anc[N][K], lvl[N];
namespace lca {
 void init(int u = 1, int p = 0, int d
       = 0) {
   lvl[u] = d;
   anc[u][0] = p;
   for (int i = 1; i < K; i++)</pre>
     anc[u][i] = anc[anc[u][i - 1]][i -
   for (auto v : adj[u])
     if (v != p)
       init(v, u, d + 1);
 int getAnc(int u, int k) {
   for (int i = 0; u and i < K; i++)</pre>
     if ((k >> i) & 1)
       u = anc[u][i];
   return u;
 }
 int lca(int u, int v) {
   if (lvl[u] < lvl[v]) swap(u, v);</pre>
   u = getAnc(u, lvl[u] - lvl[v]);
   if (u == v) return u;
   for (int i = K - 1; ~i; i--)
     if (anc[u][i] != anc[v][i])
       u = anc[u][i], v = anc[v][i];
   return anc[u][0];
 int dist(int u, int v) {
   return lvl[u] + lvl[v] - 2 * lvl[lca
        (u, v)];
 }
};
struct euler_tour {
 int time = 0;
 tree &T;
 int n:
 vector<int> start, finish, level, par
  euler_tour(tree &T, int root = 0)
     : T(T), n(T.n), start(n), finish(n|}
          ), level(n), par(n) {
   time = 0;
   call(root);
 void call(int node, int p = -1) {
   if (p != -1) level[node] = level[p]
        + 1:
   start[node] = time++;
   for (int e : T[node])
     if (e != p) call(e, node);
   par[node] = p;
   finish[node] = time++;
```

```
bool isAncestor(int node, int par) {
   return start[par] <= start[node] and</pre>
         finish[par] >= finish[node];
 int subtreeSize(int node) { return
      finish[node] - start[node] + 1 >> |}
       1: }
tree virtual_tree(vector<int> &nodes,
    lca_table &table, euler_tour &tour)
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
           start[x] < tour.start[y]; });</pre>
 int n = nodes.size();
  for (int i = 0; i + 1 < n; i++)</pre>
   nodes.push_back(table.lca(nodes[i],
        nodes[i + 1]));
 sort(nodes.begin(), nodes.end());
 nodes.erase(unique(nodes.begin(),
      nodes.end()), nodes.end());
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
 n = nodes.size();
 stack<int> st;
 st.push(0);
 tree ans(n);
 for (int i = 1; i < n; i++) {</pre>
   while (!tour.isAncestor(nodes[i],
       nodes[st.top()])) st.pop();
   ans.addEdge(st.top(), i);
   st.push(i);
 return ans;
set<int> getCenters(tree &T) {
 int n = T.n;
 vector<int> deg(n), q;
 set<int> s;
 for (int i = 0; i < n; i++) {</pre>
   deg[i] = T[i].size();
   if (deg[i] == 1) q.push_back(i);
   s.insert(i);
 }
 for (vector<int> t; s.size() > 2; q =
       t) {
   for (auto x : q) {
     for (auto e : T[x])
       if (--deg[e] == 1) t.push_back(e
           );
     s.erase(x);
   }
 return s;
5.2 SCC
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
vector<int> order, component;
// tp = 0 ,finding topo order, tp = 1 ,
     reverse edge traversal
```

void dfs(int u, int tp = 0) {

```
vis[u] = true;
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
 reverse(order.begin(), order.end());
 for (int i : order) {
   if (!vis[i]) {
     // one component is found
     dfs(i, 1), component.clear();
 }
5.3 Euler Tour on Edge
```

10

```
// for simplicity, G[idx] contains the
                                   adjacency list of a node
start[x] < tour.start[y]; }); // while G(e) is a reference to the e-
                                   th edge.
                              const int N = 2e5 + 5;
                              int in[N], out[N], fwd[N], bck[N];
                              int t = 0;
                              void dfs(graph &G, int node, int par) {
                                out[node] = t:
                                for (int e : G[node]) {
                                  int v = G(e).to(node);
                                  if (v == par) continue;
                                  fwd[e] = t++;
                                  dfs(G, v, node);
                                  bck[e] = t++;
                                in[node] = t - 1;
                              void init(graph &G, int node) {
                                dfs(G, node, node);
```

5.4 LCA In O(1)

```
/* LCA in O(1)
 * depth calculates weighted distance
 * level calculates distance by number
     of edges
 * Preprocessing in NlongN */
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG
      [1] = 0;
  for (int i = 2; i <= n; i++) LOG[i] =</pre>
       LOG[i >> 1] + 1;
  for (int i = 0; i <= n; i++) adj[i].</pre>
      clear();
```

```
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
  if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v])</pre>
       return v;
 int 1 = LOG[id[v] - id[u] + 1];
 int p1 = id[u], p2 = id[v] - (1 << 1)</pre>
 int d1 = level[table[1][p1]], d2 =
      level[table[1][p2]];
  if (d1 < d2)
   return par[table[1][p1]];
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time;
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table</pre>
      [0][i] = a[i];
 for (int 1 = 0; 1 < L - 1; 1++) {</pre>
   for (int i = 1; i <= n; i++) {</pre>
     table[1 + 1][i] = table[1][i];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 <<</pre>
           1)]] < level[table[1][i]];</pre>
     if (C1 && C2) table[l + 1][i] =
          table[1][i + (1 << 1)];
   }
 }
}
```

HLD5.5

```
const int N = 1e6 + 7;
template <typename DT>
struct Segtree {
 // write lazy segtree here
```

```
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
  depth[u] = depth[p] + 1;
  for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz
         = subsz;
   SZ += subsz;
 return SZ;
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
  if (heavy[u] != -1) decompose(heavy[u
      ], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
 }
int query(int a, int b) {
 int ret = 0;
  for (; head[a] != head[b]; b = par[
      head[b]]) {
   if (depth[head[a]] > depth[head[b]])
        swap(a, b);
   ret += tree.query(1, 0, cnt, pos[
       head[b]], pos[b]);
  if (depth[a] > depth[b]) swap(a, b);
  ret += tree.query(1, 0, cnt, pos[a],
      pos[b]);
 return ret;
```

5.6 Dinic Max Flow

flow; };

```
/**
Implementation of Dinic's algorithm
    with optional scaling
Source: Chilli (https://codeforces.com/
    blog/entry/66006)
Complexity: O(ans*E) or O(V^2E) without
     scaling, O(VE log(U)) with scaling
Scaling performs much better in worst
    case, but has much higher constant
    factor
To enable scaling, call maxFlow(true)
Everything 0-indexed
*/
namespace Dinic {
   typedef long long LL;
   const int N = 5005, K = 60; /// N >
       no of nodes, K >= max bits in
       capacity
   const LL INF = 1e18;
   struct Edge { int frm, to; LL cap,
```

```
int s, t, n;
int level[N], ptr[N];
vector<Edge> edges;
vector<int> adj[N];
void init(int nodes) {
   n = nodes;
   for (int i=0; i<n; i++) adj[i].</pre>
        clear():
   edges.clear();
}
/// For adding undirected Edge (u, v
    , c) call addEdge(u, v, c, c);
int addEdge(int a, int b, LL cap, LL
     revcap = 0) {
   edges.push_back({a, b, cap, 0});
   edges.push_back({b, a, revcap,
   adj[a].push_back(edges.size()-2)
   adj[b].push_back(edges.size()-1)
   return edges.size()-2;
bool bfs(LL lim) {
   fill(level, level+n, -1);
   level[s] = 0;
   queue<int> q;
   q.push(s);
   while (!q.empty() && level[t] ==
         -1) {
       int v = q.front();
       q.pop();
       for (int id: adj[v]) {
           Edge e = edges[id];
           if (level[e.to] == -1 &&
               e.cap - e.flow >= lim
              q.push(e.to);
              level[e.to] = level[v
                   ] + 1;
           }
       }
   }
   return level[t] != -1;
LL dfs(int v, LL flow) {
   if (v == t || !flow)
                             return
       flow:
   for (; ptr[v] < adj[v].size();</pre>
       ptr[v]++) {
       int eid = adj[v][ptr[v]];
       Edge &e = edges[eid];
       if (level[e.to] != level[v]
           + 1) continue;
       if (LL pushed = dfs(e.to,
           min(flow, e.cap - e.flow
           ))) {
           e.flow += pushed;
           edges[eid^1].flow -=
               pushed;
           return pushed;
       }
   }
   return 0;
LL maxFlow(int source, int sink,
    bool SCALING = false) {
   s = source, t = sink;
```

```
long long flow = 0;
   for (LL lim = SCALING ? (1LL <<</pre>
        K) : 1; \lim > 0; \lim >>= 1)
       while (bfs(lim)) {
           fill(ptr, ptr+n, 0);
           while (LL pushed = dfs(s,
                INF)) flow += pushed
       }
   }
   return flow;
}
bool leftOfMinCut(int x) {return
    level[x] != -1;
/// Only works for undirected graph,
     Make sure to add UNDIRECTED
    edges. (u, v, c, c)
/// returns n by n matrix flow, st
    flow[i][j] = maxFlow
/// tree holds the edges of a gomory
    -hu tree of the graph
vector<vector<LL>> allPairMaxFlow(
    vector<Edge> &tree) {
   tree.clear();
   vector<vector<LL>> flow(n,
        vector<LL> (n, INF));
   vector<int> par(n);
   for (int i=1; i<n; i++) {</pre>
       for (auto &e: edges) e.flow
            = 0;
       LL f = maxFlow(i, par[i]);
       tree.push_back({i, par[i], f
           });
       for (int j=i+1; j<n; j++)</pre>
           if (par[j] == par[i] &&
               leftOfMinCut(j)) par[
               j] = i;
       flow[i][par[i]] = flow[par[i
            ]][i] = f;
       for (int j=0; j<i; j++)</pre>
           if (j != par[i]) flow[i
               ][j] = flow[j][i] =
               min(f, flow[par[i]][j
               ]);
   return flow;
}
```

5.7 Min Cost Max Flow

}

```
mt19937 rnd(chrono::steady_clock::now()
    .time_since_epoch().count());
const LL inf = 1e9;
struct edge {
 int v, rev;
 LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(
         cost), flow(0) {}
};
struct mcmf {
 int src, sink, n;
 vector<int> par, idx, Q;
 vector<bool> inq;
```

```
vector<LL> dis;
vector<vector<edge>> g;
mcmf() {}
mcmf(int src, int sink, int n)
    : src(src),
     sink(sink),
     n(n),
     par(n),
     idx(n),
     inq(n),
     dis(n),
     g(n),
     Q(10000005) {} // use Q(n) if
          not using random
void add_edge(int u, int v, LL cap,
    LL cost, bool directed = true) {
  edge _u = edge(v, g[v].size(), cap,
      cost);
  edge _v = edge(u, g[u].size(), 0, -
      cost);
  g[u].pb(_u);
  g[v].pb(_v);
  if (!directed) add_edge(v, u, cap,
      cost, true);
bool spfa() {
  for (int i = 0; i < n; i++) {</pre>
   dis[i] = inf, inq[i] = false;
  int f = 0, 1 = 0;
  dis[src] = 0, par[src] = -1, Q[1++]
      = src, inq[src] = true;
  while (f < 1) {</pre>
    int u = Q[f++];
    for (int i = 0; i < g[u].size(); i</pre>
        ++) {
     edge &e = g[u][i];
     if (e.cap <= e.flow) continue;</pre>
     if (dis[e.v] > dis[u] + e.cost)
       dis[e.v] = dis[u] + e.cost;
       par[e.v] = u, idx[e.v] = i;
       if (!inq[e.v]) inq[e.v] = true
            , Q[1++] = e.v;
       // if (!inq[e.v]) {
       // inq[e.v] = true;
       // if (f && rnd() & 7) Q[--f]
             = e.v;
       // else Q[1++] = e.v;
       // }
     }
    inq[u] = false;
 return (dis[sink] != inf);
pair<LL, LL> solve() {
  LL mincost = 0, maxflow = 0;
  while (spfa()) {
    LL bottleneck = inf;
    for (int u = par[sink], v = idx[
        sink]; u != -1; v = idx[u], u
        = par[u]) {
     edge &e = g[u][v];
     bottleneck = min(bottleneck, e.
          cap - e.flow);
   }
   for (int u = par[sink], v = idx[
        sink; u != -1; v = idx[u], u |_{void init(graph &G)} {
```

```
= par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck
     }
     mincost += bottleneck * dis[sink],
          maxflow += bottleneck;
   return make_pair(mincost, maxflow);
 }
// want to minimize cost and don't care
     about flow
// add edge from sink to dummy sink (
    cap = inf, cost = 0)
// add edge from source to sink (cap =
    inf, cost = 0)
// run mcmf, cost returned is the
    minimum cost
```

```
Bridge Tree
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G,
    int par = -1, int d = 0) {
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1)
       if (low[to] > depth[node]) {
        bridges.emplace_back(node, to)
        components.push_back({});
         for (int x = -1; x != to; x =
             st.top(), st.pop())
          components.back().push_back(
               st.top());
       }
     low[node] = min(low[node], low[to
         ]);
   }
 }
 if (par == -1) {
   components.push_back({});
   while (!st.empty()) components.back
        ().push_back(st.top()), st.pop
 }
graph &create_tree() {
 for (auto &comp : components) {
   int idx = tree.addNode();
   for (auto &e : comp) id[e] = idx;
 for (auto &[l, r] : bridges) tree.
      addEdge(id[1], id[r]);
 return tree;
```

```
int n = G.n;
depth.assign(n, -1), id.assign(n, -1)
     , low.resize(n);
for (int i = 0; i < n; i++)</pre>
 if (depth[i] == -1) find_bridges(i,
```

5.9 Tree Isomorphism

```
LL Hash(int u, int p) {
 vector<LL> childrenHash;
 for (auto v : adj[u]) if (v != p)
   childrenHash.add(Hash(v, u));
 sort(all(childrenHash));
 LL nodeHash = 0;
 for (int i = 0; i < childrenHash.size</pre>
      (); i++)
   nodeHash = (nodeHash + childrenHash[
        i] * bigmod(SEED, i, MOD)) %
       MOD;
 return nodeHash;
}
```

6 Math

```
6.1 Linear Sieve
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5],
     cnt[N + 5], POW[N + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] =
      1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   for (auto p : primes) {
     if (p * i > N \text{ or } p > spf[i]) break
     prime[p * i] = false, spf[p * i] = |map<LL, int> factorize(LL n) {
     if (i % p == 0) {
       phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] +
           1) * (cnt[i] + 2),
              cnt[p * i] = cnt[i] + 1;
       SOD[p * i] = SOD[i] / SOD[POW[i]
           ]] * (SOD[POW[i]] + p * POW[
           i]),
              POW[p * i] = p * POW[i];
       break:
     } else {
       phi[p * i] = phi[p] * phi[i];
       NOD[p * i] = NOD[p] * NOD[i],
           cnt[p * i] = 1;
       SOD[p * i] = SOD[p] * SOD[i],
           POW[p * i] = p;
     }
   }
 }
```

6.2 Pollard Rho

LL mul(LL a, LL b, LL mod) {

```
return (__int128)a * b % mod;
 // LL ans = a * b - mod * (LL) (1.L / L)
       mod * a * b);
  // return ans + mod * (ans < 0) - mod
       * (ans >= (LL) mod);
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(
      num, num, mod))
   if (pow & 1) ans = mul(ans, num, mod
        );
 return ans;
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (</pre>
      n | 1) == 3;
 LL a[] = \{2, 325, 9375, 28178,
      450775, 9780504, 1795265022};
 LL s = \__builtin\_ctzll(n - 1), d = n
      >> s;
 for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x %
        n and i--; p = mul(p, p, n)
   if (p != n - 1 and i != s) return
        false;
 return true;
LL get_factor(LL n) {
 auto f = [\&](LL x) \{ return mul(x, x,
       n) + 1; };
 LL x = 0, y = 0, t = 0, prod = 2, i =
 for (; t++ % 40 or gcd(prod, n) == 1;
       x = f(x), y = f(f(y))) {
    (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) -
       min(x, y), n) ? q : prod;
 return gcd(prod, n);
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11,
      13, 17, 19, 23, 29, 31, 37, 41,
                     43, 47, 53, 59,
                         61, 67, 71,
                         73, 79, 83,
                         89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
  auto _factor = [&](LL n, auto &
      factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
```

```
}
};
_factor(n, _factor);
return res;
```

6.3 Chinese Remainder Theorem

13

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
 if (b == 0)
   return {1, 0, a};
  else {
   auto [x, y, g] = EGCD(b, a \% b);
   return {y, x - a / b * y, g};
// given modulo equations, will apply
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1;
  for (auto &[v, m] : v) { // value %
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1,
   V += x * (v - V) / g % (m / g) * M,
       M *= m / g;
   V = (V \% M + M) \% M;
 return make_pair(V, M);
```

6.4 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
 memset(mob, -1, sizeof mob);
 mob[1] = 1;
 for (int i = 2; i < N; i++)</pre>
   if (mob[i]) {
     for (int j = i + i; j < N; j += i)
           mob[j] -= mob[i];
```

6.5 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
  assert((n & (n - 1)) == 0), N = n;
  perm = vector<int>(N, 0);
  for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
  wp[0] = wp[1] = vector < CD > (N);
  for (int i = 0; i < N; i++) {</pre>
    wp[0][i] = CD(cos(2 * PI * i / N),
        sin(2 * PI * i / N));
    wp[1][i] = CD(cos(2 * PI * i / N), -
        sin(2 * PI * i / N));
```

```
}
                                           vector<LL> anyMod(const vector<LL> &a,
                                                                                       for (LL res = 2; res <= p; res++) {</pre>
}
                                               const vector<LL> &b) {
                                                                                          bool ok = true;
void fft(vector<CD> &v, bool invert =
                                             int n = 1;
                                                                                          for (LL i = 0; i < factor.size() &&</pre>
                                             while (n < a.size() + b.size()) n <<=</pre>
    false) {
                                                                                              ok; i++)
 if (v.size() != perm.size())
                                                                                           ok &= Pow(res, phi / factor[i]) !=
                                                  1:
      precalculate(v.size());
                                             vector<CD> al(n), ar(n), bl(n), br(n)
                                                                                                 1;
 for (int i = 0; i < N; i++)</pre>
                                                                                          if (ok) return res;
   if (i < perm[i]) swap(v[i], v[perm[i</pre>
                                             for (int i = 0; i < a.size(); i++) al</pre>
                                                                                        }
                                                                                        return -1;
        ]]);
                                                 [i] = a[i] % M / B, ar[i] = a[i]
 for (int len = 2; len <= N; len *= 2)</pre>
                                                 % M % B;
                                             for (int i = 0; i < b.size(); i++) bl</pre>
                                                                                      void prepare(LL n) {
   for (int i = 0, d = N / len; i < N;</pre>
                                                 [i] = b[i] % M / B, br[i] = b[i]
                                                                                        LL sz = abs(31 - \_builtin\_clz(n));
        i += len) {
                                                 % M % B;
                                                                                        LL r = Pow(g, (MOD - 1) / n);
     for (int j = 0, idx = 0; j < len /
                                                                                        inv_n = Pow(n, MOD - 2);
                                             pairfft(al, ar);
                                             pairfft(bl, br);
           2; j++, idx += d) {
                                                                                        w[0] = w[n] = 1;
       CD x = v[i + j];
                                                      fft(al); fft(ar); fft(bl);
                                                                                        for (LL i = 1; i < n; i++) w[i] = (w[</pre>
                                             //
       CD y = wp[invert][idx] * v[i + j]
                                                 fft(br);
                                                                                            i - 1] * r) % MOD;
                                             for (int i = 0; i < n; i++) {</pre>
                                                                                        for (LL i = 1; i < n; i++)</pre>
             + len / 2];
       v[i + j] = x + y;
                                              CD 11 = (al[i] * bl[i]), lr = (al[i]
                                                                                          rev[i] = (rev[i >> 1] >> 1) | ((i &
       v[i + j + len / 2] = x - y;
                                                                                              1) << (sz - 1);
                                                    * br[i]);
     }
                                              CD rl = (ar[i] * bl[i]), rr = (ar[i]|}
   }
                                                                                      void NTT(LL *a, LL n, LL dir = 0) {
                                                    * br[i]);
 }
                                              al[i] = 11;
                                                                                        for (LL i = 1; i < n - 1; i++)</pre>
                                                                                          if (i < rev[i]) swap(a[i], a[rev[i</pre>
 if (invert) {
                                              ar[i] = lr;
                                              bl[i] = rl;
   for (int i = 0; i < N; i++) v[i] /=</pre>
                                                                                              ]]);
                                              br[i] = rr;
                                                                                        for (LL m = 2; m <= n; m <<= 1) {</pre>
                                                                                          for (LL i = 0; i < n; i += m) {</pre>
                                             pairfft(al, ar, true);
                                                                                           for (LL j = 0; j < (m >> 1); j++)
void pairfft(vector<CD> &a, vector<CD>
                                             pairfft(bl, br, true);
    &b, bool invert = false) {
                                                      fft(al, true); fft(ar, true)
                                                                                             LL \&u = a[i + j], \&v = a[i + j +
 int N = a.size();
                                                 ; fft(bl, true); fft(br, true);
                                                                                                   (m >> 1)];
 vector<CD> p(N);
                                             vector<LL> ans(n);
                                                                                             LL t = v * w[dir ? n - n / m * j]
 for (int i = 0; i < N; i++) p[i] = a[</pre>
                                             for (int i = 0; i < n; i++) {</pre>
                                                                                                   : n / m * j] % MOD;
                                              LL right = round(br[i].real()), left
                                                                                             v = u - t < 0 ? u - t + MOD : u
      i] + b[i] * CD(0, 1);
                                                    = round(al[i].real());
 fft(p, invert);
                                                                                                  - t;
 p.push_back(p[0]);
                                                                                             u = u + t >= MOD ? u + t - MOD :
 for (int i = 0; i < N; i++) {</pre>
                                              LL mid = round(round(bl[i].real()) +
                                                                                                   u + t;
   if (invert) {
                                                    round(ar[i].real()));
     a[i] = CD(p[i].real(), 0);
                                              ans[i] = ((left \% M) * B * B + (mid)
                                                                                         }
     b[i] = CD(p[i].imag(), 0);
                                                   % M) * B + right) % M;
                                                                                        }
                                                                                        if (dir)
   } else {
     a[i] = (p[i] + conj(p[N - i])) *
                                            return ans;
                                                                                          for (LL i = 0; i < n; i++) a[i] = (
                                                                                              inv_n * a[i]) % MOD;
          CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) *
                                           6.6 NTT
          CD(0, -0.5);
                                                                                      vector<LL> mul(vector<LL> p, vector<LL>
   }
                                           const LL N = 1 << 18;
                                                                                        LL n = p.size(), m = q.size();
 }
                                           const LL MOD = 786433;
}
                                                                                        LL t = n + m - 1, sz = 1;
vector<LL> multiply(const vector<LL> &a |vector<LL> P[N];
                                                                                        while (sz < t) sz <<= 1;</pre>
    , const vector<LL> &b) {
                                                                                        prepare(sz);
                                           LL rev[N], w[N | 1], a[N], b[N], inv_n,
 int n = 1;
 while (n < a.size() + b.size()) n <<= |LL Pow(LL b, LL p)| {
                                                                                        for (LL i = 0; i < n; i++) a[i] = p[i
       1;
                                                                                            ];
                                            LL ret = 1:
 vector<CD> fa(a.begin(), a.end()), fb
                                                                                        for (LL i = 0; i < m; i++) b[i] = q[i</pre>
                                             while (p) {
      (b.begin(), b.end());
                                              if (p & 1) ret = (ret * b) % MOD;
                                                                                            ];
 fa.resize(n);
                                                                                        for (LL i = n; i < sz; i++) a[i] = 0;</pre>
                                              b = (b * b) % MOD;
 fb.resize(n);
                                                                                        for (LL i = m; i < sz; i++) b[i] = 0;</pre>
                                              p >>= 1;
           fft(fa); fft(fb);
 //
                                            }
 pairfft(fa, fb);
                                                                                        NTT(a, sz);
                                            return ret;
 for (int i = 0; i < n; i++) fa[i] =</pre>
                                                                                        NTT(b, sz);
      fa[i] * fb[i];
                                                                                        for (LL i = 0; i < sz; i++) a[i] = (a
                                          LL primitive_root(LL p) {
 fft(fa, true);
                                                                                            [i] * b[i]) % MOD;
                                             vector<LL> factor;
                                             LL phi = p - 1, n = phi;
 vector<LL> ans(n);
                                                                                        NTT(a, sz, 1);
  for (int i = 0; i < n; i++) ans[i] =</pre>
                                             for (LL i = 2; i * i <= n; i++) {</pre>
      round(fa[i].real());
                                                                                        vector<LL> c(a, a + sz);
                                              if (n % i) continue;
                                                                                        while (c.size() && c.back() == 0) c.
 return ans;
                                              factor.emplace_back(i);
}
                                                                                            pop_back();
                                              while (n \% i == 0) n /= i;
const int M = 1e9 + 7, B = sqrt(M) + 1;
                                                                                        return c;
                                            if (n > 1) factor.emplace_back(n);
```

}

};

String

7.1 **Aho Corasick**

```
const int sg = 26, N = 1e3 + 9;
struct aho_corasick {
 struct node {
   node *link, *out, *par;
   bool leaf:
   LL val;
   int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '
       $', int len = 0)
       : par(par), p_ch(p_ch), len(len)
     val = leaf = cnt = last = 0;
     link = out = NULL;
 };
 vector<node> trie;
 node* root:
 aho_corasick() {
   trie.reserve(N), trie.emplace_back()
   root = &trie[0];
   root->link = root->out = root;
 inline int f(char c) { return c - 'a'
      ; }
 inline node* add_node(node* par =
      NULL, char p_ch = '$', int len =
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
 void add_str(const string& s, LL val
      = 1) {
   node* now = root;
   for (char c : s) {
     int i = f(c);
     if (!now->to[i]) now->to[i] =
     now = now -> to[i];
   now->leaf = true, now->val++;
 }
 void push_links() {
   queue<node*> q;
   for (q.push(root); q.empty(); q.pop
        ()) {
     node *cur = q.front(), *link = cur |}
          ->link;
     cur->out = link->leaf ? link :
         link->out;
     int idx = 0;
     for (auto& next : cur->to) {
       if (next != NULL) {
         next->link = cur != root ?
             link->to[idx++] : root;
         q.push(next);
       } else
         next = link->to[idx++];
   cur->val += link->val;
```

```
7.2 Double hash
                               // define +, -, * for (PLL, LL) and (
                                   PLL, PLL), % for (PLL, PLL);
                               PLL base(1949313259, 1997293877);
                               PLL mod(2091573227, 2117566807);
                               PLL power(PLL a, LL p) {
                                PLL ans = PLL(1, 1);
                                 for(; p; p >>= 1, a = a * a % mod) {
                                    if(p \& 1) ans = ans * a % mod;
                                 return ans;
                               PLL inverse(PLL a) { return power(a, (
                                   mod.ff - 1) * (mod.ss - 1) - 1); }
                               PLL inv_base = inverse(base);
                               PLL val;
                               vector<PLL> P;
                               void hash_init(int n) {
                                P.resize(n + 1);
                                P[0] = PLL(1, 1);
                                for (int i = 1; i <= n; i++) P[i] = (
                                     P[i - 1] * base) % mod;
                              PLL append(PLL cur, char c) { return (
                                   cur * base + c) % mod; }
                               /// prepends c to string with size k
                              PLL prepend(PLL cur, int k, char c) {
                                   return (P[k] * c + cur) % mod; }
                               /// replaces the i-th (0-indexed)
                                   character from right from a to b;
                               PLL replace(PLL cur, int i, char a,
                                   char b) {
                                 cur = (cur + P[i] * (b - a)) \% mod;
                                return (cur + mod) % mod;
                               /// Erases c from the back of the
                                   string
add_node(now, c, now->len + 1) PLL pop_back(PLL hash, char c) {
                                return (((hash - c) * inv_base) % mod
                                      + mod) % mod;
                               /// Erases c from front of the string
                                   with size len
                               PLL pop_front(PLL hash, int len, char c
                                   ) {
                                return ((hash - P[len - 1] * c) % mod
                                      + mod) % mod;
                               /// concatenates two strings where
                                   length of the right is k
                               PLL concat(PLL left, PLL right, int k)
                                   { return (left * P[k] + right) %
                                   mod; }
                               /// Calculates hash of string with size
                                    len repeated cnt times
                               /// This is O(\log n). For O(1), pre-
                                   calculate inverses
                               PLL repeat(PLL hash, int len, LL cnt) {
                                PLL mul = (P[len * cnt] - 1) *
                                     inverse(P[len] - 1);
                                 mul = (mul % mod + mod) % mod;
                                PLL ret = (hash * mul) % mod;
```

```
if (P[len].ff == 1) ret.ff = hash.ff
      * cnt:
  if (P[len].ss == 1) ret.ss = hash.ss
      * cnt:
 return ret;
LL get(PLL hash) { return ((hash.ff <<
    32) ^ hash.ss); }
struct hashlist {
 int len:
  vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.
        resize(len + 2, PLL(0, 0));
   for (int i = 1; i <= len; i++) H[i]</pre>
        = append(H[i - 1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i]
        = append(R[i + 1], s[i - 1]);
  /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[1 - 1] * P[r - 1 +
         1]) % mod + mod) % mod;
 PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 +
         1]) % mod + mod) % mod;
 PLL concat_range_hash(int 11, int r1,
       int 12, int r2) {
   return concat(range_hash(l1, r1),
        range_hash(12, r2), r2 - 12 +
        1);
  PLL concat_reverse_hash(int 11, int
      r1, int 12, int r2) {
   return concat(reverse_hash(12, r2),
        reverse_hash(l1, r1), r1 - l1 +
         1);
 }
};
```

7.3 $_{ m KMP}$

```
vector<int> prefix_function(string s) {
 int n = (int)s.length();
 vector<int> pi(n);
 for (int i = 1; i < n; i++) {</pre>
   int j = pi[i-1];
   while (j > 0 \&\& s[i] != s[j])
     j = pi[j-1];
   if (s[i] == s[j])
     j++;
   pi[i] = j;
 return pi;
```

Manacher's

```
vector<int> d1(n);
// d[i] = number of palindromes taking
    s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 1 : min(d1[1 + r -
      i], r - i + 1);
```

```
while (0 \le i - k \&\& i + k \le n \&\& s[i]
       - k] == s[i + k]) k++;
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking
    s[i-1] and s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 0 : min(d2[1 + r -
      i + 1], r - i + 1);
 while (0 \le i - k - 1 \&\& i + k \le n \&\&
       s[i - k - 1] == s[i + k]) k++;
 d2[i] = k--;
  if (i + k > r) l = i - k - 1, r = i +
       k;
```

7.5 String Match FFT

```
//find occurrences of t in s where '?'s
     are automatically matched with any
//res[i + m - 1] = sum_j = 0 to m - 1_{s[}
    i + j] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
 vector\langle int \rangle s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[</pre>
      i] == '?' ? 0 : s[i] - 'a' + 1;
      //assign any non zero number for
      non '?'s
 for(int i = 0; i < n; i++) s2[i] = s1
      [i] * s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1</pre>
      [i] * s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[</pre>
      i] == '?' ? 0 : t[i] - 'a' + 1;
 for(int i = 0; i < m; i++) t2[i] = t1</pre>
      [i] * t1[i];
 for(int i = 0; i < m; i++) t3[i] = t1</pre>
      [i] * t2[i];
 reverse(t1.begin(), t1.end());
 reverse(t2.begin(), t2.end());
 reverse(t3.begin(), t3.end());
 vector<int> s1t3 = multiply(s1, t3);
 vector<int> s2t2 = multiply(s2, t2);
 vector<int> s3t1 = multiply(s3, t1);
 vector<int> res(n);
 for(int i = 0; i < n; i++) res[i] =</pre>
      s1t3[i] - s2t2[i] * 2 + s3t1[i];
 vector<int> oc;
  for(int i = m - 1; i < n; i++) if(res</pre>
      [i] == 0) oc.push_back(i - m + 1)
 return oc;
}
```

7.6 Suffix Array

```
#include <bits/stdc++.h>
using namespace std;
 O(|S| + |alphabet|) Suffix Array
 LIM := max{s[i]} + 2
```

```
if ((!sl[a] and sl[a - 1]) or (!sl
void inducedSort (const vector <int> &
                                                    [b] and sl[b - 1])) {
                                                 flag = !(!sl[a] and sl[a - 1]
    vec, int val_range, vector <int> &
    SA, const vector <int> &sl, const
                                                      and !sl[b] and sl[b - 1]);
    vector <int> &lms_idx) {
                                                      break:
 vector <int> l(val_range, 0), r(
                                               }
      val_range, 0);
 for (int c : vec) {
                                             SA[j] = flag ? ++cur : cur;
   ++r[c]; if (c + 1 < val_range) ++l[c
                                           }
         + 1]:
                                           for (int i = 0; i < lms_idx.size();</pre>
                                                ++i) lms_vec[i] = SA[lms_idx[i]];
 partial_sum(1.begin(), 1.end(), 1.
                                            if (cur + 1 < lms_idx.size()) {</pre>
                                             auto lms_SA = suffixArray(lms_vec,
      begin());
                                                  cur + 1);
 partial_sum(r.begin(), r.end(), r.
                                             for (int i = 0; i < lms_idx.size();</pre>
      begin()):
 fill(SA.begin(), SA.end(), -1);
                                                  ++i) new_lms_idx[i] = lms_idx[
 for (int i = lms_idx.size() - 1; i >=
                                                  lms_SA[i]];
       0; --i) SA[--r[vec[lms_idx[i]]]]
       = lms_idx[i];
                                            inducedSort(vec, val_range, SA, s1,
 for (int i : SA) if (i > 0 and sl[i -
                                                new_lms_idx); return SA;
       1]) SA[l[vec[i-1]]++] = i-1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
                                          vector <int> getSuffixArray (const
                                              string &s, const int LIM = 128) {
 partial_sum(r.begin(), r.end(), r.
                                            vector <int> vec(s.size() + 1);
      begin());
 for (int k = SA.size() - 1, i = SA[k
                                            copy(begin(s), end(s), begin(vec));
      ]; k; --k, i = SA[k]) {
                                                vec.back() = '$';
   if (i and !sl[i - 1]) SA[--r[vec[i -
                                           auto ret = suffixArray(vec, LIM);
         1]]] = i - 1;
                                           ret.erase(ret.begin()); return ret;
                                          // build RMQ on it to get LCP of any
vector <int> suffixArray (const vector
                                              two suffix
                                          vector <int> getLCParray (const string
    <int> &vec, int val_range) {
 const int n = vec.size();
                                              &s, const vector <int> &SA) {
 vector <int> sl(n), SA(n), lms_idx;
                                            int n = s.size(), k = 0;
 for (int i = n - 2; i \ge 0; --i) {
                                            vector <int> lcp(n), rank(n);
   sl[i] = vec[i] > vec[i + 1] or (vec[i])
                                            for (int i = 0; i < n; ++i) rank[SA[i</pre>
       i] == vec[i + 1] and sl[i + 1])
                                                ]] = i;
                                            for (int i = 0; i < n; ++i, k ? --k :</pre>
   if (sl[i] and !sl[i + 1]) lms_idx.
                                                 0) {
                                             if (rank[i] == n - 1) {
        emplace_back(i + 1);
                                               k = 0; continue;
 reverse(lms_idx.begin(), lms_idx.end
                                             int j = SA[rank[i] + 1];
     ());
 inducedSort(vec, val_range, SA, s1,
                                             while (i + k < n \text{ and } j + k < n \text{ and } s
                                                  [i + k] == s[j + k]) ++k;
     lms idx):
 vector <int> new_lms_idx(lms_idx.size
                                             lcp[rank[i]] = k;
      ()), lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {</pre>
                                           lcp[n - 1] = 0; return lcp;
   if (SA[i] > 0 and !sl[SA[i]] and sl[|}
       SA[i] - 1] new_lms_idx[k++] =
       SA[i];
                                          int main() {
                                           string s; cin >> s;
 int cur = 0; SA[n - 1] = 0;
                                            for (const int i : getSuffixArray(s))
                                                 printf("%d ", i);
 for (int k = 1; k < new_lms_idx.size</pre>
                                           puts("");
      (); ++k) {
   int i = new_lms_idx[k - 1], j =
                                           return 0;
       new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
                                              \operatorname{Trie}
     SA[j] = ++cur; continue;
                                          template<int sz>
   bool flag = 0;
                                          struct Trie {
   for (int a = i + 1, b = j + 1; ; ++a
                                           Trie() : id(1) {
        , ++b) {
                                             memset(endMark, 0, sizeof endMark);
     if (vec[a] ^ vec[b]) {
                                             for_each(all(trie), [](vector<int> &
       flag = 1; break;
                                                  v) { v.assign(sz, 0); });
```

```
void insert(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false;
     cur = trie[cur][val];
   }
   return endMark[cur];
 }
private:
 int id, endMark[100005];
 vector<int> trie[100005];
```

7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     1 = r = i:
     while (r < n \&\& s[r] == s[r - 1])
         r++:
     z[i] = r - 1, r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k];
       1 = i;
       while (r < n \&\& s[r] == s[r - 1]
           ]) r++;
       z[i] = r - 1, r--;
   }
 }
 return z;
```

Extra

8.1 Make File

```
.PHONY: clean
D ?= 0
ifeq ($(D), 1)
CXXFLAGS=-std=c++17 -g -DLOCAL -Wall -
     Wextra -Wpedantic -Wshadow -
     Wformat=2 -Wfloat-equal -
     Wconversion -Wlogical-op -Wshift-
     overflow=2 -Wduplicated-cond -
     Wcast-qual -Wcast-align -Wno-
     variadic-macros -D_GLIBCXX_DEBUG -
     D_GLIBCXX_DEBUG_PEDANTIC -
     fsanitize=address -fsanitize=
     undefined -fno-sanitize-recover -
     fstack-protector -fsanitize-
```

```
address-use-after-scope
else
CXXFLAGS=-02 -std=c++17 -DLOCAL
endif
all:
Qecho 'Error: No file given.'
@echo 'Cleaning executables...'
Ofind . ! -name stress.sh -type f -
     executable -delete
@echo 'Executables cleaned.'
8.2 Stress Tester
make $1
make $2
make $3
# $1 is actual code
# $2 is good code
# $3 is generator
for ((i = 1;; i++)); do
 echo 'Test #'$i
 timeout 5s ./$3 $RANDOM > in
 timeout 5s./\$1 < in > out
 timeout 5s./$2 < in > ans
 diff ans out || break
```

8.3 Sublime Build

Auto import & Compile

```
"shell_cmd": "make $file_base_name &&
     timeout 5s ./$file_base_name <</pre>
    in > out",
"working_dir": "$file_path",
"selector": "source.cpp"
```

:autocmd BufNewFile *.cpp Or ~/template

8.4 vimrc

```
nnoremap <F4> :!xclip -o -sel clip > ~/
    cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip
    > ~/cp/in <CR><CR>
nnoremap <F6> :!xclip -sel clip % <CR><
   CR>
inoremap <F6> <ESC>:!xclip -sel clip %
    <CR><CR>
autocmd filetype cpp nnoremap <F9>
    wa \| !make %:r && timeout 5s ./%:r set nocompatible
     < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F9> <ESC
    >:wa \| !make %:r && timeout 5s
    ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp nnoremap <F10> :
    wa \l !make clean && make %:r D=1
    && ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F10> <
    ESC>:wa \| !make clean && make %:r
    D=1 \&\& ./%:r < ^/cp/in > ^/cp/out <
    CR>
```

```
autocmd filetype python nnoremap <F9> :
    wa \| !python % < ~^/cp/in> ~^/cp/out
autocmd filetype python inoremap <F9> <
    ESC>:wa \| !python % < ~/cp/in> ~/
    cp/out<CR>
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'),
     col('.')-1, 1) == ")" ? "\<Right>"
     : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'),
     col('.')-1, 1) == "}" ? "\<Right>"
     : "}"
inoremap [ []<left>
inoremap <expr> ] strpart(getline('.'),
     col('.')-1, 1) == "]" ? "\<Right>"
     : "]"
inoremap <expr> " strpart(getline('.'),
     col('.')-1, 1) == "\"" ? "\<Right>
    " : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'),
     col('.')-1, 1) == "\'" ? "\<Right
    >" : "\'\\'\<left>"
inoremap <expr> <CR> <sid>
    insert newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col(
      '.')-2, 2)
  return stridx('{}[]', pair) % 2 == 0
      && strlen(pair) == 2 ? "\<CR>\<
      ESC>\0" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>
    insert_space()
function s:insert_space() abort
  let pair = strpart(getline('.'), col(
      '.')-2, 2)
  return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<space
      >\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
 let pair = strpart(getline('.'), col('
     .')-2, 2)
 return stridx('(){}[]''', pair) %
     2 == 0 && strlen(pair) == 2 ? "\<
     del>\<bs>" : "\<bs>"
endfunction
                            " be
    iMproved, required
filetype on
                           " required
filetype plugin on
filetype plugin indent on
syntax on
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
```

set expandtab

```
18
```

```
set softtabstop=2
set smartindent
set smarttab
set autoindent
set cindent
set noerrorbells
set ruler
set guifont=*
set backspace=indent,eol,start
set ignorecase
set incsearch
set nowrap
set hlsearch
" bubt site
" set termguicolors
set foldmethod=indent
set nofoldenable
set cursorline
set laststatus=2
set showcmd
set wildmenu
if !has('nvim')
 set clipboard=unnamedplus
if !has('nvim')
 set ttymouse=xterm2
\verb"endif"
nnoremap <S-j> :m .+1<CR>==
nnoremap \langle S-k \rangle : m .-2 \langle CR \rangle ==
vnoremap <S-j> :m '>+1<CR>gv==gv
```

Equations and Formulas

9.1 Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1$$

and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

the size n+1 factors.

intersecting chords.

n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

Number of permutations of $1, \ldots, n$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n=3, these permutations are 132, 213, 231, 312 and 321.

9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count If $P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$, then, permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$
where, $S(0,0) = 1, S(n,0) = S(0,n) = 0$

$$\sum_{k=0}^{n} S(n,k) = n!$$

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k) \text{ , then,}$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 i_k.$$

9.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$ where S(0,0) = 1, S(n,0) = S(0,n) = 0 $S(n,2) = 2^{n-1} - 1 S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets,

with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = \sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$ $kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1) \qquad \sum_{k=1}^n x^{\gcd(k,n)} = \sum_{d|n} x^d \cdot \phi\left(\frac{n}{d}\right)$ The number of ways to completely parenthesize n+1 factors. Stirling numbers of the second kind, delayed $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$ Denote the n objects to partition by the noted $S^d(n,k)$, to be the number of ways $\phi(d)$ The number of triangulations of a convex to partition the integers 1, 2, ..., n into k polygon with n+2 sides (i.e. the number of nonempty subsets such that all elements in nonempty subsets such that all elements in $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right)$ partitions of polygon into disjoint triangles each subset have pairwise distance at least by using the diagonals). The number of ways to connect the 2n subset, it is required that $|i-j| \ge d$. It $\frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$ points on a circle to form n disjoint i.e. non-has been shown that these numbers satisfy, The number of rooted full binary trees with 9.4 Other Combinatorial Identities

 $\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n}$ $n, r \in N, n > r, \sum_{i=1}^{n} {i \choose r} = {n+1 \choose r+1}$

If
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

f
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

Different Math Formulas

Picks Theorem: A = i + b/2 - 1 $d(i) = (i-1) \times$ Deragements: (d(i-1) + d(i-2))

$$\frac{n}{ab} \quad - \quad \left\{\frac{b\prime n}{a}\right\} \quad - \quad \left\{\frac{a\prime n}{b}\right\} \quad + \quad 1$$

9.6 GCD and LCM

if m is any integer, then $gcd(a + m \cdot b, b) =$ gcd(a,b)

The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$ $\gcd(a_1,b)\cdot\gcd(a_2,b).$

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).For non-negative integers a and b, where aand b are not both zero, $gcd(n^a-1, n^b-1) =$ $n^{\gcd(a,b)} - 1$

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\sum_{i=1}^{n} [\gcd(i, n) = k] = \phi\left(\frac{n}{k}\right)$$

with each subset containing at least
$$r$$
 elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$ Denote the n objects to partition by the integers $1,2,\ldots,n$. Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers $1,2,\ldots,n$ into k nonempty subsets such that all elements in each subset have pairwise distance at least d . That is, for any integers i and j in a given subset, it is required that $|i-j| \ge d$. It has been shown that these numbers satisfy, $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ 0 . A Other Combinatorial Identities $\binom{n}{k} = \frac{n}{k}\binom{n+i}{k-1}$ $\binom{n}{k} = \frac{n}{k}\binom{n+i}{k-1}$ $\binom{n}{k} = \binom{n}{k} \cdot Q(k)$, then, $\binom{n}{k} = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$, then, $\binom{n}{k} = \binom{n}{k} \cdot P(k)$

19