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All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :20000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,
    ssse3,sse4,popcnt,abm,mmx,avx,tune=
    native")
#include <ext/pb_ds/assoc_container.hpp</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns
        iterator to the kth largest
        element counting from 0
   //order_of_key(val) --> returns the
        number of items in a set that
        are strictly smaller than our
template <typename DT>
using ordered_set = tree <DT, null_type</pre>
    , less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef LOCAL
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\',' << x <<</pre>
     '\'';}
void show(const string &x) {cerr << '\"</pre>
    ' << x << '\"';}
void show(bool x) {cerr << (x ? "true"</pre>
    : "false");}
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f';</pre>
    show(x.first); cerr << ", "; show(x | LT lazy[4 * N];
    .second); cerr << '}'; }
template<typename T>
void show(T x) {int f = 0; cerr << "{";</pre>
     for (auto &i: x) cerr << (f++ ? ",</pre>
     " : ""), show(i); cerr << "}";}
void debug_out(string s) {
   cerr << '\n';</pre>
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(),
         '), s.end());
   cerr << "
                  "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl:
   if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " <<</pre>
    __LINE__ << endl; debug_out(#x, x);
     cerr << endl;</pre>
#else
#define dbg(x...)
```

```
#endif
const int RANDOM = chrono::
    high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) ^ x;
unsigned hash_combine(unsigned a,
    unsigned b) { return a * 31 + b; }
struct chash {
 int operator()(int x) const { return
     hash_f(x); }
typedef gp_hash_table<int, int, chash>
    gp;
gp table;
mt19937_64 rng(atoi(argv[1]));
long long random(long long 1, long long
    r) {
   uniform_int_distribution<long long>
       dist(1, r);
   return dist(rng);
2
    Data Structure
```

2.1 Segment Tree

```
const int N = 1000006;
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N];
int L, R;
void pull(int s, int e, int node) {
 val[node] = val[node << 1] + val[node</pre>
       << 1 | 1];
void apply(const LT &U, int s, int e,
    int node) {
 val[node] += (e - s + 1) * U;
 lazy[node] += U;
void reset(int node) { lazy[node] =
    None; }
DT merge(const DT &a, const DT &b) {
   return a + b; }
DT get(int s, int e, int node) { return |}
     val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
 apply(lazy[node], s, s + e >> 1, node
       << 1):
 apply(lazy[node], s + e + 2 \gg 1, e,
     node << 1 | 1);
 reset(node);
void build(int s, int e, vector<DT> &v,
     int node = 1) {
 int m = s + e >> 1;
```

```
if (s == e) {
    val[node] = v[s];
   return;
  build(s, m, v, node * 2);
  build(m + 1, e, v, node * 2 + 1);
  pull(s, e, node);
void update(int S, int E, LT uval, int
    s = L, int e = R, int node = 1) {
  if (S > E) return;
  if (S == s and E == e) \{
    apply(uval, s, e, node);
    return;
  push(s, e, node);
  int m = s + e >> 1;
  update(S, min(m, E), uval, s, m, node
       * 2);
  update(max(S, m + 1), E, uval, m + 1,
       e, node *2 + 1);
  pull(s, e, node);
DT query(int S, int E, int s = L, int e
     = R, int node = 1) {
  if (S > E) return I;
  if (s == S and e == E) return get(s,
      e, node);
  push(s, e, node);
  int m = s + e >> 1;
  DT L = query(S, min(m, E), s, m, node
       * 2);
  DT R = query(max(S, m + 1), E, m + 1,
       e, node * 2 + 1);
  return merge(L, R);
void init(int _L, int _R, vector<DT> &v
  L = L, R = R;
  build(L, R, v);
```

2.2Persistent Segment Tree

```
struct Node {
   int 1 = 0, r = 0, val = 0;
tr[20 * N];
int ptr = 0;
int build(int st, int en) {
   int u = ++ptr;
   if (st == en) return u;
   int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   1 = build(st, mid);
   r = build(mid + 1, en);
   val = tr[l].val + tr[r].val;
   return u;
int update(int pre, int st, int en, int
     idx, int v) {
   int u = ++ptr;
   tr[u] = tr[pre];
   if (st == en) {
       tr[u].val += v;
       return u;
   int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   if (idx <= mid) {</pre>
```

```
r = tr[pre].r;
       1 = update(tr[pre].1, st, mid,
           idx, v);
   } else {
       1 = tr[pre].1;
       r = update(tr[pre].r, mid + 1,
           en, idx, v);
   tr[u].val = tr[l].val + tr[r].val;
   return u;
// finding the kth elelment in a range
int query(int left, int right, int st,
    int en, int k) {
   if (st == en) return st;
   int cnt = tr[tr[right].1].val - tr[
        tr[left].1].val;
   int mid = (st + en) / 2;
   if (cnt >= k) return query(tr[left].
        1, tr[right].1, st, mid, k);
   else return query(tr[left].r, tr[
       right].r, mid + 1, en, k - cnt)
int V[N], root[N], a[N];
int main() {
   map<int, int> mp; int n, q;
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >>
       a[i], mp[a[i]];
   int c = 0;
   for (auto x : mp) mp[x.first] = ++c,
        V[c] = x.first;
   root[0] = build(1, n);
   for (int i = 1; i <= n; i++) {</pre>
       root[i] = update(root[i - 1], 1,
            n, mp[a[i]], 1);
   while (q--) {
       int 1, r, k; cin >> 1 >> r >> k;
            1++, k++;
       cout << V[query(root[l - 1],</pre>
           root[r], 1, n, k)] << '\n';
   }
   return 0;
```

2.3 Lazy Persistent Segment Tree

```
struct node {
   int val;
   int lazy;
   int left, right;
} nodes[MAXN];
int sz = 0;
int a[MAXN];
int build(int s, int e) {
   int curr = sz;
   sz++;
   if(s == e) {
       nodes[curr].left = nodes[curr].
           right = -1;
       nodes[curr].val = a[s];
       nodes[curr].lazy = 0;
       return curr;
   int m = (s + e) / 2;
   nodes[curr].left = build(s, m);
   nodes[curr].right = build(m + 1, e);
```

```
nodes[curr].val = nodes[nodes[curr].
        left].val + nodes[nodes[curr].
        right].val;
   return curr;
int update(int prev, int s, int e, int
    1, int r, int val) {
    if(s > r or e < 1) return prev;</pre>
   int curr = sz;
   sz++;
   if(s \ge 1 \text{ and } e \le r) {
       nodes[curr].left = nodes[prev].
            left;
       nodes[curr].right = nodes[prev].
            right;
       nodes[curr].val = nodes[prev].
            val + val*(e-s+1);
       nodes[curr].lazy = nodes[prev].
            lazy + val;
       return curr;
   int m = (s + e) / 2;
   nodes[curr].left = update(nodes[prev
        ].left, s, m, l, r, val);
   nodes[curr].right = update(nodes[
        prev].right, m+1, e, l, r, val)
   nodes[curr].lazy = nodes[prev].lazy;
   nodes[curr].val = nodes[nodes[curr].
        left].val + nodes[nodes[curr].
        right].val + nodes[curr].lazy*(
        e-s+1);
   return curr;
11 query(int nd, int s, int e, int 1,
    int r, 11 c) {
    if(s > r or e < 1) return 0;</pre>
    if(s >= 1 and e <= r) return c*(e-s</pre>
        +1) + nodes[nd].val;
   int m = (s+e)/2;
   c += nodes[nd].lazy;
   return query(nodes[nd].left, s, m, 1
        , r, c) + query(nodes[nd].right
        , m+1, e, l, r, c);
```

2.4 Implicit Segment Tree

```
struct node {
   ll val;
   ll lazy;
   int left, right;
ነ:
node nodes[4*1000006];
int sz;
void reset(int nd) {
   nodes[nd].left = nodes[nd].right =
   nodes[nd].val = nodes[nd].lazy = 0;
int getLeft(int nd) {
   if(nodes[nd].left == -1) {
       nodes[nd].left = sz++;
       reset(nodes[nd].left);
   return nodes[nd].left;
int getRight(int nd) {
    if(nodes[nd].right == -1) {
       nodes[nd].right = sz++;
```

```
reset(nodes[nd].right);
   }
   return nodes[nd].right;
void update(int nd, int s, int e, int l
    , int r, ll val) {
   if(e < 1 or s > r) return;
   if(s \ge 1 and e \le r) {
       nodes[nd].val += val*(e-s + 1);
       nodes[nd].lazy += val;
       return;
   }
   int m = (s + e) / 2;
   update(getLeft(nd), s, m, l ,r, val)
   update(getRight(nd), m+1, e, 1, r,
   nodes[nd].val = nodes[nodes[nd].left
        ].val + nodes[nodes[nd].right].
        val + nodes[nd].lazy * (e-s+1);
ll query(int nd, int s, int e, int l,
    int r, ll c) {
   if(s > r or e < 1) return 0;</pre>
   if(s \ge 1 \text{ and } e \le r) {
       return nodes[nd].val + c*(e-s+1)
   int m = (s+e)/2;
   c += nodes[nd].lazy;
   return query(getLeft(nd), s, m, 1, r
        , c) + query(getRight(nd), m+1,
         e, 1, r, c);
```

DSU With Rollbacks 2.5

```
struct Rollback_DSU {
 int n;
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n,
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par
       [node])
     ; // no path compression
   return node;
 void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)])
        swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back())
     par[op.back().second] = op.back().
         second;
     sz[op.back().first] -= sz[op.back
         ().second];
 }
```

```
2.6 BIT-2D
```

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j \leftarrow -j & j)
       bit[x][j] += val;
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j \& -j) ans
       += bit[x][j];
 return ans;
}
int get(int x1, int y1, int x2, int y2)
 return get(x2, y2) - get(x1 - 1, y2)
      - get(x2, y1 - 1) + get(x1 - 1,
      y1 - 1);
```

2.7 Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];
vector<LL> merge(vector<LL> v1, vector<
    LL> v2) {
 LL i = 0, j = 0;
 vector<LL> ret;
  while (i < v1.size() || j < v2.size()</pre>
      ) {
   if (i == v1.size()) {
     ret.push_back(v2[j]);
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
     i++;
   } else {
     if (v1[i] < v2[j]) {</pre>
       ret.push_back(v1[i]);
       i++;
     } else {
       ret.push_back(v2[j]);
       j++;
   }
 return ret;
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 }
 LL leftNode = 2 * node, rightNode = 2
       * node + 1:
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
```

```
Tree[node] = merge(Tree[leftNode],
      Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1,
    LL r, LL k) {
 if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin
        (), Tree[node].end(), k) -
          Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, 1, r,
        query(rightNode, mid + 1, ed, 1
            , r, k);
```

2.8 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs =
int arr[N], freq[2 * N], cnt[2 * N], id
    [N], ans[N];
struct query {
 int 1, r, t, L, R;
 query(int l = 1, int r = 0, int t =
      1, int id = -1)
     : 1(1), r(r), t(t), L(1 / sz), R(r)
           / sz) {}
  bool operator<(const query &rhs)</pre>
      const {
   return (L < rhs.L) or (L == rhs.L
       and R < rhs.R) or
          (L == rhs.L and R == rhs.R
              and t < rhs.t);</pre>
 }
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
 cnt[freq[arr[idx]]]++;
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[ int A[MAXN][MAXN];</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
```

```
arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0);
int main() {
  int n, q;
  cin >> n >> q;
  int counter = 0;
  map<int, int> M;
  for (int i = 1; i <= n; i++) {</pre>
    cin >> arr[i];
    if (!M[arr[i]]) M[arr[i]] = ++
        counter:
    arr[i] = M[arr[i]];
  iota(id, id + N, 0);
  while (q--) {
    int tp, x, y;
    cin >> tp >> x >> y;
    if (tp == 1)
     Q[++qi] = query(x, y, ui);
    else {
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = \{x, y, arr[x]\};
     arr[x] = y;
  }
  t = ui;
  cnt[0] = 3 * n;
  sort(id + 1, id + qi + 1, [&](int x,
      int y) { return Q[x] < Q[y]; });
  for (int i = 1; i <= qi; i++) {</pre>
    int x = id[i];
    while (Q[x].t > t) apply(++t);
    while (Q[x].t < t) undo(t--);
    while (Q[x].1 < 1) add(--1);
    while (Q[x].r > r) add(++r);
    while (Q[x].1 > 1) remove(1++);
    while (Q[x].r < r) remove(r--);
   ans[x] = mex();
 for (int i = 1; i <= qi; i++) cout <<</pre>
       ans[i] << '\n';
```

2.9 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
   for (int q = 1; (1 << q) <= N; q++)
```

```
int add = 1 << (q - 1);
     for (int j = 1; j + add <= N; j++)</pre>
       M[i][j][0][q] = max(M[i][j][0][q]
             - 1], M[i][j + add][0][q -
            1]);
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q
          ++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p]
              - 1][q], M[i + add][j][p -
               1] [q]);
     }
   }
 }
// returns max of all A[i][j], where x1
    =i<=x2 and y1<=j<=y2
int Query(int x1, int y1, int x2, int
    y2) {
  int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 \ll kX;
 int addY = 1 << kY;</pre>
 int ret1 = max(M[x1][y1][kX][kY], M[
      x1][y2 - addY + 1][kX][kY]);
  int ret2 = max(M[x2 - addX + 1][y1][
      kX][kY],
                M[x2 - addX + 1][y2 -
                    addY + 1][kX][kY]);
 return max(ret1, ret2);
```

2.10 Sparse Table

```
// O-based indexing, query finds in
    range [first, last]
#define lg(x) (31 - __builtin_clz(x))
const int N = 1e5 + 1;
const int K = lg(N);
int a[N], tr[N][K + 1];
namespace sparse_table {
 int f(int p1, int p2) { return min(p1
      , p2); }
 void build() {
   for(int i = 0; i < n; i++)</pre>
     tr[i][0] = a[i];
   for(int j = 1; j <= K; j++) {</pre>
     for(int i = 0; i + (1<<j) <= n; i</pre>
          ++)
       tr[i][j] = f(tr[i][j - 1], tr[i
            + (1<<(j - 1))][j - 1]);
   }
 }
 int query(int 1, int r) {
   int d = lg(r - l + 1);
   return f(table[1][d], table[r - (1<<</pre>
        d) + 1][d]);
 }
}
```

DP

3.1 Convex Hull Trick

```
struct line {
    11 m, c;
    line() {}
    line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
    vector<line> lines;
    int ptr = 0;
    convex_hull_trick() {}
    bool bad(line a, line b, line c) {
         return 1.0 * (c.c - a.c) * (a.m - b.
                    m) < 1.0 * (b.c - a.c) * (a.m - a.c) * (a.c) * (
                      c.m);
    }
    void add(line L) {
         int sz = lines.size();
         while (sz \ge 2 \&\& bad(lines[sz - 2],
                        lines[sz - 1], L)) {
              lines.pop_back();
         lines.pb(L);
    11 get(int idx, int x) { return (111
                * lines[idx].m * x + lines[idx].c
                ); }
    11 query(int x) {
         if (lines.empty()) return 0;
         if (ptr >= lines.size()) ptr = lines
                     .size() - 1;
          while (ptr < lines.size() - 1 && get
                     (ptr, x) > get(ptr + 1, x)) ptr
                     ++;
         return get(ptr, x);
   }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
    fastio;
    int t;
    cin >> t;
    while (t--) {
         int n, a, b, c;
         cin >> n >> a >> b >> c;
         for (int i = 1; i <= n; i++) cin >>
                     sum[i];
         for (int i = 1; i <= n; i++) dp[i] =</pre>
                       0, sum[i] += sum[i - 1];
         convex_hull_trick cht;
          cht.add(line(0, 0));
          for (int pos = 1; pos <= n; pos++) {</pre>
              dp[pos] = cht.query(sum[pos]) - 1
                         11 * a * sqr(sum[pos]) - c;
              cht.add(line(211 * a * sum[pos],
                          dp[pos] - a * sqr(sum[pos])));
         ll ans = (-111 * dp[n]);
         ans += (111 * sum[n] * b);
         cout << ans << "\n";
              Dynamic CHT
3.2
```

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 IS_QUERY = -(1LL << 62);</pre>
struct line {
  11 m. b:
  mutable function <const line*()> succ
  bool operator < (const line &rhs)</pre>
      const {
    if (rhs.b != IS_QUERY) return m <</pre>
        rhs.m;
    const line *s = succ();
    if (!s) return 0;
    11 x = rhs.m;
    return b - s -> b < (s -> m - m) * x
 }
};
struct HullDynamic : public multiset <</pre>
    line> {
  bool bad (iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y -> m == z -> m && y -> b
          \langle = z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x
        -> m && y -> b <= x -> b;
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z
        -> m - y -> m) >= 1.0 * (y -> b)
         -z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 }
  void insert_line (ll m, ll b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) ==
         end() ? 0 : &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(
        y))) erase(next(y));
    while (y != begin() && bad(prev(y)))
          erase(prev(y));
  11 eval (ll x) {
    auto 1 = *lower_bound((line) {x,
        IS_QUERY});
    return 1.m * x + 1.b;
 }
int main() {
  HullDynamic hull;
  hull.insert_line(1, 1);
  hull.insert_line(-1, 1);
  cout << hull.eval(69) << endl;</pre>
  cout << hull.eval(420) << endl;</pre>
  return 0;
```

3.3 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c)
LL calc(line L, LL x) { return 1LL * L.
    m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft =
      NULL, node *rt = NULL)
     : L(line(m, c)), lft(lft), rt(rt)
};
struct LiChao {
 node *root;
 LiChao() { root = new node(); }
 void update(node *now, int L, int R,
      line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(
        lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
           node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo;
     if (now->lft == NULL) now->lft =
         new node():
     update(now->lft, L, mid, hi);
 }
 LL query(node *now, int L, int R, LL
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x <= mid)</pre>
     return max(calc(now->L, x), query(
         now->lft, L, mid, x));
   else
     return max(calc(now->L, x), query(
          now->rt, mid + 1, R, x));
 }
};
```

Geometry

```
4.1 Point
typedef double Tf;
typedef double Ti; /// use long long
    for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ?</pre>
    0 : (x < 0 ? -1 : 1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y)
```

```
Point operator+(const Point& u) const
                  { return Point(x + u.x, y + u.y) | Point normal(Point a) {
                ; }
     Point operator-(const Point& u) const
                  { return Point(x - u.x, y - u.y) |}
     Point operator*(const LL u) const {
                return Point(x * u, y * u); }
     Point operator*(const Tf u) const {
                return Point(x * u, y * u); }
     Point operator/(const Tf u) const {
                return Point(x / u, y / u); }
     bool operator==(const Point& u) const
          return dcmp(x - u.x) == 0 \&\& dcmp(y)
                     -u.y) == 0;
      bool operator!=(const Point& u) const
                   { return !(*this == u); }
     bool operator<(const Point& u) const</pre>
          return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                    -u.x) == 0 && dcmp(y - u.y) <
                    0);
 Ti dot(Point a, Point b) { return a.x *
              b.x + a.y * b.y; }
 Ti cross(Point a, Point b) { return a.x |};
              * b.y - a.y * b.x; }
 Tf length(Point a) { return sqrt(dot(a,
              a)): }
 Ti sqLength(Point a) { return dot(a, a)
            ; }
 Tf distance(Point a, Point b) { return
           length(a - b); }
 Tf angle(Point u) { return atan2(u.y, u
            .x); }
  // returns angle between oa, ob in (-PI
            , PI]
 Tf angleBetween(Point a, Point b) {
     Tf ans = angle(b) - angle(a);
     return ans <= -PI ? ans + 2 * PI : (
                ans > PI ? ans - 2 * PI : ans);
  // Rotate a ccw by rad radians, Tf Ti
 Point rotate(Point a, Tf rad) {
     return Point(a.x * cos(rad) - a.y *
                sin(rad),
                                   a.x * sin(rad) + a.y * cos
                                              (rad));
 // rotate a ccw by angle th with cos(th
           ) = co \&\& sin(th) = si, tf ti same
 Point rotatePrecise(Point a, Tf co, Tf
      return Point(a.x * co - a.y * si, a.y
                  * co + a.x * si);
 Point rotate90(Point a) { return Point
            (-a.y, a.x); }
 // scales vector a by s such that
           length of a becomes s, Tf Ti same
Point scale(Point a, Tf s) { return a /
```

length(a) * s; }

```
// returns an unit vector perpendicular
     to vector a, Tf Ti same
 Tf 1 = length(a);
 return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if
    on ab && -1 if right of ab
int orient(Point a, Point b, Point c) {
     return dcmp(cross(b - a, c - a));
/// Use as sort(v.begin(), v.end(),
    polarComp(0, dir))
/// Polar comparator around O starting
    at direction dir
struct polarComp {
 Point O, dir;
 polarComp(Point 0 = Point(0, 0),
      Point dir = Point(1, 0): O(0),
      dir(dir) {}
 bool half(Point p) {
   return dcmp(cross(dir, p)) < 0 ||</pre>
          (dcmp(cross(dir, p)) == 0 \&\&
              dcmp(dot(dir, p)) > 0);
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) <</pre>
        make_tuple(half(q), cross(p, q)
struct Segment {
 Point a, b;
 Segment(Point aa, Point bb) : a(aa),
     b(bb) {}
typedef Segment Line;
struct Circle {
 Point o;
 Tf r;
 Circle(Point o = Point(0, 0), Tf r =
      0) : o(o), r(r) {}
  // returns true if point p is in ||
      on the circle
 bool contains(Point p) { return dcmp(
      sqLength(p - o) - r * r) \le 0;
  // returns a point on the circle rad
     radians away from +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value
   return Point(o.x + cos(rad) * r, o.y
        + sin(rad) * r);
 // area of a circular sector with
      central angle rad
 Tf area(Tf rad = PI + PI) { return
      rad * r * r / 2; }
  // area of the circular sector cut by
       a chord with central angle alpha
 Tf sector(Tf alpha) { return r * r *
      0.5 * (alpha - sin(alpha)); }
```

4.2 Linear

```
// **** LINE LINE INTERSECTION START
// returns true if point p is on
    segment s
```

```
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p))
      == 0 && dcmp(dot(s.a - p, s.b - p
// returns true if segment p && q touch |}
     or intersect
bool segmentsIntersect(Segment p,
    Segment q) {
 if (onSegment(p.a, q) || onSegment(p.
      b, q)) return true;
 if (onSegment(q.a, p) || onSegment(q.
      b, p)) return true;
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 &&
      dcmp(c3) * dcmp(c4) < 0;
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.
      a)) == 0;
// lines are represented as a ray from
    a point: (point, vector)
// returns false if two lines (p, v) &&
     (q, w) are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference, Tf Ti
     Same
bool lineLineIntersection(Point p,
    Point v, Point q, Point w, Point& o
 if (dcmp(cross(v, w)) == 0) return
      false:
 Point u = p - q;
 o = p + v * (cross(w, u) / cross(v, w)
      ));
 return true;
// returns false if two lines p && q
    are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference
bool lineLineIntersection(Line p, Line
    q, Point& o) {
 return lineLineIntersection(p.a, p.b
      - p.a, q.a, q.b - q.a, o);
// returns the distance from point a to |}
     line 1
// **** LINE LINE INTERSECTION FINISH
    ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(1.b - 1.a, p - 1.a)
      / length(1.b - 1.a));
// returns the shortest distance from
    point a to segment s
Tf distancePointSegment(Point p,
    Segment s) {
 if (s.a == s.b) return length(p - s.a
      );
 Point v1 = s.b - s.a, v2 = p - s.a,
      v3 = p - s.b;
 if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
```

```
else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
   return abs(cross(v1, v2) / length(v1
        )):
// returns the shortest distance from
    segment p to segment q
Tf distanceSegmentSegment(Segment p,
    Segment q) {
  if (segmentsIntersect(p, q)) return
      0;
 Tf ans = distancePointSegment(p.a, q)
 ans = min(ans, distancePointSegment(p
      .b, q));
 ans = min(ans, distancePointSegment(q
      .a, p));
  ans = min(ans, distancePointSegment(q
      .b, p));
 return ans;
// returns the projection of point p on
     line 1, Tf Ti Same
Point projectPointLine(Point p, Line 1)
 Point v = 1.b - 1.a;
 return 1.a + v * ((Tf)dot(v, p - 1.a)
       / dot(v, v));
4.3 Circular
// Extremely inaccurate for finding
    near touches
// compute intersection of line 1 with
    circle c
// The intersections are given in order
     of the ray (l.a, l.b), Tf Ti same
vector<Point> circleLineIntersection(
    Circle c, Line 1) {
 vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B *
       B - A * C;
  if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D
      + EPS)) / A);
  if (D > EPS) ret.push_back(1.a + b *
      (-B + sqrt(D)) / A);
 return ret;
// signed area of intersection of
    circle(c.o, c.r) &&
 // triangle(c.o, s.a, s.b) [cross(a-o,
    b-o)/21
Tf circleTriangleIntersectionArea(
    Circle c, Segment s) {
 using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
 if (dcmp(distancePointSegment(c.o, s)
       - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b -
         c.o) * (c.r * c.r) / 2.0;
  // triangle
  if (dcmp(OA - c.r) <= 0 && dcmp(OB -
      c.r) \ll 0
```

```
return cross(c.o - s.b, s.a - s.b) /
         2.0:
 // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect =
      circleLineIntersection(c, s);
 return circleTriangleIntersectionArea
      (c, Segment(s.a, Sect[0])) +
        circleTriangleIntersectionArea(
            c, Segment(Sect[0], Sect
            [1])) +
        circleTriangleIntersectionArea(
            c, Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c
    .r) && simple polyson(p[])
Tf circlePolyIntersectionArea(Circle c,
     Polygon p) {
 Tf res = 0;
 int n = p.size();
 for (int i = 0; i < n; ++i)</pre>
        circleTriangleIntersectionArea(
        c, Segment(p[i], p[(i + 1) % n
        ]));
 return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
    ----> -2
// interior tangents (d = R - r)
    ----> -1
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r)
    ) ----> 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
int circleCirclePosition(Circle c1,
    Circle c2) {
 Tf d = length(c1.o - c2.o);
 int in = dcmp(d - abs(c1.r - c2.r)),
      ex = dcmp(d - (c1.r + c2.r));
 return in < 0 ? -2 : in == 0 ? -1 :
      ex == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points
    between two circles c1 && c2, Tf Ti
vector<Point> circleCircleIntersection(
    Circle c1, Circle c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o);
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return
       ret;
 if (dcmp(abs(c1.r - c2.r) - d) > 0)
      return ret;
 Point v = c2.o - c1.o;
 Tf co = (c1.r * c1.r + sqLength(v) -
      c2.r * c2.r) / (2 * c1.r * length
      (v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co,
       -si), c1.r) + c1.o;
 Point p2 = scale(rotatePrecise(v, co,
       si), c1.r) + c1.o;
```

```
int circleCircleTangencyPoints(Circle
                                                                                     for (int i = 0; i < n; i++) A[i] = A[</pre>
 ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
                                              c1, Circle c2, vector<Point> &a,
                                                                                          i + 1] - A[i];
                                                                                     for (int i = 0; i < m; i++) B[i] = B[</pre>
                                                                      vector<
 return ret;
                                                                                          i + 1] - B[i];
                                                                           Point>
// intersection area between two
                                                                           &b) {
                                           a.clear(), b.clear();
                                                                                     Polygon C(n + m + 1);
    circles c1, c2
Tf circleCircleIntersectionArea(Circle
                                           int cnt = 0;
                                                                                     C[0] = A.back() + B.back();
    c1, Circle c2) {
                                           if (dcmp(c1.r - c2.r) < 0) {
                                                                                     merge(A.begin(), A.end() - 1, B.begin
 Point AB = c2.o - c1.o;
                                             swap(c1, c2);
                                                                                          (), B.end() - 1, C.begin() + 1,
 Tf d = length(AB);
                                             swap(a, b);
                                                                                           polarComp(Point(0, 0), Point(0,
 if (d >= c1.r + c2.r) return 0;
                                                                                               -1))):
 if (d + c1.r <= c2.r) return PI * c1.</pre>
                                           Tf d2 = sqLength(c1.o - c2.o);
                                                                                     for (int i = 1; i < C.size(); i++) C[</pre>
                                                                                          i] = C[i] + C[i - 1];
      r * c1.r;
                                           Tf rdif = c1.r - c2.r, rsum = c1.r +
 if (d + c2.r <= c1.r) return PI * c2.</pre>
                                                c2.r;
                                                                                     C.pop_back();
                                           if (dcmp(d2 - rdif * rdif) < 0)</pre>
                                                                                     return C;
      r * c2.r:
                                                return 0;
 Tf alpha1 = acos((c1.r * c1.r + d * d
                                           if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
                                                                                   // finds the rectangle with minimum
       -c2.r*c2.r) / (2.0*c1.r*d
                                                ) == 0) return -1;
                                                                                        area enclosing a convex polygon and
      ));
                                                                                   // the rectangle with minimum perimeter
 Tf alpha2 = acos((c2.r * c2.r + d * d
                                           Tf base = angle(c2.o - c1.o);
                                                                                         enclosing a convex polygon
       -c1.r*c1.r) / (2.0*c2.r*d)
                                           if (dcmp(d2 - rdif * rdif) == 0) {
                                                                                   // Tf Ti Same
                                                                                   pair<Tf, Tf>
      ));
                                             a.push_back(c1.point(base));
 return c1.sector(2 * alpha1) + c2.
                                             b.push_back(c2.point(base));
                                                                                        \verb|rotatingCalipersBoundingBox(const|\\
      sector(2 * alpha2);
                                             cnt++;
                                                                                        Polygon &p) {
                                             return cnt;
                                                                                     using Linear::distancePointLine;
// returns tangents from a point p to
                                                                                     int n = p.size();
                                                                                     int 1 = 1, r = 1, j = 1;
    circle c, Tf Ti same
                                                                                     Tf area = 1e100;
vector<Point> pointCircleTangents(Point
                                           Tf ang = acos((c1.r - c2.r) / sqrt(d2)
     p, Circle c) {
                                                                                     Tf perimeter = 1e100;
 vector<Point> ret;
                                           a.push_back(c1.point(base + ang));
                                                                                     for (int i = 0; i < n; i++) {</pre>
                                                                                       Point v = (p[(i + 1) \% n] - p[i]) /
 Point u = c.o - p;
                                           b.push_back(c2.point(base + ang));
                                                                                           length(p[(i + 1) % n] - p[i]);
 Tf d = length(u);
                                           cnt++;
 if (d < c.r)
                                           a.push_back(c1.point(base - ang));
                                                                                       while (dcmp(dot(v, p[r % n] - p[i])
                                           b.push_back(c2.point(base - ang));
                                                                                           - dot(v, p[(r + 1) % n] - p[i])
 else if (dcmp(d - c.r) == 0) {
                                                                                           ) < 0)
                                           cnt++;
   ret = {rotate(u, PI / 2)};
                                                                                         r++;
 } else {
                                           if (dcmp(d2 - rsum * rsum) == 0) {
                                                                                       while (j < r \mid | dcmp(cross(v, p[j %
   Tf ang = asin(c.r / d);
                                             a.push_back(c1.point(base));
                                                                                           n] - p[i]) -
                                                                                                           cross(v, p[(j +
   ret = {rotate(u, -ang), rotate(u,
                                             b.push_back(c2.point(PI + base));
        ang)};
                                             cnt++;
                                                                                                               1) % n] - p[
                                           } else if (dcmp(d2 - rsum * rsum) >
                                                                                                               i])) < 0)
 return ret;
                                                0) {
                                                                                         j++;
}
                                             Tf ang = acos((c1.r + c2.r) / sqrt(
                                                                                       while (1 < j ||
                                                                                             dcmp(dot(v, p[1 % n] - p[i])
// returns the points on tangents that
                                                 d2)):
    touches the circle, Tf Ti Same
                                                                                                  - dot(v, p[(1 + 1) % n] -
                                             a.push_back(c1.point(base + ang));
vector<Point> pointCircleTangencyPoints
                                             b.push_back(c2.point(PI + base + ang
                                                                                                   p[i])) > 0)
    (Point p, Circle c) {
                                                                                         1++;
 Point u = p - c.o;
                                                                                       Tf w = dot(v, p[r \% n] - p[i]) - dot
                                             cnt++;
 Tf d = length(u);
                                             a.push_back(c1.point(base - ang));
                                                                                           (v, p[1 % n] - p[i]);
 if (d < c.r)
                                             b.push_back(c2.point(PI + base - ang
                                                                                       Tf h = distancePointLine(p[j % n],
   return {};
                                                                                           Line(p[i], p[(i + 1) % n]));
 else if (dcmp(d - c.r) == 0)
                                             cnt++;
                                                                                       area = min(area, w * h);
                                           }
   return {c.o + u};
                                                                                       perimeter = min(perimeter, 2 * w + 2)
 else {
                                           return cnt;
                                                                                            * h);
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
                                                                                     return make_pair(area, perimeter);
                                               Convex
                                         4.4
   return {c.o + rotate(u, -ang), c.o +
        rotate(u, ang));
                                                                                   // returns the left side of polygon u
                                          /// minkowski sum of two polygons in O(
                                                                                        after cutting it by ray a->b
                                                                                   Polygon cutPolygon(Polygon u, Point a,
                                         Polygon minkowskiSum(Polygon A, Polygon
// for two circles c1 && c2, returns
                                                                                        Point b) {
                                               B) {
    two list of points a && b
                                                                                     using Linear::lineLineIntersection;
                                           int n = A.size(), m = B.size();
// such that a[i] is on c1 && b[i] is
                                           rotate(A.begin(), min_element(A.begin
                                                                                     using Linear::onSegment;
    c2 && for every i
                                                (), A.end()), A.end());
// Line(a[i], b[i]) is a tangent to
                                           rotate(B.begin(), min_element(B.begin
                                                                                     Polygon ret;
                                                                                     int n = u.size();
    both circles
                                                (), B.end()), B.end());
// CAUTION: a[i] = b[i] in case they
                                                                                     for (int i = 0; i < n; i++) {</pre>
    touch \mid -1 for c1 = c2
                                                                                       Point c = u[i], d = u[(i + 1) \% n];
                                           A.push_back(A[0]);
                                           B.push_back(B[0]);
```

```
if (dcmp(cross(b - a, c - a)) >= 0)
       ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0)
     Point t;
     lineLineIntersection(a, b - a, c,
          d - c, t);
     if (onSegment(t, Segment(c, d)))
         ret.push_back(t);
   }
 }
 return ret;
\ensuremath{//} returns true if point p is in or on
    triangle abc
bool pointInTriangle(Point a, Point b,
    Point c, Point p) {
 return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >=
      0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no
    three collinear points
// returns inside = -1, on = 0, outside
int pointInConvexPolygon(const Polygon
    &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p -
       pt[0])) > 0)
     lo = mid;
   else
     hi = mid;
 }
 bool in = pointInTriangle(pt[0], pt[
      lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p
       - pt[lo - 1])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p -
      pt[lo])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) %
       n], p - pt[(hi + 1) % n])) == 0)
   return 0;
 return -1;
}
// Extreme Point for a direction is the
     farthest point in that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly,
    Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1)
       % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 +
            n) \% n], u)) >= 0) {
     return c;
```

```
bool a_up = dcmp(dot(poly[(a + 1) %
       n] - poly[a], u)) >= 0;
   bool c_up = dcmp(dot(poly[(c + 1) %
       n] - poly[c], u)) >= 0;
   bool a_above_c = dcmp(dot(poly[a] -
       poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 if (dcmp(dot(poly[a] - poly[(a + 1) % // Calculate [ACW, CW] tangent pair
       n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n | constexpr int CW = -1, ACW = 1;
         ) % n], u)) > 0)
   return a;
 return b % n;
// For a convex polygon p and a line 1,
     returns a list of segments
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i
   ], p[(i + 1) modulo |p|])
// #1 If a segment is collinear with
    the line, only that is returned
// #2 Else if 1 goes through i'th point
    , the i'th segment is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(
    const Polygon &p, Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v))
 int rt = extremePoint(p, rotate90(v)
      * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(l.a, l.b, p[(idx - 1 + n)
        % n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
```

```
for (int i = 0; i < 2; ++i) {</pre>
    int lo = i ? rt : lf;
    int hi = i ? lf : rt;
    int olo = i ? ort : olf;
    while (true) {
      int gap = (hi - lo + n) \% n;
      if (gap < 2) break;</pre>
     int mid = (lo + gap / 2) % n;
     int omid = orient(l.a, l.b, p[mid
          ]);
     if (!omid) {
       lo = mid;
       break;
      if (omid == olo)
       lo = mid;
      else
       hi = mid;
   ret.push_back(lo);
  return ret;
    from an external point
bool isGood(Point u, Point v, Point Q,
    int dir) {
  return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q,
     int dir) {
  return orient(Q, u, v) == dir ? u : v
Point pointPolyTangent(const Polygon &
    pt, Point Q, int dir, int lo, int
    hi) {
  while (hi - lo > 1) {
    int mid = (lo + hi) / 2;
    bool pvs = isGood(pt[mid], pt[mid -
        1], Q, dir);
    bool nxt = isGood(pt[mid], pt[mid +
        1], Q, dir);
    if (pvs && nxt) return pt[mid];
    if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q,
           dir, mid + 1, hi);
     Point p2 = pointPolyTangent(pt, Q,
           dir, lo, mid - 1);
     return better(p1, p2, Q, dir);
    if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) ==
          dir)
       hi = mid - 1;
      else if (better(pt[lo], pt[hi], Q,
           dir) == pt[lo])
       hi = mid - 1;
      else
       lo = mid + 1;
    if (!nxt) {
      if (orient(Q, pt[mid], pt[lo]) ==
```

```
lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++)
      ret = better(ret, pt[i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(
    const Polygon &pt, Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt,
      Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q
      , CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &
    poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0
        && (ret.empty() || b != ret.
       back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p
     of n vertices
Tf signedPolygonArea(const Polygon &p)
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() -</pre>
   ret += cross(p[i] - p[0], p[i + 1] -
        p[0]);
 return ret / 2;
}
// given a polygon p of n vertices,
    generates the convex hull in in CCW
// Tested on https://acm.timus.ru/
    problem.aspx?space=1&num=1185
// Caution: when all points are
    colinear AND removeRedundant ==
    false
// output will be contain duplicate
    points (from upper hull) at back
Polygon convexHull(Polygon p, bool
    removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
```

```
p.erase(unique(p.begin(), p.end()), p
      .end());
  int n = p.size();
  Polygon ch(n + n);
  int m = 0; // preparing lower hull
  for (int i = 0; i < n; i++) {</pre>
    while (m > 1 &&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 1]))
               <= check)
     m--;
    ch[m++] = p[i];
  int k = m; // preparing upper hull
  for (int i = n - 2; i >= 0; i--) {
    while (m > k \&\&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 2]))
               <= check)
    ch[m++] = p[i];
  if (n > 1) m--;
  ch.resize(m);
  return ch;
// returns inside = -1, on = 0, outside
int pointInPolygon(const Polygon &p,
    Point o) {
  using Linear::onSegment;
  int wn = 0, n = p.size();
  for (int i = 0; i < n; i++) {</pre>
    int j = (i + 1) \% n;
    if (onSegment(o, Segment(p[i], p[j])
        ) || o == p[i]) return 0;
    int k = dcmp(cross(p[j] - p[i], o -
        p[i]));
    int d1 = dcmp(p[i].y - o.y);
    int d2 = dcmp(p[j].y - o.y);
    if (k > 0 && d1 \le 0 && d2 > 0) wn
    if (k < 0 && d2 <= 0 && d1 > 0) wn
  }
 return wn ? -1 : 1;
// Given a simple polygon p, and a line
     1, returns (x, y)
// x = longest segment of 1 in p, y =
    total length of l in p.
pair<Tf, Tf> linePolygonIntersection(
    Line 1, const Polygon &p) {
  using Linear::lineLineIntersection;
  int n = p.size();
  vector<pair<Tf, int>> ev;
  for (int i = 0; i < n; ++i) {</pre>
    Point a = p[i], b = p[(i + 1) \% n],
        z = p[(i - 1 + n) \% n];
    int ora = orient(1.a, 1.b, a), orb =
         orient(1.a, 1.b, b),
       orz = orient(l.a, l.b, z);
    if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(
            d, 0);
       // else // Point Touch
```

```
} else if (orz)
     ev.emplace_back(d, orz);
   else if (orb)
     ev.emplace_back(d, orb);
 } else if (ora == -orb) {
   Point ins;
   lineLineIntersection(1, Line(a, b)
        , ins);
   ev.emplace_back(dot(ins - 1.a, 1.b
         - 1.a), 0);
}
sort(ev.begin(), ev.end());
Tf ans = 0, len = 0, last = 0, tot =
    0:
bool active = false;
int sign = 0;
for (auto &qq : ev) {
 int tp = qq.second;
 Tf d = qq.first; /// current Segment
       is (last, d)
                  /// On Border
 if (sign) {
   len += d - last;
   tot += d - last;
   ans = max(ans, len);
   if (tp != sign) active = !active;
   sign = 0;
 } else {
   if (active) { /// Strictly Inside
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
   if (tp == 0)
     active = !active;
   else
     sign = tp;
 last = d;
 if (!active) len = 0;
ans /= length(l.b - l.a);
tot /= length(l.b - l.a);
return {ans, tot};
```

4.6 Point Redifined

```
namespace geo {
   #define x real()
   #define y imag()
   #define setx(i) real(i)
   #define sety(i) imag(i)
   typedef complex<double> point;
   typedef point vec2;
   double dist(point a, point b) {
       return abs(b-a); }
   double dot(vec2 a, vec2 b) { return
        (conj(a) * b).x; }
   double cross(vec2 a, vec2 b) {
       return (conj(a) * b).y; }
   vec2 rotate90(vec2 a) { return point
        (-a.y, a.x); }
   double area(vec2 a, vec2 b) { return
        abs(cross(a,b)); }
   // double area(point a, point b,
       point c) { return area( getVec2
        (a,b), getVec2(a,c)); }
```

5 Graph

5.1 LCA, ETT, VT

```
#define lg(n) (31 - __builtin_clz(n))
const int N = 1e5 + 1;
const int K = lg(N) + 1;
vector<int> adj[N];
int anc[N][K], lvl[N];
namespace lca {
 void init(int u = 1, int p = 0, int d
       = 0) {
   lvl[u] = d;
   anc[u][0] = p;
   for (int i = 1; i < K; i++)</pre>
     anc[u][i] = anc[anc[u][i - 1]][i -
   for (auto v : adj[u])
     if (v != p)
       init(v, u, d + 1);
 int getAnc(int u, int k) {
   for (int i = 0; u and i < K; i++)</pre>
     if ((k >> i) & 1)
       u = anc[u][i];
   return u;
 }
 int lca(int u, int v) {
   if (lvl[u] < lvl[v]) swap(u, v);</pre>
   u = getAnc(u, lvl[u] - lvl[v]);
   if (u == v) return u;
   for (int i = K - 1; ~i; i--)
     if (anc[u][i] != anc[v][i])
       u = anc[u][i], v = anc[v][i];
   return anc[u][0];
 int dist(int u, int v) {
   return lvl[u] + lvl[v] - 2 * lvl[lca
        (u, v)];
 }
};
struct euler_tour {
 int time = 0;
 tree &T;
 int n:
 vector<int> start, finish, level, par
  euler_tour(tree &T, int root = 0)
     : T(T), n(T.n), start(n), finish(n|}
          ), level(n), par(n) {
   time = 0;
   call(root);
 void call(int node, int p = -1) {
   if (p != -1) level[node] = level[p]
        + 1:
   start[node] = time++;
   for (int e : T[node])
     if (e != p) call(e, node);
   par[node] = p;
   finish[node] = time++;
```

```
bool isAncestor(int node, int par) {
   return start[par] <= start[node] and</pre>
         finish[par] >= finish[node];
 int subtreeSize(int node) { return
      finish[node] - start[node] + 1 >> |}
       1: }
tree virtual_tree(vector<int> &nodes,
    lca_table &table, euler_tour &tour)
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
           start[x] < tour.start[y]; });</pre>
 int n = nodes.size();
  for (int i = 0; i + 1 < n; i++)</pre>
   nodes.push_back(table.lca(nodes[i],
        nodes[i + 1]));
 sort(nodes.begin(), nodes.end());
 nodes.erase(unique(nodes.begin(),
      nodes.end()), nodes.end());
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
 n = nodes.size();
 stack<int> st;
 st.push(0);
 tree ans(n);
 for (int i = 1; i < n; i++) {</pre>
   while (!tour.isAncestor(nodes[i],
       nodes[st.top()])) st.pop();
   ans.addEdge(st.top(), i);
   st.push(i);
 return ans;
set<int> getCenters(tree &T) {
 int n = T.n;
 vector<int> deg(n), q;
 set<int> s;
 for (int i = 0; i < n; i++) {</pre>
   deg[i] = T[i].size();
   if (deg[i] == 1) q.push_back(i);
   s.insert(i);
 }
 for (vector<int> t; s.size() > 2; q =
       t) {
   for (auto x : q) {
     for (auto e : T[x])
       if (--deg[e] == 1) t.push_back(e
           );
     s.erase(x);
   }
 return s;
5.2 SCC
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
vector<int> order, component;
// tp = 0 ,finding topo order, tp = 1 ,
     reverse edge traversal
```

void dfs(int u, int tp = 0) {

```
vis[u] = true;
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
 reverse(order.begin(), order.end());
 for (int i : order) {
   if (!vis[i]) {
     // one component is found
     dfs(i, 1), component.clear();
 }
5.3 Euler Tour on Edge
```

10

```
// for simplicity, G[idx] contains the
                                   adjacency list of a node
start[x] < tour.start[y]; }); // while G(e) is a reference to the e-
                                   th edge.
                              const int N = 2e5 + 5;
                              int in[N], out[N], fwd[N], bck[N];
                              int t = 0;
                              void dfs(graph &G, int node, int par) {
                                out[node] = t:
                                for (int e : G[node]) {
                                  int v = G(e).to(node);
                                  if (v == par) continue;
                                  fwd[e] = t++;
                                  dfs(G, v, node);
                                  bck[e] = t++;
                                in[node] = t - 1;
                              void init(graph &G, int node) {
                                dfs(G, node, node);
```

5.4 LCA In O(1)

```
/* LCA in O(1)
 * depth calculates weighted distance
 * level calculates distance by number
     of edges
 * Preprocessing in NlongN */
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG
      [1] = 0;
  for (int i = 2; i <= n; i++) LOG[i] =</pre>
       LOG[i >> 1] + 1;
  for (int i = 0; i <= n; i++) adj[i].</pre>
      clear();
```

```
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
  if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v])</pre>
       return v;
 int 1 = LOG[id[v] - id[u] + 1];
 int p1 = id[u], p2 = id[v] - (1 << 1)</pre>
 int d1 = level[table[1][p1]], d2 =
      level[table[1][p2]];
  if (d1 < d2)
   return par[table[1][p1]];
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time;
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table</pre>
      [0][i] = a[i];
 for (int 1 = 0; 1 < L - 1; 1++) {</pre>
   for (int i = 1; i <= n; i++) {</pre>
     table[1 + 1][i] = table[1][i];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 <<</pre>
           1)]] < level[table[1][i]];</pre>
     if (C1 && C2) table[l + 1][i] =
          table[1][i + (1 << 1)];
   }
 }
}
```

HLD5.5

```
const int N = 1e6 + 7;
template <typename DT>
struct Segtree {
 // write lazy segtree here
```

```
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
  depth[u] = depth[p] + 1;
  for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz
         = subsz;
   SZ += subsz;
 return SZ;
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
  if (heavy[u] != -1) decompose(heavy[u
      ], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
 }
int query(int a, int b) {
 int ret = 0;
  for (; head[a] != head[b]; b = par[
      head[b]]) {
   if (depth[head[a]] > depth[head[b]])
        swap(a, b);
   ret += tree.query(1, 0, cnt, pos[
       head[b]], pos[b]);
  if (depth[a] > depth[b]) swap(a, b);
  ret += tree.query(1, 0, cnt, pos[a],
      pos[b]);
 return ret;
```

5.6 Dinic Max Flow

flow; };

```
/**
Implementation of Dinic's algorithm
    with optional scaling
Source: Chilli (https://codeforces.com/
    blog/entry/66006)
Complexity: O(ans*E) or O(V^2E) without
     scaling, O(VE log(U)) with scaling
Scaling performs much better in worst
    case, but has much higher constant
    factor
To enable scaling, call maxFlow(true)
Everything 0-indexed
*/
namespace Dinic {
   typedef long long LL;
   const int N = 5005, K = 60; /// N >
       no of nodes, K >= max bits in
       capacity
   const LL INF = 1e18;
   struct Edge { int frm, to; LL cap,
```

```
int s, t, n;
int level[N], ptr[N];
vector<Edge> edges;
vector<int> adj[N];
void init(int nodes) {
   n = nodes;
   for (int i=0; i<n; i++) adj[i].</pre>
        clear():
   edges.clear();
}
/// For adding undirected Edge (u, v
    , c) call addEdge(u, v, c, c);
int addEdge(int a, int b, LL cap, LL
     revcap = 0) {
   edges.push_back({a, b, cap, 0});
   edges.push_back({b, a, revcap,
   adj[a].push_back(edges.size()-2)
   adj[b].push_back(edges.size()-1)
   return edges.size()-2;
bool bfs(LL lim) {
   fill(level, level+n, -1);
   level[s] = 0;
   queue<int> q;
   q.push(s);
   while (!q.empty() && level[t] ==
         -1) {
       int v = q.front();
       q.pop();
       for (int id: adj[v]) {
           Edge e = edges[id];
           if (level[e.to] == -1 &&
               e.cap - e.flow >= lim
              q.push(e.to);
              level[e.to] = level[v
                   ] + 1;
           }
       }
   }
   return level[t] != -1;
LL dfs(int v, LL flow) {
   if (v == t || !flow)
                             return
       flow:
   for (; ptr[v] < adj[v].size();</pre>
       ptr[v]++) {
       int eid = adj[v][ptr[v]];
       Edge &e = edges[eid];
       if (level[e.to] != level[v]
           + 1) continue;
       if (LL pushed = dfs(e.to,
           min(flow, e.cap - e.flow
           ))) {
           e.flow += pushed;
           edges[eid^1].flow -=
               pushed;
           return pushed;
       }
   }
   return 0;
LL maxFlow(int source, int sink,
    bool SCALING = false) {
   s = source, t = sink;
```

```
long long flow = 0;
   for (LL lim = SCALING ? (1LL <<</pre>
        K) : 1; \lim > 0; \lim >>= 1)
       while (bfs(lim)) {
           fill(ptr, ptr+n, 0);
           while (LL pushed = dfs(s,
                INF)) flow += pushed
       }
   }
   return flow;
}
bool leftOfMinCut(int x) {return
    level[x] != -1;
/// Only works for undirected graph,
     Make sure to add UNDIRECTED
    edges. (u, v, c, c)
/// returns n by n matrix flow, st
    flow[i][j] = maxFlow
/// tree holds the edges of a gomory
    -hu tree of the graph
vector<vector<LL>> allPairMaxFlow(
    vector<Edge> &tree) {
   tree.clear();
   vector<vector<LL>> flow(n,
        vector<LL> (n, INF));
   vector<int> par(n);
   for (int i=1; i<n; i++) {</pre>
       for (auto &e: edges) e.flow
            = 0;
       LL f = maxFlow(i, par[i]);
       tree.push_back({i, par[i], f
           });
       for (int j=i+1; j<n; j++)</pre>
           if (par[j] == par[i] &&
               leftOfMinCut(j)) par[
               j] = i;
       flow[i][par[i]] = flow[par[i
            ]][i] = f;
       for (int j=0; j<i; j++)</pre>
           if (j != par[i]) flow[i
               ][j] = flow[j][i] =
               min(f, flow[par[i]][j
               ]);
   return flow;
}
```

5.7 Min Cost Max Flow

}

```
mt19937 rnd(chrono::steady_clock::now()
    .time_since_epoch().count());
const LL inf = 1e9;
struct edge {
 int v, rev;
 LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(
         cost), flow(0) {}
};
struct mcmf {
 int src, sink, n;
 vector<int> par, idx, Q;
 vector<bool> inq;
```

```
vector<LL> dis;
vector<vector<edge>> g;
mcmf() {}
mcmf(int src, int sink, int n)
    : src(src),
     sink(sink),
     n(n),
     par(n),
     idx(n),
     inq(n),
     dis(n),
     g(n),
     Q(10000005) {} // use Q(n) if
          not using random
void add_edge(int u, int v, LL cap,
    LL cost, bool directed = true) {
  edge _u = edge(v, g[v].size(), cap,
      cost);
  edge _v = edge(u, g[u].size(), 0, -
      cost);
  g[u].pb(_u);
  g[v].pb(_v);
  if (!directed) add_edge(v, u, cap,
      cost, true);
bool spfa() {
  for (int i = 0; i < n; i++) {</pre>
   dis[i] = inf, inq[i] = false;
  int f = 0, 1 = 0;
  dis[src] = 0, par[src] = -1, Q[1++]
      = src, inq[src] = true;
  while (f < 1) {</pre>
    int u = Q[f++];
    for (int i = 0; i < g[u].size(); i</pre>
        ++) {
     edge &e = g[u][i];
     if (e.cap <= e.flow) continue;</pre>
     if (dis[e.v] > dis[u] + e.cost)
       dis[e.v] = dis[u] + e.cost;
       par[e.v] = u, idx[e.v] = i;
       if (!inq[e.v]) inq[e.v] = true
            , Q[1++] = e.v;
       // if (!inq[e.v]) {
       // inq[e.v] = true;
       // if (f && rnd() & 7) Q[--f]
             = e.v;
       // else Q[1++] = e.v;
       // }
     }
    inq[u] = false;
 return (dis[sink] != inf);
pair<LL, LL> solve() {
  LL mincost = 0, maxflow = 0;
  while (spfa()) {
    LL bottleneck = inf;
    for (int u = par[sink], v = idx[
        sink]; u != -1; v = idx[u], u
        = par[u]) {
     edge &e = g[u][v];
     bottleneck = min(bottleneck, e.
          cap - e.flow);
   }
   for (int u = par[sink], v = idx[
        sink; u != -1; v = idx[u], u |_{void init(graph &G)} {
```

```
= par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck
     }
     mincost += bottleneck * dis[sink],
          maxflow += bottleneck;
   return make_pair(mincost, maxflow);
 }
// want to minimize cost and don't care
     about flow
// add edge from sink to dummy sink (
    cap = inf, cost = 0)
// add edge from source to sink (cap =
    inf, cost = 0)
// run mcmf, cost returned is the
    minimum cost
```

```
Bridge Tree
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G,
    int par = -1, int d = 0) {
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1)
       if (low[to] > depth[node]) {
        bridges.emplace_back(node, to)
        components.push_back({});
         for (int x = -1; x != to; x =
             st.top(), st.pop())
          components.back().push_back(
               st.top());
       }
     low[node] = min(low[node], low[to
         ]);
   }
 }
 if (par == -1) {
   components.push_back({});
   while (!st.empty()) components.back
        ().push_back(st.top()), st.pop
 }
graph &create_tree() {
 for (auto &comp : components) {
   int idx = tree.addNode();
   for (auto &e : comp) id[e] = idx;
 for (auto &[l, r] : bridges) tree.
      addEdge(id[1], id[r]);
 return tree;
```

```
int n = G.n;
depth.assign(n, -1), id.assign(n, -1)
        , low.resize(n);
for (int i = 0; i < n; i++)
   if (depth[i] == -1) find_bridges(i,
        G);
}</pre>
```

5.9 Tree Isomorphism

```
LL Hash(int u, int p) {
  vector<LL> childrenHash;
  for (auto v : adj[u]) if (v != p)
     childrenHash.add(Hash(v, u));
  sort(all(childrenHash));
  LL nodeHash = 0;
  for (int i = 0; i < childrenHash.size
     (); i++)
  nodeHash = (nodeHash + childrenHash[
        i] * bigmod(SEED, i, MOD)) %
     MOD;
  return nodeHash;
}</pre>
```

6 Math

6.1 Linear Sieve

```
const int N = 10000000;
vector<int> lp(N+1);
vector<int> pr;

for (int i=2; i <= N; ++i) {
    if (lp[i] == 0) {
        lp[i] = i;
        pr.push_back(i);
    }
    for (int j = 0; i * pr[j] <= N; ++j)
        {
        lp[i * pr[j]] = pr[j];
        if (pr[j] == lp[i]) {
            break;
        }
    }
}</pre>
```

6.2 Pollard Rho

```
LL mul(LL a, LL b, LL mod) {
 return (__int128)a * b % mod;
  // LL ans = a * b - mod * (LL) (1.L /
       mod * a * b);
  // return ans + mod * (ans < 0) - mod
       * (ans >= (LL) mod);
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(
      num, num, mod))
   if (pow & 1) ans = mul(ans, num, mod
       );
 return ans;
}
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (
      n \mid 1) == 3;
 LL a[] = \{2, 325, 9375, 28178,
      450775, 9780504, 1795265022};
  LL s = \__builtin_ctzll(n - 1), d = n
      >> s;
 for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
```

```
for (; p != 1 and p != n - 1 and x %
        n and i--; p = mul(p, p, n))
   if (p != n - 1 \text{ and } i != s) \text{ return}
        false;
 return true;
LL get_factor(LL n) {
 auto f = [\&](LL x) \{ return mul(x, x,
       n) + 1; };
 LL x = 0, y = 0, t = 0, prod = 2, i =
       2, q;
 for (; t++ % 40 or gcd(prod, n) == 1;
       x = f(x), y = f(f(y))) {
    (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) -
        min(x, y), n) ? q : prod;
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11,
      13, 17, 19, 23, 29, 31, 37, 41,
                     43, 47, 53, 59,
                         61, 67, 71,
                         73, 79, 83,
                          89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
 auto _factor = [&](LL n, auto &
      _factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
 };
  _factor(n, _factor);
 return res;
```

6.3 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
   if (b == 0)
     return {1, 0, a};
   else {
     auto [x, y, g] = EGCD(b, a % b);
     return {y, x - a / b * y, g};
   }
}
// given modulo equations, will apply
   CRT
PLL CRT(vector<PLL> &v) {
   LL V = 0, M = 1;
   for (auto &[v, m] : v) { // value %
        mod
     auto [x, y, g] = EGCD(M, m);
     if ((v - V) % g != 0) return {-1,
```

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6.4 Mobius Function

using CD = complex<double>;

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i)
            mob[j] -= mob[i];
    }
}</pre>
```

6.5 FFT

```
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N),
        sin(2 * PI * i / N));
   wp[1][i] = CD(cos(2 * PI * i / N), -
        sin(2 * PI * i / N));
void fft(vector<CD> &v, bool invert =
    false) {
  if (v.size() != perm.size())
      precalculate(v.size());
  for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i</pre>
        ]]);
  for (int len = 2; len <= N; len *= 2)</pre>
   for (int i = 0, d = N / len; i < N;</pre>
        i += len) {
     for (int j = 0, idx = 0; j < len /
           2; j++, idx += d) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j]
             + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
  if (invert) {
   for (int i = 0; i < N; i++) v[i] /=</pre>
```

```
}
}
void pairfft(vector<CD> &a, vector<CD>
    &b, bool invert = false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[</pre>
      i] + b[i] * CD(0, 1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) *
         CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) *
         CD(0, -0.5);
   }
 }
}
, const vector<LL> &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<=</pre>
 vector<CD> fa(a.begin(), a.end()), fb
      (b.begin(), b.end());
 fa.resize(n);
 fb.resize(n);
          fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] =</pre>
     fa[i] * fb[i];
 fft(fa, true);
 vector<LL> ans(n);
 for (int i = 0; i < n; i++) ans[i] =
     round(fa[i].real());
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a,
    const vector<LL> &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<=</pre>
      1:
 vector<CD> al(n), ar(n), bl(n), br(n)
 for (int i = 0; i < a.size(); i++) al</pre>
      [i] = a[i] % M / B, ar[i] = a[i]
     % M % B;
 for (int i = 0; i < b.size(); i++) bl</pre>
      [i] = b[i] % M / B, br[i] = b[i]
     % M % B;
 pairfft(al, ar);
 pairfft(bl, br);
          fft(al); fft(ar); fft(bl);
     fft(br);
 for (int i = 0; i < n; i++) {</pre>
   CD ll = (al[i] * bl[i]), lr = (al[i]
   CD rl = (ar[i] * bl[i]), rr = (ar[i])
        * br[i]);
   al[i] = 11;
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr;
```

```
pairfft(al, ar, true);
pairfft(bl, br, true);
         fft(al, true); fft(ar, true)
    ; fft(bl, true); fft(br, true);
vector<LL> ans(n);
for (int i = 0; i < n; i++) {</pre>
 LL right = round(br[i].real()), left
       = round(al[i].real());
 LL mid = round(round(bl[i].real()) +
       round(ar[i].real()));
 ans[i] = ((left \% M) * B * B + (mid
      % M) * B + right) % M;
return ans:
```

```
6.6 NTT
const LL N = 1 << 18;</pre>
const LL MOD = 786433;
vector<LL> P[N];
LL Pow(LL b, LL p) {
 LL ret = 1;
  while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 }
 return ret;
LL primitive_root(LL p) {
  vector<LL> factor;
  LL phi = p - 1, n = phi;
  for (LL i = 2; i * i <= n; i++) {
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
  if (n > 1) factor.emplace_back(n);
  for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() &&</pre>
     ok &= Pow(res, phi / factor[i]) != auto insert = [&](int x) {
           1;
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - \_builtin\_clz(n));
  LL r = Pow(g, (MOD - 1) / n);
  inv_n = Pow(n, MOD - 2);
  w[0] = w[n] = 1;
  for (LL i = 1; i < n; i++) w[i] = (w[ ];
      i - 1] * r) % MOD;
  for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i &
        1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)</pre>
   if (i < rev[i]) swap(a[i], a[rev[i</pre>
  for (LL m = 2; m <= n; m <<= 1) {</pre>
```

for (LL i = 0; i < n; i += m) {</pre>

```
for (LL j = 0; j < (m >> 1); j++)
       LL &u = a[i + j], &v = a[i + j +
            (m >> 1)];
       LL t = v * w[dir ? n - n / m * j]
            : n / m * j] % MOD;
       v = u - t < 0 ? u - t + MOD : u
            - t:
       u = u + t \ge MOD ? u + t - MOD :
            u + t;
   }
 }
 if (dir)
   for (LL i = 0; i < n; i++) a[i] = (
        inv_n * a[i]) % MOD;
vector<LL> mul(vector<LL> p, vector<LL>
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;</pre>
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
 for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
 for (LL i = 0; i < sz; i++) a[i] = (a</pre>
      [i] * b[i]) % MOD;
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
 while (c.size() && c.back() == 0) c.
     pop_back();
 return c;
```

6.7 XOR Basis

```
int basis[21] = {}, sz = 0;
 while (x) {
   int i = __builtin_ctz(x & -x);
   if (not basis[i]) {
     basis[i] = x;
     sz++;
     return x;
   x ^= basis[i];
 return 0;
```

String

Aho Corasick

```
const int sg = 26, N = 1e3 + 9;
struct aho_corasick {
 struct node {
   node *link, *out, *par;
   bool leaf;
   LL val;
```

```
int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '
        $', int len = 0)
       : par(par), p_ch(p_ch), len(len)
     val = leaf = cnt = last = 0;
     link = out = NULL;
   }
 };
 vector<node> trie;
 node* root;
 aho_corasick() {
   trie.reserve(N), trie.emplace_back()
   root = &trie[0];
   root->link = root->out = root;
 }
 inline int f(char c) { return c - 'a'
      ; }
 inline node* add_node(node* par =
      NULL, char p_ch = '$', int len =
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
 void add_str(const string& s, LL val
      = 1) {
   node* now = root;
   for (char c : s) {
     int i = f(c);
     if (!now->to[i]) now->to[i] =
         add_node(now, c, now->len + 1)
     now = now->to[i];
   now->leaf = true, now->val++;
 void push_links() {
   queue<node*> q;
   for (q.push(root); q.empty(); q.pop
        ()) {
     node *cur = q.front(), *link = cur
         ->link:
     cur->out = link->leaf ? link :
         link->out;
     int idx = 0;
     for (auto& next : cur->to) {
       if (next != NULL) {
         next->link = cur != root ?
             link->to[idx++] : root;
         q.push(next);
       } else
         next = link->to[idx++];
     }
   cur->val += link->val;
};
    Double hash
```

```
// define +, -, * for (PLL, LL) and (
    PLL, PLL), % for (PLL, PLL);
PLL base(1949313259, 1997293877);
namespace Hashing {
    using LL = long long;
    using PLL = pair<LL,LL>;
    #define ff first
```

```
#define ss second
const PLL M = {2091573227,
                    ///Should be
    2117566807};
    large primes
const LL base = 1259;
    /Should be larger than alphabet
     size
const int N = 1e6+7;
    /Highest length of string
PLL operator+ (const PLL& a, LL x)
    {return {a.ff + x, a.ss + x};}
PLL operator- (const PLL& a, LL x)
    {return {a.ff - x, a.ss - x};}
PLL operator* (const PLL& a, LL x)
    {return {a.ff * x, a.ss * x};}
PLL operator+ (const PLL& a, PLL x)
    {return {a.ff + x.ff, a.ss + x.
PLL operator- (const PLL& a, PLL x)
    {return {a.ff - x.ff, a.ss - x.
PLL operator* (const PLL& a, PLL x)
    {return {a.ff * x.ff, a.ss * x.
    ss};}
PLL operator% (const PLL& a, PLL m)
    {return {a.ff % m.ff, a.ss % m.
ostream& operator << (ostream& os, PLL
     hash) {
   return os<<"("<<hash.ff<<", "<<
       hash.ss<<")";
}
PLL pb[N];
              ///powers of base mod
///Call pre before everything
void hashPre() {
   pb[0] = \{1,1\};
   for (int i=1; i<N; i++)</pre>
                             pb[i] =
         (pb[i-1] * base)%M;
///Calculates hashes of all prefixes
     of s including empty prefix
vector<PLL> hashList(string s) {
   int n = s.size();
   vector<PLL> ans(n+1);
   ans[0] = \{0,0\};
   for (int i=1; i<=n; i++) ans[i]</pre>
        = (ans[i-1] * base + s[i-1])
        %M;
   return ans;
}
///Calculates hash of substring s[l
    ..r] (1 indexed)
PLL substringHash(const vector<PLL>
    &hashlist, int 1, int r) {
   return (hashlist[r]+(M-hashlist[
        l-1])*pb[r-l+1])%M;
///Calculates Hash of a string
PLL Hash (string s) {
   PLL ans = \{0,0\};
   for (int i=0; i<s.size(); i++)</pre>
        ans=(ans*base + s[i])%M;
   return ans;
///Tested on https://toph.co/p/
    palindromist
///appends c to string
```

PLL append(PLL cur, char c) {

```
return (cur*base + c)%M;
}
///Tested on https://toph.co/p/
    palindromist
///prepends c to string with size k
PLL prepend(PLL cur, int k, char c)
   return (pb[k]*c + cur)%M;
///Tested on https://toph.co/p/
    chikongunia
///replaces the i-th (0-indexed)
    character from right from a to
PLL replace(PLL cur, int i, char a,
    char b) {
   return cur + pb[i] * (M+b-a)%M;
///Erases c from front of the string
     with size len
PLL pop_front(PLL hash, int len,
    char c) {
   return (hash + pb[len-1]*(M-c))%
///Tested on https://toph.co/p/
    palindromist
///concatenates two strings where
    length of the right is k
PLL concat(PLL left, PLL right, int
    k) {
   return (left*pb[k] + right)%M;
PLL power (const PLL& a, LL p) {
   if (p==0) return {1,1};
   PLL ans = power(a, p/2);
   ans = (ans * ans) %M;
   if (p%2)
             ans = (ans*a)%M;
   return ans;
PLL inverse(PLL a) {
   if (M.ss == 1) return power(a, M
        .ff-2);
   return power(a, (M.ff-1)*(M.ss
       -1)-1);
///Erases c from the back of the
    string
PLL invb = inverse({base, base});
PLL pop_back(PLL hash, char c) {
   return ((hash-c+M)*invb)%M;
///Tested on https://toph.co/p/
    palindromist
///Calculates hash of string with
    size len repeated cnt times
///This is O(\log n). For O(1), pre-
    calculate inverses
PLL repeat(PLL hash, int len, LL cnt
    ) {
   PLL mul = ((pb[len*cnt]-1+M) *
        inverse(pb[len]-1+M))%M;
   PLL ans = (hash*mul);
   if (pb[len].ff == 1) ans.ff =
       hash.ff*cnt:
   if (pb[len].ss == 1) ans.ss =
       hash.ss*cnt;
   return ans%M;
}
```

```
struct pair_hash {
   inline std::size_t operator()(
       const std::pair<LL,LL> & v)
       const {
       return v.first*31+v.second;
    }
};
```

7.3 KMP

```
vector<int> prefix_function(string s) {
  int n = (int)s.length();
  vector<int> pi(n);
  for (int i = 1; i < n; i++) {
    int j = pi[i-1];
    while (j > 0 && s[i] != s[j])
        j = pi[j-1];
    if (s[i] == s[j])
        j++;
    pi[i] = j;
  }
  return pi;
}
```

7.4 Manacher's

```
vector<int> d1(n);
// d[i] = number of palindromes taking
    s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i /*
    ++) {
 int k = (i > r) ? 1 : min(d1[1 + r -
     i], r - i + 1);
 while (0 <= i - k && i + k < n && s[i
       -k] == s[i + k]) k++;
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n):
// d[i] = number of palindromes taking
    s[i-1] and s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 0 : min(d2[1 + r -
      i + 1], r - i + 1);
 while (0 \le i - k - 1 \&\& i + k \le n \&\&
       s[i - k - 1] == s[i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i +
```

7.5 String Match FFT

```
//find occurrences of t in s where '?'s
     are automatically matched with any
     character
//res[i + m - 1] = sum_j=0 to m - 1_{s[}
    i + j] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
 vector < int > s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[</pre>
      i] == '?' ? 0 : s[i] - 'a' + 1;
      //assign any non zero number for
      non '?'s
  for(int i = 0; i < n; i++) s2[i] = s1</pre>
 for(int i = 0; i < n; i++) s3[i] = s1</pre>
      [i] * s2[i];
```

```
vector<int> t1(m), t2(m), t3(m);
for(int i = 0; i < m; i++) t1[i] = t[</pre>
    i] == '?' ? 0 : t[i] - 'a' + 1;
for(int i = 0; i < m; i++) t2[i] = t1</pre>
    [i] * t1[i];
for(int i = 0; i < m; i++) t3[i] = t1</pre>
    [i] * t2[i];
reverse(t1.begin(), t1.end());
reverse(t2.begin(), t2.end());
reverse(t3.begin(), t3.end());
vector<int> s1t3 = multiply(s1, t3);
vector<int> s2t2 = multiply(s2, t2);
vector<int> s3t1 = multiply(s3, t1);
vector<int> res(n);
for(int i = 0; i < n; i++) res[i] =</pre>
    s1t3[i] - s2t2[i] * 2 + s3t1[i];
vector<int> oc;
for(int i = m - 1; i < n; i++) if(res</pre>
    [i] == 0) oc.push_back(i - m + 1)
return oc;
```

7.6 Suffix Array

```
#include <bits/stdc++.h>
using namespace std;
 O(|S| + |alphabet|) Suffix Array
LIM := max{s[i]} + 2
void inducedSort (const vector <int> &
    vec, int val_range, vector <int> &
    SA, const vector <int> &sl, const
    vector <int> &lms_idx) {
 vector <int> l(val_range, 0), r(
     val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c
        + 1]:
 partial_sum(l.begin(), l.end(), l.
     begin());
 partial_sum(r.begin(), r.end(), r.
     begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i >=
      0; --i) SA[--r[vec[lms_idx[i]]]]
      = lms_idx[i];
 for (int i : SA) if (i > 0 and sl[i -
       1]) SA[1[vec[i-1]]++] = i-1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
 partial_sum(r.begin(), r.end(), r.
     begin());
 for (int k = SA.size() - 1, i = SA[k]
     ]; k; --k, i = SA[k]) {
   if (i and !sl[i - 1]) SA[--r[vec[i -
        1]]] = i - 1;
vector <int> suffixArray (const vector
    <int> &vec, int val_range) {
 const int n = vec.size();
 vector <int> sl(n), SA(n), lms_idx;
```

for (int i = n - 2; $i \ge 0$; --i) {

```
sl[i] = vec[i] > vec[i + 1] or (vec[i])
        i] == vec[i + 1] and sl[i + 1])
   if (sl[i] and !sl[i + 1]) lms_idx.
        emplace_back(i + 1);
 reverse(lms_idx.begin(), lms_idx.end
 inducedSort(vec, val_range, SA, s1,
      lms_idx);
 vector <int> new_lms_idx(lms_idx.size
      ()), lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {</pre>
   if (SA[i] > 0 and !sl[SA[i]] and sl[
        SA[i] - 1]) new_lms_idx[k++] =
        SA[i];
 int cur = 0; SA[n - 1] = 0;
 for (int k = 1; k < new_lms_idx.size</pre>
      (); ++k) {
   int i = new_lms_idx[k - 1], j =
        new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
     SA[j] = ++cur; continue;
   bool flag = 0;
   for (int a = i + 1, b = j + 1; ; ++a
        , ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
     }
     if ((!sl[a] and sl[a - 1]) or (!sl
          [b] and sl[b - 1])) {
       flag = !(!sl[a] and sl[a - 1]
           and !sl[b] and sl[b - 1]);
           break:
     }
   SA[j] = flag ? ++cur : cur;
 for (int i = 0; i < lms_idx.size();</pre>
      ++i) lms_vec[i] = SA[lms_idx[i]];
 if (cur + 1 < lms_idx.size()) {</pre>
   auto lms_SA = suffixArray(lms_vec,
       cur + 1);
   for (int i = 0; i < lms_idx.size();</pre>
        ++i) new_lms_idx[i] = lms_idx[
        lms_SA[i]];
 inducedSort(vec, val_range, SA, sl,
      new_lms_idx); return SA;
vector <int> getSuffixArray (const
    string &s, const int LIM = 128) {
 vector <int> vec(s.size() + 1);
 copy(begin(s), end(s), begin(vec));
      vec.back() = '$';
 auto ret = suffixArray(vec, LIM);
 ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any
    two suffix
vector <int> getLCParray (const string
    &s, const vector <int> &SA) {
 int n = s.size(), k = 0;
 vector <int> lcp(n), rank(n);
```

```
17
```

```
for (int i = 0; i < n; ++i) rank[SA[i</pre>
      ]] = i;
 for (int i = 0; i < n; ++i, k ? --k :</pre>
       0) {
   if (rank[i] == n - 1) {
     k = 0; continue;
   int j = SA[rank[i] + 1];
   while (i + k < n \text{ and } j + k < n \text{ and } s
         [i + k] == s[j + k]) ++k;
   lcp[rank[i]] = k;
 }
 lcp[n - 1] = 0; return lcp;
int main() {
 string s; cin >> s;
  for (const int i : getSuffixArray(s))
       printf("%d ", i);
 puts("");
 return 0;
     \operatorname{Trie}
template<int sz>
```

```
struct Trie {
 Trie() : id(1) {
   memset(endMark, 0, sizeof endMark);
   for_each(all(trie), [](vector<int> &
       v) { v.assign(sz, 0); });
 void insert(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false:
     cur = trie[cur][val];
   return endMark[cur];
 }
 int id, endMark[100005];
 vector<int> trie[100005];
```

7.8 Z Algo

```
vector<int> calcz(string s) {
  int n = s.size();
  vector<int> z(n);
  int 1 = 0, r = 0;
  for (int i = 1; i < n; i++) {
    if (i > r) {
        1 = r = i;
        while (r < n && s[r] == s[r - 1])
        r++;
    z[i] = r - 1, r--;</pre>
```

8 Extra

8.1 Stress Tester

```
g++ -02 -std=c++17 "$1".cpp -o $1
g++ -02 -std=c++17 "$2".cpp -o $2
g++ -02 -std=c++17 "$3".cpp -o $3
# $1 is actual code
# $2 is good code
# $3 is generator
for ((i = 1;; i++)); do
 echo 'Test #'$i
 timeout 5s ./$3 $RANDOM > in
  timeout 5s./\$1 < in > out
  timeout 5s./$2 < in > ans
 diff -i ans out > diff.out
  if [ $? -ne 0 ]; then
   cat in
   cat diff.out
   break
 fi
done
```

8.2 Sublime Build

8.3 vimrc

" Auto import & Compile

```
:autocmd BufNewFile *.cpp Or ~/template
    .cpp
nnoremap <F4> :!xclip -o -sel clip > ~/
   cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip
     > ~/cp/in <CR><CR>
nnoremap <F6> :!xclip -sel clip % <CR><
   CR>
inoremap <F6> <ESC>:!xclip -sel clip %
    <CR><CR>
autocmd filetype cpp nnoremap <F9>
    wa \| !g++ -02 -std=c++17 % -o %:r
    && timeout 5s ./%:r < ^{\sim}/cp/in> ^{\sim}/cp
autocmd filetype cpp inoremap <F9> <ESC set softtabstop=2
    >:wa \| !g++ -02 -std=c++17 % -o %: set smartindent
    r && timeout 5s ./%:r < ~/cp/in> ~/ set smarttab
```

```
cp/out<CR>
autocmd filetype cpp nnoremap <F10> :
    wa \| !make %:r D=1 && ./%:r < ~/cp
    /in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F10> <
    ESC>:wa \| !make clean && make %:r
    D=1 \&\& ./\%:r < ^/cp/in > ^/cp/out <
    CR>
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'),
     col('.')-1, 1) == ")" ? "\<Right>"
     : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'),
     col('.')-1, 1) == "}" ? "\<Right>"
     : "}"
inoremap [ [] < left >
inoremap <expr> ] strpart(getline('.'),
     col('.')-1, 1) == "]" ? "\<Right>"
     : "]"
inoremap <expr> " strpart(getline('.'),
     col('.')-1, 1) == "\"" ? "\<Right>
    " : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'),
    col('.')-1, 1) == "\'" ? "\<Right
    >" : "\'\\'\<left>"
inoremap <expr> <CR> <sid>
    insert_newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col(
      '.')-2, 2)
  return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<CR>\<
      ESC>\0" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>
    insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col(
      (,,,) -2, 2)
  return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<space
      >\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('
     .')-2, 2)
 return stridx('(){}[]''', pair) %
     2 == 0 && strlen(pair) == 2 ? "\<
     del>\<bs>" : "\<bs>"
endfunction
set nocompatible
    iMproved, required
                            " required
filetype on
filetype plugin on
filetype plugin indent on
syntax on
set splitright splitbelow
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
set expandtab
```

```
18
```

```
set autoindent
set cindent
set noerrorbells
set ruler
set guifont=*
set backspace=indent,eol,start
" set ignorecase
set incsearch
set nowrap
set hlsearch
" set termguicolors
set foldmethod=indent
set nofoldenable
" set cursorline
set laststatus=2
set showcmd
set wildmenu
" colorscheme torte
if !has('nvim')
  set clipboard=unnamedplus
endif
if !has('nvim')
 set ttymouse=xterm2
endif
nnoremap \langle S-j \rangle : m .+1 \langle CR \rangle ==
nnoremap \langle S-k \rangle : m .-2 \langle CR \rangle ==
vnoremap <S-j> :m '>+1<CR>gv==gv
vnoremap \langle S-k \rangle : m \langle -2 \langle CR \rangle gv == gv
nnoremap <A-h> <C-w>h
nnoremap <A-j> <C-w>j
nnoremap <A-k> <C-w>k
nnoremap <A-l> <C-w>l
let mapleader = ','
map <leader>cp :50 vsplit in<CR>:split
     out<CR><C-w>h
```

Equations and Formulas

9.1 Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1$$

and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

the size n+1 factors.

partitions of polygon into disjoint triangles each subset have pairwise distance at least

points on a circle to form n disjoint i.e. non-has been shown that these numbers satisfy, intersecting chords.

The number of rooted full binary trees with 9.4 Other Combinatorial Identities n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

Number of permutations of $1, \ldots, n$ that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n=3, these permutations are 132, 213, 231, 312 and 321.

9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count If $P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$, then, permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$
where, $S(0,0) = 1, S(n,0) = S(0,n) = 0$

$$\sum_{k=0}^{n} S(n,k) = n!$$

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k) \text{ , then,}$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 i_k.$$

9.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$ where S(0,0) = 1, S(n,0) = S(0,n) = 0 $S(n,2) = 2^{n-1} - 1 S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets,

with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = \sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$ $kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1) \qquad \sum_{k=1}^n x^{\gcd(k,n)} = \sum_{d|n} x^d \cdot \phi\left(\frac{n}{d}\right)$ Denote the n objects to partition by the noted $S^d(n,k)$, to be the number of ways $\phi(d)$ by using the diagonals). The number of ways to connect the 2n subset, it is required that $|i-j| \ge d$. It $\frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$

 $\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n}$ $n, r \in N, n > r, \sum_{i=1}^{n} {i \choose r} = {n+1 \choose r+1}$

If
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

f
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

Different Math Formulas

Picks Theorem: A = i + b/2 - 1 $d(i) = (i-1) \times$ Deragements: (d(i-1) + d(i-2))

$$\frac{n}{ab} \quad - \quad \left\{\frac{b\prime n}{a}\right\} \quad - \quad \left\{\frac{a\prime n}{b}\right\} \quad + \quad 1$$

9.6 GCD and LCM

if m is any integer, then $gcd(a + m \cdot b, b) =$ gcd(a,b)

The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$ $\gcd(a_1,b)\cdot\gcd(a_2,b).$

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).For non-negative integers a and b, where aand b are not both zero, $gcd(n^a-1, n^b-1) =$ $n^{\gcd(a,b)} - 1$

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\sum_{i=1}^{n} [\gcd(i, n) = k] = \phi\left(\frac{n}{k}\right)$$

9 Equations and Formulas 9.1 Catalan Numbers
$$C_n = \frac{1}{n+1} \binom{2n}{n} C_0 = 1, C_1 = \begin{cases} with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) = \frac{1}{n} \sum_{n=1}^{n} \frac{1}{n} \frac{1}{n} \cdot \phi(\frac{n}{d}) = \frac{1}{n} \sum_{n=1}^{n} \frac{1}{n} \frac{1}{n} \cdot \phi(\frac{n}{d}) = \frac{1}{n} \sum_{n=1}^{n} \frac{1}{n} \cdot \phi(\frac{n}{d}) = \frac{1}{n}$$$

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