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All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,
    ssse3,sse4,popcnt,abm,mmx,avx,tune=
    native")
#include <ext/pb_ds/assoc_container.hpp</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns
        iterator to the kth largest
        element counting from 0
   //order_of_key(val) --> returns the
        number of items in a set that
        are strictly smaller than our
template <typename DT>
using ordered_set = tree <DT, null_type</pre>
    , less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef SFT
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\',' << x <<</pre>
     '\'';}
void show(const string &x) {cerr << '\"</pre>
    ' << x << '\"';}
void show(bool x) {cerr << (x ? "true"</pre>
    : "false");}
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f';</pre>
    show(x.first); cerr << ", "; show(x void reset(int node) { lazy[node] =</pre>
    .second); cerr << '}'; }
template<typename T>
void show(T x) {int f = 0; cerr << "{";</pre>
     for (auto &i: x) cerr << (f++ ? ",</pre>
     " : ""), show(i); cerr << "}";}
void debug_out(string s) {
   cerr << '\n';</pre>
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(),
         '), s.end());
   cerr << "
                "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl:
   if(sizeof...(v)) debug_out(s, v...);
#define debug(x...) cerr << "LINE: " <<</pre>
     __LINE__ << endl; debug_out(#x, x) }
    ; cerr << endl;
#else
#define debug(x...)
#endif
```

```
const int RANDOM = chrono::
    high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) ^x;
unsigned hash_combine(unsigned a,
    unsigned b) { return a * 31 + b; }
struct chash {
 int operator()(int x) const { return
     hash_f(x); }
typedef gp_hash_table<int, int, chash>
    gp;
gp table;
    Data Structure
2.1 Segment Tree
const int N = 1000006;
```

```
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N];
LT lazy[4 * N];
int L, R;
void pull(int s, int e, int node) {
 val[node] = val[node << 1] + val[node</pre>
       << 1 | 1];
void apply(const LT &U, int s, int e,
    int node) {
 val[node] += (e - s + 1) * U;
 lazy[node] += U;
    None; }
DT merge(const DT &a, const DT &b) {
    return a + b; }
DT get(int s, int e, int node) { return int a[MAXN];
     val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
  apply(lazy[node], s, s + e >> 1, node
       << 1);
 apply(lazy[node], s + e + 2 \gg 1, e,
      node << 1 | 1);
 reset(node);
void build(int s, int e, vector<DT> &v,
     int node = 1) {
  int m = s + e >> 1;
  if (s == e) {
   val[node] = v[s];
   return;
 build(s, m, v, node * 2);
 build(m + 1, e, v, node * 2 + 1);
 pull(s, e, node);
void update(int S, int E, LT uval, int
    s = L, int e = R, int node = 1) {
 if (S > E) return;
```

if (S == s and E == e){

apply(uval, s, e, node);

```
return;
  }
  push(s, e, node);
  int m = s + e \gg 1;
  update(S, min(m, E), uval, s, m, node
       * 2):
  update(max(S, m + 1), E, uval, m + 1,
       e, node *2 + 1);
  pull(s, e, node);
DT query(int S, int E, int s = L, int e
     = R, int node = 1) {
  if (S > E) return I;
  if (s == S and e == E) return get(s,
      e, node);
  push(s, e, node);
  int m = s + e >> 1;
  DT L = query(S, min(m, E), s, m, node
       * 2);
  DT R = query(max(S, m + 1), E, m + 1,
       e, node *2 + 1);
  return merge(L, R);
void init(int _L, int _R, vector<DT> &v
 L = _L, R = _R;
  build(L, R, v);
```

2.2 Persistent Segment Tree

```
struct Node {
  Node *1, *r;
  int sum;
  Node(int val) : 1(nullptr), r(nullptr
       ), sum(val) {}
  Node(Node* 1, Node* r) : 1(1), r(r),
      sum(0) {
    if (1) sum += 1->sum;
    if (r) sum += r->sum;
  }
};
Node* root[MAXN];
Node* Build(int bg, int ed) {
  if (bg == ed) return new Node(a[bg]);
  int mid = (bg + ed) / 2;
  return new Node (Build (bg, mid), Build
       (mid + 1, ed));
int Query(Node* v, int bg, int ed, int
    1, int r) {
  if (1 > ed || r < bg) return 0;</pre>
  if (1 <= bg && ed <= r) return v->sum
  int mid = (bg + ed) / 2;
  return Query(v->1, bg, mid, 1, r) +
       Query(v->r, mid + 1, ed, l, r);
Node* Update(Node* v, int bg, int ed,
     int pos, int new_val) {
  if (bg == ed) return new Node(v->sum
       + new_val);
  int mid = (bg + ed) / 2;
```

if (pos <= mid)</pre>

```
node *root;
 implicit_segtree() {}
 implicit_segtree(int n) {
   root = new node(n);
 }
 void update(node *now, int L, int R,
      int idx, int val) {
    if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
         node(mid - L + 1);
   if (now->rt == NULL) now->rt = new
       node(R - mid);
   if (idx <= mid) update(now->lft, L,
       mid, idx, val);
   else update(now->rt, mid + 1, R, idx
       , val);
   now->val = (now->lft)->val + (now->
       rt)->val;
 }
 int query(node *now, int L, int R,
      int k) {
   if (L == R) return L;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new|}
         node(mid - L + 1);
   if (now->rt == NULL) now->rt = new
       node(R - mid);
   if (k <= (now->lft)->val) return
        query(now->lft, L, mid, k);
   else return query(now->rt, mid + 1,
       R, k - (now->lft)->val);
 }
};
```

```
2.4 Centroid Decomposition
// problem: Xenia and Tree
// Centroid Tree Implementation
#define lg(n) (31 - __builtin_clz(n))

const int N = 1e5 + 1;
const int K = lg(N) + 1;

int n;
vector<int> adj[N];
int anc[N][K], lvl[N], par[N], sz[N],
    vis[N];
int minD[N];
```

```
// Call dfs at the very outset
 void dfs(int u = 1, int p = 0) {
   sz[u] = 1:
   for (auto v : adj[u])
     if (v != p) {
       dfs_size(v, u);
       sz[u] += sz[v];
 }
 int findCentroid(int u, int p) {
   int total = sz[u];
   for (auto v : adj[u])
     if (v != p and not vis[v] and 2 *
         sz[v] > total) {
       sz[u] = total - sz[v];
       sz[v] = total;
      return findCentroid(v, u);
     }
   return u;
 int query(int u) {
   int ans = 1e6;
   for (int i = u; i; i = par[i])
     ans = min(ans, minD[i] + lca::dist
         (i, u));
   return ans;
 }
 void update(int u) {
   for (int i = u; i; i = par[i])
     minD[i] = min(minD[i], lca::dist(i
         , u));
 int decompose(int u, int p) {
   u = findCentroid(u, p);
   vis[u] = 1;
   for (auto v : adj[u])
     if (not vis[v])
       par[decompose(v, u)] = u;
   return u;
int main() {
 cin.tie(NULL)->sync_with_stdio(false)
 fill(all(minD), 1e9);
 int m;
 cin >> n >> m;
 for (int u, v, i = 1; i < n; i++) {</pre>
   cin >> u >> v;
   adj[u].push_back(v);
   adj[v].push_back(u);
 lca::init();
 CD::dfs();
 CD::decompose(1, 0);
 CD::update(1);
 int t, v;
 while (m--) {
   cin >> t >> v;
   if (t == 1) CD::update(v);
   else cout << CD::query(v) << '\n';</pre>
```

```
return 0;
}
```

2.5 DSU With Rollbacks

```
struct Rollback_DSU {
 int n;
 vector<int> par, sz;
  vector<pair<int, int>> op;
  Rollback_DSU(int n) : par(n), sz(n,
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par
        [node])
     ; // no path compression
   return node;
  void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)])
         swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back())
     par[op.back().second] = op.back().
     sz[op.back().first] -= sz[op.back
          ().second];
 }
};
```

2.6 BIT-2D

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j += -j & j)
        bit[x][j] += val;
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j & -j) ans
        += bit[x][j];
 return ans;
int get(int x1, int y1, int x2, int y2)
 return get(x2, y2) - get(x1 - 1, y2)
      - get(x2, y1 - 1) + get(x1 - 1,
      y1 - 1);
```

2.7 Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];

vector<LL> merge(vector<LL> v1, vector<
    LL> v2) {
  LL i = 0, j = 0;
  vector<LL> ret;
```

```
while (i < v1.size() || j < v2.size()</pre>
   if (i == v1.size()) {
     ret.push_back(v2[j]);
     j++;
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
   } else {
     if (v1[i] < v2[j]) {</pre>
       ret.push_back(v1[i]);
     } else {
       ret.push_back(v2[j]);
       j++;
 }
 return ret;
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 }
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
 Tree[node] = merge(Tree[leftNode],
      Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1,
    LL r, LL k) {
  if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin
        (), Tree[node].end(), k) -
          Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, 1, r,
       k) +
        query(rightNode, mid + 1, ed, 1
             , r, k);
```

2.8 MO with Update

```
bool operator<(const query &rhs)</pre>
      const {
   return (L < rhs.L) or (L == rhs.L
        and R < rhs.R) or
          (L == rhs.L and R == rhs.R
              and t < rhs.t);</pre>
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
 cnt[freq[arr[idx]]]++;
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].last;
 if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
 assert(0);
int main() {
 int n, q;
 cin >> n >> q;
 int counter = 0;
 map<int, int> M;
 for (int i = 1; i <= n; i++) {</pre>
   cin >> arr[i];
   if (!M[arr[i]]) M[arr[i]] = ++
        counter;
   arr[i] = M[arr[i]];
 }
 iota(id, id + N, 0);
 while (q--) {
   int tp, x, y;
   cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = {x, y, arr[x]};
     arr[x] = y;
 }
 t = ui;
 cnt[0] = 3 * n;
```

2.9 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
    for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
    for (int q = 1; (1 << q) <= N; q++)</pre>
      int add = 1 << (q - 1);
     for (int j = 1; j + add <= N; j++)</pre>
       M[i][j][0][q] = max(M[i][j][0][q]
             - 1], M[i][j + add][0][q -
            1]);
     }
   }
 }
  for (int p = 1; (1 << p) <= N; p++) {</pre>
    int add = 1 << (p - 1);
    for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q
          ++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p]
              - 1][q], M[i + add][j][p -
               1][q]);
     }
   }
 }
// returns max of all A[i][j], where x1
    =i<=x2 and y1<=j<=y2
```

int Query(int x1, int y1, int x2, int

int kX = log2(x2 - x1 + 1);

int kY = log2(y2 - y1 + 1);

int addX = 1 << kX;</pre>

int addY = 1 << kY;</pre>

y2) {

2.10 Sparse Table

```
// tested by: https://judge.yosupo.jp/
    problem/staticrmq
// O-based indexing, query finds in
    range [first, last]
#define lg(x) (31 - __builtin_clz(x))
const int N = 1e5 + 1;
const int K = lg(N);
int a[N], tr[N][K + 1];
namespace sparse_table {
 int f(int p1, int p2) { return min(p1
      , p2); }
 void build() {
   for(int i = 0; i < n; i++)</pre>
     tr[i][0] = a[i];
   for(int j = 1; j <= K; j++) {</pre>
     for(int i = 0; i + (1<<j) <= n; i</pre>
       tr[i][j] = f(tr[i][j - 1], tr[i
            + (1<<(j - 1))][j - 1]);
 }
 int query(int 1, int r) {
   int d = lg(r - 1 + 1);
   return f(table[1][d], table[r - (1<<||
        d) + 1][d]);
 }
}
```

3 DP

3.1 Convex Hull Trick

```
struct line {
           ll m, c;
           line() {}
          line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
           vector<line> lines;
           int ptr = 0;
           convex_hull_trick() {}
           bool bad(line a, line b, line c) {
                      return 1.0 * (c.c - a.c) * (a.m - b.
                                                m) < 1.0 * (b.c - a.c) * (a.m - a.c) * (a.c) * (
                                                       c.m);
           void add(line L) {
                     int sz = lines.size();
                      while (sz >= 2 && bad(lines[sz - 2],
                                                       lines[sz - 1], L)) {
                                lines.pop_back();
                      lines.pb(L);
```

```
ll get(int idx, int x) { return (111
      * lines[idx].m * x + lines[idx].c
      ): }
 ll query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines
        .size() - 1;
   while (ptr < lines.size() - 1 && get
        (ptr, x) > get(ptr + 1, x)) ptr
   return get(ptr, x);
 }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio;
 int t;
 cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >>
        sum[i]:
   for (int i = 1; i <= n; i++) dp[i] =</pre>
         0, sum[i] += sum[i - 1];
   convex_hull_trick cht;
    cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 1
         11 * a * sqr(sum[pos]) - c;
     cht.add(line(211 * a * sum[pos],
         dp[pos] - a * sqr(sum[pos]));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

3.2 Dynamic CHT

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 IS_QUERY = -(1LL << 62);</pre>
struct line {
 ll m, b;
  mutable function <const line*()> succ
  bool operator < (const line &rhs)</pre>
      const {
    if (rhs.b != IS_QUERY) return m <</pre>
        rhs.m;
   const line *s = succ();
   if (!s) return 0;
   11 x = rhs.m;
   return b - s -> b < (s -> m - m) * x
 }
};
struct HullDynamic : public multiset <</pre>
    line> {
```

bool bad (iterator y) {

```
auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
     return y -> m == z -> m && y -> b
          \langle = z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x
        -> m && y -> b <= x -> b;
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z
        -> m - y -> m) >= 1.0 * (y -> b)
         -z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 void insert_line (ll m, ll b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) ==
        end() ? 0 : &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(
        y))) erase(next(y));
    while (y != begin() && bad(prev(y)))
         erase(prev(y));
 ll eval (ll x) {
    auto 1 = *lower_bound((line) {x,
        IS_QUERY});
    return 1.m * x + 1.b;
 }
int main() {
 HullDynamic hull;
 hull.insert_line(1, 1);
 hull.insert_line(-1, 1);
  cout << hull.eval(69) << endl;</pre>
  cout << hull.eval(420) << endl;</pre>
  return 0:
```

3.3 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c)
LL calc(line L, LL x) { return 1LL * L.
    m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft =
      NULL, node *rt = NULL)
     : L(line(m, c)), lft(lft), rt(rt)
};
struct LiChao {
 node *root;
 LiChao() { root = new node(); }
 void update(node *now, int L, int R,
      line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(
        lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
```

```
if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
          node():
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo;
     if (now->lft == NULL) now->lft =
         new node();
     update(now->lft, L, mid, hi);
   }
 LL query(node *now, int L, int R, LL
      x) {
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x \le mid)
     return max(calc(now->L, x), query(
         now->lft, L, mid, x));
   else
     return max(calc(now->L, x), query(
         now->rt, mid + 1, R, x));
 }
};
```

SOS DP 3.4

```
for(int i = 0; i<(1<<N); ++i)</pre>
F[i] = A[i];
for(int i = 0;i < N; ++i) for(int mask</pre>
    = 0; mask < (1<<N); ++mask){
 if(mask & (1<<i))</pre>
 F[mask] += F[mask^(1<<i)];
```

Geometry

4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long
    for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ?</pre>
    0 : (x < 0 ? -1 : 1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y)
 Point operator+(const Point& u) const
      { return Point(x + u.x, y + u.y)
 Point operator-(const Point& u) const
      { return Point(x - u.x, y - u.y)
      ; }
 Point operator*(const LL u) const {
      return Point(x * u, y * u); }
 Point operator*(const Tf u) const {
     return Point(x * u, y * u); }
 Point operator/(const Tf u) const {
      return Point(x / u, y / u); }
 bool operator==(const Point& u) const
   return dcmp(x - u.x) == 0 && dcmp(y)
        - u.y) == 0;
 bool operator!=(const Point& u) const
       { return !(*this == u); }
```

```
bool operator<(const Point& u) const</pre>
        return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                  -u.x) == 0 && dcmp(y - u.y) <
   }
Ti dot(Point a, Point b) { return a.x *
            b.x + a.y * b.y; }
Ti cross(Point a, Point b) { return a.x | struct Segment {
            * b.y - a.y * b.x; }
Tf length(Point a) { return sqrt(dot(a,
            a)); }
Ti sqLength(Point a) { return dot(a, a)
          ; }
Tf distance(Point a, Point b) { return
          length(a - b); }
Tf angle(Point u) { return atan2(u.y, u
// returns angle between oa, ob in (-PI
Tf angleBetween(Point a, Point b) {
   Tf ans = angle(b) - angle(a);
    return ans <= -PI ? ans + 2 * PI : (
              ans > PI ? ans - 2 * PI : ans);
// Rotate a ccw by rad radians, Tf Ti
Point rotate(Point a, Tf rad) {
    return Point(a.x * cos(rad) - a.y *
              sin(rad),
                                a.x * sin(rad) + a.y * cos
                                           (rad)):
// rotate a ccw by angle th with cos(th
         ) = co \&\& sin(th) = si, tf ti same
Point rotatePrecise(Point a, Tf co, Tf
         si) {
    return Point(a.x * co - a.y * si, a.y
                * co + a.x * si);
Point rotate90(Point a) { return Point
          (-a.y, a.x); }
// scales vector a by s such that
         length of a becomes s, Tf Ti same
Point scale(Point a, Tf s) { return a /
            length(a) * s; }
 // returns an unit vector perpendicular
            to vector a, Tf Ti same
Point normal(Point a) {
   Tf l = length(a);
    return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if
          on ab && -1 if right of ab
int orient(Point a, Point b, Point c) {
            return dcmp(cross(b - a, c - a));
/// Use as sort(v.begin(), v.end(),
          polarComp(0, dir))
/// Polar comparator around O starting
          at direction dir
struct polarComp {
   Point O, dir;
    polarComp(Point 0 = Point(0, 0),
              Point dir = Point(1, 0)) : O(0),
              dir(dir) {}
    bool half(Point p) {
        return dcmp(cross(dir, p)) < 0 ||</pre>
                                                                                                 bool linesParallel(Line p, Line q) {
```

```
(dcmp(cross(dir, p)) == 0 \&\&
              dcmp(dot(dir, p)) > 0);
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) <</pre>
       make_tuple(half(q), cross(p, q)
 Point a, b;
 Segment(Point aa, Point bb) : a(aa),
     b(bb) {}
typedef Segment Line;
struct Circle {
 Point o;
 Tf r;
 Circle(Point o = Point(0, 0), Tf r =
      0) : o(o), r(r) {}
 // returns true if point p is in ||
     on the circle
 bool contains(Point p) { return dcmp(
      sqLength(p - o) - r * r) <= 0; }
 // returns a point on the circle rad
     radians away from +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value
   return Point(o.x + cos(rad) * r, o.y
        + sin(rad) * r);
 // area of a circular sector with
      central angle rad
 Tf area(Tf rad = PI + PI) { return
     rad * r * r / 2; }
 // area of the circular sector cut by
      a chord with central angle alpha
 Tf sector(Tf alpha) { return r * r *
     0.5 * (alpha - sin(alpha)); }
```

4.2 Linear

```
// **** LINE LINE INTERSECTION START
// returns true if point p is on
   segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p))
     == 0 && dcmp(dot(s.a - p, s.b - p
     )) <= 0;</pre>
// returns true if segment p && q touch
     or intersect
bool segmentsIntersect(Segment p,
   Segment q) {
 if (onSegment(p.a, q) || onSegment(p.
     b, q)) return true;
 if (onSegment(q.a, p) || onSegment(q.
     b, p)) return true;
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 &&
     dcmp(c3) * dcmp(c4) < 0;
```

```
return dcmp(cross(p.b - p.a, q.b - q. // returns the projection of point p on
                                                                                          ]));
                                              line 1, Tf Ti Same
                                                                                    return abs(res);
      a)) == 0:
                                         |Point projectPointLine(Point p, Line 1) |}
// lines are represented as a ray from
                                           Point v = 1.b - 1.a;
    a point: (point, vector)
                                                                                  // interior
                                           return 1.a + v * ((Tf)dot(v, p - 1.a)
// returns false if two lines (p, v) &&
                                                                                       ---> -2
     (q, w) are parallel or collinear
                                                / dot(v, v));
// true otherwise, intersection point
                                                                                       ----> -1
    is stored at o via reference, Tf Ti
                                                                                   // concentric
                                                                                                      (d = 0)
     Same
                                         4.3 Circular
                                                                                  // secants
                                                                                       ) ----> 0
bool lineLineIntersection(Point p,
                                          // Extremely inaccurate for finding
    Point v, Point q, Point w, Point& o
                                             near touches
    ) {
                                                                                       ----> 1
                                         // compute intersection of line 1 with
 if (dcmp(cross(v, w)) == 0) return
                                                                                   // exterior
                                             circle c
                                                                                       ----> 2
      false:
                                          // The intersections are given in order
 Point u = p - q;
                                              of the ray (l.a, l.b), Tf Ti same
 o = p + v * (cross(w, u) / cross(v, w)
                                                                                       Circle c2) {
                                         vector<Point> circleLineIntersection(
                                                                                    Tf d = length(c1.o - c2.o);
      ));
                                             Circle c, Line 1) {
 return true;
                                           vector<Point> ret;
                                           Point b = 1.b - 1.a, a = 1.a - c.o;
// returns false if two lines p && q
                                           Tf A = dot(b, b), B = dot(a, b);
    are parallel or collinear
                                           Tf C = dot(a, a) - c.r * c.r, D = B *
// true otherwise, intersection point
                                                B - A * C;
    is stored at o via reference
                                           if (D < -EPS) return ret;</pre>
bool lineLineIntersection(Line p, Line
                                           ret.push_back(l.a + b * (-B - sqrt(D
    q, Point& o) {
                                                                                        same
                                               + EPS)) / A);
 return lineLineIntersection(p.a, p.b
                                           if (D > EPS) ret.push_back(1.a + b *
      - p.a, q.a, q.b - q.a, o);
                                                                                       Circle c1, Circle c2) {
                                               (-B + sqrt(D)) / A);
                                                                                    vector<Point> ret;
                                           return ret;
// returns the distance from point a to |
                                                                                    Tf d = length(c1.o - c2.o);
     line 1
                                                                                    if (dcmp(d) == 0) return ret;
                                         // signed area of intersection of
// **** LINE LINE INTERSECTION FINISH
                                             circle(c.o, c.r) &&
    ****
                                                                                          ret;
                                          // triangle(c.o, s.a, s.b) [cross(a-o,
Tf distancePointLine(Point p, Line 1) {
                                             b-o)/2
 return abs(cross(1.b - 1.a, p - 1.a)
                                                                                        return ret;
                                         Tf circleTriangleIntersectionArea(
      / length(1.b - 1.a));
                                             Circle c, Segment s) {
                                                                                    Point v = c2.o - c1.o;
                                           using Linear::distancePointSegment;
// returns the shortest distance from
                                           Tf OA = length(c.o - s.a);
    point a to segment s
                                           Tf OB = length(c.o - s.b);
Tf distancePointSegment(Point p,
                                                                                         (v));
                                           // sector
    Segment s) {
                                           if (dcmp(distancePointSegment(c.o, s)
 if (s.a == s.b) return length(p - s.a
                                                -c.r) >= 0)
                                                                                          -si), c1.r) + c1.o;
      );
                                             return angleBetween(s.a - c.o, s.b -
 Point v1 = s.b - s.a, v2 = p - s.a,
                                                  c.o) * (c.r * c.r) / 2.0;
      v3 = p - s.b;
                                                                                          si), c1.r) + c1.o;
                                           // triangle
 if (dcmp(dot(v1, v2)) < 0)</pre>
                                           if (dcmp(OA - c.r) <= 0 && dcmp(OB -
   return length(v2);
                                                                                    ret.push_back(p1);
                                               c.r) <= 0)
 else if (dcmp(dot(v1, v3)) > 0)
                                             return cross(c.o - s.b, s.a - s.b) /
   return length(v3);
                                                                                    return ret;
                                                  2.0:
                                           // three part: (A, a) (a, b) (b, B)
   return abs(cross(v1, v2) / length(v1
                                           vector<Point> Sect =
       ));
                                                                                       circles c1, c2
                                               circleLineIntersection(c, s);
                                           return circleTriangleIntersectionArea
// returns the shortest distance from
                                                                                       c1, Circle c2) {
                                               (c, Segment(s.a, Sect[0])) +
                                                                                    Point AB = c2.o - c1.o;
    segment p to segment q
                                                 circleTriangleIntersectionArea(
Tf distanceSegmentSegment(Segment p,
                                                                                    Tf d = length(AB);
                                                      c, Segment(Sect[0], Sect
    Segment q) {
                                                      [1])) +
 if (segmentsIntersect(p, q)) return
                                                 circleTriangleIntersectionArea(
      0;
                                                                                        r * c1.r;
                                                      c, Segment(Sect[1], s.b));
 Tf ans = distancePointSegment(p.a, q)
                                         // area of intersecion of circle(c.o, c
 ans = min(ans, distancePointSegment(p
                                              .r) && simple polyson(p[])
      .b, q));
                                         Tf circlePolyIntersectionArea(Circle c,
 ans = min(ans, distancePointSegment(q
                                              Polygon p) {
                                                                                         ));
      .a, p));
                                           Tf res = 0;
 ans = min(ans, distancePointSegment(q
                                           int n = p.size();
      .b, p));
                                           for (int i = 0; i < n; ++i)
 return ans:
                                                                                         )):
                                                 circleTriangleIntersectionArea(
```

c, Segment(p[i], p[(i + 1) % n

```
// locates circle c2 relative to c1
                      (d < R - r)
// interior tangents (d = R - r)
                     (R - r < d < R + r)
// exterior tangents (d = R + r)
                      (d > R + r)
int circleCirclePosition(Circle c1,
 int in = dcmp(d - abs(c1.r - c2.r)),
      ex = dcmp(d - (c1.r + c2.r));
 return in < 0 ? -2 : in == 0 ? -1 :
      ex == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points
    between two circles c1 && c2, Tf Ti
vector<Point> circleCircleIntersection(
 if (dcmp(c1.r + c2.r - d) < 0) return
 if (dcmp(abs(c1.r - c2.r) - d) > 0)
 Tf co = (c1.r * c1.r + sqLength(v) -
      c2.r * c2.r) / (2 * c1.r * length
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co,
 Point p2 = scale(rotatePrecise(v, co,
 if (p1 != p2) ret.push_back(p2);
// intersection area between two
Tf circleCircleIntersectionArea(Circle
 if (d >= c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.</pre>
 if (d + c2.r <= c1.r) return PI * c2.</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d
       -c2.r*c2.r) / (2.0*c1.r*d
 Tf alpha2 = acos((c2.r * c2.r + d * d))
       -c1.r*c1.r) / (2.0*c2.r*d)
 return c1.sector(2 * alpha1) + c2.
      sector(2 * alpha2);
```

```
// returns tangents from a point p to
    circle c, Tf Ti same
vector<Point> pointCircleTangents(Point
     p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
   ;
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u,
        ang)};
 return ret;
// returns the points on tangents that
    touches the circle, Tf Ti Same
vector<Point> pointCircleTangencyPoints
    (Point p, Circle c) {
 Point u = p - c.o;
 Tf d = length(u);
 if (d < c.r)
   return {};
 else if (dcmp(d - c.r) == 0)
   return {c.o + u};
 else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o +
        rotate(u, ang)};
 }
// for two circles c1 && c2, returns
    two list of points a && b
// such that a[i] is on c1 && b[i] is
    c2 && for every i
// Line(a[i], b[i]) is a tangent to
    both circles
// CAUTION: a[i] = b[i] in case they
    touch \mid -1 for c1 = c2
int circleCircleTangencyPoints(Circle
    c1, Circle c2, vector<Point> &a,
                            vector<
                                 Point>
                                 &b) {
 a.clear(), b.clear();
 int cnt = 0;
 if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r +
      c2.r;
 if (dcmp(d2 - rdif * rdif) < 0)</pre>
      return 0:
 if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
      ) == 0) return -1;
 Tf base = angle(c2.o - c1.o);
 if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++:
   return cnt;
```

```
Tf ang = acos((c1.r - c2.r) / sqrt(d2)
 a.push_back(c1.point(base + ang));
 b.push_back(c2.point(base + ang));
 a.push_back(c1.point(base - ang));
 b.push_back(c2.point(base - ang));
 if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++;
 } else if (dcmp(d2 - rsum * rsum) >
   Tf ang = acos((c1.r + c2.r) / sqrt(
       d2)):
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang
   cnt++;
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang
   cnt++:
 }
 return cnt;
    Convex
/// minkowski sum of two polygons in O(
```

```
Polygon minkowskiSum(Polygon A, Polygon
     B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin
      (), A.end()), A.end());
 rotate(B.begin(), min_element(B.begin
      (), B.end()), B.end());
 A.push_back(A[0]);
 B.push_back(B[0]);
 for (int i = 0; i < n; i++) A[i] = A[</pre>
      i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[</pre>
      i + 1] - B[i];
 Polygon C(n + m + 1);
 C[0] = A.back() + B.back();
 merge(A.begin(), A.end() - 1, B.begin
      (), B.end() - 1, C.begin() + 1,
       polarComp(Point(0, 0), Point(0,
           -1)));
 for (int i = 1; i < C.size(); i++) C[</pre>
      i] = C[i] + C[i - 1];
 C.pop_back();
 return C;
 / finds the rectangle with minimum
    area enclosing a convex polygon and
// the rectangle with minimum perimeter
     enclosing a convex polygon
// Tf Ti Same
pair<Tf, Tf>
    rotatingCalipersBoundingBox(const
    Polygon &p) {
 using Linear::distancePointLine;
 int n = p.size();
```

int l = 1, r = 1, j = 1;

```
Tf area = 1e100;
  Tf perimeter = 1e100;
  for (int i = 0; i < n; i++) {</pre>
   Point v = (p[(i + 1) \% n] - p[i]) /
        length(p[(i + 1) % n] - p[i]);
   while (dcmp(dot(v, p[r % n] - p[i])
        - dot(v, p[(r + 1) \% n] - p[i])
        ) < 0)
     r++;
   while (j < r || dcmp(cross(v, p[j %</pre>
        n] - p[i]) -
                       cross(v, p[(j +
                            1) % n] - p[
                            i])) < 0)
     j++;
   while (1 < j ||
          dcmp(dot(v, p[1 % n] - p[i])
               - dot(v, p[(1 + 1) % n] -
               p[i])) > 0)
     1++;
   Tf w = dot(v, p[r \% n] - p[i]) - dot
        (v, p[1 % n] - p[i]);
   Tf h = distancePointLine(p[j % n],
        Line(p[i], p[(i + 1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2)
         * h);
 return make_pair(area, perimeter);
// returns the left side of polygon u
    after cutting it by ray a->b
Polygon cutPolygon(Polygon u, Point a,
    Point b) {
  using Linear::lineLineIntersection;
  using Linear::onSegment;
  Polygon ret;
  int n = u.size();
  for (int i = 0; i < n; i++) {</pre>
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0)
        ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0)
     Point t:
     lineLineIntersection(a, b - a, c,
          d - c, t);
     if (onSegment(t, Segment(c, d)))
          ret.push_back(t);
 }
 return ret;
^{\prime}/ returns true if point p is in or on
    triangle abc
bool pointInTriangle(Point a, Point b,
    Point c, Point p) {
  return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >=
      0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no
    three collinear points
// returns inside = -1, on = 0, outside
int pointInConvexPolygon(const Polygon
    &pt, Point p) {
  int n = pt.size();
```

```
assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p -
       pt[0])) > 0)
     lo = mid;
   else
     hi = mid;
 bool in = pointInTriangle(pt[0], pt[
      lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p
       - pt[lo - 1])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p -
      pt[lo])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) %
       n], p - pt[(hi + 1) % n])) == 0)
   return 0;
 return -1;
// Extreme Point for a direction is the
     farthest point in that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly,
    Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1)
       % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 +
            n) % n], u)) >= 0) {
     return c;
   bool a_up = dcmp(dot(poly[(a + 1) %
       n] - poly[a], u)) >= 0;
   bool c_up = dcmp(dot(poly[(c + 1) %
       n] - poly[c], u)) >= 0;
   bool a_above_c = dcmp(dot(poly[a] -
       poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c:
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c:
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 if (dcmp(dot(poly[a] - poly[(a + 1) %
       n], u)) > 0 &&
         ) % n], u)) > 0)
                                           return orient(Q, u, v) != -dir;
   return a;
```

```
return b % n;
                                   // For a convex polygon p and a line l,
                                        returns a list of segments
                                   // of p that touch or intersect line 1.
                                   // the i'th segment is considered (p[i
                                        ], p[(i + 1) modulo |p|])
                                    // #1 If a segment is collinear with
                                        the line, only that is returned
                                    // #2 Else if l goes through i'th point
                                        , the i'th segment is added
                                    // Complexity: O(lg |p|)
                                   vector<int> lineConvexPolyIntersection(
                                        const Polygon &p, Line 1) {
                                     assert((int)p.size() >= 3);
                                     assert(1.a != 1.b);
                                     int n = p.size();
                                     vector<int> ret;
                                     Point v = 1.b - 1.a;
                                     int lf = extremePoint(p, rotate90(v))
                                     int rt = extremePoint(p, rotate90(v)
                                         * Ti(-1)):
                                     int olf = orient(l.a, l.b, p[lf]);
                                     int ort = orient(l.a, l.b, p[rt]);
                                     if (!olf || !ort) {
                                       int idx = (!olf ? lf : rt);
                                       if (orient(l.a, l.b, p[(idx - 1 + n)
                                            % n]) == 0)
                                         ret.push_back((idx - 1 + n) \% n);
                                         ret.push_back(idx);
                                       return ret:
                                     if (olf == ort) return ret;
                                     for (int i = 0; i < 2; ++i) {
                                       int lo = i ? rt : lf;
                                       int hi = i ? lf : rt;
                                       int olo = i ? ort : olf;
                                       while (true) {
                                         int gap = (hi - lo + n) \% n;
                                         if (gap < 2) break;</pre>
                                         int mid = (lo + gap / 2) % n;
                                         int omid = orient(l.a, l.b, p[mid
                                             ]);
                                         if (!omid) {
                                          lo = mid;
                                          break;
                                         if (omid == olo)
                                           lo = mid;
                                         else
                                          hi = mid;
                                       ret.push_back(lo);
                                     return ret;
                                    // Calculate [ACW, CW] tangent pair
                                        from an external point
                                   constexpr int CW = -1, ACW = 1;
dcmp(dot(poly[a] - poly[(a - 1 + n bool isGood(Point u, Point v, Point Q,
                                        int dir) {
```

```
Point better (Point u, Point v, Point Q,
     int dir) {
 return orient(Q, u, v) == dir ? u : v
Point pointPolyTangent(const Polygon &
    pt, Point Q, int dir, int lo, int
    hi) {
  while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid -
        1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid +
        1], Q, dir);
   if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q,
          dir, mid + 1, hi);
     Point p2 = pointPolyTangent(pt, Q,
          dir, lo, mid - 1);
     return better(p1, p2, Q, dir);
   if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) ==
         dir)
       hi = mid - 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) ==
         dir)
       lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
 }
 Point ret = pt[lo];
  for (int i = lo + 1; i <= hi; i++)</pre>
      ret = better(ret, pt[i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(
    const Polygon &pt, Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt,
      Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q
      , CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
4.5 Polygon
```

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &
    poly) {
 Polygon ret;
```

```
int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0
        && (ret.empty() || b != ret.
        back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p
     of n vertices
Tf signedPolygonArea(const Polygon &p)
    {
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() -</pre>
      1; i++)
   ret += cross(p[i] - p[0], p[i + 1] -
        p[0]);
 return ret / 2;
// given a polygon p of n vertices,
    generates the convex hull in in CCW
// Tested on https://acm.timus.ru/
    problem.aspx?space=1&num=1185
// Caution: when all points are
    colinear AND removeRedundant ==
// output will be contain duplicate
    points (from upper hull) at back
Polygon convexHull(Polygon p, bool
    removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p
      .end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
          dcmp(cross(ch[m - 1] - ch[m -
               2], p[i] - ch[m - 1]))
              <= check)
     m--:
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i >= 0; i--) {
   while (m > k \&\&
          dcmp(cross(ch[m - 1] - ch[m -
               2], p[i] - ch[m - 2]))
              <= check)
     m--:
   ch[m++] = p[i];
 if (n > 1) m--;
 ch.resize(m);
 return ch;
// returns inside = -1, on = 0, outside
int pointInPolygon(const Polygon &p,
    Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
```

```
int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])
       ) || o == p[i]) return 0;
   int k = dcmp(cross(p[j] - p[i], o -
       p[i]));
   int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn
   if (k < 0 && d2 <= 0 && d1 > 0) wn
 return wn ? -1 : 1;
// Given a simple polygon p, and a line
     1, returns (x, y)
// x = longest segment of 1 in p, y =
    total length of 1 in p.
pair<Tf, Tf> linePolygonIntersection(
    Line 1, const Polygon &p) {
 using Linear::lineLineIntersection;
 int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {</pre>
   Point a = p[i], b = p[(i + 1) \% n],
       z = p[(i - 1 + n) \% n];
   int ora = orient(l.a, l.b, a), orb =
        orient(l.a, l.b, b),
       orz = orient(l.a, l.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(
           d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins;
     lineLineIntersection(1, Line(a, b)
         , ins);
     ev.emplace_back(dot(ins - 1.a, 1.b
          - 1.a), 0);
 }
 sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot =
 bool active = false;
 int sign = 0;
 for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment
        is (last, d)
                    /// On Border
   if (sign) {
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
   } else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
     if (tp == 0)
```

```
active = !active;
else
    sign = tp;
}
last = d;
if (!active) len = 0;
}
ans /= length(1.b - 1.a);
tot /= length(1.b - 1.a);
return {ans, tot};
}
```

5 Graph

5.1 LCA, ETT, VT

```
struct lca_table {
 tree &T;
 int n, LOG = 20;
 vector<vector<int>> anc;
 vector<int> level;
 void setupLifting(int node, int par)
   for (int v : T[node])
     if (v != par) {
       anc[v][0] = node, level[v] =
           level[node] + 1;
       for (int k = 1; k < LOG; k++)</pre>
           anc[v][k] = anc[anc[v][k -
           1]][k - 1];
       setupLifting(v, node);
 lca_table(tree &T, int root = 0) : T(
     T), n(T.n) {
   LOG = 33 - \_builtin_clz(n);
   anc.assign(n, vector<int>(LOG, root)
   level.resize(n);
   setupLifting(root, root);
 int lca(int u, int v) {
   if (level[u] > level[v]) swap(u, v);
   for (int k = LOG - 1; ~k; k--)
     if (level[u] + (1 << k) <= level[v</pre>
         ]) v = anc[v][k];
   if (u == v) return u;
   for (int k = LOG - 1; ~k; k--)
     if (anc[u][k] != anc[v][k]) u =
         anc[u][k], v = anc[v][k];
   return anc[u][0];
 int getAncestor(int node, int ht) {
   for (int k = 0; k < LOG; k++)
     if (ht & (1 << k)) node = anc[node</pre>
         ][k];
   return node;
 int distance(int u, int v) {
   int g = lca(u, v);
   return level[u] + level[v] - 2 *
       level[g];
 }
struct euler_tour {
 int time = 0;
 tree &T;
 int n;
 vector<int> start, finish, level, par
 euler_tour(tree &T, int root = 0)
```

```
: T(T), n(T.n), start(n), finish(n 5.2 Euler Tour on Edge
          ), level(n), par(n) {
   time = 0:
   call(root);
 void call(int node, int p = -1) {
   if (p != -1) level[node] = level[p]
        + 1;
   start[node] = time++;
   for (int e : T[node])
     if (e != p) call(e, node);
   par[node] = p;
   finish[node] = time++;
 }
 bool isAncestor(int node, int par) {
   return start[par] <= start[node] and</pre>
         finish[par] >= finish[node];
 int subtreeSize(int node) { return
      finish[node] - start[node] + 1 >>
       1; }
};
tree virtual_tree(vector<int> &nodes,
    lca_table &table, euler_tour &tour)
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
           start[x] < tour.start[y]; });</pre>
  int n = nodes.size();
 for (int i = 0; i + 1 < n; i++)</pre>
   nodes.push_back(table.lca(nodes[i],
        nodes[i + 1]));
 sort(nodes.begin(), nodes.end());
 nodes.erase(unique(nodes.begin(),
      nodes.end()), nodes.end());
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.
           start[x] < tour.start[y]; });</pre>
 n = nodes.size():
 stack<int> st;
 st.push(0);
 tree ans(n);
 for (int i = 1; i < n; i++) {</pre>
   while (!tour.isAncestor(nodes[i],
        nodes[st.top()])) st.pop();
   ans.addEdge(st.top(), i);
   st.push(i);
 return ans;
set<int> getCenters(tree &T) {
 int n = T.n;
 vector<int> deg(n), q;
 set<int> s;
 for (int i = 0; i < n; i++) {</pre>
   deg[i] = T[i].size();
   if (deg[i] == 1) q.push_back(i);
   s.insert(i):
 for (vector<int> t; s.size() > 2; q =
   for (auto x : q) {
     for (auto e : T[x])
       if (--deg[e] == 1) t.push_back(e
           ):
     s.erase(x);
   }
 return s;
```

```
// for simplicity, G[idx] contains the
    adjacency list of a node
^{\prime\prime} while G(e) is a reference to the e-
    th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0;
void dfs(graph &G, int node, int par) {
 out[node] = t;
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node);
   bck[e] = t++;
 in[node] = t - 1;
void init(graph &G, int node) {
 dfs(G, node, node);
5.3 HLD
const int N = 1e6 + 7;
template <typename DT>
struct Segtree {
 // write lazy segtree here
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
 depth[u] = depth[p] + 1;
 for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
   if (subsz > mxsz) heavy[u] = v, mxsz
        = subsz;
   SZ += subsz;
 }
 return SZ;
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u
      ], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
int query(int a, int b) {
 int ret = 0;
 for (; head[a] != head[b]; b = par[
      head[b]]) {
```

if (depth[head[a]] > depth[head[b]])

ret += tree.query(1, 0, cnt, pos[

if (depth[a] > depth[b]) swap(a, b);

swap(a, b);

head[b]], pos[b]);

```
ret += tree.query(1, 0, cnt, pos[a],
      pos[b]);
 return ret;
    SCC
5.4
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
int root[N];
vector<int> adj[N], adjr[N], adj_scc[N
vector<int> order, component, roots;
// tp = 0 ,finding topo order, tp = 1 ,
     reverse edge traversal
void dfs(int u, int tp = 0) {
 vis[u] = true;
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
 reverse(order.begin(), order.end());
 for (int i : order) {
   if (!vis[i]) {
     dfs(i, 1);
     int rt = component[0];
     for (auto j : component)
       root[j] = rt;
     roots.push_back(rt);
     component.clear();
 for (int u = 1; u <= n; u++)</pre>
   for (auto v : adj[u]) {
     int rtu = root[u],
         rtv = root[v];
     if (rtu != rtv)
       adj_scc[rtu].push_back(rtv);
5.5 LCA in O(1)
/* LCA in O(1)
```

```
* depth calculates weighted distance
 * level calculates distance by number
     of edges
 * Preprocessing in NlongN */
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG
      [1] = 0;
```

```
for (int i = 2; i <= n; i++) LOG[i] = /// flow with demand(lower bound) only
       LOG[i >> 1] + 1;
 for (int i = 0; i <= n; i++) adj[i].</pre>
      clear():
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
 if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v])</pre>
       return v:
 int 1 = LOG[id[v] - id[u] + 1];
  int p1 = id[u], p2 = id[v] - (1 << 1)</pre>
  int d1 = level[table[1][p1]], d2 =
      level[table[1][p2]];
 if (d1 < d2)
   return par[table[1][p1]];
 else
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[
      1] * 2));
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time:
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table</pre>
      [0][i] = a[i];
 for (int 1 = 0; 1 < L - 1; 1++) {</pre>
   for (int i = 1; i <= n; i++) {</pre>
     table[1 + 1][i] = table[1][i];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 <<</pre>
           1)]] < level[table[1][i]];</pre>
     if (C1 && C2) table[l + 1][i] =
          table[1][i + (1 << 1)];
   }
 }
}
     Dinic Max Flow
                                          Ti maxFlow(int s, int t) {
```

```
for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u
    1))
// add_edge(u, new sink, sum(out_demand
    [u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then
     demand satisfied
// flow in every edge i = demand[i] + e |}
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
 edge(int v, int u, Ti cap) : v(v), u(
      u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.
      empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level
          [ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q
            .push(ed.u);
   }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].</pre>
      size(); cid++) {
   int id = adj[v][cid];
   auto &ed = edges[id];
   if (level[v] + 1 != level[ed.u] ||
        ed.cap - ed.flow < 1) continue;</pre>
   Ti tr = dfs(ed.u, min(pushed, ed.cap
         - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr;
   edges[id ^ 1].flow -= tr;
   return tr;
 }
 return 0;
void init(int nodes) {
 m = 0, n = nodes;
 for (int i = 0; i < n; i++) level[i]</pre>
      = -1, ptr[i] = 0, adj[i].clear();
void addEdge(int v, int u, Ti cap) {
 edges.emplace_back(v, u, cap), adj[v
      ].push_back(m++);
 edges.emplace_back(u, v, 0), adj[u].
      push_back(m++);
```

```
Ti f = 0;
for (auto &ed : edges) ed.flow = 0;
for (; bfs(s, t); memset(level, -1, n
     * 4)) {
 for (memset(ptr, 0, n * 4); Ti
      pushed = dfs(s, INF, t); f +=
      pushed)
}
return f;
```

```
5.7 Min Cost Max Flow
mt19937 rnd(chrono::steady_clock::now()
    .time_since_epoch().count());
const LL inf = 1e9;
struct edge {
  int v, rev;
  LL cap, cost, flow;
  edge() {}
  edge(int v, int rev, LL cap, LL cost)
      : v(v), rev(rev), cap(cap), cost(
          cost), flow(0) {}
};
struct mcmf {
  int src, sink, n;
  vector<int> par, idx, Q;
  vector<bool> inq;
  vector<LL> dis;
  vector<vector<edge>> g;
  mcmf() {}
  mcmf(int src, int sink, int n)
      : src(src),
        sink(sink),
       n(n),
       par(n),
       idx(n),
       inq(n),
       dis(n),
       g(n),
       Q(10000005) {} // use Q(n) if
            not using random
  void add_edge(int u, int v, LL cap,
      LL cost, bool directed = true) {
    edge _u = edge(v, g[v].size(), cap,
        cost);
    edge _v = edge(u, g[u].size(), 0, -
        cost);
    g[u].pb(_u);
    g[v].pb(_v);
    if (!directed) add_edge(v, u, cap,
        cost, true);
  bool spfa() {
    for (int i = 0; i < n; i++) {</pre>
     dis[i] = inf, inq[i] = false;
    int f = 0, 1 = 0;
    dis[src] = 0, par[src] = -1, Q[1++]
        = src, inq[src] = true;
    while (f < 1) {</pre>
      int u = Q[f++];
      for (int i = 0; i < g[u].size(); i</pre>
          ++) {
        edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
        if (dis[e.v] > dis[u] + e.cost)
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
```

```
if (!inq[e.v]) inq[e.v] = true
             , Q[1++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
         // if (f && rnd() & 7) Q[--f]
              = e.v;
         // else Q[1++] = e.v;
        // }
      }
     inq[u] = false;
   return (dis[sink] != inf);
 pair<LL, LL> solve() {
   LL mincost = 0, maxflow = 0;
   while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
         = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.
           cap - e.flow);
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
         = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck
     }
     mincost += bottleneck * dis[sink],
          maxflow += bottleneck;
   return make_pair(mincost, maxflow);
};
// want to minimize cost and don't care
     about flow
// add edge from sink to dummy sink (
    cap = inf, cost = 0)
// add edge from source to sink (cap =
    inf, cost = 0)
// run mcmf, cost returned is the
    minimum cost
```

5.8 Bridge Tree

```
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G,
    int par = -1, int d = 0) {
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1)
       if (low[to] > depth[node]) {
         bridges.emplace_back(node, to)
         components.push_back({});
```

```
for (int x = -1; x != to; x =
              st.top(), st.pop())
           components.back().push_back(
                st.top());
       }
      }
      low[node] = min(low[node], low[to
    }
  }
  if (par == -1) {
    components.push_back({});
    while (!st.empty()) components.back
         ().push_back(st.top()), st.pop
         ();
  }
graph &create_tree() {
  for (auto &comp : components) {
    int idx = tree.addNode();
    for (auto &e : comp) id[e] = idx;
  }
  for (auto &[l, r] : bridges) tree.
      addEdge(id[l], id[r]);
  return tree;
void init(graph &G) {
  int n = G.n;
  depth.assign(n, -1), id.assign(n, -1)
       , low.resize(n);
  for (int i = 0; i < n; i++)</pre>
    if (depth[i] == -1) find_bridges(i,
        G);
```

5.9Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
 }
 if (!cnt) return 1;
 sort(vs.begin(), vs.end());
 string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] =
      ++ind;
 int ret = mp[s];
 return ret;
```

Math

6.1 Linear Sieve

```
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5],
      cnt[N + 5], POW[N + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
```

```
fill(prime + 2, prime + N + 1, 1);
SOD[1] = NOD[1] = phi[1] = spf[1] =
for (LL i = 2; i <= N; i++) {</pre>
 if (prime[i]) {
   primes.push_back(i), spf[i] = i;
   phi[i] = i - 1;
   NOD[i] = 2, cnt[i] = 1;
   SOD[i] = i + 1, POW[i] = i;
 for (auto p : primes) {
   if (p * i > N or p > spf[i]) break
   prime[p * i] = false, spf[p * i] =
   if (i % p == 0) {
     phi[p * i] = p * phi[i];
     NOD[p * i] = NOD[i] / (cnt[i] +
          1) * (cnt[i] + 2),
            cnt[p * i] = cnt[i] + 1;
     SOD[p * i] = SOD[i] / SOD[POW[i
          ]] * (SOD[POW[i]] + p * POW[
          i]),
            POW[p * i] = p * POW[i];
     break;
   } else {
     phi[p * i] = phi[p] * phi[i];
     NOD[p * i] = NOD[p] * NOD[i],
         cnt[p * i] = 1;
     SOD[p * i] = SOD[p] * SOD[i],
         POW[p * i] = p;
   }
 }
}
```

```
6.2 Pollard Rho
LL mul(LL a, LL b, LL mod) {
  return (__int128)a * b % mod;
  // LL ans = a * b - mod * (LL) (1.L /
       mod * a * b);
  // return ans + mod * (ans < 0) - mod
       * (ans >= (LL) mod);
LL bigmod(LL num, LL pow, LL mod) {
  LL ans = 1;
  for (; pow > 0; pow >>= 1, num = mul(
      num, num, mod))
    if (pow & 1) ans = mul(ans, num, mod
        );
  return ans;
bool is_prime(LL n) {
  if (n < 2 or n % 6 % 4 != 1) return (</pre>
      n \mid 1) == 3;
  LL a[] = \{2, 325, 9375, 28178,
      450775, 9780504, 1795265022};
  LL s = \_builtin\_ctzll(n - 1), d = n
      >> s;
  for (LL x : a) {
   LL p = bigmod(x % n, d, n), i = s;
    for (; p != 1 and p != n - 1 and x %
         n and i--; p = mul(p, p, n)
    if (p != n - 1 \text{ and } i != s) \text{ return}
        false;
  return true;
LL get_factor(LL n) {
```

```
auto f = [&](LL x) { return mul(x, x, |void mobius() {
      n) + 1; };
 LL x = 0, y = 0, t = 0, prod = 2, i =
       2, q;
 for (; t++ % 40 or gcd(prod, n) == 1;
       x = f(x), y = f(f(y))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) -
       min(x, y), n) ? q : prod;
 }
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11,
      13, 17, 19, 23, 29, 31, 37, 41,
                     43, 47, 53, 59,
                         61, 67, 71,
                         73, 79, 83,
                         89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
 auto _factor = [&](LL n, auto &
      _factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
   }
 }:
 _factor(n, _factor);
 return res;
```

6.3 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
 if (b == 0)
   return {1, 0, a};
 else {
   auto [x, y, g] = EGCD(b, a \% b);
   return {y, x - a / b * y, g};
 }
// given modulo equations, will apply
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1;
 for (auto &[v, m] : v) { // value %
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1,
   V += x * (v - V) / g % (m / g) * M,
       M *= m / g;
   V = (V \% M + M) \% M;
 }
 return make_pair(V, M);
```

6.4 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
```

```
memset(mob, -1, sizeof mob);
 mob[1] = 1;
 for (int i = 2; i < N; i++)</pre>
   if (mob[i]) {
     for (int j = i + i; j < N; j += i)
           mob[j] -= mob[i];
6.5 FFT
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N),
        sin(2 * PI * i / N));
   wp[1][i] = CD(cos(2 * PI * i / N), -
        sin(2 * PI * i / N));
 }
void fft(vector<CD> &v, bool invert =
    false) {
 if (v.size() != perm.size())
      precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i</pre>
        ]]);
 for (int len = 2; len <= N; len *= 2)</pre>
   for (int i = 0, d = N / len; i < N;</pre>
        i += len) {
     for (int j = 0, idx = 0; j < len /
           2; j++, idx += d) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j]
             + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /=</pre>
void pairfft(vector<CD> &a, vector<CD>
    &b, bool invert = false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[</pre>
      i] + b[i] * CD(0, 1);
 fft(p, invert);
 p.push_back(p[0]);
```

for (int i = 0; i < N; i++) {</pre>

```
if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) *
          CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) *
          CD(0, -0.5);
 }
vector<LL> multiply(const vector<LL> &a
    , const vector<LL> &b) {
 int n = 1;
  while (n < a.size() + b.size()) n <<=
  vector<CD> fa(a.begin(), a.end()), fb
      (b.begin(), b.end());
  fa.resize(n);
  fb.resize(n);
  //
           fft(fa); fft(fb);
  pairfft(fa, fb);
  for (int i = 0; i < n; i++) fa[i] =</pre>
      fa[i] * fb[i];
 fft(fa, true);
  vector<LL> ans(n);
  for (int i = 0; i < n; i++) ans[i] =</pre>
      round(fa[i].real());
  return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a,
    const vector<LL> &b) {
 int n = 1:
  while (n < a.size() + b.size()) n <<=</pre>
  vector<CD> al(n), ar(n), bl(n), br(n)
  for (int i = 0; i < a.size(); i++) al</pre>
      [i] = a[i] % M / B, ar[i] = a[i]
      % M % B;
  for (int i = 0; i < b.size(); i++) bl</pre>
      [i] = b[i] % M / B, br[i] = b[i]
      % M % B;
  pairfft(al, ar);
 pairfft(bl, br);
           fft(al); fft(ar); fft(bl);
      fft(br);
  for (int i = 0; i < n; i++) {</pre>
   CD ll = (al[i] * bl[i]), lr = (al[i]
         * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i]
         * br[i]);
   al[i] = 11;
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr;
 pairfft(al, ar, true);
 pairfft(bl, br, true);
           fft(al, true); fft(ar, true)
      ; fft(bl, true); fft(br, true);
  vector<LL> ans(n);
  for (int i = 0; i < n; i++) {</pre>
   LL right = round(br[i].real()), left
         = round(al[i].real());
   LL mid = round(round(bl[i].real()) +
         round(ar[i].real()));
```

```
return ans;
6.6 NTT
const LL N = 1 << 18;</pre>
const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n,
LL Pow(LL b, LL p) {
 LL ret = 1;
  while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) % MOD;
   p >>= 1;
 }
 return ret;
LL primitive_root(LL p) {
  vector<LL> factor;
  LL phi = p - 1, n = phi;
  for (LL i = 2; i * i <= n; i++) {</pre>
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
  }
 if (n > 1) factor.emplace_back(n);
  for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() &&</pre>
        ok; i++)
     ok &= Pow(res, phi / factor[i]) != struct aho_corasick {
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - __builtin_clz(n));
  LL r = Pow(g, (MOD - 1) / n);
  inv_n = Pow(n, MOD - 2);
  w[0] = w[n] = 1;
  for (LL i = 1; i < n; i++) w[i] = (w[</pre>
      i - 1] * r) % MOD;
  for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i &
        1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)</pre>
   if (i < rev[i]) swap(a[i], a[rev[i</pre>
        ]]);
  for (LL m = 2; m <= n; m <<= 1) {</pre>
   for (LL i = 0; i < n; i += m) {</pre>
     for (LL j = 0; j < (m >> 1); j++)
       LL &u = a[i + j], &v = a[i + j +
             (m >> 1)];
       LL t = v * w[dir ? n - n / m * j]
             : n / m * j] % MOD;
       v = u - t < 0 ? u - t + MOD : u
       u = u + t >= MOD ? u + t - MOD :
             u + t;
   }
```

ans[i] = ((left % M) * B * B + (mid)

% M) * B + right) % M;

```
if (dir)
   for (LL i = 0; i < n; i++) a[i] = (</pre>
        inv_n * a[i]) % MOD;
vector<LL> mul(vector<LL> p, vector<LL>
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i
      ];
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
  for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
 for (LL i = 0; i < sz; i++) a[i] = (a
      [i] * b[i]) % MOD;
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
  while (c.size() && c.back() == 0) c.
      pop_back();
 return c;
```

for (char c : s) {

now = now->to[i];

if (!now->to[i]) now->to[i] =

now->leaf = true, now->val++;

for (q.push(root); q.empty(); q.pop

cur->out = link->leaf ? link :

for (auto& next : cur->to) {

if (next != NULL) {

q.push(next);

cur->val += link->val;

node *cur = q.front(), *link = cur

next->link = cur != root ?

next = link->to[idx++];

link->to[idx++] : root;

add_node(now, c, now->len + 1)

int i = f(c);

void push_links() {

queue<node*> q;

()) {

->link:

int idx = 0;

} else

}

}

link->out;

```
7.2 Double hash
// define +, -, * for (PLL, LL) and (
    PLL, PLL), % for (PLL, PLL);
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p & 1) ans = ans * a % mod;
 return ans;
PLL inverse(PLL a) { return power(a, (
    mod.ff - 1) * (mod.ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = PLL(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (</pre>
      P[i - 1] * base) % mod;
PLL append(PLL cur, char c) { return (
    cur * base + c) % mod; }
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) {
    return (P[k] * c + cur) % mod; }
/// replaces the i-th (0-indexed)
    character from right from a to b;
PLL replace(PLL cur, int i, char a,
    char b) {
  cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
```

String

Aho Corasick

```
const int sg = 26, N = 1e3 + 9;
 struct node {
   node *link, *out, *par;
   bool leaf;
   LL val;
   int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '
       $', int len = 0)
       : par(par), p_ch(p_ch), len(len)
     val = leaf = cnt = last = 0;
     link = out = NULL;
   }
 };
 vector<node> trie;
 node* root:
 aho corasick() {
   trie.reserve(N), trie.emplace_back()
   root = &trie[0];
   root->link = root->out = root;
 inline int f(char c) { return c - 'a'
 inline node* add_node(node* par =
     NULL, char p_ch = '$', int len =
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
 void add_str(const string& s, LL val
     = 1) {
   node* now = root;
```

```
/// Erases c from the back of the
    string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod
       + mod) % mod;
/// Erases c from front of the string
    with size len
PLL pop_front(PLL hash, int len, char c
    ) {
 return ((hash - P[len - 1] * c) % mod
       + mod) % mod;
/// concatenates two strings where
    length of the right is k
PLL concat(PLL left, PLL right, int k)
    { return (left * P[k] + right) %
    mod: }
/// Calculates hash of string with size
     len repeated cnt times
/// This is O(\log n). For O(1), pre-
    calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) *
      inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff
 if (P[len].ss == 1) ret.ss = hash.ss
      * cnt;
 return ret;
LL get(PLL hash) { return ((hash.ff <<
    32) ^ hash.ss); }
struct hashlist {
 int len;
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.
       resize(len + 2, PLL(0, 0));
   for (int i = 1; i <= len; i++) H[i]</pre>
       = append(H[i - 1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i]
       = append(R[i + 1], s[i - 1]);
 /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - 1 +
        1]) % mod + mod) % mod;
 PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 +
        1]) % mod + mod) % mod;
 PLL concat_range_hash(int 11, int r1,
       int 12, int r2) {
   return concat(range_hash(l1, r1),
       range_hash(12, r2), r2 - 12 +
 PLL concat_reverse_hash(int l1, int
      r1, int 12, int r2) {
   return concat(reverse_hash(12, r2),
       reverse_hash(l1, r1), r1 - l1 +
        1);
 }
```

7.4 Manacher's

```
vector<int> d1(n);
// d[i] = number of palindromes taking
    s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i /*
    ++) {
 int k = (i > r) ? 1 : min(d1[1 + r -
      i], r - i + 1);
 while (0 <= i - k && i + k < n && s[i
       - k] == s[i + k]) k++;
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking
    s[i-1] and s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 0 : min(d2[1 + r -
      i + 1], r - i + 1);
 while (0 <= i - k - 1 && i + k < n &&
      s[i - k - 1] == s[i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i +
```

7.5 String Match FFT

```
//find occurrences of t in s where '?'s
     are automatically matched with any
     character
//res[i + m - 1] = sum_j=0 to m - 1_{s[}
    i + j] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
 vector<int> s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[</pre>
      i] == '?' ? 0 : s[i] - 'a' + 1;
      //assign any non zero number for
      non '?'s
 for(int i = 0; i < n; i++) s2[i] = s1</pre>
      [i] * s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1</pre>
      [i] * s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[</pre>
      i] == '?' ? 0 : t[i] - 'a' + 1;
 for(int i = 0; i < m; i++) t2[i] = t1</pre>
      [i] * t1[i];
  for(int i = 0; i < m; i++) t3[i] = t1</pre>
      [i] * t2[i];
 reverse(t1.begin(), t1.end());
```

```
reverse(t2.begin(), t2.end());
reverse(t3.begin(), t3.end());
vector<int> s1t3 = multiply(s1, t3);
vector<int> s2t2 = multiply(s2, t2);
vector<int> s3t1 = multiply(s3, t1);
vector<int> res(n);
for(int i = 0; i < n; i++) res[i] =
    s1t3[i] - s2t2[i] * 2 + s3t1[i];
vector<int> oc;
for(int i = m - 1; i < n; i++) if(res
    [i] == 0) oc.push_back(i - m + 1)
    ;
return oc;
}</pre>
```

7.6 Suffix Array

```
#include <bits/stdc++.h>
using namespace std;
 O(|S| + |alphabet|) Suffix Array
 LIM := \max\{s[i]\} + 2
void inducedSort (const vector <int> &
    vec, int val_range, vector <int> &
    SA, const vector <int> &sl, const
    vector <int> &lms_idx) {
 vector <int> 1(val_range, 0), r(
      val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c
         + 1];
 partial_sum(l.begin(), l.end(), l.
      begin());
 partial_sum(r.begin(), r.end(), r.
      begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i >=
       0; --i) SA[--r[vec[lms_idx[i]]]]
       = lms_idx[i];
 for (int i : SA) if (i > 0 and sl[i -
       1]) SA[1[vec[i-1]]++] = i-1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
 partial_sum(r.begin(), r.end(), r.
      begin());
 for (int k = SA.size() - 1, i = SA[k
      ]; k; --k, i = SA[k]) {
   if (i and !sl[i - 1]) SA[--r[vec[i -
        1]]] = i - 1;
vector <int> suffixArray (const vector
    <int> &vec, int val_range) {
 const int n = vec.size();
```

vector <int> sl(n), SA(n), lms_idx;

for (int i = n - 2; $i \ge 0$; --i) {

sl[i] = vec[i] > vec[i + 1] or (vec[i])

if (sl[i] and !sl[i + 1]) lms_idx.

reverse(lms_idx.begin(), lms_idx.end

emplace_back(i + 1);

());

i] == vec[i + 1] and sl[i + 1])

```
inducedSort(vec, val_range, SA, s1,
      lms idx):
 vector <int> new_lms_idx(lms_idx.size
      ()), lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {
   if (SA[i] > 0 \text{ and } !sl[SA[i]] \text{ and } sl[const int MX = 1e6+7, K = 20;
        SA[i] - 1]) new_lms_idx[k++] =
 int cur = 0; SA[n - 1] = 0;
 for (int k = 1; k < new_lms_idx.size</pre>
      (); ++k) {
   int i = new_lms_idx[k - 1], j =
       new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
     SA[j] = ++cur; continue;
   bool flag = 0;
   for (int a = i + 1, b = j + 1; ; ++a
        , ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
     if ((!sl[a] and sl[a - 1]) or (!sl
          [b] and sl[b - 1])) {
       flag = !(!sl[a] and sl[a - 1]
            and !sl[b] and sl[b - 1]);
            break;
     }
   }
   SA[j] = flag ? ++cur : cur;
 }
 for (int i = 0; i < lms_idx.size();</pre>
      ++i) lms_vec[i] = SA[lms_idx[i]];
 if (cur + 1 < lms_idx.size()) {</pre>
   auto lms_SA = suffixArray(lms_vec,
        cur + 1);
   for (int i = 0; i < lms_idx.size();</pre>
        ++i) new_lms_idx[i] = lms_idx[
        lms_SA[i]];
 }
 inducedSort(vec, val_range, SA, sl,
      new_lms_idx); return SA;
vector <int> getSuffixArray (const
    string &s, const int LIM = 128) {
 vector <int> vec(s.size() + 1);
  copy(begin(s), end(s), begin(vec));
      vec.back() = '$';
  auto ret = suffixArray(vec, LIM);
 ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any
    two suffix
vector <int> getLCParray (const string
    &s, const vector <int> &SA) {
 int n = s.size(), k = 0;
 vector <int> lcp(n), rank(n);
  for (int i = 0; i < n; ++i) rank[SA[i</pre>
      ]] = i;
 for (int i = 0; i < n; ++i, k ? --k :</pre>
       0) {
   if (rank[i] == n - 1) {
     k = 0; continue;
   int j = SA[rank[i] + 1];
   while (i + k < n \text{ and } j + k < n \text{ and } s
        [i + k] == s[j + k]) ++k;
```

```
lcp[rank[i]] = k;
 lcp[n - 1] = 0; return lcp;
int lg[MX];
void pre()
   lg[1] = 0;
   for (int i=2; i<MX; i++)</pre>
       lg[i] = lg[i/2]+1;
struct RMQ{
   int N;
   VI v[K];
   RMQ(const VI &a) {
       N = a.size();
       v[0] = a;
       for (int k = 0; (1<<(k+1)) <= N;
            k++) {
           v[k+1].resize(N);
           for (int i = 0; i-1+(1<<(k
               +1)) < N; i++) {
              v[k+1][i] = min(v[k][i],
                   v[k][i+(1<< k)]);
           }
       }
   }
   int findMin(int i, int j) {
       int k = lg[j-i+1];
       return min(v[k][i], v[k][j
           +1-(1<<k)]);
int main() {
 string s; cin >> s;
 for (const int i : getSuffixArray(s))
       printf("%d ", i);
 puts("");
 return 0;
```

7.7Trie

```
template<int sz>
struct Trie {
 Trie() : id(1) {
   memset(endMark, 0, sizeof endMark);
   for_each(all(trie), [](vector<int> &
       v) { v.assign(sz, 0); });
 void insert(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0;
```

```
for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false:
     cur = trie[cur][val];
   return endMark[cur];
 }
private:
 int id, endMark[100005];
 vector<int> trie[100005];
```

7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     l = r = i;
     while (r < n \&\& s[r] == s[r - 1])
     z[i] = r - 1, r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
       1 = i;
       while (r < n \&\& s[r] == s[r - 1]
            ]) r++;
       z[i] = r - 1, r--;
     }
   }
 }
 return z;
```

Extra 8

Make File

```
.PHONY: clean
D ?= 0
ifeq ($(D), 1)
CXXFLAGS=-std=c++2a -g -DMUSHFIQ_DEBUG
      -DDBG_MACRO_NO_WARNING -Wall -
     Wextra -Wpedantic -Wshadow -
     Wformat=2 -Wfloat-equal -
     Wconversion -Wlogical-op -Wshift-
     overflow=2 -Wduplicated-cond -
     Wcast-qual -Wcast-align -Wno-
     variadic-macros -D_GLIBCXX_DEBUG -
     D_GLIBCXX_DEBUG_PEDANTIC -
     fsanitize=address -fsanitize=
     undefined -fno-sanitize-recover -
     fstack-protector -fsanitize-
     address-use-after-scope
else
CXXFLAGS=-02 -std=c++2a
endif
@echo 'Error: No file given.'
clean:
@echo 'Cleaning executables...'
@find . ! -name stress.sh -type f -
     executable -delete
@echo 'Executables cleaned.'
```

```
8.2
    Stress Tester
make $1
make $2
make $3
# $1 is actual code
# $2 is good code
# $3 is generator
for i in 'seq 1 1000'; do
 echo 'Test #'$i
 timeout 5s ./$3 $RANDOM > in
 timeout 5s ./$1 < in> out
 timeout 5s ./$2 < in> ans
 diff ans out || break
done
8.3
     Sublime Build
{
  "shell_cmd": "make $file_base_name &&
      timeout 5s ./$file_base_name <</pre>
      in > out",
 "working_dir": "$file_path",
 "selector": "source.cpp"
```

```
8.4
     Vimrc
nmap <F5> :NERDTreeToggle<CR>
" nnoremap <silent> <expr> <F5> g:
    NERDTree.IsOpen() ? "\:
    NERDTreeClose<CR>" : bufexists(
    expand('%')) ? "\:NERDTreeFind<CR>"
     : "\:NERDTree<CR>"
" Auto import & Compile
:autocmd BufNewFile *.cpp Or ~/template
    .cpp
nnoremap <F4> :!xclip -o -sel clip > ~/
    cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip
     > ~/cp/in <CR><CR>
nnoremap <F6> :!xclip -sel clip % <CR><
inoremap <F6> <ESC>:!xclip -sel clip %
    <CR><CR>
map <F3> :50 vsplit in<CR>:split out<CR
    ><C-w>h
autocmd filetype cpp nnoremap <F9>
    wa \| !make %:r && timeout 5s ./%:r
     < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F9> <ESC
    >:wa \| !make %:r && timeout 5s
    ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp nnoremap <F10>
    wa \| !make clean && make %:r D=1
    && ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F10> <
    ESC>:wa \| !make clean && make %:r
    D=1 && ./%:r < ^/cp/in > ^/cp/out <
    CR>
autocmd filetype python nnoremap <F9> :
    wa \| !python % < ~/cp/in> ~/cp/out
```

autocmd filetype python inoremap <F9> <

ESC>:wa \| !python % < ~/cp/in> ~/

set cindent

```
cp/out<CR>
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'),
     col('.')-1, 1) == ")" ? "\<Right>"
     : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'),
     col('.')-1, 1) == "}" ? "\<Right>"
     : "}"
inoremap [ []<left>
inoremap <expr> ] strpart(getline('.'),
     col('.')-1, 1) == "]" ? "\<Right>"
inoremap <expr> " strpart(getline('.'),
     col('.')-1, 1) == "\"" ? "\<Right>
    " : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'),
     col('.')-1, 1) == "\'" ? "\<Right
    >" : "\',\',\<left>"
inoremap <expr> <CR> <sid>
    insert newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col(
      (., .)^{-2}, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<CR>\< nnoremap <S-k> :m .-2<CR>==
      ESC>\0" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>
   insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col(
      (., .) -2, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<space</pre>
      >\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('
     .')-2, 2)
return stridx('(){}[]''', pair) %
     2 == 0 && strlen(pair) == 2 ? "\<
     del>\<bs>" : "\<bs>"
endfunction
set nocompatible
    iMproved, required
                           " required
filetype on
filetype plugin on
filetype plugin indent on
syntax on
set splitright splitbelow
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
set expandtab
set softtabstop=2
set smartindent
set smarttab
set autoindent
```

```
set noerrorbells
set ruler
set guifont=*
set backspace=indent,eol,start
set ignorecase
set incsearch
set nowrap
 set hlsearch
" bubt site
 " set termguicolors
set foldmethod=indent
set nofoldenable
 " set cursorline
 set laststatus=2
set showcmd
set wildmenu
 " colorscheme torte
if !has('nvim')
  set clipboard=unnamedplus
endif
if !has('nvim')
 set ttymouse=xterm2
endif
nnoremap \langle S-j \rangle : m .+1 \langle CR \rangle ==
 vnoremap <S-j> :m '>+1<CR>gv==gv
vnoremap <S-k> :m '<-2<CR>gv==gv
filetype plugin indent on " required
```

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1$$

1 and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

the size n+1 factors.

intersecting chords.

n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

patterns of length 3); that is, the number of $\binom{n+k+1}{k}$ permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321.

9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count of $P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$, then, permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$

where, $S(0,0) = 1, S(n,0) = S(0,n) = 0$

$$\sum_{k=0}^{n} S(n,k) = n!$$

 $\sum_{k=0}^{k=0}$ The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x$$

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 i_k.$$

Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$ where S(0,0) = 1, S(n,0) = S(0,n) = 0 $S(n,2) = 2^{n-1} - 1$ $S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, $\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$

9 Equations and Formulas
9.1 Catalan Numbers $C_n = \frac{1}{n+1} \binom{2n}{n} \quad C_0 = 1, C_1 = \begin{cases} \text{with each subset containing at least } r \text{ elements. It is denoted by } S_r(n,k) \text{ and obeys the recurrence relation.} \\ kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \\ kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \end{cases}$ The number of ways to completely parenthesize n+1 factors.

Denote the n objects to partition by the integers $1,2,\ldots,n$. Define the reduced Stirling numbers of the second kind, defined n+1 factors.

Denote the n objects to partition by the integers $1,2,\ldots,n$. Define the reduced Stirling numbers of the second kind, defined n+1 factors. noted $S^d(n,k)$, to be the number of ways $\phi(d)$ The number of triangulations of a convex polygon with n+2 sides (i.e. the number of nonempty subsets such that all elements in $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right)$ partitions of polygon into disjoint triangles each subset have pairwise distance at least by using the diagonals).

The number of ways to connect the 2n subset, it is required that $|i-j| \ge d$. It $\left| \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d) \right|$ points on a circle to form n disjoint i.e. non- has been shown that these numbers satisfy, The number of rooted full binary trees with 9.4 Other Combinatorial Identities

n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children. Number of permutations of
$$1, \ldots, n$$
 that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For $n=3$, these permutations are 132, 213, 231, 312 and 321.

1. 2. Stirling Numbers First Kind.

If
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$
 where, $S(0,0) = 1, S(n,0) = S(0,n) = 0$ If $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

9.5 Different Math Formulas

Picks Theorem: A = i + b/2 - 1 $x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{i=1}^{n} S(n,k)x^{k}$ Deragements : $d(i) = (i-1) \times (i-1) + d(i-2)$ |(d(i-1)+d(i-2))|

$$\frac{n}{ab}$$
 - $\left\{\frac{b'n}{a}\right\}$ - $\left\{\frac{a'n}{b}\right\}$ + 1

9.6 GCD and LCM

if m is any integer, then $gcd(a + m \cdot b, b) =$ gcd(a,b)

The gcd is a multiplicative function in the following sense: if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$ $\gcd(a_1,b)\cdot\gcd(a_2,b).$

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).For non-negative integers a and b, where aand b are not both zero, $gcd(n^a-1, n^b-1) =$

and
$$b$$
 are not both zero, go $n^{\gcd(a,b)} - 1$ $\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$

$$\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$$

with each subset containing at least
$$r$$
 elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k)=kS_r(n,k)+\binom{n}{r-1}S_r(n-r+1,k-1)$
Denote the n objects to partition by the integers $1,2,\ldots,n$. Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers $1,2,\ldots,n$ into k nonempty subsets such that all elements in each subset have pairwise distance at least d . That is, for any integers i and j in a given subsets, it is required that $|i-j| \geq d$. It has been shown that these numbers satisfy, $S^d(n,k) = S(n-d+1,k-d+1), n \geq k \geq d$ i and i and