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# All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :20000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,
    ssse3,sse4,popcnt,abm,mmx,avx,tune=
    native")
#include <ext/pb_ds/assoc_container.hpp</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns
        iterator to the kth largest
        element counting from 0
   //order_of_key(val) --> returns the
        number of items in a set that
        are strictly smaller than our
template <typename DT>
using ordered_set = tree <DT, null_type</pre>
    , less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef SFT
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\',' << x <<</pre>
     '\'';}
void show(const string &x) {cerr << '\"</pre>
    ' << x << '\"';}
void show(bool x) {cerr << (x ? "true"</pre>
    : "false");}
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f';</pre>
    show(x.first); cerr << ", "; show(x void reset(int node) { lazy[node] =</pre>
    .second); cerr << '}'; }
template<typename T>
void show(T x) {int f = 0; cerr << "{";</pre>
     for (auto &i: x) cerr << (f++ ? ",</pre>
     " : ""), show(i); cerr << "}";}
void debug_out(string s) {
   cerr << '\n';</pre>
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(),
         '), s.end());
   cerr << "
                "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl:
   if(sizeof...(v)) debug_out(s, v...);
#define debug(x...) cerr << "LINE: " <<</pre>
     __LINE__ << endl; debug_out(#x, x) }
    ; cerr << endl;
#else
#define debug(x...)
#endif
```

```
const int RANDOM = chrono::
    high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) ^x;
unsigned hash_combine(unsigned a,
    unsigned b) { return a * 31 + b; }
struct chash {
 int operator()(int x) const { return
     hash_f(x); }
typedef gp_hash_table<int, int, chash>
    gp;
gp table;
    Data Structure
2.1 Segment Tree
const int N = 1000006;
```

```
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N];
LT lazy[4 * N];
int L, R;
void pull(int s, int e, int node) {
 val[node] = val[node << 1] + val[node</pre>
       << 1 | 1];
void apply(const LT &U, int s, int e,
    int node) {
 val[node] += (e - s + 1) * U;
 lazy[node] += U;
    None; }
DT merge(const DT &a, const DT &b) {
    return a + b; }
DT get(int s, int e, int node) { return int a[MAXN];
     val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
  apply(lazy[node], s, s + e >> 1, node
       << 1);
 apply(lazy[node], s + e + 2 \gg 1, e,
      node << 1 | 1);
 reset(node);
void build(int s, int e, vector<DT> &v,
     int node = 1) {
  int m = s + e >> 1;
  if (s == e) {
   val[node] = v[s];
   return;
 build(s, m, v, node * 2);
 build(m + 1, e, v, node * 2 + 1);
 pull(s, e, node);
void update(int S, int E, LT uval, int
    s = L, int e = R, int node = 1) {
 if (S > E) return;
```

if (S == s and E == e){

apply(uval, s, e, node);

```
return;
  }
  push(s, e, node);
  int m = s + e \gg 1;
  update(S, min(m, E), uval, s, m, node
       * 2):
  update(max(S, m + 1), E, uval, m + 1,
       e, node *2 + 1);
  pull(s, e, node);
DT query(int S, int E, int s = L, int e
     = R, int node = 1) {
  if (S > E) return I;
  if (s == S and e == E) return get(s,
      e, node);
  push(s, e, node);
  int m = s + e >> 1;
  DT L = query(S, min(m, E), s, m, node
       * 2);
  DT R = query(max(S, m + 1), E, m + 1,
       e, node *2 + 1);
  return merge(L, R);
void init(int _L, int _R, vector<DT> &v
 L = _L, R = _R;
  build(L, R, v);
```

## 2.2 Persistent Segment Tree

```
struct Node {
  Node *1, *r;
  int sum;
  Node(int val) : 1(nullptr), r(nullptr
       ), sum(val) {}
  Node(Node* 1, Node* r) : 1(1), r(r),
      sum(0) {
    if (1) sum += 1->sum;
    if (r) sum += r->sum;
  }
};
Node* root[MAXN];
Node* Build(int bg, int ed) {
  if (bg == ed) return new Node(a[bg]);
  int mid = (bg + ed) / 2;
  return new Node (Build (bg, mid), Build
       (mid + 1, ed));
int Query(Node* v, int bg, int ed, int
    1, int r) {
  if (1 > ed || r < bg) return 0;</pre>
  if (1 <= bg && ed <= r) return v->sum
  int mid = (bg + ed) / 2;
  return Query(v->1, bg, mid, 1, r) +
       Query(v->r, mid + 1, ed, l, r);
Node* Update(Node* v, int bg, int ed,
     int pos, int new_val) {
  if (bg == ed) return new Node(v->sum
       + new_val);
  int mid = (bg + ed) / 2;
```

if (pos <= mid)</pre>

```
node *root;
 implicit_segtree() {}
 implicit_segtree(int n) {
   root = new node(n);
 }
 void update(node *now, int L, int R,
      int idx, int val) {
    if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new
         node(mid - L + 1);
   if (now->rt == NULL) now->rt = new
       node(R - mid);
   if (idx <= mid) update(now->lft, L,
       mid, idx, val);
   else update(now->rt, mid + 1, R, idx
       , val);
   now->val = (now->lft)->val + (now->
       rt)->val;
 }
 int query(node *now, int L, int R,
      int k) {
   if (L == R) return L;
   int mid = L + (R - L) / 2;
   if (now->lft == NULL) now->lft = new|}
         node(mid - L + 1);
   if (now->rt == NULL) now->rt = new
       node(R - mid);
   if (k <= (now->lft)->val) return
        query(now->lft, L, mid, k);
   else return query(now->rt, mid + 1,
       R, k - (now->lft)->val);
 }
};
```

```
2.4 Centroid Decomposition
// problem: Xenia and Tree
// Centroid Tree Implementation
#define lg(n) (31 - __builtin_clz(n))

const int N = 1e5 + 1;
const int K = lg(N) + 1;

int n;
vector<int> adj[N];
int anc[N][K], lvl[N], par[N], sz[N],
    vis[N];
int minD[N];
```

```
// Call dfs at the very outset
 void dfs(int u = 1, int p = 0) {
   sz[u] = 1:
   for (auto v : adj[u])
     if (v != p) {
       dfs_size(v, u);
       sz[u] += sz[v];
 }
 int findCentroid(int u, int p) {
   int total = sz[u];
   for (auto v : adj[u])
     if (v != p and not vis[v] and 2 *
         sz[v] > total) {
       sz[u] = total - sz[v];
       sz[v] = total;
      return findCentroid(v, u);
     }
   return u;
 int query(int u) {
   int ans = 1e6;
   for (int i = u; i; i = par[i])
     ans = min(ans, minD[i] + lca::dist
         (i, u));
   return ans;
 }
 void update(int u) {
   for (int i = u; i; i = par[i])
     minD[i] = min(minD[i], lca::dist(i
         , u));
 int decompose(int u, int p) {
   u = findCentroid(u, p);
   vis[u] = 1;
   for (auto v : adj[u])
     if (not vis[v])
       par[decompose(v, u)] = u;
   return u;
int main() {
 cin.tie(NULL)->sync_with_stdio(false)
 fill(all(minD), 1e9);
 int m;
 cin >> n >> m;
 for (int u, v, i = 1; i < n; i++) {</pre>
   cin >> u >> v;
   adj[u].push_back(v);
   adj[v].push_back(u);
 lca::init();
 CD::dfs();
 CD::decompose(1, 0);
 CD::update(1);
 int t, v;
 while (m--) {
   cin >> t >> v;
   if (t == 1) CD::update(v);
   else cout << CD::query(v) << '\n';</pre>
```

```
return 0;
}
```

### 2.5 DSU With Rollbacks

```
struct Rollback_DSU {
 int n;
 vector<int> par, sz;
  vector<pair<int, int>> op;
  Rollback_DSU(int n) : par(n), sz(n,
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par
        [node])
     ; // no path compression
   return node;
  void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)])
         swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back())
     par[op.back().second] = op.back().
     sz[op.back().first] -= sz[op.back
          ().second];
 }
};
```

# 2.6 BIT-2D

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j += -j & j)
        bit[x][j] += val;
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j & -j) ans
        += bit[x][j];
 return ans;
int get(int x1, int y1, int x2, int y2)
 return get(x2, y2) - get(x1 - 1, y2)
      - get(x2, y1 - 1) + get(x1 - 1,
      y1 - 1);
```

## 2.7 Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];

vector<LL> merge(vector<LL> v1, vector<
    LL> v2) {
  LL i = 0, j = 0;
  vector<LL> ret;
```

```
while (i < v1.size() || j < v2.size()</pre>
   if (i == v1.size()) {
     ret.push_back(v2[j]);
     j++;
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
   } else {
     if (v1[i] < v2[j]) {</pre>
       ret.push_back(v1[i]);
     } else {
       ret.push_back(v2[j]);
       j++;
 }
 return ret;
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 }
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
 Tree[node] = merge(Tree[leftNode],
      Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1,
    LL r, LL k) {
  if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin
        (), Tree[node].end(), k) -
          Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, 1, r,
       k) +
        query(rightNode, mid + 1, ed, 1
             , r, k);
```

## 2.8 MO with Update

```
bool operator<(const query &rhs)</pre>
      const {
   return (L < rhs.L) or (L == rhs.L
        and R < rhs.R) or
          (L == rhs.L and R == rhs.R
              and t < rhs.t);</pre>
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
 cnt[freq[arr[idx]]]++;
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].last;
 if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
 assert(0);
int main() {
 int n, q;
 cin >> n >> q;
 int counter = 0;
 map<int, int> M;
 for (int i = 1; i <= n; i++) {</pre>
   cin >> arr[i];
   if (!M[arr[i]]) M[arr[i]] = ++
        counter;
   arr[i] = M[arr[i]];
 }
 iota(id, id + N, 0);
 while (q--) {
   int tp, x, y;
   cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = {x, y, arr[x]};
     arr[x] = y;
 }
 t = ui;
 cnt[0] = 3 * n;
```

## 2.9 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
    for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
    for (int q = 1; (1 << q) <= N; q++)</pre>
      int add = 1 << (q - 1);
     for (int j = 1; j + add <= N; j++)</pre>
       M[i][j][0][q] = max(M[i][j][0][q]
             - 1], M[i][j + add][0][q -
            1]);
     }
   }
 }
  for (int p = 1; (1 << p) <= N; p++) {</pre>
    int add = 1 << (p - 1);
    for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q
          ++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p]
              - 1][q], M[i + add][j][p -
               1][q]);
     }
   }
 }
// returns max of all A[i][j], where x1
    =i<=x2 and y1<=j<=y2
```

int Query(int x1, int y1, int x2, int

int kX = log2(x2 - x1 + 1);

int kY = log2(y2 - y1 + 1);

int addX = 1 << kX;</pre>

int addY = 1 << kY;</pre>

y2) {

# 2.10 Sparse Table

```
// tested by: https://judge.yosupo.jp/
    problem/staticrmq
// O-based indexing, query finds in
    range [first, last]
#define lg(x) (31 - __builtin_clz(x))
const int N = 1e5 + 1;
const int K = lg(N);
int a[N], tr[N][K + 1];
namespace sparse_table {
 int f(int p1, int p2) { return min(p1
      , p2); }
 void build() {
   for(int i = 0; i < n; i++)</pre>
     tr[i][0] = a[i];
   for(int j = 1; j <= K; j++) {</pre>
     for(int i = 0; i + (1<<j) <= n; i</pre>
       tr[i][j] = f(tr[i][j - 1], tr[i
            + (1<<(j - 1))][j - 1]);
 }
 int query(int 1, int r) {
   int d = lg(r - 1 + 1);
   return f(table[1][d], table[r - (1<<||
        d) + 1][d]);
 }
}
```

# 3 DP

# 3.1 Convex Hull Trick

```
struct line {
           ll m, c;
           line() {}
          line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
           vector<line> lines;
           int ptr = 0;
           convex_hull_trick() {}
           bool bad(line a, line b, line c) {
                      return 1.0 * (c.c - a.c) * (a.m - b.
                                                m) < 1.0 * (b.c - a.c) * (a.m - a.c) * (a.c) * (
                                                       c.m);
           void add(line L) {
                     int sz = lines.size();
                      while (sz >= 2 && bad(lines[sz - 2],
                                                       lines[sz - 1], L)) {
                                lines.pop_back();
                      lines.pb(L);
```

```
ll get(int idx, int x) { return (111
      * lines[idx].m * x + lines[idx].c
      ): }
 ll query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines
        .size() - 1;
   while (ptr < lines.size() - 1 && get
        (ptr, x) > get(ptr + 1, x)) ptr
   return get(ptr, x);
 }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio;
 int t;
 cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >>
        sum[i]:
   for (int i = 1; i <= n; i++) dp[i] =</pre>
         0, sum[i] += sum[i - 1];
   convex_hull_trick cht;
    cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 1
         11 * a * sqr(sum[pos]) - c;
     cht.add(line(211 * a * sum[pos],
         dp[pos] - a * sqr(sum[pos]));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

# 3.2 Dynamic CHT

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 IS_QUERY = -(1LL << 62);</pre>
struct line {
 ll m, b;
  mutable function <const line*()> succ
  bool operator < (const line &rhs)</pre>
      const {
    if (rhs.b != IS_QUERY) return m <</pre>
        rhs.m;
   const line *s = succ();
   if (!s) return 0;
   11 x = rhs.m;
   return b - s -> b < (s -> m - m) * x
 }
};
struct HullDynamic : public multiset <</pre>
    line> {
```

bool bad (iterator y) {

```
auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
     return y -> m == z -> m && y -> b
          \langle = z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x
        -> m && y -> b <= x -> b;
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z
        -> m - y -> m) >= 1.0 * (y -> b)
         -z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 void insert_line (ll m, ll b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) ==
        end() ? 0 : &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(
        y))) erase(next(y));
    while (y != begin() && bad(prev(y)))
         erase(prev(y));
 ll eval (ll x) {
    auto 1 = *lower_bound((line) {x,
        IS_QUERY});
    return 1.m * x + 1.b;
 }
int main() {
 HullDynamic hull;
 hull.insert_line(1, 1);
 hull.insert_line(-1, 1);
  cout << hull.eval(69) << endl;</pre>
  cout << hull.eval(420) << endl;</pre>
  return 0:
```

# 3.3 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c)
LL calc(line L, LL x) { return 1LL * L.
    m * x + L.c; }
struct node {
 LL m, c;
 line L;
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft =
      NULL, node *rt = NULL)
     : L(line(m, c)), lft(lft), rt(rt)
};
struct LiChao {
 node *root;
 LiChao() { root = new node(); }
 void update(node *now, int L, int R,
      line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(
        lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
```

```
if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new
          node():
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo;
     if (now->lft == NULL) now->lft =
         new node();
     update(now->lft, L, mid, hi);
   }
 LL query(node *now, int L, int R, LL
      x) {
   if (now == NULL) return -inf;
   int mid = L + (R - L) / 2;
   if (x \le mid)
     return max(calc(now->L, x), query(
         now->lft, L, mid, x));
   else
     return max(calc(now->L, x), query(
         now->rt, mid + 1, R, x));
 }
};
```

### SOS DP 3.4

```
for(int i = 0; i<(1<<N); ++i)</pre>
F[i] = A[i];
for(int i = 0;i < N; ++i) for(int mask</pre>
    = 0; mask < (1<<N); ++mask){
 if(mask & (1<<i))</pre>
 F[mask] += F[mask^(1<<i)];
```

## Geometry

## 4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long
    for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ?</pre>
    0 : (x < 0 ? -1 : 1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y)
 Point operator+(const Point& u) const
      { return Point(x + u.x, y + u.y)
 Point operator-(const Point& u) const
      { return Point(x - u.x, y - u.y)
      ; }
 Point operator*(const LL u) const {
      return Point(x * u, y * u); }
 Point operator*(const Tf u) const {
     return Point(x * u, y * u); }
 Point operator/(const Tf u) const {
      return Point(x / u, y / u); }
 bool operator==(const Point& u) const
   return dcmp(x - u.x) == 0 && dcmp(y)
        - u.y) == 0;
 bool operator!=(const Point& u) const
       { return !(*this == u); }
```

```
bool operator<(const Point& u) const</pre>
        return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                  -u.x) == 0 && dcmp(y - u.y) <
   }
Ti dot(Point a, Point b) { return a.x *
            b.x + a.y * b.y; }
Ti cross(Point a, Point b) { return a.x | struct Segment {
            * b.y - a.y * b.x; }
Tf length(Point a) { return sqrt(dot(a,
            a)); }
Ti sqLength(Point a) { return dot(a, a)
          ; }
Tf distance(Point a, Point b) { return
          length(a - b); }
Tf angle(Point u) { return atan2(u.y, u
// returns angle between oa, ob in (-PI
Tf angleBetween(Point a, Point b) {
   Tf ans = angle(b) - angle(a);
    return ans <= -PI ? ans + 2 * PI : (
              ans > PI ? ans - 2 * PI : ans);
// Rotate a ccw by rad radians, Tf Ti
Point rotate(Point a, Tf rad) {
    return Point(a.x * cos(rad) - a.y *
              sin(rad),
                                a.x * sin(rad) + a.y * cos
                                           (rad)):
// rotate a ccw by angle th with cos(th
         ) = co \&\& sin(th) = si, tf ti same
Point rotatePrecise(Point a, Tf co, Tf
         si) {
    return Point(a.x * co - a.y * si, a.y
                * co + a.x * si);
Point rotate90(Point a) { return Point
          (-a.y, a.x); }
// scales vector a by s such that
         length of a becomes s, Tf Ti same
Point scale(Point a, Tf s) { return a /
            length(a) * s; }
 // returns an unit vector perpendicular
            to vector a, Tf Ti same
Point normal(Point a) {
   Tf l = length(a);
    return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if
          on ab && -1 if right of ab
int orient(Point a, Point b, Point c) {
            return dcmp(cross(b - a, c - a));
/// Use as sort(v.begin(), v.end(),
          polarComp(0, dir))
/// Polar comparator around O starting
          at direction dir
struct polarComp {
   Point O, dir;
    polarComp(Point 0 = Point(0, 0),
              Point dir = Point(1, 0)) : O(0),
              dir(dir) {}
    bool half(Point p) {
        return dcmp(cross(dir, p)) < 0 ||</pre>
                                                                                                 bool linesParallel(Line p, Line q) {
```

```
(dcmp(cross(dir, p)) == 0 \&\&
              dcmp(dot(dir, p)) > 0);
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) <</pre>
       make_tuple(half(q), cross(p, q)
 Point a, b;
 Segment(Point aa, Point bb) : a(aa),
     b(bb) {}
typedef Segment Line;
struct Circle {
 Point o;
 Tf r;
 Circle(Point o = Point(0, 0), Tf r =
      0) : o(o), r(r) {}
 // returns true if point p is in ||
     on the circle
 bool contains(Point p) { return dcmp(
      sqLength(p - o) - r * r) <= 0; }
 // returns a point on the circle rad
     radians away from +X CCW
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value
   return Point(o.x + cos(rad) * r, o.y
        + sin(rad) * r);
 // area of a circular sector with
      central angle rad
 Tf area(Tf rad = PI + PI) { return
     rad * r * r / 2; }
 // area of the circular sector cut by
      a chord with central angle alpha
 Tf sector(Tf alpha) { return r * r *
     0.5 * (alpha - sin(alpha)); }
```

# 4.2 Linear

```
// **** LINE LINE INTERSECTION START
// returns true if point p is on
   segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p))
     == 0 && dcmp(dot(s.a - p, s.b - p
     )) <= 0;</pre>
// returns true if segment p && q touch
     or intersect
bool segmentsIntersect(Segment p,
   Segment q) {
 if (onSegment(p.a, q) || onSegment(p.
     b, q)) return true;
 if (onSegment(q.a, p) || onSegment(q.
     b, p)) return true;
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 &&
     dcmp(c3) * dcmp(c4) < 0;
```

```
return dcmp(cross(p.b - p.a, q.b - q. // returns the projection of point p on
                                                                                          ]));
                                              line 1, Tf Ti Same
                                                                                    return abs(res);
      a)) == 0:
                                         |Point projectPointLine(Point p, Line 1) |}
// lines are represented as a ray from
                                           Point v = 1.b - 1.a;
    a point: (point, vector)
                                                                                  // interior
                                           return 1.a + v * ((Tf)dot(v, p - 1.a)
// returns false if two lines (p, v) &&
                                                                                       ---> -2
     (q, w) are parallel or collinear
                                                / dot(v, v));
// true otherwise, intersection point
                                                                                       ----> -1
    is stored at o via reference, Tf Ti
                                                                                   // concentric
                                                                                                      (d = 0)
     Same
                                         4.3 Circular
                                                                                  // secants
                                                                                       ) ----> 0
bool lineLineIntersection(Point p,
                                          // Extremely inaccurate for finding
    Point v, Point q, Point w, Point& o
                                             near touches
    ) {
                                                                                       ----> 1
                                         // compute intersection of line 1 with
 if (dcmp(cross(v, w)) == 0) return
                                                                                   // exterior
                                             circle c
                                                                                       ----> 2
      false:
                                          // The intersections are given in order
 Point u = p - q;
                                              of the ray (l.a, l.b), Tf Ti same
 o = p + v * (cross(w, u) / cross(v, w)
                                                                                       Circle c2) {
                                         vector<Point> circleLineIntersection(
                                                                                    Tf d = length(c1.o - c2.o);
      ));
                                             Circle c, Line 1) {
 return true;
                                           vector<Point> ret;
                                           Point b = 1.b - 1.a, a = 1.a - c.o;
// returns false if two lines p && q
                                           Tf A = dot(b, b), B = dot(a, b);
    are parallel or collinear
                                           Tf C = dot(a, a) - c.r * c.r, D = B *
// true otherwise, intersection point
                                                B - A * C;
    is stored at o via reference
                                           if (D < -EPS) return ret;</pre>
bool lineLineIntersection(Line p, Line
                                           ret.push_back(l.a + b * (-B - sqrt(D
    q, Point& o) {
                                                                                        same
                                               + EPS)) / A);
 return lineLineIntersection(p.a, p.b
                                           if (D > EPS) ret.push_back(1.a + b *
      - p.a, q.a, q.b - q.a, o);
                                                                                       Circle c1, Circle c2) {
                                               (-B + sqrt(D)) / A);
                                                                                    vector<Point> ret;
                                           return ret;
// returns the distance from point a to |
                                                                                    Tf d = length(c1.o - c2.o);
     line 1
                                                                                    if (dcmp(d) == 0) return ret;
                                         // signed area of intersection of
// **** LINE LINE INTERSECTION FINISH
                                             circle(c.o, c.r) &&
    ****
                                                                                          ret;
                                          // triangle(c.o, s.a, s.b) [cross(a-o,
Tf distancePointLine(Point p, Line 1) {
                                             b-o)/2
 return abs(cross(1.b - 1.a, p - 1.a)
                                                                                        return ret;
                                         Tf circleTriangleIntersectionArea(
      / length(1.b - 1.a));
                                             Circle c, Segment s) {
                                                                                    Point v = c2.o - c1.o;
                                           using Linear::distancePointSegment;
// returns the shortest distance from
                                           Tf OA = length(c.o - s.a);
    point a to segment s
                                           Tf OB = length(c.o - s.b);
Tf distancePointSegment(Point p,
                                                                                         (v));
                                           // sector
    Segment s) {
                                           if (dcmp(distancePointSegment(c.o, s)
 if (s.a == s.b) return length(p - s.a
                                                -c.r) >= 0)
                                                                                          -si), c1.r) + c1.o;
      );
                                             return angleBetween(s.a - c.o, s.b -
 Point v1 = s.b - s.a, v2 = p - s.a,
                                                  c.o) * (c.r * c.r) / 2.0;
      v3 = p - s.b;
                                                                                          si), c1.r) + c1.o;
                                           // triangle
 if (dcmp(dot(v1, v2)) < 0)</pre>
                                           if (dcmp(OA - c.r) <= 0 && dcmp(OB -
   return length(v2);
                                                                                    ret.push_back(p1);
                                               c.r) <= 0)
 else if (dcmp(dot(v1, v3)) > 0)
                                             return cross(c.o - s.b, s.a - s.b) /
   return length(v3);
                                                                                    return ret;
                                                  2.0:
                                           // three part: (A, a) (a, b) (b, B)
   return abs(cross(v1, v2) / length(v1
                                           vector<Point> Sect =
       ));
                                                                                       circles c1, c2
                                               circleLineIntersection(c, s);
                                           return circleTriangleIntersectionArea
// returns the shortest distance from
                                                                                       c1, Circle c2) {
                                               (c, Segment(s.a, Sect[0])) +
                                                                                    Point AB = c2.o - c1.o;
    segment p to segment q
                                                 circleTriangleIntersectionArea(
Tf distanceSegmentSegment(Segment p,
                                                                                    Tf d = length(AB);
                                                      c, Segment(Sect[0], Sect
    Segment q) {
                                                      [1])) +
 if (segmentsIntersect(p, q)) return
                                                 circleTriangleIntersectionArea(
      0;
                                                                                        r * c1.r;
                                                      c, Segment(Sect[1], s.b));
 Tf ans = distancePointSegment(p.a, q)
                                         // area of intersecion of circle(c.o, c
 ans = min(ans, distancePointSegment(p
                                              .r) && simple polyson(p[])
      .b, q));
                                         Tf circlePolyIntersectionArea(Circle c,
 ans = min(ans, distancePointSegment(q
                                              Polygon p) {
                                                                                         ));
      .a, p));
                                           Tf res = 0;
 ans = min(ans, distancePointSegment(q
                                           int n = p.size();
      .b, p));
                                           for (int i = 0; i < n; ++i)
 return ans:
                                                                                         )):
                                                 circleTriangleIntersectionArea(
```

c, Segment(p[i], p[(i + 1) % n

```
// locates circle c2 relative to c1
                      (d < R - r)
// interior tangents (d = R - r)
                     (R - r < d < R + r)
// exterior tangents (d = R + r)
                      (d > R + r)
int circleCirclePosition(Circle c1,
 int in = dcmp(d - abs(c1.r - c2.r)),
      ex = dcmp(d - (c1.r + c2.r));
 return in < 0 ? -2 : in == 0 ? -1 :
      ex == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points
    between two circles c1 && c2, Tf Ti
vector<Point> circleCircleIntersection(
 if (dcmp(c1.r + c2.r - d) < 0) return
 if (dcmp(abs(c1.r - c2.r) - d) > 0)
 Tf co = (c1.r * c1.r + sqLength(v) -
      c2.r * c2.r) / (2 * c1.r * length
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co,
 Point p2 = scale(rotatePrecise(v, co,
 if (p1 != p2) ret.push_back(p2);
// intersection area between two
Tf circleCircleIntersectionArea(Circle
 if (d >= c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.</pre>
 if (d + c2.r <= c1.r) return PI * c2.</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d
       -c2.r*c2.r) / (2.0*c1.r*d
 Tf alpha2 = acos((c2.r * c2.r + d * d))
       -c1.r*c1.r) / (2.0*c2.r*d)
 return c1.sector(2 * alpha1) + c2.
      sector(2 * alpha2);
```

```
// returns tangents from a point p to
    circle c, Tf Ti same
vector<Point> pointCircleTangents(Point
     p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
   ;
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u,
        ang)};
 return ret;
// returns the points on tangents that
    touches the circle, Tf Ti Same
vector<Point> pointCircleTangencyPoints
    (Point p, Circle c) {
 Point u = p - c.o;
 Tf d = length(u);
 if (d < c.r)
   return {};
 else if (dcmp(d - c.r) == 0)
   return {c.o + u};
 else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o +
        rotate(u, ang)};
 }
// for two circles c1 && c2, returns
    two list of points a && b
// such that a[i] is on c1 && b[i] is
    c2 && for every i
// Line(a[i], b[i]) is a tangent to
    both circles
// CAUTION: a[i] = b[i] in case they
    touch \mid -1 for c1 = c2
int circleCircleTangencyPoints(Circle
    c1, Circle c2, vector<Point> &a,
                            vector<
                                 Point>
                                 &b) {
 a.clear(), b.clear();
 int cnt = 0;
 if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r +
      c2.r;
 if (dcmp(d2 - rdif * rdif) < 0)</pre>
      return 0:
 if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
      ) == 0) return -1;
 Tf base = angle(c2.o - c1.o);
 if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++:
   return cnt;
```

```
Tf ang = acos((c1.r - c2.r) / sqrt(d2)
 a.push_back(c1.point(base + ang));
 b.push_back(c2.point(base + ang));
 a.push_back(c1.point(base - ang));
 b.push_back(c2.point(base - ang));
 if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++;
 } else if (dcmp(d2 - rsum * rsum) >
   Tf ang = acos((c1.r + c2.r) / sqrt(
       d2)):
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang
   cnt++;
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang
   cnt++:
 }
 return cnt;
    Convex
/// minkowski sum of two polygons in O(
```

```
Polygon minkowskiSum(Polygon A, Polygon
     B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin
      (), A.end()), A.end());
 rotate(B.begin(), min_element(B.begin
      (), B.end()), B.end());
 A.push_back(A[0]);
 B.push_back(B[0]);
 for (int i = 0; i < n; i++) A[i] = A[</pre>
      i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[</pre>
      i + 1] - B[i];
 Polygon C(n + m + 1);
 C[0] = A.back() + B.back();
 merge(A.begin(), A.end() - 1, B.begin
      (), B.end() - 1, C.begin() + 1,
       polarComp(Point(0, 0), Point(0,
           -1)));
 for (int i = 1; i < C.size(); i++) C[</pre>
      i] = C[i] + C[i - 1];
 C.pop_back();
 return C;
 / finds the rectangle with minimum
    area enclosing a convex polygon and
// the rectangle with minimum perimeter
     enclosing a convex polygon
// Tf Ti Same
pair<Tf, Tf>
    rotatingCalipersBoundingBox(const
    Polygon &p) {
 using Linear::distancePointLine;
 int n = p.size();
```

int l = 1, r = 1, j = 1;

```
Tf area = 1e100;
  Tf perimeter = 1e100;
  for (int i = 0; i < n; i++) {</pre>
   Point v = (p[(i + 1) \% n] - p[i]) /
        length(p[(i + 1) % n] - p[i]);
   while (dcmp(dot(v, p[r % n] - p[i])
        - dot(v, p[(r + 1) \% n] - p[i])
        ) < 0)
     r++;
   while (j < r || dcmp(cross(v, p[j %</pre>
        n] - p[i]) -
                       cross(v, p[(j +
                            1) % n] - p[
                            i])) < 0)
     j++;
   while (1 < j ||
          dcmp(dot(v, p[1 % n] - p[i])
               - dot(v, p[(1 + 1) % n] -
               p[i])) > 0)
     1++;
   Tf w = dot(v, p[r \% n] - p[i]) - dot
        (v, p[1 % n] - p[i]);
   Tf h = distancePointLine(p[j % n],
        Line(p[i], p[(i + 1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2)
         * h);
 return make_pair(area, perimeter);
// returns the left side of polygon u
    after cutting it by ray a->b
Polygon cutPolygon(Polygon u, Point a,
    Point b) {
  using Linear::lineLineIntersection;
  using Linear::onSegment;
  Polygon ret;
  int n = u.size();
  for (int i = 0; i < n; i++) {</pre>
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0)
        ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0)
     Point t:
     lineLineIntersection(a, b - a, c,
          d - c, t);
     if (onSegment(t, Segment(c, d)))
          ret.push_back(t);
 }
 return ret;
^{\prime}/ returns true if point p is in or on
    triangle abc
bool pointInTriangle(Point a, Point b,
    Point c, Point p) {
  return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >=
      0 &&
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no
    three collinear points
// returns inside = -1, on = 0, outside
int pointInConvexPolygon(const Polygon
    &pt, Point p) {
  int n = pt.size();
```

```
assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p -
       pt[0])) > 0)
     lo = mid;
   else
     hi = mid;
 bool in = pointInTriangle(pt[0], pt[
      lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p
       - pt[lo - 1])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p -
      pt[lo])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) %
       n], p - pt[(hi + 1) % n])) == 0)
   return 0;
 return -1;
// Extreme Point for a direction is the
     farthest point in that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly,
    Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1)
       % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 +
            n) % n], u)) >= 0) {
     return c;
   bool a_up = dcmp(dot(poly[(a + 1) %
       n] - poly[a], u)) >= 0;
   bool c_up = dcmp(dot(poly[(c + 1) %
       n] - poly[c], u)) >= 0;
   bool a_above_c = dcmp(dot(poly[a] -
       poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c:
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c:
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 if (dcmp(dot(poly[a] - poly[(a + 1) %
       n], u)) > 0 &&
         ) % n], u)) > 0)
                                           return orient(Q, u, v) != -dir;
   return a;
```

```
return b % n;
                                   // For a convex polygon p and a line l,
                                        returns a list of segments
                                   // of p that touch or intersect line 1.
                                   // the i'th segment is considered (p[i
                                        ], p[(i + 1) modulo |p|])
                                    // #1 If a segment is collinear with
                                        the line, only that is returned
                                    // #2 Else if l goes through i'th point
                                        , the i'th segment is added
                                    // Complexity: O(lg |p|)
                                   vector<int> lineConvexPolyIntersection(
                                        const Polygon &p, Line 1) {
                                     assert((int)p.size() >= 3);
                                     assert(1.a != 1.b);
                                     int n = p.size();
                                     vector<int> ret;
                                     Point v = 1.b - 1.a;
                                     int lf = extremePoint(p, rotate90(v))
                                     int rt = extremePoint(p, rotate90(v)
                                         * Ti(-1)):
                                     int olf = orient(l.a, l.b, p[lf]);
                                     int ort = orient(l.a, l.b, p[rt]);
                                     if (!olf || !ort) {
                                       int idx = (!olf ? lf : rt);
                                       if (orient(l.a, l.b, p[(idx - 1 + n)
                                            % n]) == 0)
                                         ret.push_back((idx - 1 + n) \% n);
                                         ret.push_back(idx);
                                       return ret:
                                     if (olf == ort) return ret;
                                     for (int i = 0; i < 2; ++i) {
                                       int lo = i ? rt : lf;
                                       int hi = i ? lf : rt;
                                       int olo = i ? ort : olf;
                                       while (true) {
                                         int gap = (hi - lo + n) \% n;
                                         if (gap < 2) break;</pre>
                                         int mid = (lo + gap / 2) % n;
                                         int omid = orient(l.a, l.b, p[mid
                                             ]);
                                         if (!omid) {
                                          lo = mid;
                                          break;
                                         if (omid == olo)
                                           lo = mid;
                                         else
                                          hi = mid;
                                       ret.push_back(lo);
                                     return ret;
                                    // Calculate [ACW, CW] tangent pair
                                        from an external point
                                   constexpr int CW = -1, ACW = 1;
dcmp(dot(poly[a] - poly[(a - 1 + n bool isGood(Point u, Point v, Point Q,
                                        int dir) {
```

```
Point better (Point u, Point v, Point Q,
     int dir) {
 return orient(Q, u, v) == dir ? u : v
Point pointPolyTangent(const Polygon &
    pt, Point Q, int dir, int lo, int
    hi) {
  while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid -
        1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid +
        1], Q, dir);
   if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q,
          dir, mid + 1, hi);
     Point p2 = pointPolyTangent(pt, Q,
          dir, lo, mid - 1);
     return better(p1, p2, Q, dir);
   if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) ==
         dir)
       hi = mid - 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) ==
         dir)
       lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
 }
 Point ret = pt[lo];
  for (int i = lo + 1; i <= hi; i++)</pre>
      ret = better(ret, pt[i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(
    const Polygon &pt, Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt,
      Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q
      , CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
4.5 Polygon
```

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &
    poly) {
 Polygon ret;
```

```
int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0
        && (ret.empty() || b != ret.
       back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p
     of n vertices
Tf signedPolygonArea(const Polygon &p)
    {
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() -</pre>
      1; i++)
   ret += cross(p[i] - p[0], p[i + 1] -
        p[0]);
 return ret / 2;
}
// given a polygon p of n vertices,
    generates the convex hull in in CCW
// Tested on https://acm.timus.ru/
    problem.aspx?space=1&num=1185
// Caution: when all points are
    colinear AND removeRedundant ==
    false
// output will be contain duplicate
    points (from upper hull) at back
Polygon convexHull(Polygon p, bool
    removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p
      .end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
          dcmp(cross(ch[m - 1] - ch[m -
               2], p[i] - ch[m - 1]))
              <= check)
     m--:
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i >= 0; i--) {
   while (m > k \&\&
          dcmp(cross(ch[m - 1] - ch[m -
               2], p[i] - ch[m - 2]))
              <= check)
     m--:
   ch[m++] = p[i];
 if (n > 1) m--;
 ch.resize(m);
 return ch;
// returns inside = -1, on = 0, outside
int pointInPolygon(const Polygon &p,
    Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
```

```
int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])
       ) || o == p[i]) return 0;
   int k = dcmp(cross(p[j] - p[i], o -
       p[i]));
   int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn
   if (k < 0 && d2 <= 0 && d1 > 0) wn
 return wn ? -1 : 1;
// Given a simple polygon p, and a line
     1, returns (x, y)
// x = longest segment of 1 in p, y =
    total length of 1 in p.
pair<Tf, Tf> linePolygonIntersection(
    Line 1, const Polygon &p) {
 using Linear::lineLineIntersection;
 int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {</pre>
   Point a = p[i], b = p[(i + 1) \% n],
       z = p[(i - 1 + n) \% n];
   int ora = orient(l.a, l.b, a), orb =
        orient(l.a, l.b, b),
       orz = orient(l.a, l.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(
           d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins;
     lineLineIntersection(1, Line(a, b)
          , ins);
     ev.emplace_back(dot(ins - 1.a, 1.b
          - 1.a), 0);
 }
 sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot =
 bool active = false;
 int sign = 0;
 for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment
         is (last, d)
                    /// On Border
   if (sign) {
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
   } else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
```

if (tp == 0)

```
active = !active;
   else
     sign = tp;
 last = d;
 if (!active) len = 0;
ans /= length(l.b - l.a);
tot /= length(1.b - 1.a);
return {ans, tot};
```

# Graph

```
5.1 LCA, ETT, VT
#define lg(n) (31 - __builtin_clz(n))
const int N = 1e5 + 1;
const int K = lg(N) + 1;
vector<int> adj[N];
int anc[N][K], lvl[N];
namespace lca {
 void init(int u = 1, int p = 0, int d
       = 0) {
   lvl[u] = d;
   anc[u][0] = p;
   for (int i = 1; i < K; i++)</pre>
     anc[u][i] = anc[anc[u][i - 1]][i -
           1];
   for (auto v : adj[u])
     if (v != p)
       init(v, u, d + 1);
 int getAnc(int u, int k) {
   for (int i = 0; u and i < K; i++)</pre>
     if ((k >> i) & 1)
       u = anc[u][i];
   return u;
 int lca(int u, int v) {
   if (lvl[u] < lvl[v]) swap(u, v);</pre>
   u = getAnc(u, lvl[u] - lvl[v]);
   if (u == v) return u;
   for (int i = K - 1; ~i; i--)
     if (anc[u][i] != anc[v][i])
       u = anc[u][i], v = anc[v][i];
   return anc[u][0];
 int dist(int u, int v) {
   return lvl[u] + lvl[v] - 2 * lvl[lca
        (u, v)];
struct euler_tour {
 int time = 0;
 tree &T;
 int n;
 vector<int> start, finish, level, par
 euler_tour(tree &T, int root = 0)
     : T(T), n(T.n), start(n), finish(n
         ), level(n), par(n) {
```

time = 0;

call(root);

```
void call(int node, int p = -1) {
                                         // for simplicity, G[idx] contains the
                                              adjacency list of a node
   if (p != -1) level[node] = level[p]
                                                                                         pos[b]);
        + 1;
                                          // while G(e) is a reference to the e-
                                                                                     return ret;
   start[node] = time++;
                                              th edge.
   for (int e : T[node])
                                         const int N = 2e5 + 5;
                                                                                   5.4
                                                                                         SCC
                                         int in[N], out[N], fwd[N], bck[N];
     if (e != p) call(e, node);
   par[node] = p;
                                         int t = 0;
                                                                                    typedef long long LL;
   finish[node] = time++;
                                         void dfs(graph &G, int node, int par) {
                                                                                   const LL N = 1e6 + 7;
 }
                                           out[node] = t;
 bool isAncestor(int node, int par) {
                                           for (int e : G[node]) {
                                                                                   bool vis[N];
   return start[par] <= start[node] and</pre>
                                             int v = G(e).to(node);
                                                                                   int root[N];
         finish[par] >= finish[node];
                                             if (v == par) continue;
                                             fwd[e] = t++;
 int subtreeSize(int node) { return
                                             dfs(G, v, node);
      finish[node] - start[node] + 1 >>
                                             bck[e] = t++;
       1; }
                                           in[node] = t - 1;
tree virtual_tree(vector<int> &nodes,
                                          void init(graph &G, int node) {
                                                                                     vis[u] = true;
    lca_table &table, euler_tour &tour)
                                           t = 0;
                                           dfs(G, node, node);
 sort(nodes.begin(), nodes.end(),
                                                                                     for (int v : ad[u])
      [&](int x, int y) { return tour.
          start[x] < tour.start[y]; });</pre>
 int n = nodes.size():
                                         5.3 HLD
 for (int i = 0; i + 1 < n; i++)
                                          const int N = 1e6 + 7;
   nodes.push_back(table.lca(nodes[i],
                                                                                    int main() {
                                         template <typename DT>
        nodes[i + 1]));
                                         struct Segtree {
 sort(nodes.begin(), nodes.end());
                                           // write lazy segtree here
 nodes.erase(unique(nodes.begin(),
                                         };
      nodes.end()), nodes.end());
                                         Segtree<int> tree(N);
 sort(nodes.begin(), nodes.end(),
                                         vector<int> adj[N];
      [&](int x, int y) { return tour.
                                         int depth[N], par[N], pos[N];
          start[x] < tour.start[y]; });</pre>
                                                                                       if (!vis[i]) {
                                         int head[N], heavy[N], cnt;
 n = nodes.size():
                                                                                         dfs(i, 1);
 stack<int> st;
                                          int dfs(int u, int p) {
 st.push(0);
                                           int SZ = 1, mxsz = 0, heavyc;
 tree ans(n);
                                                                                           root[j] = rt;
                                           depth[u] = depth[p] + 1;
 for (int i = 1; i < n; i++) {</pre>
   while (!tour.isAncestor(nodes[i],
                                           for (auto v : adj[u]) {
        nodes[st.top()])) st.pop();
                                             if (v == p) continue;
   ans.addEdge(st.top(), i);
                                                                                     }
                                             par[v] = u;
   st.push(i);
                                             int subsz = dfs(v, u);
 }
                                             if (subsz > mxsz) heavy[u] = v, mxsz
 return ans;
                                                   = subsz;
                                             SZ += subsz;
                                           }
set<int> getCenters(tree &T) {
                                                                                         if (rtu != rtv)
                                           return SZ;
 int n = T.n;
 vector<int> deg(n), q;
                                          void decompose(int u, int h) {
 set<int> s;
                                           head[u] = h, pos[u] = ++cnt;
 for (int i = 0; i < n; i++) {</pre>
                                           if (heavy[u] != -1) decompose(heavy[u
   deg[i] = T[i].size();
                                                                                    5.5 LCA in O(1)
                                                ], h);
   if (deg[i] == 1) q.push_back(i);
                                                                                    /* LCA in O(1)
   s.insert(i);
                                           for (int v : adj[u]) {
                                             if (v == par[u]) continue;
 for (vector<int> t; s.size() > 2; q =
                                             if (v != heavy[u]) decompose(v, v);
                                                                                         of edges
       t) {
                                           }
   for (auto x : q) {
                                                                                   LL depth[N];
     for (auto e : T[x])
                                          int query(int a, int b) {
                                                                                   int level[N];
       if (--deg[e] == 1) t.push_back(e
                                           int ret = 0;
                                           for (; head[a] != head[b]; b = par[
     s.erase(x);
                                                head[b]]) {
   }
                                             if (depth[head[a]] > depth[head[b]])
 }
                                                                                   vector<PII> adj[N];
                                                   swap(a, b);
 return s;
                                             ret += tree.query(1, 0, cnt, pos[
                                                 head[b]], pos[b]);
```

if (depth[a] > depth[b]) swap(a, b);

Euler Tour on Edge

```
ret += tree.query(1, 0, cnt, pos[a],
vector<int> adj[N], adjr[N], adj_scc[N
vector<int> order, component, roots;
// tp = 0 ,finding topo order, tp = 1 ,
     reverse edge traversal
void dfs(int u, int tp = 0) {
 if (tp) component.push_back(u);
 auto& ad = (tp ? adjr : adj);
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
 reverse(order.begin(), order.end());
 for (int i : order) {
     int rt = component[0];
     for (auto j : component)
     roots.push_back(rt);
     component.clear();
 for (int u = 1; u <= n; u++)</pre>
   for (auto v : adj[u]) {
     int rtu = root[u],
         rtv = root[v];
       adj_scc[rtu].push_back(rtv);
 * depth calculates weighted distance
 * level calculates distance by number
 * Preprocessing in NlongN */
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG
      [1] = 0;
```

```
for (int i = 2; i <= n; i++) LOG[i] = /// flow with demand(lower bound) only
       LOG[i >> 1] + 1;
 for (int i = 0; i <= n; i++) adj[i].</pre>
      clear():
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
 if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v])</pre>
       return v:
 int 1 = LOG[id[v] - id[u] + 1];
  int p1 = id[u], p2 = id[v] - (1 << 1)</pre>
  int d1 = level[table[1][p1]], d2 =
      level[table[1][p2]];
 if (d1 < d2)
   return par[table[1][p1]];
 else
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[
      1] * 2));
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time:
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table</pre>
      [0][i] = a[i];
 for (int 1 = 0; 1 < L - 1; 1++) {</pre>
   for (int i = 1; i <= n; i++) {</pre>
     table[1 + 1][i] = table[1][i];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 <<</pre>
           1)]] < level[table[1][i]];</pre>
     if (C1 && C2) table[l + 1][i] =
          table[1][i + (1 << 1)];
   }
 }
}
     Dinic Max Flow
                                          Ti maxFlow(int s, int t) {
```

```
for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u
    1))
// add_edge(u, new sink, sum(out_demand
    [u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then
     demand satisfied
// flow in every edge i = demand[i] + e |}
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
 edge(int v, int u, Ti cap) : v(v), u(
      u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.
      empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level
          [ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q
            .push(ed.u);
   }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].</pre>
      size(); cid++) {
   int id = adj[v][cid];
   auto &ed = edges[id];
   if (level[v] + 1 != level[ed.u] ||
        ed.cap - ed.flow < 1) continue;</pre>
   Ti tr = dfs(ed.u, min(pushed, ed.cap
         - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr;
   edges[id ^ 1].flow -= tr;
   return tr;
 }
 return 0;
void init(int nodes) {
 m = 0, n = nodes;
 for (int i = 0; i < n; i++) level[i]</pre>
      = -1, ptr[i] = 0, adj[i].clear();
void addEdge(int v, int u, Ti cap) {
 edges.emplace_back(v, u, cap), adj[v
      ].push_back(m++);
 edges.emplace_back(u, v, 0), adj[u].
      push_back(m++);
```

```
Ti f = 0;
for (auto &ed : edges) ed.flow = 0;
for (; bfs(s, t); memset(level, -1, n
     * 4)) {
 for (memset(ptr, 0, n * 4); Ti
      pushed = dfs(s, INF, t); f +=
      pushed)
}
return f;
```

```
5.7 Min Cost Max Flow
mt19937 rnd(chrono::steady_clock::now()
    .time_since_epoch().count());
const LL inf = 1e9;
struct edge {
  int v, rev;
  LL cap, cost, flow;
  edge() {}
  edge(int v, int rev, LL cap, LL cost)
      : v(v), rev(rev), cap(cap), cost(
          cost), flow(0) {}
};
struct mcmf {
  int src, sink, n;
  vector<int> par, idx, Q;
  vector<bool> inq;
  vector<LL> dis;
  vector<vector<edge>> g;
  mcmf() {}
  mcmf(int src, int sink, int n)
      : src(src),
        sink(sink),
       n(n),
       par(n),
       idx(n),
       inq(n),
       dis(n),
       g(n),
       Q(10000005) {} // use Q(n) if
            not using random
  void add_edge(int u, int v, LL cap,
      LL cost, bool directed = true) {
    edge _u = edge(v, g[v].size(), cap,
        cost);
    edge _v = edge(u, g[u].size(), 0, -
        cost);
    g[u].pb(_u);
    g[v].pb(_v);
    if (!directed) add_edge(v, u, cap,
        cost, true);
  bool spfa() {
    for (int i = 0; i < n; i++) {</pre>
     dis[i] = inf, inq[i] = false;
    int f = 0, 1 = 0;
    dis[src] = 0, par[src] = -1, Q[1++]
        = src, inq[src] = true;
    while (f < 1) {</pre>
      int u = Q[f++];
      for (int i = 0; i < g[u].size(); i</pre>
          ++) {
        edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
        if (dis[e.v] > dis[u] + e.cost)
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
```

```
if (!inq[e.v]) inq[e.v] = true
             , Q[1++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
         // if (f && rnd() & 7) Q[--f]
              = e.v;
         // else Q[1++] = e.v;
        // }
      }
     inq[u] = false;
   return (dis[sink] != inf);
 pair<LL, LL> solve() {
   LL mincost = 0, maxflow = 0;
   while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
         = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.
           cap - e.flow);
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
         = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck
     }
     mincost += bottleneck * dis[sink],
          maxflow += bottleneck;
   return make_pair(mincost, maxflow);
};
// want to minimize cost and don't care
     about flow
// add edge from sink to dummy sink (
    cap = inf, cost = 0)
// add edge from source to sink (cap =
    inf, cost = 0)
// run mcmf, cost returned is the
    minimum cost
```

# 5.8 Bridge Tree

```
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G,
    int par = -1, int d = 0) {
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1)
       if (low[to] > depth[node]) {
         bridges.emplace_back(node, to)
         components.push_back({});
```

```
for (int x = -1; x != to; x =
              st.top(), st.pop())
           components.back().push_back(
                st.top());
       }
      }
      low[node] = min(low[node], low[to
    }
  }
  if (par == -1) {
    components.push_back({});
    while (!st.empty()) components.back
         ().push_back(st.top()), st.pop
         ();
  }
graph &create_tree() {
  for (auto &comp : components) {
    int idx = tree.addNode();
    for (auto &e : comp) id[e] = idx;
  }
  for (auto &[l, r] : bridges) tree.
      addEdge(id[l], id[r]);
  return tree;
void init(graph &G) {
  int n = G.n;
  depth.assign(n, -1), id.assign(n, -1)
       , low.resize(n);
  for (int i = 0; i < n; i++)</pre>
    if (depth[i] == -1) find_bridges(i,
        G);
```

### 5.9Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
 }
 if (!cnt) return 1;
 sort(vs.begin(), vs.end());
 string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] =
      ++ind;
 int ret = mp[s];
 return ret;
```

## Math

# 6.1 Linear Sieve

```
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5],
      cnt[N + 5], POW[N + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
```

```
fill(prime + 2, prime + N + 1, 1);
SOD[1] = NOD[1] = phi[1] = spf[1] =
for (LL i = 2; i <= N; i++) {</pre>
 if (prime[i]) {
   primes.push_back(i), spf[i] = i;
   phi[i] = i - 1;
   NOD[i] = 2, cnt[i] = 1;
   SOD[i] = i + 1, POW[i] = i;
 for (auto p : primes) {
   if (p * i > N or p > spf[i]) break
   prime[p * i] = false, spf[p * i] =
   if (i % p == 0) {
     phi[p * i] = p * phi[i];
     NOD[p * i] = NOD[i] / (cnt[i] +
          1) * (cnt[i] + 2),
            cnt[p * i] = cnt[i] + 1;
     SOD[p * i] = SOD[i] / SOD[POW[i
          ]] * (SOD[POW[i]] + p * POW[
          i]),
            POW[p * i] = p * POW[i];
     break;
   } else {
     phi[p * i] = phi[p] * phi[i];
     NOD[p * i] = NOD[p] * NOD[i],
         cnt[p * i] = 1;
     SOD[p * i] = SOD[p] * SOD[i],
         POW[p * i] = p;
   }
 }
}
```

```
6.2 Pollard Rho
LL mul(LL a, LL b, LL mod) {
  return (__int128)a * b % mod;
  // LL ans = a * b - mod * (LL) (1.L /
       mod * a * b);
  // return ans + mod * (ans < 0) - mod
       * (ans >= (LL) mod);
LL bigmod(LL num, LL pow, LL mod) {
  LL ans = 1;
  for (; pow > 0; pow >>= 1, num = mul(
      num, num, mod))
    if (pow & 1) ans = mul(ans, num, mod
        );
  return ans;
bool is_prime(LL n) {
  if (n < 2 or n % 6 % 4 != 1) return (</pre>
      n \mid 1) == 3;
  LL a[] = \{2, 325, 9375, 28178,
      450775, 9780504, 1795265022};
  LL s = \_builtin\_ctzll(n - 1), d = n
      >> s;
  for (LL x : a) {
   LL p = bigmod(x % n, d, n), i = s;
    for (; p != 1 and p != n - 1 and x %
         n and i--; p = mul(p, p, n)
    if (p != n - 1 \text{ and } i != s) \text{ return}
        false;
  return true;
LL get_factor(LL n) {
```

```
auto f = [&](LL x) { return mul(x, x, |void mobius() {
      n) + 1; };
 LL x = 0, y = 0, t = 0, prod = 2, i =
       2, q;
 for (; t++ % 40 or gcd(prod, n) == 1;
       x = f(x), y = f(f(y))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) -
       min(x, y), n) ? q : prod;
 }
 return gcd(prod, n);
map<LL, int> factorize(LL n) {
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11,
      13, 17, 19, 23, 29, 31, 37, 41,
                     43, 47, 53, 59,
                         61, 67, 71,
                         73, 79, 83,
                         89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
 auto _factor = [&](LL n, auto &
      _factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
   }
 }:
 _factor(n, _factor);
 return res;
```

# 6.3 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
 if (b == 0)
   return {1, 0, a};
 else {
   auto [x, y, g] = EGCD(b, a \% b);
   return {y, x - a / b * y, g};
 }
// given modulo equations, will apply
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1;
 for (auto &[v, m] : v) { // value %
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1,
   V += x * (v - V) / g % (m / g) * M,
       M *= m / g;
   V = (V \% M + M) \% M;
 }
 return make_pair(V, M);
```

# 6.4 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
```

```
memset(mob, -1, sizeof mob);
 mob[1] = 1;
 for (int i = 2; i < N; i++)</pre>
   if (mob[i]) {
     for (int j = i + i; j < N; j += i)
           mob[j] -= mob[i];
6.5 FFT
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N),
        sin(2 * PI * i / N));
   wp[1][i] = CD(cos(2 * PI * i / N), -
        sin(2 * PI * i / N));
 }
void fft(vector<CD> &v, bool invert =
    false) {
 if (v.size() != perm.size())
      precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i</pre>
        ]]);
 for (int len = 2; len <= N; len *= 2)</pre>
   for (int i = 0, d = N / len; i < N;</pre>
        i += len) {
     for (int j = 0, idx = 0; j < len /
           2; j++, idx += d) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j]
             + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /=</pre>
void pairfft(vector<CD> &a, vector<CD>
    &b, bool invert = false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[</pre>
      i] + b[i] * CD(0, 1);
 fft(p, invert);
 p.push_back(p[0]);
```

for (int i = 0; i < N; i++) {</pre>

```
if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) *
          CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) *
          CD(0, -0.5);
 }
vector<LL> multiply(const vector<LL> &a
    , const vector<LL> &b) {
 int n = 1;
  while (n < a.size() + b.size()) n <<=
  vector<CD> fa(a.begin(), a.end()), fb
      (b.begin(), b.end());
  fa.resize(n);
  fb.resize(n);
  //
           fft(fa); fft(fb);
  pairfft(fa, fb);
  for (int i = 0; i < n; i++) fa[i] =</pre>
      fa[i] * fb[i];
 fft(fa, true);
  vector<LL> ans(n);
  for (int i = 0; i < n; i++) ans[i] =</pre>
      round(fa[i].real());
  return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a,
    const vector<LL> &b) {
 int n = 1:
  while (n < a.size() + b.size()) n <<=</pre>
  vector<CD> al(n), ar(n), bl(n), br(n)
  for (int i = 0; i < a.size(); i++) al</pre>
      [i] = a[i] % M / B, ar[i] = a[i]
      % M % B;
  for (int i = 0; i < b.size(); i++) bl</pre>
      [i] = b[i] % M / B, br[i] = b[i]
      % M % B;
  pairfft(al, ar);
 pairfft(bl, br);
           fft(al); fft(ar); fft(bl);
      fft(br);
  for (int i = 0; i < n; i++) {</pre>
   CD ll = (al[i] * bl[i]), lr = (al[i]
         * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i]
         * br[i]);
   al[i] = 11;
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr;
 pairfft(al, ar, true);
 pairfft(bl, br, true);
           fft(al, true); fft(ar, true)
      ; fft(bl, true); fft(br, true);
  vector<LL> ans(n);
  for (int i = 0; i < n; i++) {</pre>
   LL right = round(br[i].real()), left
         = round(al[i].real());
   LL mid = round(round(bl[i].real()) +
         round(ar[i].real()));
```

```
return ans;
6.6 NTT
const LL N = 1 << 18;</pre>
const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n,
LL Pow(LL b, LL p) {
 LL ret = 1;
  while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) % MOD;
   p >>= 1;
 }
 return ret;
LL primitive_root(LL p) {
  vector<LL> factor;
  LL phi = p - 1, n = phi;
  for (LL i = 2; i * i <= n; i++) {</pre>
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
  }
 if (n > 1) factor.emplace_back(n);
  for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() &&</pre>
        ok; i++)
     ok &= Pow(res, phi / factor[i]) != struct aho_corasick {
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - __builtin_clz(n));
  LL r = Pow(g, (MOD - 1) / n);
  inv_n = Pow(n, MOD - 2);
  w[0] = w[n] = 1;
  for (LL i = 1; i < n; i++) w[i] = (w[</pre>
      i - 1] * r) % MOD;
  for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i &
        1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)</pre>
   if (i < rev[i]) swap(a[i], a[rev[i</pre>
        ]]);
  for (LL m = 2; m <= n; m <<= 1) {</pre>
   for (LL i = 0; i < n; i += m) {</pre>
     for (LL j = 0; j < (m >> 1); j++)
       LL &u = a[i + j], &v = a[i + j +
             (m >> 1)];
       LL t = v * w[dir ? n - n / m * j]
             : n / m * j] % MOD;
       v = u - t < 0 ? u - t + MOD : u
       u = u + t >= MOD ? u + t - MOD :
             u + t;
   }
```

ans[i] = ((left % M) \* B \* B + (mid)

% M) \* B + right) % M;

```
if (dir)
   for (LL i = 0; i < n; i++) a[i] = (</pre>
        inv_n * a[i]) % MOD;
vector<LL> mul(vector<LL> p, vector<LL>
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i
      ];
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
  for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
 for (LL i = 0; i < sz; i++) a[i] = (a
      [i] * b[i]) % MOD;
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
  while (c.size() && c.back() == 0) c.
      pop_back();
 return c;
```

for (char c : s) {

now = now->to[i];

if (!now->to[i]) now->to[i] =

now->leaf = true, now->val++;

for (q.push(root); q.empty(); q.pop

cur->out = link->leaf ? link :

for (auto& next : cur->to) {

if (next != NULL) {

q.push(next);

cur->val += link->val;

node \*cur = q.front(), \*link = cur

next->link = cur != root ?

next = link->to[idx++];

link->to[idx++] : root;

add\_node(now, c, now->len + 1)

int i = f(c);

void push\_links() {

queue<node\*> q;

()) {

->link:

int idx = 0;

} else

}

}

link->out;

```
7.2 Double hash
// define +, -, * for (PLL, LL) and (
    PLL, PLL), % for (PLL, PLL);
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p & 1) ans = ans * a % mod;
 return ans;
PLL inverse(PLL a) { return power(a, (
    mod.ff - 1) * (mod.ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = PLL(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (</pre>
      P[i - 1] * base) % mod;
PLL append(PLL cur, char c) { return (
    cur * base + c) % mod; }
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) {
    return (P[k] * c + cur) % mod; }
/// replaces the i-th (0-indexed)
    character from right from a to b;
PLL replace(PLL cur, int i, char a,
    char b) {
  cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
```

# String

## **Aho Corasick**

```
const int sg = 26, N = 1e3 + 9;
 struct node {
   node *link, *out, *par;
   bool leaf;
   LL val;
   int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '
       $', int len = 0)
       : par(par), p_ch(p_ch), len(len)
     val = leaf = cnt = last = 0;
     link = out = NULL;
   }
 };
 vector<node> trie;
 node* root:
 aho corasick() {
   trie.reserve(N), trie.emplace_back()
   root = &trie[0];
   root->link = root->out = root;
 inline int f(char c) { return c - 'a'
 inline node* add_node(node* par =
     NULL, char p_ch = '$', int len =
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
 void add_str(const string& s, LL val
     = 1) {
   node* now = root;
```

```
/// Erases c from the back of the
    string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod
       + mod) % mod;
/// Erases c from front of the string
    with size len
PLL pop_front(PLL hash, int len, char c
    ) {
 return ((hash - P[len - 1] * c) % mod
       + mod) % mod;
/// concatenates two strings where
    length of the right is k
PLL concat(PLL left, PLL right, int k)
    { return (left * P[k] + right) %
    mod: }
/// Calculates hash of string with size
     len repeated cnt times
/// This is O(\log n). For O(1), pre-
    calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) *
      inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff
 if (P[len].ss == 1) ret.ss = hash.ss
      * cnt;
 return ret;
LL get(PLL hash) { return ((hash.ff <<
    32) ^ hash.ss); }
struct hashlist {
 int len;
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.
       resize(len + 2, PLL(0, 0));
   for (int i = 1; i <= len; i++) H[i]</pre>
       = append(H[i - 1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i]
       = append(R[i + 1], s[i - 1]);
 /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - 1 +
        1]) % mod + mod) % mod;
 PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 +
        1]) % mod + mod) % mod;
 PLL concat_range_hash(int 11, int r1,
       int 12, int r2) {
   return concat(range_hash(l1, r1),
       range_hash(12, r2), r2 - 12 +
 PLL concat_reverse_hash(int l1, int
      r1, int 12, int r2) {
   return concat(reverse_hash(12, r2),
       reverse_hash(l1, r1), r1 - l1 +
        1);
 }
```

### 7.4 Manacher's

```
vector<int> d1(n);
// d[i] = number of palindromes taking
    s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i /*
    ++) {
 int k = (i > r) ? 1 : min(d1[1 + r -
      i], r - i + 1);
 while (0 <= i - k && i + k < n && s[i
       - k] == s[i + k]) k++;
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking
    s[i-1] and s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 0 : min(d2[1 + r -
      i + 1], r - i + 1);
 while (0 <= i - k - 1 && i + k < n &&
      s[i - k - 1] == s[i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i +
```

### 7.5 String Match FFT

```
//find occurrences of t in s where '?'s
     are automatically matched with any
     character
//res[i + m - 1] = sum_j=0 to m - 1_{s[}
    i + j] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
    string &t) {
 int n = s.size(), m = t.size();
 vector<int> s1(n), s2(n), s3(n);
 for(int i = 0; i < n; i++) s1[i] = s[</pre>
      i] == '?' ? 0 : s[i] - 'a' + 1;
      //assign any non zero number for
      non '?'s
 for(int i = 0; i < n; i++) s2[i] = s1</pre>
      [i] * s1[i];
 for(int i = 0; i < n; i++) s3[i] = s1</pre>
      [i] * s2[i];
 vector<int> t1(m), t2(m), t3(m);
 for(int i = 0; i < m; i++) t1[i] = t[</pre>
      i] == '?' ? 0 : t[i] - 'a' + 1;
 for(int i = 0; i < m; i++) t2[i] = t1</pre>
      [i] * t1[i];
  for(int i = 0; i < m; i++) t3[i] = t1</pre>
      [i] * t2[i];
 reverse(t1.begin(), t1.end());
```

```
reverse(t2.begin(), t2.end());
reverse(t3.begin(), t3.end());
vector<int> s1t3 = multiply(s1, t3);
vector<int> s2t2 = multiply(s2, t2);
vector<int> s3t1 = multiply(s3, t1);
vector<int> res(n);
for(int i = 0; i < n; i++) res[i] =
    s1t3[i] - s2t2[i] * 2 + s3t1[i];
vector<int> oc;
for(int i = m - 1; i < n; i++) if(res
    [i] == 0) oc.push_back(i - m + 1)
    ;
return oc;
}</pre>
```

### 7.6 Suffix Array

```
#include <bits/stdc++.h>
using namespace std;
 O(|S| + |alphabet|) Suffix Array
 LIM := \max\{s[i]\} + 2
void inducedSort (const vector <int> &
    vec, int val_range, vector <int> &
    SA, const vector <int> &sl, const
    vector <int> &lms_idx) {
 vector <int> 1(val_range, 0), r(
      val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c
         + 1];
 partial_sum(l.begin(), l.end(), l.
      begin());
 partial_sum(r.begin(), r.end(), r.
      begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i >=
       0; --i) SA[--r[vec[lms_idx[i]]]]
       = lms_idx[i];
 for (int i : SA) if (i > 0 and sl[i -
       1]) SA[1[vec[i-1]]++] = i-1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
 partial_sum(r.begin(), r.end(), r.
      begin());
 for (int k = SA.size() - 1, i = SA[k
      ]; k; --k, i = SA[k]) {
   if (i and !sl[i - 1]) SA[--r[vec[i -
        1]]] = i - 1;
vector <int> suffixArray (const vector
    <int> &vec, int val_range) {
 const int n = vec.size();
```

vector <int> sl(n), SA(n), lms\_idx;

for (int i = n - 2;  $i \ge 0$ ; --i) {

sl[i] = vec[i] > vec[i + 1] or (vec[i])

if (sl[i] and !sl[i + 1]) lms\_idx.

reverse(lms\_idx.begin(), lms\_idx.end

emplace\_back(i + 1);

());

i] == vec[i + 1] and sl[i + 1])

```
inducedSort(vec, val_range, SA, s1,
      lms idx):
 vector <int> new_lms_idx(lms_idx.size
      ()), lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {
   if (SA[i] > 0 \text{ and } !sl[SA[i]] \text{ and } sl[const int MX = 1e6+7, K = 20;
        SA[i] - 1]) new_lms_idx[k++] =
 int cur = 0; SA[n - 1] = 0;
 for (int k = 1; k < new_lms_idx.size</pre>
      (); ++k) {
   int i = new_lms_idx[k - 1], j =
       new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
     SA[j] = ++cur; continue;
   bool flag = 0;
   for (int a = i + 1, b = j + 1; ; ++a
        , ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
     if ((!sl[a] and sl[a - 1]) or (!sl
          [b] and sl[b - 1])) {
       flag = !(!sl[a] and sl[a - 1]
            and !sl[b] and sl[b - 1]);
            break;
     }
   }
   SA[j] = flag ? ++cur : cur;
 }
 for (int i = 0; i < lms_idx.size();</pre>
      ++i) lms_vec[i] = SA[lms_idx[i]];
 if (cur + 1 < lms_idx.size()) {</pre>
   auto lms_SA = suffixArray(lms_vec,
        cur + 1);
   for (int i = 0; i < lms_idx.size();</pre>
        ++i) new_lms_idx[i] = lms_idx[
        lms_SA[i]];
 }
 inducedSort(vec, val_range, SA, s1,
      new_lms_idx); return SA;
vector <int> getSuffixArray (const
    string &s, const int LIM = 128) {
 vector <int> vec(s.size() + 1);
  copy(begin(s), end(s), begin(vec));
      vec.back() = '$';
  auto ret = suffixArray(vec, LIM);
 ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any
    two suffix
vector <int> getLCParray (const string
    &s, const vector <int> &SA) {
 int n = s.size(), k = 0;
 vector <int> lcp(n), rank(n);
  for (int i = 0; i < n; ++i) rank[SA[i</pre>
      ]] = i;
 for (int i = 0; i < n; ++i, k ? --k :</pre>
       0) {
   if (rank[i] == n - 1) {
     k = 0; continue;
   int j = SA[rank[i] + 1];
   while (i + k < n \text{ and } j + k < n \text{ and } s
        [i + k] == s[j + k]) ++k;
```

```
lcp[rank[i]] = k;
 lcp[n - 1] = 0; return lcp;
int lg[MX];
void pre()
   lg[1] = 0;
   for (int i=2; i<MX; i++)</pre>
       lg[i] = lg[i/2]+1;
struct RMQ{
   int N;
   VI v[K];
   RMQ(const VI &a) {
       N = a.size();
       v[0] = a;
       for (int k = 0; (1<<(k+1)) <= N;
            k++) {
           v[k+1].resize(N);
           for (int i = 0; i-1+(1<<(k
               +1)) < N; i++) {
              v[k+1][i] = min(v[k][i],
                   v[k][i+(1<< k)]);
           }
       }
   }
   int findMin(int i, int j) {
       int k = lg[j-i+1];
       return min(v[k][i], v[k][j
           +1-(1<<k)]);
int main() {
 string s; cin >> s;
 for (const int i : getSuffixArray(s))
       printf("%d ", i);
 puts("");
 return 0;
```

### 7.7Trie

```
template<int sz>
struct Trie {
 Trie() : id(1) {
   memset(endMark, 0, sizeof endMark);
   for_each(all(trie), [](vector<int> &
       v) { v.assign(sz, 0); });
 void insert(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0;
```

```
for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false:
     cur = trie[cur][val];
   return endMark[cur];
 }
private:
 int id, endMark[100005];
 vector<int> trie[100005];
```

## 7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     l = r = i;
     while (r < n \&\& s[r] == s[r - 1])
     z[i] = r - 1, r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
       1 = i;
       while (r < n \&\& s[r] == s[r - 1]
            ]) r++;
       z[i] = r - 1, r--;
     }
   }
 }
 return z;
```

### Extra 8

# Make File

```
.PHONY: clean
D ?= 0
ifeq ($(D), 1)
CXXFLAGS=-std=c++2a -g -DMUSHFIQ_DEBUG
      -DDBG_MACRO_NO_WARNING -Wall -
     Wextra -Wpedantic -Wshadow -
     Wformat=2 -Wfloat-equal -
     Wconversion -Wlogical-op -Wshift-
     overflow=2 -Wduplicated-cond -
     Wcast-qual -Wcast-align -Wno-
     variadic-macros -D_GLIBCXX_DEBUG -
     D_GLIBCXX_DEBUG_PEDANTIC -
     fsanitize=address -fsanitize=
     undefined -fno-sanitize-recover -
     fstack-protector -fsanitize-
     address-use-after-scope
else
CXXFLAGS=-02 -std=c++2a
endif
@echo 'Error: No file given.'
clean:
@echo 'Cleaning executables...'
@find . ! -name stress.sh -type f -
     executable -delete
@echo 'Executables cleaned.'
```

```
8.2
    Stress Tester
make $1
make $2
make $3
# $1 is actual code
# $2 is good code
# $3 is generator
for i in 'seq 1 1000'; do
 echo 'Test #'$i
 timeout 5s ./$3 $RANDOM > in
 timeout 5s ./$1 < in> out
 timeout 5s ./$2 < in> ans
 diff ans out || break
done
8.3
     Sublime Build
{
  "shell_cmd": "make $file_base_name &&
      timeout 5s ./$file_base_name <</pre>
      in > out",
 "working_dir": "$file_path",
 "selector": "source.cpp"
```

```
8.4
     Vimrc
nmap <F5> :NERDTreeToggle<CR>
" nnoremap <silent> <expr> <F5> g:
    NERDTree.IsOpen() ? "\:
    NERDTreeClose<CR>" : bufexists(
    expand('%')) ? "\:NERDTreeFind<CR>"
     : "\:NERDTree<CR>"
" Auto import & Compile
:autocmd BufNewFile *.cpp Or ~/template
    .cpp
nnoremap <F4> :!xclip -o -sel clip > ~/
    cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip
     > ~/cp/in <CR><CR>
nnoremap <F6> :!xclip -sel clip % <CR><
inoremap <F6> <ESC>:!xclip -sel clip %
    <CR><CR>
map <F3> :50 vsplit in<CR>:split out<CR
    ><C-w>h
autocmd filetype cpp nnoremap <F9>
    wa \| !make %:r && timeout 5s ./%:r
     < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F9> <ESC
    >:wa \| !make %:r && timeout 5s
    ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp nnoremap <F10>
    wa \| !make clean && make %:r D=1
    && ./%:r < ~/cp/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F10> <
    ESC>:wa \| !make clean && make %:r
    D=1 && ./%:r < ^/cp/in > ^/cp/out <
    CR>
autocmd filetype python nnoremap <F9> :
    wa \| !python % < ~/cp/in> ~/cp/out
```

autocmd filetype python inoremap <F9> <

ESC>:wa \| !python % < ~/cp/in> ~/

set cindent

```
cp/out<CR>
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'),
     col('.')-1, 1) == ")" ? "\<Right>"
     : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'),
     col('.')-1, 1) == "}" ? "\<Right>"
     : "}"
inoremap [ []<left>
inoremap <expr> ] strpart(getline('.'),
     col('.')-1, 1) == "]" ? "\<Right>"
inoremap <expr> " strpart(getline('.'),
     col('.')-1, 1) == "\"" ? "\<Right>
    " : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'),
     col('.')-1, 1) == "\'" ? "\<Right
    >" : "\',\',\<left>"
inoremap <expr> <CR> <sid>
    insert newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col(
      (., .)^{-2}, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<CR>\< nnoremap <S-k> :m .-2<CR>==
      ESC>\0" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>
   insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col(
      (., .) -2, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<space</pre>
      >\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('
     .')-2, 2)
return stridx('(){}[]''', pair) %
     2 == 0 && strlen(pair) == 2 ? "\<
     del>\<bs>" : "\<bs>"
endfunction
set nocompatible
    iMproved, required
                           " required
filetype on
filetype plugin on
filetype plugin indent on
syntax on
set splitright splitbelow
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
set expandtab
set softtabstop=2
set smartindent
set smarttab
set autoindent
```

```
set noerrorbells
set ruler
set guifont=*
set backspace=indent,eol,start
set ignorecase
set incsearch
set nowrap
 set hlsearch
" bubt site
 " set termguicolors
set foldmethod=indent
set nofoldenable
 " set cursorline
 set laststatus=2
set showcmd
set wildmenu
 " colorscheme torte
if !has('nvim')
  set clipboard=unnamedplus
endif
if !has('nvim')
 set ttymouse=xterm2
endif
nnoremap \langle S-j \rangle : m .+1 \langle CR \rangle ==
 vnoremap <S-j> :m '>+1<CR>gv==gv
vnoremap <S-k> :m '<-2<CR>gv==gv
filetype plugin indent on " required
```

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1$$
  
1 and  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$ 

the size n+1 factors.

intersecting chords.

n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

patterns of length 3); that is, the number of  $\binom{n+k+1}{k}$ permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321.

# 9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count of  $P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$ , then, permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$
  
where,  $S(0,0) = 1, S(n,0) = S(0,n) = 0$ 

$$\sum_{k=0}^{n} S(n,k) = n!$$

 $\sum_{k=0}^{k=0}$  The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$
 | Picks Theorem :  $A = i+b/2-1$  | Deragements :  $d(i) = (i-1) \times (d(i-1)+d(i-2))$ 

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 .... i_k.$$

# Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$ where S(0,0) = 1, S(n,0) = S(0,n) = 0  $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets,  $\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$ 

9 Equations and Formulas
9.1 Catalan Numbers  $C_n = \frac{1}{n+1} \binom{2n}{n} \quad C_0 = 1, C_1 = \begin{cases} \text{with each subset containing at least } r \text{ elements. It is denoted by } S_r(n,k) \text{ and obeys the recurrence relation.} \\ kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \\ kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \end{cases}$ The number of ways to completely parenthesize n+1 factors.

Denote the n objects to partition by the integers  $1,2,\ldots,n$ . Define the reduced Stirling numbers of the second kind, defined n+1 factors.

Denote the n objects to partition by the integers  $1,2,\ldots,n$ . Define the reduced Stirling numbers of the second kind, defined n+1 factors. noted  $S^d(n,k)$ , to be the number of ways  $\phi(d)$ The number of triangulations of a convex polygon with n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that all elements in n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the number of nonempty subsets such that n+2 sides (i.e. the nu partitions of polygon into disjoint triangles each subset have pairwise distance at least by using the diagonals).

The number of ways to connect the 2n subset, it is required that  $|i-j| \ge d$ . It  $\left| \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d) \right|$ points on a circle to form n disjoint i.e. non- has been shown that these numbers satisfy, The number of rooted full binary trees with 9.4 Other Combinatorial Identities

Number of permutations of  $1, \ldots, n$  that avoid the pattern 123 (or any of the other  $\sum_{i=0}^{k} \binom{n+i}{i}$  =  $\sum_{i=0}^{k} \binom{n+i}{n}$ 

$$|n, r \in N, n > r, \sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1}$$
If  $P(n) = \sum_{i=r}^{n} {n \choose i} \cdot Q(k)$ , then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$
 where,  $S(0,0) = 1, S(n,0) = S(0,n) = 0$  If  $P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$ , then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

9.5 Different Math Formulas

Picks Theorem: A = i + b/2 - 1|(d(i-1)+d(i-2))|

$$\frac{n}{ab}$$
 -  $\left\{\frac{b'n}{a}\right\}$  -  $\left\{\frac{a'n}{b}\right\}$  + 1

## 9.6 GCD and LCM

if m is any integer, then  $gcd(a + m \cdot b, b) =$ gcd(a,b)

The gcd is a multiplicative function in the following sense: if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$  $\gcd(a_1,b)\cdot\gcd(a_2,b).$ 

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).For non-negative integers a and b, where aand b are not both zero,  $gcd(n^a-1, n^b-1) =$ 

and 
$$b$$
 are not both zero, go  $n^{\gcd(a,b)} - 1$   $\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$ 

$$\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$$

with each subset containing at least 
$$r$$
 elements. It is denoted by  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k)=kS_r(n,k)+\binom{n}{r-1}S_r(n-r+1,k-1)$  Denote the n objects to partition by the integers  $1,2,\ldots,n$ . Define the reduced Stirling numbers of the second kind, denoted  $S^d(n,k)$ , to be the number of ways to partition the integers  $1,2,\ldots,n$  into  $k$  nonempty subsets such that all elements in each subset have pairwise distance at least  $d$ . That is, for any integers  $i$  and  $j$  in a given subset, it is required that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,  $S^d(n,k) = S(n-d+1,k-d+1), n \geq k \geq d$   $i$  and  $i$