# IUT Dhrubotara, IUT

Contents			Extra	a Stress Tester	17 17
1	All Macros 1		8.2	Sublime Build	17 17 17
2	Data Structure         1           2.1 Segment Tree         1           2.2 Persistent Segment Tree         1           2.3 Lazy Persistent Segment Tree         2           2.4 Implicit Segment Tree         2           2.5 DSU With Rollbacks         2           2.6 BIT-2D         3           2.7 Merge Sort Tree         3           2.8 MO with Update         3           2.9 SparseTable (Rectangle Query)         3           2.10 Sparse Table         4		Equa 9.1 9.2 9.3 9.4 9.5	Ations and Formulas Catalan Numbers Stirling Numbers First Kind . Stirling Numbers Second Kind Other Combinatorial Identities Different Math Formulas GCD and LCM	17 19 19 19 19 19 19
3	DP       4         3.1 Convex Hull Trick       4         3.2 Dynamic CHT       4         3.3 Li Chao Tree       5				
4	Geometry       5         4.1 Point       5         4.2 Linear       5         4.3 Circular       6         4.4 Convex       7         4.5 Polygon       9				
5	Graph         9           5.1 LCA, ETT, VT         9           5.2 SCC         10           5.3 Euler Tour on Edge         10           5.4 LCA In O(1)         10           5.5 HLD         11           5.6 Centroid Tree         11           5.7 Dinic Max Flow         11           5.8 Min Cost Max Flow         12           5.9 Bridge Tree         12           5.10 Tree Isomorphism         13           5.11 Grundy         13				
6	Math       13         6.1       Linear Sieve       13         6.2       Pollard Rho       13         6.3       Extended Euclidean       13         6.4       Chinese Remainder Theorem       14         6.5       Mobius Function       14         6.6       FFT       14         6.7       NTT       14				
7	String       15         7.1 Aho Corasick       15         7.2 Double hash       15         7.3 KMP       16         7.4 Manacher's       16         7.5 String Match FFT       16         7.6 Suffix Array       16         7.7 Trie       17         7.8 Z Algo       17				

## All Macros

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack
    :20000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,
    ssse3,sse4,popcnt,abm,mmx,avx,tune=
    native")
#include <ext/pb_ds/assoc_container.hpp</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns
        iterator to the kth largest
        element counting from 0
   //order_of_key(val) --> returns the
        number of items in a set that
        are strictly smaller than our
template <typename DT>
using ordered_set = tree <DT, null_type</pre>
    , less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef LOCAL
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\',' << x <<</pre>
     '\'';}
void show(const string &x) {cerr << '\"</pre>
    ' << x << '\"';}
void show(bool x) {cerr << (x ? "true"</pre>
    : "false");}
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f';</pre>
    show(x.first); cerr << ", "; show(x</pre>
    .second); cerr << '}'; }
template<typename T>
void show(T x) {int f = 0; cerr << "{";</pre>
     for (auto &i: x) cerr << (f++ ? ",</pre>
     " : ""), show(i); cerr << "}";}
void debug_out(string s) {
   cerr << '\n';</pre>
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(),
         '), s.end());
   cerr << "
                 "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ";
   show(t);
   cerr << endl:</pre>
   if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " <<</pre>
    __LINE__ << endl; debug_out(#x, x);
     cerr << endl;</pre>
#else
#define dbg(x...)
```

```
#endif
const int RANDOM = chrono::
    high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) ^ x;
unsigned hash_combine(unsigned a,
    unsigned b) { return a * 31 + b; }
struct chash {
 int operator()(int x) const { return
     hash_f(x); }
typedef gp_hash_table<int, int, chash>
    gp;
gp table;
mt19937_64 rng(atoi(argv[1]));
long long random(long long 1, long long
    r) {
   uniform_int_distribution<long long>
       dist(1, r);
   return dist(rng);
2
    Data Structure
```

#### 2.1 Segment Tree

```
namespace segtree {
 const int N = 1000006;
 using DT = long long;
 using LT = long long;
 constexpr DT I = 0;
 constexpr LT None = 0;
 DT val[N<<2];</pre>
 LT lz[N<<2];
 int L, R;
 void apply(int u, const LT &U, int 1,
       int r) {
   if (U != None) val[u] += (r - 1 + 1) | 2.2 Persistent Segment Tree
         * U;
   lz[u] += U;
 DT merge(const DT &a, const DT &b,
      int 1, int r) {
   return a + b;
 /* -- Do Not Touch Anything Below
     This -- */
 void push(int 1, int r, int u) {
   if(1 == r) return;
   apply(u << 1, lz[u], l, (l + r) >>
       1);
   apply(u << 1 | 1, lz[u], (1 + r + 2)
        >> 1, r);
   lz[u] = None;
 void build(int 1, int r, vector <DT>
      const &v, int u = 1) {
   lz[u] = None;
   if(1 == r) {
```

```
val[u] = v[1];
   return;
 int m = (1 + r) >> 1, lft = u << 1,
     ryt = lft | 1;
 build(1, m, v, lft);
 build(m + 1, r, v, ryt);
 val[u] = merge(val[lft], val[ryt], 1
      , r);
}
void update(int ql,int qr, LT uval,
    int 1 = L, int r = R, int u = 1)
 if (qr < 1 or ql > r) return;
 if(ql \le l and r \le qr) {
   apply(u, uval, l, r);
   return;
 push(1, r, u);
 int m = (1 + r) >> 1, lft = u << 1,
      ryt = lft | 1;
 update(ql, qr, uval, l, m, lft);
 update(ql, qr, uval, m + 1, r, ryt);
 val[u] = merge(val[lft], val[ryt], 1
DT query(int ql, int qr, int l = L,
    int r = R, int u = 1) {
 if (qr < 1 or ql > r) return I;
 if (ql <= l and r <= qr) return val[</pre>
      u];
 push(1, r, u);
 int m = (1 + r) >> 1, lft = u << 1,
     ryt = lft | 1;
 DT ansl = query(ql, qr, l, m, lft);
 DT ansr = query(ql, qr, m + 1, r,
      ryt);
 return merge(ansl, ansr, 1, r);
void init(int _L, int _R, vector <DT>
 L = _L, R = _R;
 build(L, R, v);
```

```
struct Node {
   int 1 = 0, r = 0, val = 0;
tr[20 * N];
int ptr = 0;
int build(int st, int en) {
   int u = ++ptr;
   if (st == en) return u;
   int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   1 = build(st, mid);
   r = build(mid + 1, en);
   val = tr[l].val + tr[r].val;
   return u;
int update(int pre, int st, int en, int
     idx, int v) {
   int u = ++ptr;
   tr[u] = tr[pre];
   if (st == en) {
       tr[u].val += v;
       return u;
```

```
int mid = (st + en) / 2;
   auto& [1, r, val] = tr[u];
   if (idx <= mid) {</pre>
       r = tr[pre].r;
       1 = update(tr[pre].1, st, mid,
           idx, v);
   } else {
       1 = tr[pre].1;
       r = update(tr[pre].r, mid + 1,
           en, idx, v);
   tr[u].val = tr[1].val + tr[r].val;
   return u;
// finding the kth elelment in a range
int query(int left, int right, int st,
    int en, int k) {
   if (st == en) return st;
   int cnt = tr[tr[right].1].val - tr[
        tr[left].1].val;
   int mid = (st + en) / 2;
   if (cnt >= k) return query(tr[left].
        1, tr[right].1, st, mid, k);
   else return query(tr[left].r, tr[
       right].r, mid + 1, en, k - cnt)
int V[N], root[N], a[N];
int main() {
   map<int, int> mp; int n, q;
   cin >> n >> q;
   for (int i = 1; i <= n; i++) cin >>
       a[i], mp[a[i]];
   int c = 0;
   for (auto x : mp) mp[x.first] = ++c,
         V[c] = x.first;
   root[0] = build(1, n);
   for (int i = 1; i <= n; i++) {</pre>
       root[i] = update(root[i - 1], 1,
            n, mp[a[i]], 1);
   while (q--) {
       int 1, r, k; cin >> 1 >> r >> k;
            l++, k++;
       cout << V[query(root[l - 1],</pre>
           root[r], 1, n, k)] << '\n';
   }
   return 0;
```

```
2.3 Lazy Persistent Segment Tree
struct node {
   int val;
   int lazy;
   int left, right;
} nodes[MAXN];
int sz = 0;
int a[MAXN];
int build(int s, int e) {
   int curr = sz;
   sz++;
   if(s == e) {
       nodes[curr].left = nodes[curr].
           right = -1;
       nodes[curr].val = a[s];
       nodes[curr].lazy = 0;
       return curr;
   }
```

```
int m = (s + e) / 2;
   nodes[curr].left = build(s, m);
   nodes[curr].right = build(m + 1, e);
   nodes[curr].val = nodes[nodes[curr].
       left].val + nodes[nodes[curr].
       right].val;
   return curr;
int update(int prev, int s, int e, int
    1, int r, int val) {
   if(s > r or e < 1) return prev;</pre>
   int curr = sz;
   sz++;
   if(s \ge 1 \text{ and } e \le r) {
       nodes[curr].left = nodes[prev].
       nodes[curr].right = nodes[prev].
           right;
       nodes[curr].val = nodes[prev].
           val + val*(e-s+1);
       nodes[curr].lazy = nodes[prev].
           lazy + val;
       return curr;
   int m = (s + e) / 2;
   nodes[curr].left = update(nodes[prev
        ].left, s, m, l, r, val);
   nodes[curr].right = update(nodes[
       prev].right, m+1, e, l, r, val)
   nodes[curr].lazy = nodes[prev].lazy;
   nodes[curr].val = nodes[nodes[curr].
       left].val + nodes[nodes[curr].
       right].val + nodes[curr].lazy*(
       e-s+1);
   return curr;
ll query(int nd, int s, int e, int l,
    int r, 11 c) {
   if(s > r or e < 1) return 0;
   if(s >= 1 and e <= r) return c*(e-s)
       +1) + nodes[nd].val;
   int m = (s+e)/2;
   c += nodes[nd].lazy;
   return query(nodes[nd].left, s, m, 1
        , r, c) + query(nodes[nd].right
        , m+1, e, l, r, c);
```

#### 2.4 Implicit Segment Tree

```
struct node {
   ll val;
   ll lazy;
   int left, right;
node nodes[4*1000006];
int sz;
void reset(int nd) {
   nodes[nd].left = nodes[nd].right =
   nodes[nd].val = nodes[nd].lazy = 0;
int getLeft(int nd) {
   if(nodes[nd].left == -1) {
       nodes[nd].left = sz++;
       reset(nodes[nd].left);
   return nodes[nd].left;
```

```
int getRight(int nd) {
    if(nodes[nd].right == -1) {
       nodes[nd].right = sz++;
       reset(nodes[nd].right);
    return nodes[nd].right;
void update(int nd, int s, int e, int 1
     , int r, ll val) {
    if(e < 1 or s > r) return;
    if(s \ge 1 \text{ and } e \le r) {
       nodes[nd].val += val*(e-s + 1);
       nodes[nd].lazy += val;
       return;
    int m = (s + e) / 2;
    update(getLeft(nd), s, m, l ,r, val)
    update(getRight(nd), m+1, e, 1, r,
        val);
    nodes[nd].val = nodes[nodes[nd].left
        ].val + nodes[nodes[nd].right].
        val + nodes[nd].lazy * (e-s+1);
11 query(int nd, int s, int e, int 1,
    int r, 11 c) {
    if(s > r or e < 1) return 0;</pre>
    if(s \ge 1 \text{ and } e \le r) {
       return nodes[nd].val + c*(e-s+1)
    int m = (s+e)/2;
    c += nodes[nd].lazy;
    return query(getLeft(nd), s, m, l, r
        , c) + query(getRight(nd), m+1,
         e, 1, r, c);
```

## DSU With Rollbacks

```
struct Rollback_DSU {
 int n:
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n,
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par
       [node])
     ; // no path compression
   return node;
 void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)])
        swap(x, y);
   op.emplace_back(x, y);
   par[y] = x;
   sz[x] += sz[y];
 void Undo(int t) {
   for (; op.size() > t; op.pop_back())
     par[op.back().second] = op.back().
     sz[op.back().first] -= sz[op.back
         ().second];
```

```
}
};
     BIT-2D
2.6
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
 for (; x < N; x += -x & x)
   for (int j = y; j < N; j += -j & j)
       bit[x][j] += val;
}
int get(int x, int y) {
 int ans = 0;
 for (; x; x -= x & -x)
   for (int j = y; j; j -= j & -j) ans
       += bit[x][j];
 return ans;
int get(int x1, int y1, int x2, int y2)
 return get(x2, y2) - get(x1 - 1, y2)
      - get(x2, y1 - 1) + get(x1 - 1,
      y1 - 1);
```

## Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];
vector<LL> merge(vector<LL> v1, vector<</pre>
    LL> v2) {
 LL i = 0, j = 0;
 vector<LL> ret;
  while (i < v1.size() || j < v2.size()</pre>
      ) {
   if (i == v1.size()) {
     ret.push_back(v2[j]);
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
     i++;
   } else {
     if (v1[i] < v2[j]) {</pre>
       ret.push_back(v1[i]);
       i++;
     } else {
       ret.push_back(v2[j]);
       j++;
     }
   }
 return ret;
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
```

```
Tree[node] = merge(Tree[leftNode],
      Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1,
    LL r, LL k) {
 if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin
        (), Tree[node].end(), k) -
          Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2
       * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, 1, r,
        query(rightNode, mid + 1, ed, 1
            , r, k);
```

### 2.8 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs =
int arr[N], freq[2 * N], cnt[2 * N], id
    [N], ans[N];
struct query {
 int 1, r, t, L, R;
 query(int l = 1, int r = 0, int t =
      1, int id = -1)
     : l(1), r(r), t(t), L(1 / sz), R(r)
           / sz) {}
 bool operator<(const query &rhs)</pre>
      const {
   return (L < rhs.L) or (L == rhs.L
       and R < rhs.R) or
          (L == rhs.L and R == rhs.R
              and t < rhs.t);</pre>
 }
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]++;
 cnt[freq[arr[idx]]]++;
void remove(int idx) {
 --cnt[freq[arr[idx]]];
 freq[arr[idx]]--;
 cnt[freq[arr[idx]]]++;
void apply(int t) {
 const bool f = 1 <= Up[t].idx and Up[ // Supports Rectangular Query</pre>
      t].idx <= r;
 if (f) remove(Up[t].idx);
 arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
 const bool f = 1 <= Up[t].idx and Up[</pre>
      t].idx <= r;
```

```
if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].last;
 if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0);
int main() {
 int n, q;
  cin >> n >> q;
  int counter = 0;
 map<int, int> M;
 for (int i = 1; i <= n; i++) {
   cin >> arr[i];
   if (!M[arr[i]]) M[arr[i]] = ++
        counter:
   arr[i] = M[arr[i]];
  iota(id, id + N, 0);
  while (q--) {
   int tp, x, y;
   cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
   else {
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = \{x, y, arr[x]\};
     arr[x] = y;
   }
 }
 t = ui;
  cnt[0] = 3 * n;
  sort(id + 1, id + qi + 1, [\&](int x,
      int y) { return Q[x] < Q[y]; });
  for (int i = 1; i <= qi; i++) {</pre>
   int x = id[i];
   while (Q[x].t > t) apply(++t);
   while (Q[x].t < t) undo(t--);
   while (Q[x].1 < 1) add(--1);
   while (Q[x].r > r) add(++r);
   while (Q[x].1 > 1) remove(1++);
   while (Q[x].r < r) remove(r--);</pre>
   ans[x] = mex();
 for (int i = 1; i <= qi; i++) cout <<</pre>
       ans[i] << '\n';
```

## 2.9 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {</pre>
     M[i][j][0][0] = A[i][j];
```

}

```
for (int q = 1; (1 << q) <= N; q++)
     int add = 1 << (q - 1);
     for (int j = 1; j + add \le N; j++) 3
       M[i][j][0][q] = max(M[i][j][0][q]
             - 1], M[i][j + add][0][q -
            1]);
     }
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {
   int add = 1 << (p - 1);</pre>
   for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q
          ++) {
       for (int j = 1; j <= N; j++) {</pre>
         M[i][j][p][q] = max(M[i][j][p]
              - 1][q], M[i + add][j][p -
               1] [q]);
       }
     }
   }
 }
// returns max of all A[i][j], where x1
    =i<=x2 and y1<=j<=y2
int Query(int x1, int y1, int x2, int
    y2) {
 int kX = log2(x2 - x1 + 1);
 int kY = log2(y2 - y1 + 1);
 int addX = 1 << kX;</pre>
 int addY = 1 << kY;</pre>
 int ret1 = max(M[x1][y1][kX][kY], M[
      x1][y2 - addY + 1][kX][kY]);
  int ret2 = max(M[x2 - addX + 1][y1][
      kX][kY],
                M[x2 - addX + 1][y2 -
                    addY + 1][kX][kY];
 return max(ret1, ret2);
```

```
2.10 Sparse Table
// 0-based indexing, query finds in
    range [first, last]
#define lg(x) (31 - __builtin_clz(x))
const int N = 1e5 + 1;
const int K = lg(N);
int a[N], tr[N][K + 1];
namespace sparse_table {
 int f(int p1, int p2) { return min(p1
      , p2); }
 void build() {
   for(int i = 0; i < n; i++)</pre>
     tr[i][0] = a[i];
   for(int j = 1; j \le K; j++) {
     for(int i = 0; i + (1<<j) <= n; i</pre>
          ++)
       tr[i][j] = f(tr[i][j-1], tr[i]
           + (1<<(j - 1))][j - 1]);
   }
 int query(int 1, int r) {
   int d = lg(r - l + 1);
   return f(table[l][d], table[r - (1<<</pre>
        d) + 1][d]);
```

```
\mathbf{DP}
     Convex Hull Trick
3.1
struct line {
 11 m, c;
 line() {}
 line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
 vector<line> lines;
 int ptr = 0;
 convex_hull_trick() {}
 bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.
       m) < 1.0 * (b.c - a.c) * (a.m -
        c.m):
 void add(line L) {
   int sz = lines.size();
   while (sz >= 2 && bad(lines[sz - 2],
        lines[sz - 1], L)) {
     lines.pop_back();
   lines.pb(L);
 ll get(int idx, int x) { return (111
      * lines[idx].m * x + lines[idx].c
      ); }
 11 query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines
        .size() - 1;
   while (ptr < lines.size() - 1 && get</pre>
        (ptr, x) > get(ptr + 1, x)) ptr
        ++:
   return get(ptr, x);
 }
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio:
 int t;
 cin >> t;
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >>
   for (int i = 1; i <= n; i++) dp[i] =</pre>
         0, sum[i] += sum[i - 1];
   convex_hull_trick cht;
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 1
         11 * a * sqr(sum[pos]) - c;
     cht.add(line(211 * a * sum[pos],
         dp[pos] - a * sqr(sum[pos])));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
```

```
3.2 Dynamic CHT
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 IS_QUERY = -(1LL << 62);</pre>
struct line {
  ll m, b;
  mutable function <const line*()> succ
  bool operator < (const line &rhs)</pre>
      const {
    if (rhs.b != IS_QUERY) return m <</pre>
        rhs.m:
    const line *s = succ();
    if (!s) return 0;
    11 x = rhs.m;
    return b - s -> b < (s -> m - m) * x
 }
};
struct HullDynamic : public multiset <</pre>
    line> {
  bool bad (iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
     return y -> m == z -> m && y -> b
          \langle = z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x
        -> m && y -> b <= x -> b;
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z
        -> m - y -> m) >= 1.0 * (y -> b
         -z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
  void insert_line (ll m, ll b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) ==
        end() ? 0 : &*next(y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(
        y))) erase(next(y));
    while (y != begin() && bad(prev(y)))
         erase(prev(y));
  11 eval (ll x) {
    auto 1 = *lower_bound((line) {x,
        IS_QUERY});
    return 1.m * x + 1.b;
int main() {
```

HullDynamic hull;

return 0;

hull.insert\_line(1, 1);

hull.insert\_line(-1, 1);

cout << hull.eval(69) << endl;</pre>

cout << hull.eval(420) << endl;</pre>

# 3.3 Li Chao Tree struct line { LL m, c; line(LL m = 0, LL c = 0) : m(m), c(c)LL calc(line L, LL x) { return 1LL \* L. m \* x + L.c;} struct node { LL m, c; line L; node \*lft, \*rt; node(LL m = 0, LL c = 0, node \*lft =NULL, node \*rt = NULL) : L(line(m, c)), lft(lft), rt(rt) struct LiChao { node \*root; LiChao() { root = new node(); } void update(node \*now, int L, int R, line newline) { int mid = L + (R - L) / 2; line lo = now->L, hi = newline; if (calc(lo, L) > calc(hi, L)) swap( lo, hi); if (calc(lo, R) <= calc(hi, R)) {</pre> now->L = hi;return; now->L = hi;node(): update(now->rt, mid + 1, R, lo); } else { now->L = lo;if (now->lft == NULL) now->lft = new node(); update(now->lft, L, mid, hi); } } LL query(node \*now, int L, int R, LL x) { if (now == NULL) return -inf; int mid = L + (R - L) / 2; $if (x \le mid)$ return max(calc(now->L, x), query(|| now->lft, L, mid, x)); return max(calc(now->L, x), query( now->rt, mid + 1, R, x)); } };

# Geometry

#### Point

```
typedef double Tf;
typedef double Ti; /// use long long
    for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) { return abs(x) < EPS ?</pre>
    0 : (x < 0 ? -1 : 1); }
struct Point {
 Ti x, y;
```

```
) {}
                                                                                                 Point operator+(const Point& u) const
                                                                                                             { return Point(x + u.x, y + u.y)
                                                                                                           ; }
                                                                                                 Point operator-(const Point& u) const
                                                                                                             { return Point(x - u.x, y - u.y) |}
                                                                                                 Point operator*(const LL u) const {
                                                                                                           return Point(x * u, y * u); }
                                                                                                 Point operator*(const Tf u) const {
                                                                                                           return Point(x * u, y * u); }
                                                                                                 Point operator/(const Tf u) const {
                                                                                                           return Point(x / u, y / u); }
                                                                                                 bool operator==(const Point& u) const
                                                                                                     return dcmp(x - u.x) == 0 && dcmp(y)
                                                                                                               - u.y) == 0;
                                                                                                 bool operator!=(const Point& u) const
                                                                                                              { return !(*this == u); }
                                                                                                 bool operator<(const Point& u) const</pre>
                                                                                                     return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x)) \mid
                                                                                                               -u.x) == 0 && dcmp(y - u.y) <
                                                                                                               0);
                                                                                           Ti dot(Point a, Point b) { return a.x *
                                                                                                         b.x + a.y * b.y; }
if (calc(lo, mid) < calc(hi, mid)) { |Ti cross(Point a, Point b) { return a.x | };</pre>
                                                                                                         * b.y - a.y * b.x; }
    if (now->rt == NULL) now->rt = new Tf length(Point a) { return sqrt(dot(a,
                                                                                                         a)): }
                                                                                            Ti sqLength(Point a) { return dot(a, a)
                                                                                                       ; }
                                                                                            Tf distance(Point a, Point b) { return
                                                                                                      length(a - b); }
                                                                                            Tf angle(Point u) { return atan2(u.y, u
                                                                                                       .x); }
                                                                                             // returns angle between oa, ob in (-PI
                                                                                                       , PI]
                                                                                            Tf angleBetween(Point a, Point b) {
                                                                                                Tf ans = angle(b) - angle(a);
                                                                                                 return ans <= -PI ? ans + 2 * PI : (
                                                                                                           ans > PI ? ans - 2 * PI : ans);
                                                                                             ^{\prime\prime} Rotate a ccw by rad radians, Tf Ti
                                                                                            Point rotate(Point a, Tf rad) {
                                                                                                 return Point(a.x * cos(rad) - a.y *
                                                                                                           sin(rad),
                                                                                                                              a.x * sin(rad) + a.y * cos
                                                                                                                                         (rad));
                                                                                             // rotate a ccw by angle th with cos(th
                                                                                                       ) = co \&\& sin(th) = si, tf ti same
                                                                                            Point rotatePrecise(Point a, Tf co, Tf
                                                                                                 return Point(a.x * co - a.y * si, a.y
                                                                                                             * co + a.x * si);
                                                                                            Point rotate90(Point a) { return Point
                                                                                                       (-a.v, a.x); }
                                                                                                 scales vector a by s such that
                                                                                                       length of a becomes s, Tf Ti same
```

```
Point(Ti x = 0, Ti y = 0) : x(x), y(y) Point scale(Point a, Tf s) { return a /
                                             length(a) * s; }
                                        // returns an unit vector perpendicular
                                             to vector a, Tf Ti same
                                       Point normal(Point a) {
                                         Tf l = length(a);
                                         return Point(-a.y / 1, a.x / 1);
                                        // returns 1 if c is left of ab, 0 if
                                           on ab && -1 if right of ab
                                       int orient(Point a, Point b, Point c) {
                                            return dcmp(cross(b - a, c - a));
                                       /// Use as sort(v.begin(), v.end(),
                                            polarComp(0, dir))
                                        /// Polar comparator around O starting
                                            at direction dir
                                        struct polarComp {
                                         Point O, dir;
                                         polarComp(Point 0 = Point(0, 0),
                                             Point dir = Point(1, 0): O(0),
                                              dir(dir) {}
                                         bool half(Point p) {
                                           return dcmp(cross(dir, p)) < 0 ||</pre>
                                                  (dcmp(cross(dir, p)) == 0 \&\&
                                                      dcmp(dot(dir, p)) > 0);
                                         bool operator()(Point p, Point q) {
                                           return make_tuple(half(p), 0) <</pre>
                                               make_tuple(half(q), cross(p, q)
                                         }
                                        struct Segment {
                                         Point a, b;
                                         Segment(Point aa, Point bb) : a(aa),
                                             b(bb) {}
                                        typedef Segment Line;
                                        struct Circle {
                                         Point o;
                                         Tf r;
                                         Circle(Point o = Point(0, 0), Tf r =
                                             0) : o(o), r(r) {}
                                         // returns true if point p is in ||
                                             on the circle
                                         bool contains(Point p) { return dcmp(
                                             sqLength(p - o) - r * r) <= 0; }
                                         // returns a point on the circle rad
                                             radians away from +X CCW
                                         Point point(Tf rad) {
                                           static_assert(is_same<Tf, Ti>::value
                                               );
                                           return Point(o.x + cos(rad) * r, o.y
                                                + sin(rad) * r);
                                         // area of a circular sector with
                                              central angle rad
                                         Tf area(Tf rad = PI + PI) { return
                                             rad * r * r / 2; }
                                         // area of the circular sector cut by
                                              a chord with central angle alpha
                                         Tf sector(Tf alpha) { return r * r *
                                             0.5 * (alpha - sin(alpha)); }
```

# 4.2 Linear

```
// **** LINE LINE INTERSECTION START
```

```
// returns true if point p is on
    segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p))
      == 0 && dcmp(dot(s.a - p, s.b - p
// returns true if segment p && q touch |}
     or intersect
bool segmentsIntersect(Segment p,
    Segment q) {
  if (onSegment(p.a, q) || onSegment(p.
      b, q)) return true;
  \  \  \, \textbf{if} \  \, (\texttt{onSegment}(\texttt{q}.\texttt{a},\ \texttt{p})\  \, |\, | \  \, \texttt{onSegment}(\texttt{q}.
      b, p)) return true;
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 &&
      dcmp(c3) * dcmp(c4) < 0;
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.
      a)) == 0;
// lines are represented as a ray from
    a point: (point, vector)
// returns false if two lines (p, v) &&
     (q, w) are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference, Tf Ti
     Same
bool lineLineIntersection(Point p,
    Point v, Point q, Point w, Point& o
 if (dcmp(cross(v, w)) == 0) return
      false:
 Point u = p - q;
  o = p + v * (cross(w, u) / cross(v, w)
      ));
 return true;
// returns false if two lines p && q
    are parallel or collinear
// true otherwise, intersection point
    is stored at o via reference
bool lineLineIntersection(Line p, Line
    q, Point& o) {
 return lineLineIntersection(p.a, p.b
      - p.a, q.a, q.b - q.a, o);
// returns the distance from point a to |}
     line 1
// **** LINE LINE INTERSECTION FINISH
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(1.b - 1.a, p - 1.a)
      / length(1.b - 1.a));
// returns the shortest distance from
    point a to segment s
Tf distancePointSegment(Point p,
    Segment s) {
  if (s.a == s.b) return length(p - s.a
      );
 Point v1 = s.b - s.a, v2 = p - s.a,
      v3 = p - s.b;
```

```
if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
 else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
   return abs(cross(v1, v2) / length(v1
        ));
// returns the shortest distance from
    \hbox{segment $p$ to segment $q$}
Tf distanceSegmentSegment(Segment p,
    Segment q) {
  if (segmentsIntersect(p, q)) return
      0;
 Tf ans = distancePointSegment(p.a, q)
  ans = min(ans, distancePointSegment(p
      .b, q));
  ans = min(ans, distancePointSegment(q
  ans = min(ans, distancePointSegment(q
      .b, p));
 return ans;
// returns the projection of point p on
     line 1, Tf Ti Same
Point projectPointLine(Point p, Line 1)
 Point v = 1.b - 1.a;
 return 1.a + v * ((Tf)dot(v, p - 1.a)
       / dot(v, v));
4.3 Circular
// Extremely inaccurate for finding
    near touches
// compute intersection of line 1 with
    circle c
// The intersections are given in order
     of the ray (l.a, l.b), Tf Ti same
vector<Point> circleLineIntersection(
    Circle c, Line 1) {
 vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B *
       B - A * C;
  if (D < -EPS) return ret;</pre>
 ret.push_back(l.a + b * (-B - sqrt(D
      + EPS)) / A);
  if (D > EPS) ret.push_back(1.a + b *
      (-B + sqrt(D)) / A);
 return ret;
// signed area of intersection of
    circle(c.o, c.r) &&
 // triangle(c.o, s.a, s.b) [cross(a-o,
    b-o)/2
Tf circleTriangleIntersectionArea(
    Circle c, Segment s) {
 using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
 if (dcmp(distancePointSegment(c.o, s)
       - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b -
         c.o) * (c.r * c.r) / 2.0;
```

// triangle

```
if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB -
      c.r) \ll 0
   return cross(c.o - s.b, s.a - s.b) /
  // three part: (A, a) (a, b) (b, B)
  vector<Point> Sect =
      circleLineIntersection(c, s);
  return circleTriangleIntersectionArea
      (c, Segment(s.a, Sect[0])) +
        circleTriangleIntersectionArea(
             c, Segment(Sect[0], Sect
             [1])) +
        \verb|circleTriangleIntersectionArea| (
             c, Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c
    .r) && simple polyson(p[])
Tf circlePolyIntersectionArea(Circle c,
     Polygon p) {
  Tf res = 0;
  int n = p.size();
  for (int i = 0; i < n; ++i)</pre>
        circleTriangleIntersectionArea(
        c, Segment(p[i], p[(i + 1) % n
        ]));
 return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
    ---> -2
// interior tangents (d = R - r)
    ----> -1
                   (d = 0)
// concentric
                     (R - r < d < R + r)
// secants
    ) ----> 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
int circleCirclePosition(Circle c1,
    Circle c2) {
 Tf d = length(c1.o - c2.o);
  int in = dcmp(d - abs(c1.r - c2.r)),
      ex = dcmp(d - (c1.r + c2.r));
 return in < 0 ? -2 : in == 0 ? -1 :
      ex == 0 ? 1 : ex > 0 ? 2 : 0;
// compute the intersection points
    between two circles c1 && c2, Tf Ti
vector<Point> circleCircleIntersection(
    Circle c1, Circle c2) {
  vector<Point> ret;
 Tf d = length(c1.o - c2.o);
  if (dcmp(d) == 0) return ret;
  if (dcmp(c1.r + c2.r - d) < 0) return
       ret:
  if (dcmp(abs(c1.r - c2.r) - d) > 0)
      return ret;
  Point v = c2.o - c1.o;
  Tf co = (c1.r * c1.r + sqLength(v) -
      c2.r * c2.r) / (2 * c1.r * length
      (v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co,
       -si), c1.r) + c1.o;
```

```
Point p2 = scale(rotatePrecise(v, co, // CAUTION: a[i] = b[i] in case they
       si), c1.r) + c1.o;
 ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
 return ret;
// intersection area between two
    circles c1, c2
Tf circleCircleIntersectionArea(Circle
    c1, Circle c2) {
 Point AB = c2.o - c1.o;
 Tf d = length(AB);
 if (d >= c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.</pre>
      r * c1.r;
 if (d + c2.r <= c1.r) return PI * c2.</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d
       -c2.r*c2.r) / (2.0*c1.r*d
      ));
 Tf alpha2 = acos((c2.r * c2.r + d * d))
       -c1.r * c1.r) / (2.0 * c2.r * d
 return c1.sector(2 * alpha1) + c2.
      sector(2 * alpha2);
// returns tangents from a point p to
    circle c, Tf Ti same
vector<Point> pointCircleTangents(Point
     p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u,
        ang)};
 }
 return ret;
}
// returns the points on tangents that
    touches the circle, Tf Ti Same
vector<Point> pointCircleTangencyPoints
    (Point p, Circle c) {
 Point u = p - c.o;
 Tf d = length(u);
 if (d < c.r)
   return {};
 else if (dcmp(d - c.r) == 0)
                                           }
   return {c.o + u};
 else {
   Tf ang = acos(c.r / d);
                                         4.4
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o + | / / / minkowski sum of two polygons in O(
         rotate(u, ang)};
// for two circles c1 && c2, returns
    two list of points a && b
// such that a[i] is on c1 && b[i] is
    c2 && for every i
// Line(a[i], b[i]) is a tangent to
    both circles
```

```
touch \mid -1 for c1 = c2
int circleCircleTangencyPoints(Circle
    c1, Circle c2, vector<Point> &a,
                            vector<
                                 Point>
                                 &b) {
 a.clear(), b.clear();
 int cnt = 0;
 if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r +
      c2.r;
 if (dcmp(d2 - rdif * rdif) < 0)</pre>
      return 0;
 if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r)
      ) == 0) return -1;
 Tf base = angle(c2.o - c1.o);
 if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++;
   return cnt;
 Tf ang = acos((c1.r - c2.r) / sqrt(d2)
 a.push_back(c1.point(base + ang));
 b.push_back(c2.point(base + ang));
 cnt++:
 a.push_back(c1.point(base - ang));
 b.push_back(c2.point(base - ang));
 cnt++;
 if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++;
 } else if (dcmp(d2 - rsum * rsum) >
      0) {
   Tf ang = acos((c1.r + c2.r) / sqrt(
       d2)):
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang
       ));
   cnt++;
 return cnt;
     Convex
```

```
Polygon minkowskiSum(Polygon A, Polygon
     B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin
      (), A.end()), A.end());
 rotate(B.begin(), min_element(B.begin
      (), B.end()), B.end());
```

```
A.push_back(A[0]);
  B.push_back(B[0]);
  for (int i = 0; i < n; i++) A[i] = A[</pre>
      i + 1] - A[i];
  for (int i = 0; i < m; i++) B[i] = B[</pre>
      i + 1] - B[i];
  Polygon C(n + m + 1);
  C[0] = A.back() + B.back();
  merge(A.begin(), A.end() - 1, B.begin
      (), B.end() - 1, C.begin() + 1,
       polarComp(Point(0, 0), Point(0,
            -1)));
  for (int i = 1; i < C.size(); i++) C[</pre>
      i] = C[i] + C[i - 1];
  C.pop_back();
  return C;
// finds the rectangle with minimum
     area enclosing a convex polygon and
// the rectangle with minimum perimeter
     enclosing a convex polygon
// Tf Ti Same
pair<Tf, Tf>
    rotatingCalipersBoundingBox(const
    Polygon &p) {
  using Linear::distancePointLine;
  int n = p.size();
  int l = 1, r = 1, j = 1;
  Tf area = 1e100;
  Tf perimeter = 1e100;
  for (int i = 0; i < n; i++) {</pre>
    Point v = (p[(i + 1) \% n] - p[i]) /
        length(p[(i + 1) % n] - p[i]);
    while (dcmp(dot(v, p[r % n] - p[i])
        - dot(v, p[(r + 1) % n] - p[i])
        ) < 0)
     r++;
    while (j < r || dcmp(cross(v, p[j %</pre>
        n] - p[i]) -
                        cross(v, p[(j +
                            1) % n] - p[
                            i])) < 0)
      j++;
    while (1 < j ||
          dcmp(dot(v, p[1 % n] - p[i])
               - dot(v, p[(1 + 1) % n] -
               p[i])) > 0)
     1++:
    Tf w = dot(v, p[r \% n] - p[i]) - dot
        (v, p[1 % n] - p[i]);
    Tf h = distancePointLine(p[j % n],
        Line(p[i], p[(i + 1) % n]));
    area = min(area, w * h);
    perimeter = min(perimeter, 2 * w + 2)
         * h);
  return make_pair(area, perimeter);
// returns the left side of polygon u
    after cutting it by ray a->b
Polygon cutPolygon(Polygon u, Point a,
    Point b) {
  using Linear::lineLineIntersection;
  using Linear::onSegment;
  Polygon ret;
  int n = u.size();
```

for (int i = 0; i < n; i++) {</pre>

```
Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0)
       ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0)
       {
     Point t;
     lineLineIntersection(a, b - a, c,
         d - c, t);
     if (onSegment(t, Segment(c, d)))
         ret.push_back(t);
   }
 }
 return ret;
// returns true if point p is in or on
    triangle abc
bool pointInTriangle(Point a, Point b,
    Point c, Point p) {
 return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >=
        dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no
    three collinear points
// returns inside = -1, on = 0, outside
int pointInConvexPolygon(const Polygon
    &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p -
       pt[0])) > 0)
     lo = mid;
   else
     hi = mid;
 }
 bool in = pointInTriangle(pt[0], pt[
      lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p
       - pt[lo - 1])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p -
      pt[lo])) == 0) return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) %
       n], p - pt[(hi + 1) % n])) == 0)
   return 0;
 return -1;
}
// Extreme Point for a direction is the
     farthest point in that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly,
    Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1)
       % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 +
            n) % n], u)) >= 0) {
     return c;
```

```
bool a_up = dcmp(dot(poly[(a + 1) %
       n] - poly[a], u)) >= 0;
   bool c_up = dcmp(dot(poly[(c + 1) %
       n] - poly[c], u)) >= 0;
   bool a_above_c = dcmp(dot(poly[a] -
       poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
       b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
       b = c;
     else
       a = c;
 if (dcmp(dot(poly[a] - poly[(a + 1) % |// Calculate [ACW, CW] tangent pair
       n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n | constexpr int CW = -1, ACW = 1;
         ) % n], u)) > 0)
   return a;
 return b % n;
// For a convex polygon p and a line 1,
     returns a list of segments
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i
    ], p[(i + 1) modulo |p|])
// #1 If a segment is collinear with
    the line, only that is returned
// #2 Else if l goes through i'th point
    , the i'th segment is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(
    const Polygon &p, Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v))
 int rt = extremePoint(p, rotate90(v)
      * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(l.a, l.b, p[(idx - 1 + n)
        % n]) == 0)
     ret.push_back((idx - 1 + n) % n);
   else
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
```

```
for (int i = 0; i < 2; ++i) {</pre>
    int lo = i ? rt : lf;
    int hi = i ? lf : rt;
    int olo = i ? ort : olf;
    while (true) {
      int gap = (hi - lo + n) \% n;
      if (gap < 2) break;</pre>
      int mid = (lo + gap / 2) % n;
      int omid = orient(1.a, 1.b, p[mid
          ]);
      if (!omid) {
        lo = mid;
       break:
      if (omid == olo)
        lo = mid;
      else
        hi = mid;
    ret.push_back(lo);
  return ret;
     from an external point
bool isGood(Point u, Point v, Point Q,
     int dir) {
  return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q,
      int dir) {
  return orient(Q, u, v) == dir ? u : v
Point pointPolyTangent(const Polygon &
    pt, Point Q, int dir, int lo, int
    hi) {
  while (hi - lo > 1) {
    int mid = (lo + hi) / 2;
    bool pvs = isGood(pt[mid], pt[mid -
        1], Q, dir);
    bool nxt = isGood(pt[mid], pt[mid +
        1], Q, dir);
    if (pvs && nxt) return pt[mid];
    if (!(pvs || nxt)) {
      Point p1 = pointPolyTangent(pt, Q,
           dir, mid + 1, hi);
      Point p2 = pointPolyTangent(pt, Q,
           dir, lo, mid - 1);
      return better(p1, p2, Q, dir);
    if (!pvs) {
      if (orient(Q, pt[mid], pt[lo]) ==
          dir)
        hi = mid - 1;
      else if (better(pt[lo], pt[hi], Q,
           dir) == pt[lo])
       hi = mid - 1;
      else
        lo = mid + 1;
    if (!nxt) {
```

```
if (orient(Q, pt[mid], pt[lo]) ==
          dir)
       lo = mid + 1;
     else if (better(pt[lo], pt[hi], Q,
          dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
 }
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++)</pre>
      ret = better(ret, pt[i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(
    const Polygon &pt, Point Q) {
 int n = pt.size();
 Point acw_tan = pointPolyTangent(pt,
      Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q
      , CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

### 4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &
    poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = poly[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0
        && (ret.empty() || b != ret.
        back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p
     of n vertices
Tf signedPolygonArea(const Polygon &p)
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() -</pre>
      1; i++)
   ret += cross(p[i] - p[0], p[i + 1] -
        p[0]);
 return ret / 2;
}
// given a polygon p of n vertices,
    generates the convex hull in in CCW
// Tested on https://acm.timus.ru/
    problem.aspx?space=1&num=1185
// Caution: when all points are
    colinear AND removeRedundant ==
// output will be contain duplicate
    points (from upper hull) at back
Polygon convexHull(Polygon p, bool
    removeRedundant) {
 int check = removeRedundant ? 0 : -1;
```

```
sort(p.begin(), p.end());
  p.erase(unique(p.begin(), p.end()), p
      .end());
  int n = p.size();
  Polygon ch(n + n);
  int m = 0; // preparing lower hull
  for (int i = 0; i < n; i++) {</pre>
    while (m > 1 &&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 1]))
               <= check)
     m--;
    ch[m++] = p[i];
  int k = m; // preparing upper hull
  for (int i = n - 2; i >= 0; i--) {
    while (m > k \&\&
          dcmp(cross(ch[m - 1] - ch[m -
                2], p[i] - ch[m - 2]))
               <= check)
     m--;
    ch[m++] = p[i];
  if (n > 1) m--;
  ch.resize(m);
  return ch;
// returns inside = -1, on = 0, outside
int pointInPolygon(const Polygon &p,
    Point o) {
  using Linear::onSegment;
  int wn = 0, n = p.size();
  for (int i = 0; i < n; i++) {</pre>
    int j = (i + 1) \% n;
    if (onSegment(o, Segment(p[i], p[j])
        ) || o == p[i]) return 0;
    int k = dcmp(cross(p[j] - p[i], o -
        p[i]));
    int d1 = dcmp(p[i].y - o.y);
    int d2 = dcmp(p[j].y - o.y);
    if (k > 0 && d1 <= 0 && d2 > 0) wn
    if (k < 0 && d2 <= 0 && d1 > 0) wn
        --;
  return wn ? -1 : 1;
// Given a simple polygon p, and a line
     1, returns (x, y)
// x = longest segment of 1 in p, y =
    total length of 1 in p.
pair<Tf, Tf> linePolygonIntersection(
    Line 1, const Polygon &p) {
  using Linear::lineLineIntersection;
  int n = p.size();
  vector<pair<Tf, int>> ev;
  for (int i = 0; i < n; ++i) {</pre>
    Point a = p[i], b = p[(i + 1) \% n],
        z = p[(i - 1 + n) \% n];
    int ora = orient(1.a, 1.b, a), orb =
         orient(1.a, 1.b, b),
       orz = orient(l.a, l.b, z);
    if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
      if (orz && orb) {
       if (orz != orb) ev.emplace_back(
            d, 0);
```

```
// else // Point Touch
   } else if (orz)
     ev.emplace_back(d, orz);
   else if (orb)
     ev.emplace_back(d, orb);
 } else if (ora == -orb) {
   Point ins;
   lineLineIntersection(1, Line(a, b)
        , ins);
   ev.emplace_back(dot(ins - 1.a, 1.b
         - 1.a), 0);
 }
sort(ev.begin(), ev.end());
Tf ans = 0, len = 0, last = 0, tot =
bool active = false;
int sign = 0;
for (auto &qq : ev) {
 int tp = qq.second;
 Tf d = qq.first; /// current Segment
       is (last, d)
                  /// On Border
 if (sign) {
   len += d - last;
   tot += d - last;
   ans = max(ans, len);
   if (tp != sign) active = !active;
   sign = 0;
 } else {
   if (active) { /// Strictly Inside
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
   if (tp == 0)
     active = !active;
   else
     sign = tp;
 last = d;
 if (!active) len = 0;
ans /= length(l.b - l.a);
tot /= length(l.b - l.a);
return {ans, tot};
```

### Graph

#### LCA, ETT, VT

```
#define lg(n) (31 - __builtin_clz(n))
const int N = 1e5 + 1;
const int K = lg(N) + 1;
vector<int> adj[N];
int anc[N][K], lvl[N];
namespace lca {
 void init(int u = 1, int p = 0, int d
       = 0) {
   lvl[u] = d;
   anc[u][0] = p;
   for (int i = 1; i < K; i++)</pre>
     anc[u][i] = anc[anc[u][i - 1]][i -
   for (auto v : adj[u])
     if (v != p)
       init(v, u, d + 1);
  int getAnc(int u, int k) {
```

```
for (int i = 0; u and i < K; i++)</pre>
                                                                                     out[node] = t;
                                             while (!tour.isAncestor(nodes[i],
                                                 nodes[st.top()])) st.pop();
                                                                                     for (int e : G[node]) {
     if ((k >> i) & 1)
                                                                                       int v = G(e).to(node);
       u = anc[u][i];
                                             ans.addEdge(st.top(), i);
                                             st.push(i);
                                                                                       if (v == par) continue;
   return u;
                                                                                       fwd[e] = t++;
 int lca(int u, int v) {
                                           return ans;
                                                                                       dfs(G, v, node);
   if (lvl[u] < lvl[v]) swap(u, v);</pre>
                                                                                       bck[e] = t++;
   u = getAnc(u, lvl[u] - lvl[v]);
                                          set<int> getCenters(tree &T) {
   if (u == v) return u;
                                           int n = T.n;
                                                                                     in[node] = t - 1;
   for (int i = K - 1; ~i; i--)
                                           vector<int> deg(n), q;
     if (anc[u][i] != anc[v][i])
                                           set<int> s;
                                                                                    void init(graph &G, int node) {
       u = anc[u][i], v = anc[v][i];
                                           for (int i = 0; i < n; i++) {</pre>
                                             deg[i] = T[i].size();
                                                                                     dfs(G, node, node);
   return anc[u][0];
 }
                                             if (deg[i] == 1) q.push_back(i);
 int dist(int u, int v) {
                                             s.insert(i);
                                                                                    5.4 LCA In O(1)
   return lvl[u] + lvl[v] - 2 * lvl[lca
        (u, v)];
                                           for (vector<int> t; s.size() > 2; q =
                                                                                   /* LCA in O(1)
                                                                                    * depth calculates weighted distance
                                             for (auto x : q) {
                                                                                     * level calculates distance by number
struct euler_tour {
                                               for (auto e : T[x])
                                                                                         of edges
 int time = 0;
                                                 if (--deg[e] == 1) t.push_back(e
                                                                                    * Preprocessing in NlongN */
 tree &T;
                                                     );
                                                                                    LL depth[N];
 int n;
                                               s.erase(x);
                                                                                    int level[N];
 vector<int> start, finish, level, par
                                                                                    int st[N], en[N], LOG[N], par[N];
 euler_tour(tree &T, int root = 0)
                                           return s:
                                                                                    int a[N], id[N], table[L][N];
     : T(T), n(T.n), start(n), finish(n)
         ), level(n), par(n) {
                                                                                    vector<PII> adj[N];
                                          5.2 SCC
   time = 0;
                                                                                   int n, root, Time, cur;
   call(root);
                                          typedef long long LL;
 }
                                          const LL N = 1e6 + 7;
                                                                                    void init(int nodes, int root_) {
 void call(int node, int p = -1) {
                                                                                     n = nodes, root = root_, LOG[0] = LOG
   if (p != -1) level[node] = level[p]
                                          bool vis[N];
                                                                                          [1] = 0;
        + 1:
                                          vector<int> adj[N], adjr[N];
                                                                                     for (int i = 2; i <= n; i++) LOG[i] =</pre>
   start[node] = time++;
                                          vector<int> order, component;
                                                                                          LOG[i >> 1] + 1;
   for (int e : T[node])
                                          // tp = 0 ,finding topo order, tp = 1 ,
                                                                                     for (int i = 0; i <= n; i++) adj[i].</pre>
     if (e != p) call(e, node);
                                               reverse edge traversal
                                                                                          clear();
   par[node] = p;
   finish[node] = time++;
                                          void dfs(int u, int tp = 0) {
                                           vis[u] = true;
                                                                                    void addEdge(int u, int v, int w) {
 bool isAncestor(int node, int par) {
                                           if (tp) component.push_back(u);
                                                                                     adj[u].push_back(PII(v, w));
   return start[par] <= start[node] and</pre>
                                           auto& ad = (tp ? adjr : adj);
                                                                                     adj[v].push_back(PII(u, w));
         finish[par] >= finish[node];
                                           for (int v : ad[u])
                                             if (!vis[v]) dfs(v, tp);
 int subtreeSize(int node) { return
                                           if (!tp) order.push_back(u);
                                                                                    int lca(int u, int v) {
      finish[node] - start[node] + 1 >> |}
                                                                                     if (en[u] > en[v]) swap(u, v);
                                          int main() {
                                                                                     if (st[v] <= st[u] && en[u] <= en[v])</pre>
                                           for (int i = 1; i <= n; i++) {</pre>
                                                                                           return v;
tree virtual_tree(vector<int> &nodes,
                                             if (!vis[i]) dfs(i);
    lca_table &table, euler_tour &tour)
                                                                                     int 1 = LOG[id[v] - id[u] + 1];
                                           memset(vis, 0, sizeof vis);
                                                                                     int p1 = id[u], p2 = id[v] - (1 << 1)
 sort(nodes.begin(), nodes.end(),
                                           reverse(order.begin(), order.end());
      [&](int x, int y) { return tour.
                                           for (int i : order) {
                                                                                     int d1 = level[table[1][p1]], d2 =
          start[x] < tour.start[y]; });</pre>
                                             if (!vis[i]) {
                                                                                          level[table[1][p2]];
 int n = nodes.size();
                                               // one component is found
 for (int i = 0; i + 1 < n; i++)</pre>
                                               dfs(i, 1), component.clear();
                                                                                     if (d1 < d2)
   nodes.push_back(table.lca(nodes[i],
                                                                                       return par[table[1][p1]];
       nodes[i + 1]));
                                           }
 sort(nodes.begin(), nodes.end());
                                                                                       return par[table[1][p2]];
 nodes.erase(unique(nodes.begin(),
                                          5.3 Euler Tour on Edge
      nodes.end()), nodes.end());
 sort(nodes.begin(), nodes.end(),
                                          // for simplicity, G[idx] contains the
                                                                                   LL dist(int u, int v) {
      [&](int x, int y) { return tour.
                                              adjacency list of a node
                                                                                     int 1 = lca(u, v);
           start[x] < tour.start[y]; }); // while G(e) is a reference to the e-
                                                                                     return (depth[u] + depth[v] - (depth[
 n = nodes.size();
                                              th edge.
                                                                                          1] * 2));
 stack<int> st;
                                          const int N = 2e5 + 5;
 st.push(0);
                                          int in[N], out[N], fwd[N], bck[N];
 tree ans(n):
                                          int t = 0;
                                                                                    /* Euler tour */
 for (int i = 1; i < n; i++) {</pre>
                                          void dfs(graph &G, int node, int par) { void dfs(int u, int p) {
```

int query(int a, int b) {

```
st[u] = ++Time, par[u] = p;
                                            int ret = 0;
                                                head[b]]) {
 for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
                                                   swap(a, b);
   level[v] = level[u] + 1;
   dfs(v, u);
 en[u] = ++Time;
 a[++cur] = u, id[u] = cur;
                                                pos[b]);
                                           return ret;
/* RMQ */
void pre() {
                                               Centroid Tree
                                          5.6
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table</pre>
      [0][i] = a[i];
                                            sz[u] = 1;
 for (int 1 = 0; 1 < L - 1; 1++) {</pre>
                                             if (v != p) {
   for (int i = 1; i <= n; i++) {</pre>
                                               dfs_size(v, u);
     table[1 + 1][i] = table[1][i];
                                               sz[u] += sz[v];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 << int findCentroid(int u, int p) {</pre>
           1)]] < level[table[1][i]];</pre>
                                            int total = sz[u];
     if (C1 && C2) table[l + 1][i] =
          table[1][i + (1 << 1)];
   }
 }
                                               sz[v] = total;
}
                                             }
5.5 HLD
                                           return u;
const int N = 1e6 + 7;
                                          int query(int u) {
template <typename DT>
                                           int ans = 1e6;
struct Segtree {
 // write lazy segtree here
                                                  , u));
Segtree<int> tree(N);
vector<int> adj[N];
                                           return ans;
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
                                          void update(int u) {
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
                                                  u));
 depth[u] = depth[p] + 1;
 for (auto v : adj[u]) {
   if (v == p) continue;
                                            vis[u] = 1;
   par[v] = u;
   int subsz = dfs(v, u);
                                             if (not vis[v])
   if (subsz > mxsz) heavy[u] = v, mxsz
                                           return u;
         = subsz;
   SZ += subsz;
 }
                                          5.7
 return SZ;
                                          /**
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u
      ], h);
                                              blog/entry/66006)
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
```

```
for (; head[a] != head[b]; b = par[
   if (depth[head[a]] > depth[head[b]]) namespace Dinic {
   ret += tree.query(1, 0, cnt, pos[
       head[b]], pos[b]);
                                                 capacity
 if (depth[a] > depth[b]) swap(a, b);
 ret += tree.query(1, 0, cnt, pos[a],
                                                 flow; };
                                            int s, t, n;
void dfs_size(int u = 1, int p = 0) {
                                                n = nodes;
 for (auto v : adj[u])
 for (auto v : adj[u])
   if (v != p and not vis[v] and 2 * sz
        [v] > total) {
     sz[u] = total - sz[v];
     return findCentroid(v, u);
                                            }
 for (int i = u; i; i = par[i])
   ans = min(ans, minD[i] + lca::dist(i
                                                q.push(s);
 for (int i = u; i; i = par[i])
   minD[i] = min(minD[i], lca::dist(i,
int decompose(int u, int p) {
 u = findCentroid(u, p);
 for (auto v : adj[u])
     par[decompose(v, u)] = u;
                                                       }
                                                    }
                                                }
     Dinic Max Flow
Implementation of Dinic's algorithm
    with optional scaling
Source: Chilli (https://codeforces.com/
Complexity: O(ans*E) or O(V^2E) without
     scaling, O(VE log(U)) with scaling
Scaling performs much better in worst
```

case, but has much higher constant

factor

```
To enable scaling, call maxFlow(true)
Everything 0-indexed
   typedef long long LL;
   const int N = 5005, K = 60; /// N >
        no of nodes, K >= max bits in
   const LL INF = 1e18;
   struct Edge { int frm, to; LL cap,
   int level[N], ptr[N];
   vector<Edge> edges;
   vector<int> adj[N];
   void init(int nodes) {
       for (int i=0; i<n; i++) adj[i].</pre>
            clear();
       edges.clear();
   /// For adding undirected Edge (u, v
        , c) call addEdge(u, v, c, c);
   int addEdge(int a, int b, LL cap, LL
         revcap = 0) {
       edges.push_back({a, b, cap, 0});
       edges.push_back({b, a, revcap,
       adj[a].push_back(edges.size()-2)
       adj[b].push_back(edges.size()-1)
       return edges.size()-2;
   bool bfs(LL lim) {
       fill(level, level+n, -1);
       level[s] = 0;
       queue<int> q;
       while (!q.empty() && level[t] ==
             -1) {
           int v = q.front();
           q.pop();
           for (int id: adj[v]) {
              Edge e = edges[id];
               if (level[e.to] == -1 &&
                   e.cap - e.flow >= lim
                   ) {
                  q.push(e.to);
                  level[e.to] = level[v
                       ] + 1;
       return level[t] != -1;
   LL dfs(int v, LL flow) {
       if (v == t || !flow)
                                 return
       for (; ptr[v] < adj[v].size();</pre>
           ptr[v]++) {
           int eid = adj[v][ptr[v]];
           Edge &e = edges[eid];
           if (level[e.to] != level[v]
               + 1) continue;
           if (LL pushed = dfs(e.to,
               min(flow, e.cap - e.flow
               ))) {
```

```
e.flow += pushed;
           edges[eid^1].flow -=
               pushed;
           return pushed;
   }
   return 0;
}
LL maxFlow(int source, int sink,
    bool SCALING = false) {
   s = source, t = sink;
   long long flow = 0;
   for (LL lim = SCALING ? (1LL <<</pre>
        K) : 1; \lim > 0; \lim >>= 1)
       while (bfs(lim)) {
           fill(ptr, ptr+n, 0);
           while (LL pushed = dfs(s,
                INF)) flow += pushed
   }
   return flow;
bool leftOfMinCut(int x) {return
    level[x] != -1;
/// Only works for undirected graph,
     Make sure to add UNDIRECTED
    edges. (u, v, c, c)
/// returns n by n matrix flow, st
    flow[i][j] = maxFlow
/// tree holds the edges of a gomory
    -hu tree of the graph
vector<vector<LL>> allPairMaxFlow(
    vector<Edge> &tree) {
   tree.clear();
   vector<vector<LL>> flow(n,
        vector<LL> (n, INF));
   vector<int> par(n);
   for (int i=1; i<n; i++) {</pre>
       for (auto &e: edges) e.flow
       LL f = maxFlow(i, par[i]);
       tree.push_back({i, par[i], f
           });
       for (int j=i+1; j<n; j++)</pre>
           if (par[j] == par[i] &&
               leftOfMinCut(j)) par[
               j] = i;
       flow[i][par[i]] = flow[par[i
           ]][i] = f;
       for (int j=0; j<i; j++)</pre>
           if (j != par[i]) flow[i
               ][j] = flow[j][i] =
               min(f, flow[par[i]][j
               ]);
   return flow;
}
```

#### 5.8 Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now()
    .time_since_epoch().count());
const LL inf = 1e9;
struct edge {
```

```
int v, rev;
 LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(
          cost), flow(0) {}
struct mcmf {
 int src, sink, n;
 vector<int> par, idx, Q;
 vector<bool> inq;
 vector<LL> dis;
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int n)
     : src(src),
       sink(sink),
       n(n),
       par(n),
       idx(n),
       inq(n),
       dis(n),
       g(n),
       Q(10000005) {} // use Q(n) if
           not using random
 void add_edge(int u, int v, LL cap,
      LL cost, bool directed = true) {
   edge _u = edge(v, g[v].size(), cap,
   edge _v = edge(u, g[u].size(), 0, -
       cost);
   g[u].pb(_u);
   g[v].pb(_v);
   if (!directed) add_edge(v, u, cap,
       cost, true);
 bool spfa() {
   for (int i = 0; i < n; i++) {</pre>
     dis[i] = inf, inq[i] = false;
   int f = 0, 1 = 0;
   dis[src] = 0, par[src] = -1, Q[1++]
       = src, inq[src] = true;
   while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i</pre>
         ++) {
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost)
         dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
         if (!inq[e.v]) inq[e.v] = true
              , Q[1++] = e.v;
         // if (!inq[e.v]) {
         // inq[e.v] = true;
            if (f && rnd() & 7) Q[--f]
              = e.v;
             else Q[1++] = e.v;
         // }
       }
     inq[u] = false;
   return (dis[sink] != inf);
 }
 pair<LL, LL> solve() {
```

LL mincost = 0, maxflow = 0;

```
while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
         = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.
            cap - e.flow);
     for (int u = par[sink], v = idx[
         sink]; u != -1; v = idx[u], u
          = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck
     mincost += bottleneck * dis[sink],
           maxflow += bottleneck;
   return make_pair(mincost, maxflow);
 }
};
// want to minimize cost and don't care
     about flow
// add edge from sink to dummy sink (
    cap = inf, cost = 0)
// add edge from source to sink (cap =
    inf, cost = 0)
// run mcmf, cost returned is the
    minimum cost
```

#### 5.9Bridge Tree

();

```
vector<vector<int>> components;
vector<int> depth, low;
stack<int> st;
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G,
    int par = -1, int d = 0) {
 low[node] = depth[node] = d;
 st.push(node);
 for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1)
       if (low[to] > depth[node]) {
         bridges.emplace_back(node, to)
         components.push_back({});
         for (int x = -1; x != to; x =
             st.top(), st.pop())
           components.back().push_back(
               st.top());
       }
     }
     low[node] = min(low[node], low[to
         ]);
   }
 }
 if (par == -1) {
   components.push_back({});
   while (!st.empty()) components.back
        ().push_back(st.top()), st.pop
```

```
graph &create_tree() {
 for (auto &comp : components) {
   int idx = tree.addNode();
   for (auto &e : comp) id[e] = idx;
 for (auto &[1, r] : bridges) tree.
      addEdge(id[1], id[r]);
 return tree;
}
void init(graph &G) {
 int n = G.n;
 depth.assign(n, -1), id.assign(n, -1)
      , low.resize(n);
 for (int i = 0; i < n; i++)</pre>
   if (depth[i] == -1) find_bridges(i,
```

### 5.10 Tree Isomorphism

```
LL Hash(int u, int p) {
 vector<LL> childrenHash;
 for (auto v : adj[u]) if (v != p)
   childrenHash.add(Hash(v, u));
 sort(all(childrenHash));
 LL nodeHash = 0:
 for (int i = 0; i < childrenHash.size</pre>
      (); i++)
   nodeHash = (nodeHash + childrenHash[
        i] * bigmod(SEED, i, MOD)) %
        MOD:
 return nodeHash;
}
```

#### Grundy 5.11

```
single pile game-> greedy or game dp
multiple pile game and disjunctive(
    before playing, choose 1 pile) ->
    NIM game
else-> Grundy(converts n any game piles
    to n NIM piles)
```

grundy(x)->the smallest nonreachable grundy value

```
there are n pile of games and k type of
     moves.
if XOR(grundy(games)) == 0: losing
    state
else winning state
vector<int> moves, dp;
int mex(vector<int> &a) {
   set<int> b(a.begin(), a.end());
   for (int i = 0; ; ++i)
       if (!b.count(i))
          return i;
}
int grundy(int x) {
if (dp[x] != -1) return dp[x];
vector<int> reachable;
for (auto m : moves) {
 if (x - m < 0) continue;</pre>
 int val = grundy(x - m);
```

reachable.push\_back(val);

```
return dp[x] = mex(reachable);
6
    Math
6.1
   Linear Sieve
using ULL = unsigned long long;
namespace sieve{
 const int N = 1e7;
 vector <int> primes;
```

```
int spf[N+5], phi[N+5], NOD[N+5], cnt
    [N+5], POW[N+5];
bool prime[N+5];
int SOD[N+5];
void init(){
 fill(prime+2, prime+N+1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] =
 for(LL i=2;i<=N;i++){</pre>
   if(prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i-1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i+1, POW[i] = i;
   for(auto p:primes){
     if(p*i>N or p > spf[i]) break;
     prime[p*i] = false, spf[p*i] = p |map<LL, int> factorize(LL n) {
     if(i\%p == 0){
       phi[p*i]=p*phi[i];
       NOD[p*i]=NOD[i]/(cnt[i]+1)*(
           cnt[i]+2), cnt[p*i]=cnt[i
           ]+1;
       SOD[p*i]=SOD[i]/SOD[POW[i]]*(
           SOD[POW[i]]+p*POW[i]),POW[
           p*i]=p*POW[i];
       break;
     } else {
       phi[p*i]=phi[p]*phi[i];
       NOD[p*i]=NOD[p]*NOD[i], cnt[p*
           i]=1;
       SOD[p*i]=SOD[p]*SOD[i], POW[p*
           i]=p;
     }
   }
```

#### 6.2Pollard Rho

}

```
LL mul(LL a, LL b, LL mod) {
 return (__int128)a * b % mod;
 // LL ans = a * b - mod * (LL) (1.L /
      mod * a * b):
 // return ans + mod * (ans < 0) - mod
       * (ans >= (LL) mod);
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(
      num, num, mod))
   if (pow & 1) ans = mul(ans, num, mod
       );
 return ans;
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (</pre>
      n \mid 1) == 3;
```

```
LL a[] = \{2, 325, 9375, 28178,
      450775, 9780504, 1795265022};
 LL s = \__builtin\_ctzll(n - 1), d = n
      >> s;
 for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x %
        n and i--; p = mul(p, p, n))
   if (p != n - 1 and i != s) return
        false;
 return true;
LL get_factor(LL n) {
 auto f = [&](LL x) { return mul(x, x,
       n) + 1; };
 LL x = 0, y = 0, t = 0, prod = 2, i =
       2, q;
 for (; t++ % 40 or gcd(prod, n) == 1;
       x = f(x), y = f(f(y)) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) -
        min(x, y), n) ? q : prod;
 return gcd(prod, n);
 map<LL, int> res;
 if (n < 2) return res;</pre>
 LL small_primes[] = \{2, 3, 5, 7, 11,
      13, 17, 19, 23, 29, 31, 37, 41,
                     43, 47, 53, 59,
                         61, 67, 71,
                         73, 79, 83,
                         89, 97};
 for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
 auto _factor = [&](LL n, auto &
      _factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++;
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
 };
 _factor(n, _factor);
 return res;
```

#### 6.3 Extended Euclidean

```
int gcd(int a, int b, int& x, int& y) {
 x = 1, y = 0;
 int x1 = 0, y1 = 1, a1 = a, b1 = b;
 while (b1) {
   int q = a1 / b1;
   tie(x, x1) = make_tuple(x1, x - q *
       x1);
   tie(y, y1) = make_tuple(y1, y - q *
       y1);
   tie(a1, b1) = make_tuple(b1, a1 - q
 return a1;
```

```
6.4 Chinese Remainder Theorem
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
 if (b == 0)
   return {1, 0, a};
   auto [x, y, g] = EGCD(b, a \% b);
   return {y, x - a / b * y, g};
 }
// given modulo equations, will apply
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1;
 for (auto &[v, m] : v) { // value %
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1,
   V += x * (v - V) / g % (m / g) * M,
       M *= m / g;
   V = (V \% M + M) \% M;
 return make_pair(V, M);
```

#### 6.5 Mobius Function

```
const int N = 1e6 + 5;
int mob[N]:
void mobius() {
 memset(mob, -1, sizeof mob);
 mob[1] = 1;
 for (int i = 2; i < N; i++)</pre>
   if (mob[i]) {
     for (int j = i + i; j < N; j += i)
          mob[j] -= mob[i];
   }
```

#### 6.6 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
 }
 wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N),
        sin(2 * PI * i / N));
   wp[1][i] = CD(cos(2 * PI * i / N), -
        sin(2 * PI * i / N));
void fft(vector<CD> &v, bool invert =
    false) {
```

```
if (v.size() != perm.size())
      precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i</pre>
       ]]);
 for (int len = 2; len <= N; len *= 2)</pre>
   for (int i = 0, d = N / len; i < N;</pre>
        i += len) {
     for (int j = 0, idx = 0; j < len /
           2; j++, idx += d) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j]
            + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
  if (invert) {
   for (int i = 0; i < N; i++) v[i] /=</pre>
void pairfft(vector<CD> &a, vector<CD>
    &b, bool invert = false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[</pre>
      i] + b[i] * CD(0, 1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) *
          CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) *
          CD(0, -0.5);
   }
 }
vector<LL> multiply(const vector<LL> &a
    , const vector<LL> &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<= |LL Pow(LL b, LL p) {
  vector<CD> fa(a.begin(), a.end()), fb
      (b.begin(), b.end());
 fa.resize(n);
 fb.resize(n);
 //
           fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] =</pre>
      fa[i] * fb[i];
 fft(fa, true);
 vector<LL> ans(n);
 for (int i = 0; i < n; i++) ans[i] =</pre>
      round(fa[i].real());
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a,
    const vector<LL> &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<=</pre>
```

```
vector<CD> al(n), ar(n), bl(n), br(n)
for (int i = 0; i < a.size(); i++) al</pre>
    [i] = a[i] % M / B, ar[i] = a[i]
    % M % B;
for (int i = 0; i < b.size(); i++) bl</pre>
    [i] = b[i] % M / B, br[i] = b[i]
    % M % B;
pairfft(al, ar);
pairfft(bl, br);
         fft(al); fft(ar); fft(bl);
    fft(br);
for (int i = 0; i < n; i++) {</pre>
  CD 11 = (al[i] * bl[i]), lr = (al[i]
       * br[i]);
  CD rl = (ar[i] * bl[i]), rr = (ar[i]
       * br[i]);
  al[i] = 11;
  ar[i] = lr;
  bl[i] = rl;
 br[i] = rr;
pairfft(al, ar, true);
pairfft(bl, br, true);
         fft(al, true); fft(ar, true)
    ; fft(bl, true); fft(br, true);
vector<LL> ans(n);
for (int i = 0; i < n; i++) {</pre>
  LL right = round(br[i].real()), left
       = round(al[i].real());
 LL mid = round(round(bl[i].real()) +
       round(ar[i].real()));
  ans[i] = ((left \% M) * B * B + (mid)
      % M) * B + right) % M;
return ans;
```

#### 6.7 NTT

```
const LL N = 1 << 18;</pre>
 const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N \mid 1], a[N], b[N], inv_n,
  LL ret = 1;
  while (p) {
    if (p & 1) ret = (ret * b) % MOD;
    b = (b * b) \% MOD;
    p >>= 1;
  return ret;
LL primitive_root(LL p) {
  vector<LL> factor;
  LL phi = p - 1, n = phi;
  for (LL i = 2; i * i <= n; i++) {</pre>
    if (n % i) continue;
    factor.emplace_back(i);
    while (n \% i == 0) n /= i;
  if (n > 1) factor.emplace_back(n);
  for (LL res = 2; res <= p; res++) {</pre>
    bool ok = true;
    for (LL i = 0; i < factor.size() &&</pre>
         ok; i++)
```

```
ok &= Pow(res, phi / factor[i]) != struct AhoCorasick {
   if (ok) return res;
 return -1:
}
void prepare(LL n) {
 LL sz = abs(31 - \_builtin\_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
 w[0] = w[n] = 1;
 for (LL i = 1; i < n; i++) w[i] = (w[</pre>
      i - 1] * r) % MOD;
 for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i &
        1) << (sz - 1);
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i</pre>
 for (LL m = 2; m <= n; m <<= 1) {</pre>
   for (LL i = 0; i < n; i += m) {</pre>
     for (LL j = 0; j < (m >> 1); j++)
       LL &u = a[i + j], &v = a[i + j +
             (m >> 1)];
       LL t = v * w[dir ? n - n / m * j]
            : n / m * j] % MOD;
       v = u - t < 0 ? u - t + MOD : u
            - t;
       u = u + t >= MOD ? u + t - MOD :
            u + t;
     }
   }
 }
 if (dir)
   for (LL i = 0; i < n; i++) a[i] = (</pre>
        inv_n * a[i]) % MOD;
vector<LL> mul(vector<LL> p, vector<LL>
     q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;</pre>
 prepare(sz);
  for (LL i = 0; i < n; i++) a[i] = p[i
  for (LL i = 0; i < m; i++) b[i] = q[i</pre>
 for (LL i = n; i < sz; i++) a[i] = 0;
 for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
  for (LL i = 0; i < sz; i++) a[i] = (a
      [i] * b[i]) % MOD;
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
  while (c.size() && c.back() == 0) c.
      pop_back();
 return c;
```

### String

### **Aho Corasick**

```
int N, P;
 const int A = 256:
 vector<vector<int>> next;
 vector<int> link, out_link,
      end_in_pattern;
 vector<vector<int>> out;
 AhoCorasick() : N(0), P(0) { node();
     }
 int node() {
   next.emplace_back(A, 0);
   link.emplace_back(0);
   out_link.emplace_back(0);
   out.emplace_back(0);
   end_in_pattern.emplace_back(0);
   return N++;
 inline int get(char c) { return c; }
 int addPattern(const string T) {
   int u = 0;
   for (auto c : T) {
     if (!next[u][get(c)]) next[u][get(
          c)] = node();
     u = next[u][get(c)];
   out[u].push_back(P);
   end_in_pattern[u] = 1;
   return P++;
 void pushLinks() {
   queue<int> q;
   for (q.push(0); !q.empty();) {
     int u = q.front();
     q.pop();
     for (int c = 0; c < A; ++c) {
       int v = next[u][c];
       if (!v) next[u][c] = next[link[u
           ]][c];
       else {
         link[v] = u ? next[link[u]][c]
         out_link[v] = out[link[v]].
             empty() ? out_link[link[v
             ]] : link[v];
         q.push(v);
       end_in_pattern[v] |=
           end_in_pattern[out_link[v]];
     }
   }
 }
 int advance(int u, char c) {
   while (u && !next[u][get(c)]) u =
       link[u];
   u = next[u][get(c)];
   return u;
 }
};
```

#### 7.2Double hash

```
* Some well known primes:
   1949313259, 1997293877, 2091573227,
     2117566807
* Some Primes:
   100000007, 1000000009, 1000000861,
     1000099999 ( < 2^30 )
   108888881, 1111211111, 1500000001,
     1481481481 ( < 2<sup>31</sup> )
```

```
2147483647 (2^31-1),
 */
PLL base(1949313259, 1997293877);
namespace Hashing {
    using LL = long long;
    using PLL = pair<LL,LL>;
    #define ff first
    #define ss second
    const PLL M = \{1e9+7, 1e9+9\};
                                       11
        /Should be large primes
    const LL base = 1259;
                                       //
        /Should be larger than alphabet
         size
    const int N = 1e6+7;
                                      //
        /Highest length of string
    PLL operator+ (const PLL& a, LL x)
        {return \{a.ff + x, a.ss + x\};\}
    PLL operator- (const PLL& a, LL x)
        {return {a.ff - x, a.ss - x};}
    PLL operator* (const PLL& a, LL x)
        {return {a.ff * x, a.ss * x};}
    PLL operator+ (const PLL& a, PLL x)
        {return {a.ff + x.ff, a.ss + x.
        ss};}
    PLL operator- (const PLL& a, PLL x)
        {return {a.ff - x.ff, a.ss - x.
    PLL operator* (const PLL& a, PLL x)
        {return {a.ff * x.ff, a.ss * x.
        ss};}
    PLL operator% (const PLL& a, PLL m)
        {return {a.ff % m.ff, a.ss % m.
        ss};}
    ostream& operator<<(ostream& os, PLL
         hash) {
       return os<<"("<<hash.ff<<", "<<
           hash.ss<<")";
    PLL pb[N];
                  ///powers of base mod
    ///Call pre before everything
    void hashPre() {
       pb[0] = \{1,1\};
       for (int i=1; i<N; i++) pb[i] =</pre>
             (pb[i-1] * base)%M;
    ///Calculates hashes of all prefixes
         of s including empty prefix
    vector<PLL> hashList(string s) {
       int n = s.size();
       vector<PLL> ans(n+1);
       ans[0] = \{0,0\};
       for (int i=1; i<=n; i++) ans[i]</pre>
            = (ans[i-1] * base + s[i-1])
            %M:
       return ans;
    ///Calculates hash of substring s[l
        ..r] (1 indexed)
    PLL substringHash(const vector<PLL>
        &hashlist, int 1, int r) {
       return (hashlist[r]+(M-hashlist[
            l-1])*pb[r-l+1])%M;
    ///Calculates Hash of a string
    PLL Hash (string s) {
       PLL ans = \{0,0\};
       for (int i=0; i<s.size(); i++)</pre>
```

ans=(ans\*base + s[i])%M;

```
return ans;
}
///Tested on https://toph.co/p/
    palindromist
///appends c to string
PLL append(PLL cur, char c) {
   return (cur*base + c)%M;
///Tested on https://toph.co/p/
    palindromist
///prepends c to string with size k
PLL prepend(PLL cur, int k, char c)
   return (pb[k]*c + cur)%M;
}
///Tested on https://toph.co/p/
    chikongunia
///replaces the i-th (0-indexed)
    character from right from a to
PLL replace(PLL cur, int i, char a,
    char b) {
   return cur + pb[i] * (M+b-a)%M;
///Erases c from front of the string
     with size len
PLL pop_front(PLL hash, int len,
    char c) {
   return (hash + pb[len-1]*(M-c))%
}
///Tested on https://toph.co/p/
    palindromist
///concatenates two strings where
    length of the right is k
PLL concat(PLL left, PLL right, int
    k) {
   return (left*pb[k] + right)%M;
PLL power (const PLL& a, LL p) {
   if (p==0) return {1,1};
   PLL ans = power(a, p/2);
   ans = (ans * ans)%M;
   if (p\%2) ans = (ans*a)\%M;
   return ans;
PLL inverse(PLL a) {
   if (M.ss == 1) return power(a, M
   return power(a, (M.ff-1)*(M.ss
       -1)-1);
}
///Erases c from the back of the
    string
PLL invb = inverse({base, base});
PLL pop_back(PLL hash, char c) {
   return ((hash-c+M)*invb)%M;
///Tested on https://toph.co/p/
    palindromist
///Calculates hash of string with
    size len repeated cnt times
///This is O(\log n). For O(1), pre-
    calculate inverses
PLL repeat(PLL hash, int len, LL cnt
    ) {
   PLL mul = ((pb[len*cnt]-1+M) *
       inverse(pb[len]-1+M))%M;
   PLL ans = (hash*mul);
```

```
if (pb[len].ff == 1) ans.ff =
    hash.ff*cnt;
if (pb[len].ss == 1) ans.ss =
    hash.ss*cnt;
return ans%M;
}
struct pair_hash {
  inline std::size_t operator()(
    const std::pair<LL,LL> & v)
    const {
    return v.first*31+v.second;
  }
};
```

#### 7.3 KMP

```
vector<int> prefix_function(string s) {
  int n = (int)s.length();
  vector<int> pi(n);
  for (int i = 1; i < n; i++) {
    int j = pi[i-1];
    while (j > 0 && s[i] != s[j])
        j = pi[j-1];
    j += s[i] == s[j];
    pi[i] = j;
  }
  return pi;
```

# 7.4 Manacher's

```
vector<int> d1(n);
// d[i] = number of palindromes taking
    s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 1 : min(d1[1 + r -
      i], r - i + 1);
 while (0 <= i - k && i + k < n && s[i
       - k] == s[i + k]) k++;
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking
    s[i-1] and s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i
    ++) {
 int k = (i > r) ? 0 : min(d2[1 + r -
      i + 1], r - i + 1);
 while (0 \le i - k - 1 \&\& i + k \le n \&\&
       s[i - k - 1] == s[i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i +
```

#### 7.5 String Match FFT

```
//find occurrences of t in s where '?'s
    are automatically matched with any
    character
//res[i + m - 1] = sum_j=0 to m - 1_{s[
        i + j] * t[j] * (s[i + j] - t[j])
vector<int> string_matching(string &s,
        string &t) {
    int n = s.size(), m = t.size();
    vector<int> s1(n), s2(n), s3(n);
    for(int i = 0; i < n; i++) s1[i] = s[
        i] == '?' ? 0 : s[i] - 'a' + 1;
        //assign any non zero number for</pre>
```

```
non '?'s
for(int i = 0; i < n; i++) s2[i] = s1</pre>
    [i] * s1[i];
for(int i = 0; i < n; i++) s3[i] = s1</pre>
    [i] * s2[i];
vector<int> t1(m), t2(m), t3(m);
for(int i = 0; i < m; i++) t1[i] = t[</pre>
    i] == '?' ? 0 : t[i] - 'a' + 1;
for(int i = 0; i < m; i++) t2[i] = t1</pre>
    [i] * t1[i];
for(int i = 0; i < m; i++) t3[i] = t1</pre>
    [i] * t2[i];
reverse(t1.begin(), t1.end());
reverse(t2.begin(), t2.end());
reverse(t3.begin(), t3.end());
vector<int> s1t3 = multiply(s1, t3);
vector<int> s2t2 = multiply(s2, t2);
vector<int> s3t1 = multiply(s3, t1);
vector<int> res(n);
for(int i = 0; i < n; i++) res[i] =</pre>
    s1t3[i] - s2t2[i] * 2 + s3t1[i];
vector<int> oc;
for(int i = m - 1; i < n; i++) if(res</pre>
    [i] == 0) \text{ oc.push\_back}(i - m + 1)
return oc:
```

#### 7.6 Suffix Array

```
void inducedSort (const vector <int> &
    vec, int val_range, vector <int> &
    SA, const vector <int> &sl, const
    vector <int> &lms_idx) {
 vector <int> l(val_range, 0), r(
     val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c
        + 1];
 partial_sum(1.begin(), 1.end(), 1.
     begin());
 partial_sum(r.begin(), r.end(), r.
     begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i >=
      0; --i) SA[--r[vec[lms_idx[i]]]]
      = lms_idx[i];
 for (int i : SA) if (i > 0 and sl[i -
      1]) SA[l[vec[i-1]]++] = i-1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
 partial_sum(r.begin(), r.end(), r.
     begin());
 for (int k = SA.size() - 1, i = SA[k
     ]; k; --k, i = SA[k]) {
   if (i and !sl[i - 1]) SA[--r[vec[i -
        1]]] = i - 1;
```

vector <int> suffixArray (const vector

vector <int> sl(n), SA(n), lms\_idx;

for (int i = n - 2;  $i \ge 0$ ; --i) {

sl[i] = vec[i] > vec[i + 1] or (vec[i])

i] == vec[i + 1] and sl[i + 1])

<int> &vec, int val\_range) {

const int n = vec.size();

```
if (sl[i] and !sl[i + 1]) lms_idx.
        emplace_back(i + 1);
 reverse(lms_idx.begin(), lms_idx.end
      ()):
 inducedSort(vec, val_range, SA, sl,
      lms_idx);
 vector <int> new_lms_idx(lms_idx.size
      ()), lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {
   if (SA[i] > 0 and !sl[SA[i]] and sl[
        SA[i] - 1] new_lms_idx[k++] =
       SA[i];
 int cur = 0; SA[n - 1] = 0;
 for (int k = 1; k < new_lms_idx.size</pre>
      (); ++k) {
   int i = new_lms_idx[k - 1], j =
       new_lms_idx[k];
   if (vec[i] ^ vec[j]) {
     SA[j] = ++cur; continue;
   bool flag = 0;
   for (int a = i + 1, b = j + 1; ; ++a
        , ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
     if ((!sl[a] and sl[a - 1]) or (!sl
          [b] and sl[b - 1])) {
       flag = !(!sl[a] and sl[a - 1]
           and !sl[b] and sl[b - 1]);
           break;
     }
   SA[j] = flag ? ++cur : cur;
 }
 for (int i = 0; i < lms_idx.size();</pre>
      ++i) lms_vec[i] = SA[lms_idx[i]];
 if (cur + 1 < lms_idx.size()) {</pre>
   auto lms_SA = suffixArray(lms_vec,
        cur + 1);
   for (int i = 0; i < lms_idx.size();</pre>
        ++i) new_lms_idx[i] = lms_idx[
       lms_SA[i]];
 inducedSort(vec, val_range, SA, sl,
      new_lms_idx); return SA;
vector <int> getSuffixArray (const
    string &s, const int LIM = 128) {
 vector <int> vec(s.size() + 1);
 copy(begin(s), end(s), begin(vec));
      vec.back() = '$';
 auto ret = suffixArray(vec, LIM);
 ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any
    two suffix
vector <int> getLCParray (const string
    &s, const vector <int> &SA) {
 int n = s.size(), k = 0;
 vector <int> lcp(n), rank(n);
 for (int i = 0; i < n; ++i) rank[SA[i</pre>
      ]] = i;
 for (int i = 0; i < n; ++i, k ? --k :</pre>
       0) {
```

```
if (rank[i] == n - 1) {
     k = 0; continue;
   int j = SA[rank[i] + 1];
   while (i + k < n \text{ and } j + k < n \text{ and } s
        [i + k] == s[j + k]) ++k;
   lcp[rank[i]] = k;
 lcp[n - 1] = 0; return lcp;
int main() {
 string s; cin >> s;
 for (const int i : getSuffixArray(s))
       printf("%d ", i);
 puts("");
 return 0;
    Trie
template<int sz>
struct Trie {
```

#### 7.7

```
Trie() : id(1) {
   memset(endMark, 0, sizeof endMark);
   for_each(all(trie), [](vector<int> &
       v) { v.assign(sz, 0); });
 void insert(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       trie[cur][val] = id++;
     cur = trie[cur][val];
   endMark[cur] = true;
 bool search(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a';
     if (not trie[cur][val])
       return false;
     cur = trie[cur][val];
   return endMark[cur];
 }
private:
 int id, endMark[100005];
 vector<int> trie[100005];
```

#### 7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
     l = r = i;
     while (r < n \&\& s[r] == s[r - 1])
         r++;
     z[i] = r - 1, r--;
   } else {
     int k = i - 1;
     if (z[k] < r - i + 1) z[i] = z[k]; autocmd filetype cpp nnoremap <F10> :
```

```
l = i;
     while (r < n \&\& s[r] == s[r - 1]
          ]) r++;
     z[i] = r - 1, r--;
 }
}
return z;
```

#### 8 Extra

#### Stress Tester

```
# $2 is good code
# $3 is generator
g++ -02 -std=c++17 "$1".cpp -o $1
g++ -02 -std=c++17 "$2".cpp -o $2
g++ -02 -std=c++17 "$3".cpp -o $3
for ((i = 1;; i++)); do
  echo 'Test #'$i
  timeout 5s ./$3 $RANDOM > in
  timeout 5s./\$1 < in > out
  timeout 5s ./$2 < in > ans
  diff -yi ans out > diff.out
  if [ $? -ne 0 ]; then
   echo "\nInput:"
   cat in
   echo "\nDiff:"
   cat diff.out
   break
 fi
done
```

#### 8.2Sublime Build

```
"shell_cmd": "g++ -02 -std=c++17 -g -
    DLOCAL -Wall -Wextra -Wpedantic
    Wfloat-equal -Wshift-overflow=2 -
    fsanitize=address -fsanitize=
    undefined -fno-sanitize-recover
    $file_name -o $file_base_name &&
    timeout 5s ./$file_base_name < in</pre>
     > out",
"working_dir": "$file_path",
"selector": "source.cpp"
```

#### 8.3 vimrc

```
" Auto import & Compile
:autocmd BufNewFile *.cpp Or ~/template
    .cpp
nnoremap <F4> :!xclip -o -sel clip > ~/
   cp/in <CR><CR>
inoremap <F4> <ESC>:!xclip -o -sel clip
     > ~/cp/in <CR><CR>
nnoremap <F6> :!xclip -sel clip % <CR><
   CR>
inoremap <F6> <ESC>:!xclip -sel clip %
    <CR><CR>
autocmd filetype cpp nnoremap <F9>
    wa \| !g++ -02 -std=c++17 % -o %:r
    && timeout 5s ./%:r < ^{\sim}/cp/in> ^{\sim}/cp
    /out<CR>
autocmd filetype cpp inoremap <F9> <ESC
    >:wa \| !g++ -02 -std=c++17 % -o %:
    r && timeout 5s ./%:r < ~/cp/in> ~/
    cp/out<CR>
```

wa \| !make %:r D=1 && ./%:r < ~/cp

```
/in> ~/cp/out<CR>
autocmd filetype cpp inoremap <F10> <
    ESC>:wa \| !make clean && make %:r
    D=1 \&\& ./\%:r < ^/cp/in > ^/cp/out <
" Auto Completion
inoremap ( ()<left>
inoremap <expr> ) strpart(getline('.'),
     col('.')-1, 1) == ")" ? "\<Right>"
     : ")"
inoremap { {}<left>
inoremap <expr> } strpart(getline('.'), set laststatus=2
     col('.')-1, 1) == "}" ? "\<Right>"
     : "}"
inoremap [ []<left>
inoremap <expr> ] strpart(getline('.'),
     col('.')-1, 1) == "]" ? "\<Right>"
     : "]"
inoremap <expr> " strpart(getline('.'),
     col('.')-1, 1) == "\"" ? "\<Right>
    " : "\"\"\<left>"
inoremap <expr> ' strpart(getline('.'),
     col('.')-1, 1) == "\'" ? "\<Right
    >" : "\'\\'\<left>"
inoremap <expr> <CR> <sid>
    insert_newline()
function s:insert_newline() abort
 let pair = strpart(getline('.'), col(
      '.')-2, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<CR>\<
      ESC>\0" : "\<CR>"
endfunction
inoremap <expr> <space> <sid>
   insert_space()
function s:insert_space() abort
 let pair = strpart(getline('.'), col(
      '.')-2, 2)
 return stridx('(){}[]', pair) % 2 ==
      0 && strlen(pair) == 2 ? "\<space</pre>
      >\<space>\<left>" : "\<space>"
endfunction
inoremap <expr> <bs> <sid>rm_pair()
function s:rm_pair() abort
let pair = strpart(getline('.'), col('
     ., )-2, 2)
return stridx('(){}[]''', pair) %
     2 == 0 && strlen(pair) == 2 ? "\<
     del>\<bs>" : "\<bs>"
endfunction
set nocompatible
    iMproved, required
                           " required
filetype on
filetype plugin on
filetype plugin indent on
syntax on
set splitright splitbelow
set mouse=a
set number
set relativenumber
set tabstop=2
set shiftwidth=2
set expandtab
set softtabstop=2
set smartindent
set smarttab
set autoindent
set cindent
```

set noerrorbells

```
set ruler
 set guifont=*
 set backspace=indent,eol,start
 " set ignorecase
 set incsearch
 set nowrap
 set hlsearch
 " set termguicolors
set foldmethod=indent
 set nofoldenable
 " set cursorline
set showcmd
 set wildmenu
 " colorscheme torte
 if !has('nvim')
  set clipboard=unnamedplus
 if !has('nvim')
  set ttymouse=xterm2
 endif
 nnoremap \langle S-j \rangle : m .+1 \langle CR \rangle ==
 nnoremap \langle S-k \rangle : m .-2 \langle CR \rangle ==
 vnoremap <S-j> :m '>+1<CR>gv==gv
 vnoremap \langle S-k \rangle : m \langle -2 \langle CR \rangle gv == gv
 nnoremap <A-h> <C-w>h
 nnoremap <A-j> <C-w>j
nnoremap <A-k> <C-w>k
 nnoremap <A-1> <C-w>1
 let mapleader = ','
map <leader>cp :50 vsplit in<CR>:split
      out<CR><C-w>h
```

# **Equations and Formulas**

### 9.1 Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1$$
  
1 and  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$ 

the size n+1 factors.

intersecting chords.

n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

Number of permutations of  $1, \ldots, n$  that avoid the pattern 123 (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n=3, these permutations are 132, 213, 231, 312 and 321.

### 9.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count If  $P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$ , then, permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$
where,  $S(0,0) = 1, S(n,0) = S(0,n) = 0$ 

$$\sum_{k=0}^{n} S(n,k) = n!$$

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k) \text{ , then,}$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 \le i_2 \le i_k \le n} i_1 i_2 .... i_k.$$

#### 9.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$ where S(0,0) = 1, S(n,0) = S(0,n) = 0 $S(n,2) = 2^{n-1} - 1 S(n,k) \cdot k! = \text{number}$ of ways to color n nodes using colors from 1 to k such that each color is used at least

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets,

with each subset containing at least r elements. It is denoted by  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) = \sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$  $kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1) \qquad \sum_{k=1}^n x^{\gcd(k,n)} = \sum_{d|n} x^d \cdot \phi\left(\frac{n}{d}\right)$ Denote the n objects to partition by the noted  $S^d(n,k)$ , to be the number of ways  $\phi(d)$ partitions of polygon into disjoint triangles each subset have pairwise distance at least by using the diagonals). The number of ways to connect the 2n subset, it is required that  $|i-j| \ge d$ . It  $\frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$ points on a circle to form n disjoint i.e. non-has been shown that these numbers satisfy, The number of rooted full binary trees with 9.4 Other Combinatorial Identities

 $\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n}$  $n, r \in N, n > r, \sum_{i=1}^{n} {i \choose r} = {n+1 \choose r+1}$ 

If 
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If 
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

# Different Math Formulas

Picks Theorem: A = i + b/2 - 1 $d(i) = (i-1) \times$ Deragements: (d(i-1) + d(i-2))

$$\frac{n}{ab}$$
 -  $\left\{\frac{b'n}{a}\right\}$  -  $\left\{\frac{a'n}{b}\right\}$  + 1

#### 9.6 GCD and LCM

if m is any integer, then  $gcd(a + m \cdot b, b) =$ gcd(a,b)

The gcd is a multiplicative function in the following sense: if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$  $\gcd(a_1,b)\cdot\gcd(a_2,b).$ 

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$ lcm(a, gcd(b, c)) = gcd(lcm(a, b), lcm(a, c)).For non-negative integers a and b, where aand b are not both zero,  $gcd(n^a-1, n^b-1) =$  $n^{\gcd(a,b)} - 1$ 

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\sum_{i=1}^{n} [\gcd(i, n) = k] = \phi\left(\frac{n}{k}\right)$$

9.1 Catalan Numbers 
$$C_n = \frac{1}{n+1} \binom{2n}{n} \quad C_0 = 1, C_1 = \frac{1}{\ln n} \binom{2n}{n} \quad C_0 = \frac{1}{\ln n} \binom$$