

# EEE3096S: Embedded Systems II

LECTURE 15:  
ARM C PROGRAM STRUCTURE

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RECAP

## ARM: FUNCTIONS AND FRAMES

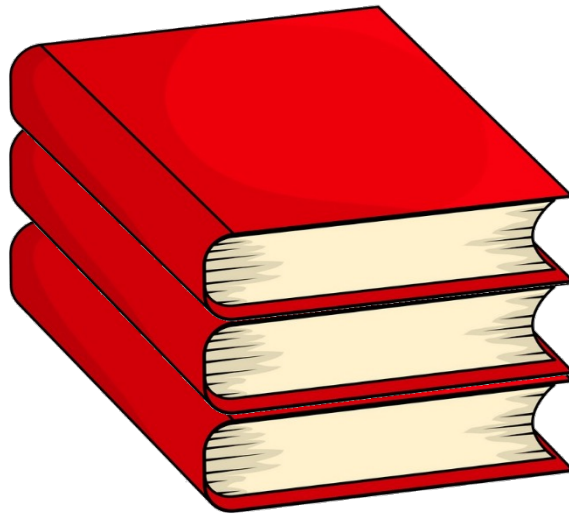
- In C, and most other executable formats, function calls are handled with **stack frames**
- These keep track of the call stack
- Special pieces of code are needed to handle these stack frames, specifically:
  - Before a function is called
  - At the start of a function (Prologue)
  - At the end of a function (Epilogue)

*For further detail:*

See text book (Patterson & Hennessy) Chapter 2, Section 2.8 " Supporting Procedures in Computer Hardware".

# Stack Data Type

A stack data type works in pretty much the way you would pile things, like books, one on top of the other to form a stack of items...



A stack data type works in the same way. It has two primary methods that works on it:

`push(x)`

adds the object or value `x` to the top of the stack

`x = pop()`

removes the object or value at the top of the stack and returns it

# Coding a Stack

You can utilize an array, and a *top* index, to construct a stack. i.e. *top* is the index of where the next item to be pushed to the stack will go. Pop uses *top* to indicate the item to pull off the stack and return.

The following code implements a simple stack:

```
// set the size of the stack
#define SIZE 100

int stack [SIZE];
int top = 0; // index for next item

void push ( int x ) {
    // place an item on top of stack
    if (top<SIZE) {
        stack[top] = x;
        top=top+1;
    }
}

int pop () {
    // remove an item from top of stack
    if (top>0) {
        top = top - 1;
        return stack[top];
    }
}

int empty () {
    // return true if the stack is empty
    return top==0;
}
```

*See next slide for full code*

# Coding a Stack – full program

```
/* The following code implements a simple integer stack in C: */
#include <stdio.h>

// set the size of the stack
#define SIZE 100

int stack [SIZE];
int top = 0; // first place to put the next item

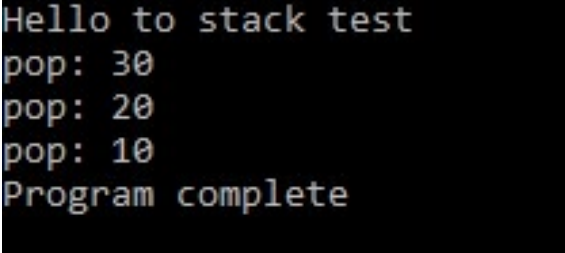
void push ( int x ) {
    // place an item on top of stack
    if (top<SIZE) {
        stack[top] = x;
        top=top+1;
    }
}

int pop () {
    // remove an item from top of stack
    if (top>0) {
        top = top - 1;
        return stack[top];
    }
}

int empty () {
    // return true if the stack is empty
    return top==0;
}

int main ( int argc, char** args )
{
    // main program to test the stack functions
    printf("Hello to stack test\n");
    push(10);
    push(20);
    push(30);
    while (!empty())
        printf("pop: %d\n", pop());
    printf("Program complete\n");
}
```

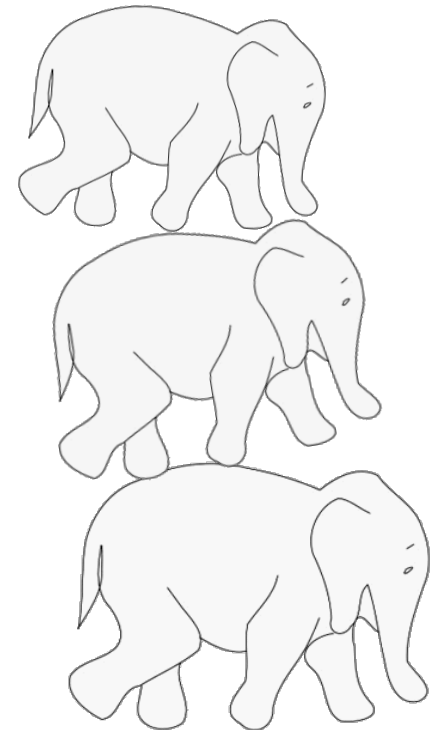
*Example run:*



```
Hello to stack test
pop: 30
pop: 20
pop: 10
Program complete
```

# WHAT IS A CALL STACK?

- When one function calls another, which calls another function, which calls another...
- The **program needs to keep track** so that it can return to the original function
- So a stack of stack frames is kept; this stack is called the **call stack**



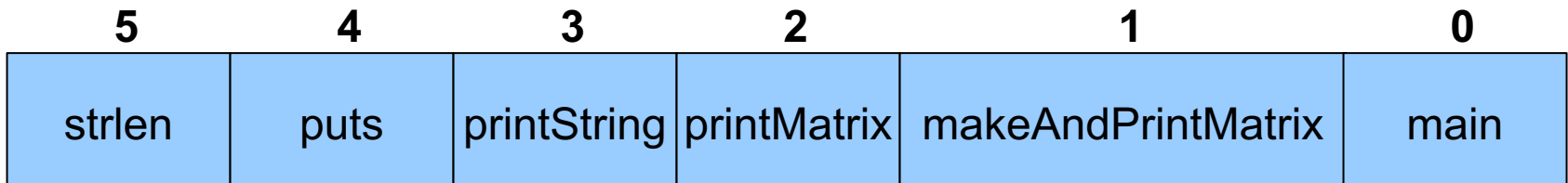
# AN EXAMPLE CALL STACK

```
#5  0xb7e92243 in strlen () from /lib/i686/cmov/libc.so.6
#4  0xb7e7cfd5 in puts () from /lib/i686/cmov/libc.so.6
#3  0x08048365 in printString (dat=0x0) at frame.c:7
#2  0x0804837e in printMatrix (pc=0xbfded4a0) at frame.c:13
#1  0x080483a2 in makeAndPrintMatrix (mdata=0xbfded4a0...) at frame.c:19
#0  0x080483db in main () at frame.c:27
```

(generated by *gdb* from a crashed program)

e.g.  
void main () {  
 makeAndPrintMatrix(0);  
}

A graphical representation of this call stack:



Top of  
call stack

Bottom of  
call stack



# C RUNTIME (CRT) ENVIRONMENT

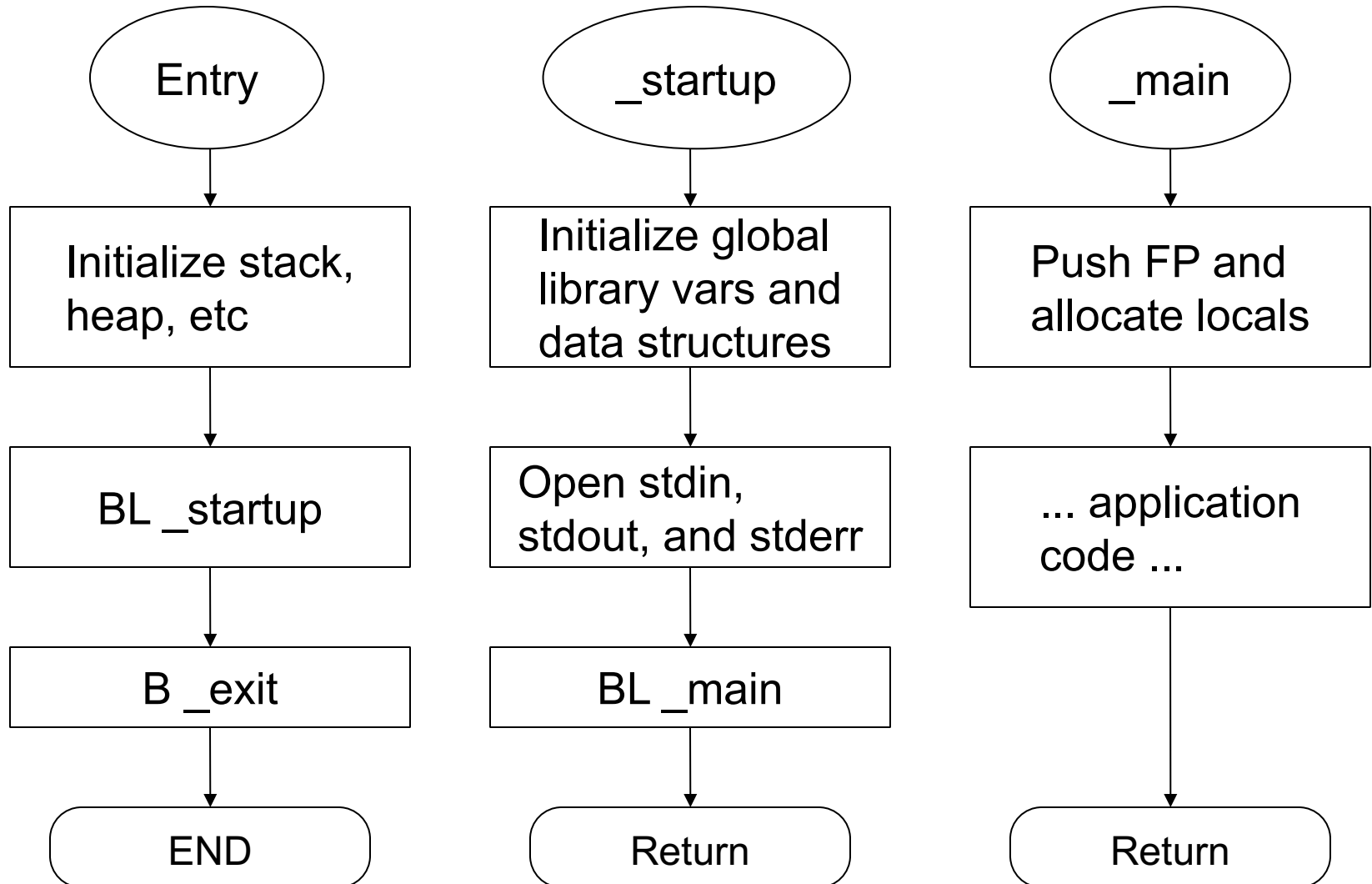
- As we now know, a **runtime environment\*** is the software structures that supports program execution
- For C programs, these are:
  - Start-up code
  - Runtime Libraries (**RTLs**) \*\*
- These RTLs provide the most basic support to programs (e.g., heap memory management, IO control)
- RTLs simplify the way programs interact with their “world” (i.e. their runtime environment)



## C START-UP (\_START OR \_STARTUP)

- The **start-up code** is the code than **transfers control to main()**. It is the code that is run even **before** the C program's main function starts.
- This code runs at the program **entry point**, which is a platform specific fixed address
- The most simple example would be  
**B \_main**
- In most cases it's more complex than this
- Often, this code is automatically inserted by the C compiler, in module **crt0**  
crt = C run time 0 → the first basic starting point

# HOW *CRT0* FITS IN

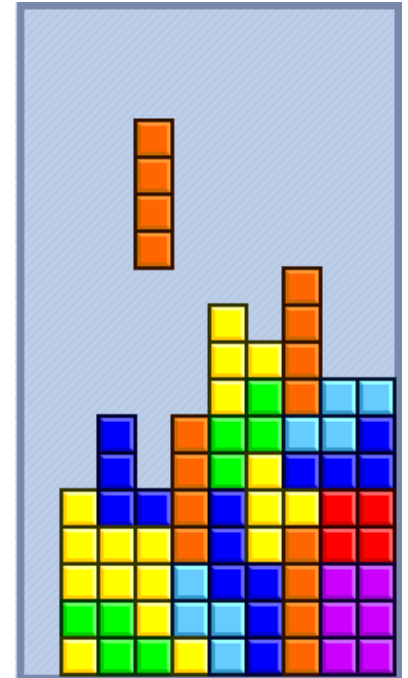


# THE C RUNTIME LIBRARY (RTL)

- The RTL provides support functions which simplify the writing of code
- For example:
  - Basic IO functions (like read and write)
  - Floating point arithmetic on platforms which don't have a floating-point unit (**FPU** — a math co-processor)
- RTLs are great, but often **too bulky** for embedded systems, so they are cut down or removed completely (e.g. *printf*, *sin*, *cos*, etc.)

# OBJECT PLACEMENT

- The C compiler generates **object code** that needs to be **positioned** in memory and integrated into a suitable **binary format**
- The **linker** is the tool for the job:
  - The linker fixes the location in memory
  - Can override linker with ASM macros (e.g. **ORG**)
- **Linker command file** controls placement of program sections
- Usually set up as a script file or add to **Makefile** for reusability



# CMake

Another option that takes Make further; as per the CMake website, it is defined as:

“CMake is an open-source, cross-platform family of tools designed to build, test and package software.”

<https://cmake.org/>

CMake can be considered an abbreviation for “cross-platform make”. The C *does not* stand for C or C++, as this tool can be used for many different languages and tools, besides C.

For some quick information, links to online tutorials and manual pages for cmake, see: <http://wiki.ee.uct.ac.za/Cmake>

# EXECUTION ENVIRONMENT MEMORY BUDGETING

- Embedded platforms have **limited** memory
- **Memory budget** = how much memory space do you need to run the program?
- Need to determine how many variables you can declare, max array sizes, max number of functions to call, etc.

# MEMORY BUDGETING

## Example:

- Using 32-byte call stack, 64-byte general stack
- 6 items/functions on call stack and each function is 4 bytes (assuming 32-bit addresses)
  - Therefore memory required is: 4 bytes x 6 = **24 bytes**
  - Thus you only have **8 bytes** of the call stack left (or space for **two functions**) before a stack overflow\* occurs

\*no, not the website

- **Stack overflow:** overflow error that occurs when a program tries to use more memory space in the call stack than has been allocated to that stack



# MEMORY BUDGETING

- Another reason for doing a memory budget is for situations where you have a small amount of RAM but a larger amount of **ROM (read-only memory)** that permanently stores instructions)
- Example:
  - 32 KB SRAM (Static RAM), directly coupled to processor
  - 512 KB external FLASH memory (ROM), with half taken up by functions (8 KB each); needs to be read one block at a time and you can't execute code directly from FLASH
  - Questions to consider:
    - How many of the functions can be in SRAM at once (so that they can be executed)?
    - How deep can the call stack become?
    - Which functions should always be in SRAM?
    - When we load one function, should we load other functions it might call?

## How to participate?



Click on the projected screen to start the question



1

Go to [woclap.com](https://www.woclap.com)

2

Enter the event code in the top banner

Event code

**ECFSRX**



If executing a Branch and Link (BL) instruction, what would be stored into the Link Register?



① SP

② PC



Click on the projected screen to start the question

③ PC + 4

④ PC + 32



The exception vector is...



1 a piece of code that is executed when an exception occurs

2 a location in memory where the exception handler is stored

3 an exception that has both magnitude and direction

4 an unexpected event that occurs from within the processor

Click on the projected screen to start the question



An Application Binary Interface (ABI) is an interface between two binary program modules.



1 True



Click on the projected screen to start the question



2 False

Serial Peripheral Interface (SPI) is said to be a serial, half-duplex, synchronous communication protocol. True or false?



1

True

0%

0



2

False

0%

0

