

Homework 5 Solution

May 20, 2007

1. Given the following information, try to prove that Garfield eats fish.

- (a) Cats like fish.
- (b) Cats eat everything they like.
- (c) Garfield is a cat.

Ans. Define the following predicates:

$isCat(x)$: x is a cat

$like(x, y)$: x like y

$eat(x, y)$: x eat y

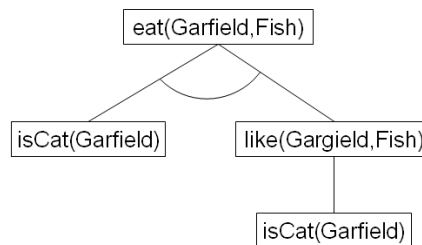
Then, by the facts above, we can define the following sentences:

$isCat(x) \Rightarrow like(x, Fish)$

$isCat(x) \wedge like(x, y) \Rightarrow eat(x, y)$

$isCat(Garfield)$

We can prove that $eat(Garfield, Fish)$ is true by backward-chaining as following:



2. Consider the following axioms, try to establish the conclusion using the axioms by applying refutation resolution.

- (a) Anyone who rides any Harley is a rough character.
- (b) Every biker rides [something that is] either a Harley or a BMW.
- (c) Anyone who rides any BMW is a yuppie.
- (d) Every yuppie is a lawyer.
- (e) Any nice girl does not date anyone who is a rough character.

- (f) Mary is a nice girl, and John is a biker.
 (g) (Conclusion) If John is not a lawyer, then Mary does not date John.

Ans. Define predicates such as:

$rides(x, y)$: x rides y
 $isBiker(x, y)$: x is a biker
 $isRoughChar(x)$: x is a rough character
 $isYuppie(x)$: x is a yuppie
 $isLawyer(x)$: x is a lawyer
 $isNiceGirl(x)$: x is a nice girl
 $isBMW(x)$: x is BMW
 $isHarley(x)$: x is Harley
 $date(x, y)$: x dates y

Then these axioms can be transferred to FOL sentences:

- (a) $\forall x \forall y rides(x, y) \wedge isHarley(y) \Rightarrow isRoughChar(x)$
 (b) $\forall x isBiker(x) \Rightarrow \exists y rides(x, y) \wedge (isBMW(y) \vee isHarley(y))$
 (c) $\forall x \forall y rides(x, y) \wedge isBMW(y) \Rightarrow isYuppie(x)$
 (d) $\forall x isYuppie(x) \Rightarrow isLawyer(x)$
 (e) $\forall x \forall y isNiceGirl(x) \Rightarrow \neg(date(x, y) \wedge isRoughChar(y))$
 (f) $isNiceGirl(Mary) \wedge isBiker(John)$
 (g) $\neg isLawyer(John) \Rightarrow \neg date(Mary, John)$

Where item (g) is what we want to prove.

Convert these sentences to CNF as:

- $$\begin{aligned} \neg rides(x, y) \vee \neg isHarley(y) \vee isRoughChar(x) & \quad (1) \\ \neg isBiker(x) \vee rides(x, F(x)) & \quad (2) \\ \neg isBiker(x) \vee isHarley(F(x)) \vee isBMW(F(x)) & \quad (3) \\ \neg rides(x, y) \vee \neg isBMW(y) \vee isYuppie(x) & \quad (4) \\ \neg isYuppie(x) \vee isLawyer(x) & \quad (5) \\ \neg isNiceGirl(x) \vee \neg date(x, y) \vee \neg isRoughChar(y) & \quad (6) \\ isNiceGirl(Mary) & \quad (7) \\ isBiker(John) & \quad (8) \\ \neg isLawyer(John) & \quad (9) \\ date(Mary, John) & \quad (10) \end{aligned}$$

The last two sentences come from the goal sentence being negated because we are going to proof by refutation using resolution.

Then, by resolution:

	Resolvent clause	Unification	
(2),(8)	$rides(John, F(John))$	$\{x/John\}$	(11)
(3),(8)	$isHarley(F(John)) \vee isBMW(F(John))$	$\{x/John\}$	(12)
(1),(11)	$\neg isHarley(F(John)) \vee isRoughChar(John)$	$\{x/John, y/F(John)\}$	(13)
(4),(11)	$\neg isBMW(F(John)) \vee isYuppie(John)$	$\{x/John, y/F(John)\}$	(14)
(6),(7)	$\neg date(Mary, y) \vee \neg isRoughChar(y)$	$\{x/Mary\}$	(15)
(10),(15)	$\neg isRoughChar(John)$	$\{y/John\}$	(16)
(5),(9)	$\neg isYuppie(John)$	$\{x/John\}$	(17)
(14),(17)	$\neg isBMW(F(John))$		(18)
(13),(16)	$\neg isHarley(F(John))$		(19)
(12),(18)	$isHarley(F(John))$		(20)
(19),(20)	$\{\}$		(21)

Thus if John is not a lawyer, Mary does not date John.