CodeBook for Competitive Programming

$MIST_EaglesExpr$

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1 Useful Tips

Big Integer C++ _int128_t

C++ FastIO

ios::sync_with_stdio(false); cin.tie(nullptr);

Python FastIO

import sys; input = sys.stdin.readline

Integer - Binary Conversion in C++

bitset < size > (val).to_string();
stoi(str, nulptr, base);

Input From File

freopen("input.txt", "r", stdin);

Python Array Input

list(map(int, input().split()))

Vim Setup

"install xclip, vim-gtk3 set nocp ai bs=2 hls ic is lbr ls=2 mouse=a nu ru sc scs smd so=3 sw=4 ts=4 rnufiletype plugin indent on syn on $\mathrm{map}\ \mathrm{gA}\ \mathrm{m'ggVG"}{+}\mathrm{y'}\ '$ inoremap {<CR> {<CR>}<Esc>ko nnoremap = mzgg=G'zvnoremap $\langle C-k \rangle : s / \langle s * \langle zs / \rangle /$ /<CR>:nohl<CR> nnoremap $\langle C-k \rangle : s / \ s * \ zs / \ / \ /$ /<CR>:nohl<CR> vnoremap $\langle C-l \rangle : s / \hat{s} \times zs / /$ //<CR>:nohl<CR> nnoremap $\langle C-l \rangle : s/^\s * \zs ///$ //<CR>:nohl<CR> autocmd FileType cpp map <F9>: w < CR > :! clear ; g++ % -DONPC-o %:r && ./%:r<CR> "autocmd FileType cpp map <F9> :w<CR> :!clear; g++ % -o %< && gnome-terminal — ./%<<CR

2 Formula

2.1 Area Formula

Rectangle Area = length * width

Square Area = Side * Side

Triangle $Area = \frac{1}{2} * length * width$

Circle $Area = \pi * radius^2$

Parallelogram Area = base * height

Pyramid Base $Area = \frac{1}{2} * base * slantHeight$ Polygon

a $Area = \frac{1}{2} |\sum_{n=1}^{n-1} (x_i y_{i+1})|$

b (Pick's formula) $Area = a + \frac{b}{2} - 1$ (for int coordinates). Here a = int points inside polygon and b = int points outside polygon.

2.2 Perimeter Formulas

Rectangle Perimeter = 2 * (length + width)

Square Perimeter = 4 * side

Triangle Perimeter = 4 * side

Circle $Perimeter = 2 * \pi * radius$

2.3 Volume Formula

Cube $Volume = side^3$

Rect Prism Volume = length * width * height

Cylinder $Volume = \pi * radius^2 * height$

Sphere $Volume = \frac{4}{3} * \pi * radius^3$

Pyramid $Volume = \frac{1}{3} * baseArea * height$

2.4 Surface Area Formula

Cube $SurfaceArea = 6 * side^2$

Rectangle Prism SurfaceArea = 2 * (length * width + length * height + width * height)

Cylinder $SurfaceArea = 2 * \pi * radius * (radius + height)$

Sphere $SurfaceArea = 4 * \pi * radius^2$

Pyramid $SurfaceArea = basearea + \frac{1}{2} * perimeterOfBase * slantHeight$

2.5 Triangles

Side Lengths a, b, c

Semi Perimeter $p = \frac{a+b+c}{2}$

Area $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumstance $R = \frac{abc}{4A}$

In Radius $r = \frac{A}{n}$

2.6 Summation Of Series

- $c^k + c^{k+1} + \dots + c^n = c^{n+1} c^k$
- $1+2+3+\ldots+n=\frac{n(n+1)}{2}$
- $1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$
- $1^3 + 2^3 + 3^3 + \dots + n^3 = (\frac{n(n+1)}{2})^2$

2.7 Miscellaneous

- $2^{100} = 2^{50} * 2^{50}$
- $\bullet \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}$
- logn! = log1 + log2 + ... + logn
- gcd(a,b) = gcd(a-b,b)
- Number of occurrence of a prime number p in n! is $\lfloor \frac{n}{p} \rfloor + \lfloor \frac{n}{p^2} \rfloor + \lfloor \frac{n}{p^3} \rfloor + \dots + 0$
- Formula of Catalan number is $\binom{2n}{n} \binom{2n}{n-1}$
- Number of divisors of $p^x q^y$ where p and q are prime is (x+1)*(y+1)
- Sum of divisors of $p^x q^y$ where p and q are prime is $(1+p+p^2+...+p^x)(1+q+q^2+...+q^y)$
- Divisibility rules for a number
 - **3** The sum of its digits is divisible by 3.
 - **4** Last two digits form a number that is divisible by 4.
 - **5** The number ends in 0 or 5.
 - 6 Divisible by both 2 and 3.
 - 7 A rule for 7 is to double the last digit, subtract it from the rest of the number, and check if the result is divisible by 7. Repeat if necessary.
 - 8 Last three digits form a number divisible by 8.
 - **9** Sum of its digits is divisible by 9
 - 11 Difference between the sum of digits of odd places and even places is divisible by 11.
- Derangement Number n person put their hat in a boy. How many way they can collect hat such that no one get their own hat $D_n = (n-1)D_{n-2} + (n-1)D_{n-1}$
- Golden Ratio $\Phi = (\frac{1+\sqrt{5}}{2}) \approx 1.618034$
- nth Fibonacci number $F_n = \frac{\Phi n (1 \Phi)n}{\sqrt{5}}$
- Stars and Bars Number of solutions of $x_1+x_2+x_3+...+x_k = n$ is $\binom{n-1}{k-1}$ where $x_i > 0$ and $\binom{n+k-1}{k-1}$ where $x_i > 0$

3 Graph Theory

All about graph.

3.1 BFS

```
void bfs(int start, int target = -1) {
   queue < int > q;
   q.push(start);
   vis[start] = true;
   while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int i : adj[u]) {
        if (!vis[i]) {
            vis[i] = true;
            q.push(i);
        }
    }
   }
}
```

3.2 DFS

```
map < int , vector < int >> adj;
map < int , int > visited , parent , level , color
   ;

void dfs(int start)
{
    visited[start]=1;
    for (auto child : adj[start])
    {
        if (!visited[child])
        {
            dfs(child);
        }
    }
    visited[start]=2;
}
```

3.3 Dijkstra Algorithm

```
void Dijkstra(int start) {
  // vector < pair < int , int >> adj[N];
  priority_queue < pair < int , int > ,
     vector<pair<int, int>>, greater<
     pair<int, int>>> pq;
  pq.push({0, start});
  while (!pq.empty()) {
    auto it = pq.top();
    pq.pop();
    int wt = it.first;
    int u = it.second;
    if (vis[u])
      continue;
    vis[u] = 1;
    for (pair<int, int> i : adj[u]) {
      int adjWt = i.second;
      int adjNode = i.first;
      if (dist[adjNode] > wt + adjWt)
        dist[adjNode] = wt + adjWt;
```

3.4 Bellman Ford

```
vector < int > dist;
vector < int > parent;
vector < vector < pair < int , int >> > adj;
// resize the vectors from main
    function
void bellmanFord(int num_of_nd, int
    src) {
  dist[src] = 0;
  for (int step = 0; step < num_of_nd;</pre>
       step++) {
    for (int i = 1; i <= num_of_nd; i</pre>
        ++) {
      for (auto it : adj[i]) {
         int u = i;
         int v = it.first;
         int wt = it.second;
         if (dist[u] != inf &&
             ((dist[u] + wt) < dist[v])
                 ) {
           if (step == num_of_nd - 1) {
             cout << "Negative cycle</pre>
                 found\n ";
                 return;
           dist[v] = dist[u] + wt;
           parent[v] = u;
        }
      }
    }
  }
  for (int i = 1;i <= num_of_nd; i++)</pre>
    cout << dist[i] << " ";
  cout << endl;</pre>
}
```

3.5 Floyed Warshall Algorithm

```
typedef double T;
typedef vector<T> VT;
typedef vector < VT > VVT;
typedef vector<int> VI;
typedef vector < VI > VVI;
bool FloydWarshall(VVT &w, VVI &prev)
   {
 int n = w.size();
 prev = VVI(n, VI(n, -1));
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {</pre>
      for (int j = 0; j < n; j++) {
        if (w[i][j] > w[i][k] + w[k][j
          w[i][j] = w[i][k] + w[k][j];
          prev[i][j] = k;
        }
      }
```

```
}
for (int i = 0; i < n; i++)
  if (w[i][i] < 0)
    return false;
return true;
}</pre>
```

3.6 Kruskal Algorithm (MST)

```
vector<pair<int, pair<int, int>>>
   Krushkal(vector<pair<int, pair<int</pre>
   , int>>> &edges, int n) {
  sort(edges.begin(), edges.end());
  vector<pair<int, pair<int, int>>>
     ans;
  DisjointSet D(n);
  for (auto it : edges) {
    if (D.findUPar(it.second.first) !=
        D.findUPar(it.second.second))
      ans.push_back({it.first, {it.
          second.first, it.second.
          second}});
      D.unionBySize(it.second.first,
          it.second.second);
 }
 return ans;
```

3.7 Prims Algorithm (MST)

```
void Prims(int start) {
  // map < int , vector < pair < int , int >>>
      adj, ans;
  priority_queue < pair < int , pair < int ,</pre>
      int>>, vector<pair<int, pair<int</pre>
      , int>>>, greater<pair<int, pair
      <int, int>>>> pq;
  pq.push({0, {start, -1}});
  while (!pq.empty()) {
    auto it = pq.top();
    pq.pop();
    int wt = it.first;
    int u = it.second.first;
    int v = it.second.second;
    if (vis[u]) continue;
    vis[u] = 1;
    if (v != -1) ans[u].push_back({v,
        wt});
    for (pair<int, int> i : adj[u]) {
      int adjWt = i.second;
      int adjNode = i.first;
      if (!vis[adjNode]) pq.push({
          adjWt, {adjNode, u}});
}
```

3.8 Strongly Connected Components

```
\verb|vector<bool>| visited; // keeps track||
                                                   int root = *min_element(begin(
   of which vertices are already
   visited
// runs depth first search starting at
// each visited vertex is appended to
   the output vector when dfs leaves
   it.
void dfs(int v, vector<vector<int>>
   const &adj, vector<int> &output) {
  visited[v] = true;
  for (auto u : adj[v])
    if (!visited[u])
                                            }
      dfs(u, adj, output);
                                            3.9 LCA
  output.push_back(v);
// input: adj -- adjacency list of G
// output: components -- the strongy
    {\it connected components in } {\it G}
// \  \, output: \  \, adj\_cond \  \, -- \  \, adjacency \  \, list
   of G^SCC (by root vertices)
                                              int n;
void scc(vector<vector<int>> const &
   adj, vector<vector<int>> &
   components, vector<vector<int>> &
   adj_cond) {
  int n = adj.size();
  components.clear(), adj_cond.clear()
  vector<int> order; // will be a
     sorted list of G's vertices by
      exit time
  visited.assign(n, false);
  // first series of depth first
     searches
  for (int i = 0; i < n; i++)</pre>
    if (!visited[i])
      dfs(i, adj, order);
  // create adjacency list of G^T
  vector < vector < int >> adj_rev(n);
  for (int v = 0; v < n; v++)
    for (int u : adj[v])
      adj_rev[u].push_back(v);
  visited.assign(n, false);
                                                     }
  reverse(order.begin(), order.end());
                                                }
  vector<int> roots(n, 0); // gives
      the root vertex of a vertex's
     SCC
  // second series of depth first
     searches
  for (auto v : order)
    if (!visited[v]) {
      std::vector<int> component;
      dfs(v, adj_rev, component);
      components.push_back(component);
```

```
component), end(component));
      for (auto u : component)
        roots[u] = root;
  // add edges to condensation graph
  adj_cond.assign(n, {});
  for (int v = 0; v < n; v++)
    for (auto u : adj[v])
      if (roots[v] != roots[u])
        adj_cond[roots[v]].push_back(
            roots[u]);
struct LCA {
  vector<int> height, euler, first,
     segtree, parent;
  vector < bool > visited;
  vector < int >> jump;
 LCA(vector < vector < int >> &adj, int
     root = 0) {
    n = adj.size();
    height.resize(n);
    first.resize(n);
    parent.resize(n);
    euler.reserve(n * 2);
    visited.assign(n, false);
    dfs(adj, root);
    int m = euler.size();
    segtree.resize(m * 4);
    build(1, 0, m - 1);
    jump.resize(n, vector<int>(32, -1)
    for(int i=0;i<n;i++) {</pre>
        jump[i][0] = parent[i];
    for(int j=1; j<20; j++) {</pre>
        for(int i=0;i<n;i++) {</pre>
            int mid = jump[i][j-1];
if(mid != -1) jump[i][j] =
                 jump[mid][j-1];
  void dfs(vector<vector<int>> &adj,
     int node, int h = 0) {
    visited[node] = true;
    height[node] = h;
    first[node] = euler.size();
    euler.push_back(node);
    for (auto to : adj[node]) {
     if (!visited[to]) {
        parent[to] = node;
        dfs(adj, to, h + 1);
```

```
euler.push_back(node);
      }
    }
  }
  void build(int node, int b, int e) {
    if (b == e) {
      segtree[node] = euler[b];
    } else {
      int mid = (b + e) / 2;
      build(node << 1, b, mid);</pre>
      build(node << 1 | 1, mid + 1, e)
      int l = segtree[node << 1], r =</pre>
          segtree[node << 1 | 1];
      segtree[node] = (height[1] <
          height[r]) ? 1 : r;
   }
  }
  int query(int node, int b, int e,
      int L, int R) {
    if (b > R || e < L)</pre>
      return -1;
    if (b >= L && e <= R)
      return segtree[node];
    int mid = (b + e) >> 1;
    int left = query(node << 1, b, mid</pre>
        , L, R);
    int right = query(node << 1 | 1,</pre>
       mid + 1, e, L, R);
    if (left == -1)
      return right;
    if (right == -1)
      return left;
    return height[left] < height[right</pre>
        ] ? left : right;
  int lca(int u, int v) {
    int left = first[u], right = first
        [v];
    if (left > right)
      swap(left, right);
    return query(1, 0, euler.size() -
        1, left, right);
  int kthParent(int u, int k) {
      for(int i=0;i<19;i++) {</pre>
          if(k & (1LL<<i)) u = jump[u</pre>
              ][i];
      }
      return u;
  }
};
3.10
      Max Flow
const int N = 505;
int capacity[N][N];
int vis[N], p[N];
```

int n, m;

```
int bfs(int s, int t) {
 memset(vis, 0, sizeof vis);
  queue < int > qu;
  qu.push(s);
  vis[s] = 1;
  while (!qu.empty()) {
    int u = qu.front();
    qu.pop();
    for (int i = 0; i \le n + m + 2; i
        ++) {
      if (capacity[u][i] > 0 && !vis[i
         ]) {
        p[i] = u;
        vis[i] = 1;
        qu.push(i);
    }
 }
  return vis[t] == 1;
int maxflow(int s, int t) {
  int cnt = 0;
  while (bfs(s, t)) {
    int cur = t;
    while (cur != s) {
      int prev = p[cur];
      capacity[prev][cur] -= 1;
      capacity[cur][prev] += 1;
      cur = prev;
    }
    cnt++;
 }
  return cnt;
}
```

4 Data Structures

Different Data Structures.

4.1 Segment Tree

```
constexpr int N = 100005;
int arr[N], seg[N];
void build(int ind, int low, int high)
  if (low == high) {
    seg[ind] = arr[low];
    return;
 int mid = (low + high) / 2;
 build(2 * ind + 1, low, mid);
  build(2 * ind + 2, mid + 1, high);
  seg[ind] = seg[2 * ind + 1] + seg[2
     * ind + 2];
int query(int ind, int low, int high,
   int 1, int r) {
  if (low \geq= 1 && high \leq= r) return
     seg[ind];
  if (low > r || high < 1) return 0;</pre>
```

```
int mid = (low + high) / 2;
                                              push(node, 1, r);
  int left = query(2 * ind + 1, low,
     mid, 1, r);
  int right = query(2 * ind + 2, mid +
                                              int mid = (1 + r) >> 1;
                                              update(node << 1, 1, mid, q1, qr,
      1, high, l, r);
  return left + right;
                                                  val);
}
                                              update(node << 1 | 1, mid + 1, r, ql
void update(int ind, int low, int high
                                                  , qr, val);
    , int node, int val) {
  if (low == high) {
                                              pull(node, 1, r);
    seg[ind] = val;
                                            }
    return;
 }
                                            int query(int node, int 1, int r, int
  int mid = (low + high) / 2;
                                                ql, int qr) {
  if (low <= node && node <= mid)</pre>
                                              if(qr < 1 || r < ql) return 0;</pre>
     update(2 * ind + 1, low, mid,
                                              if(ql <= l && r <= qr) return seg[</pre>
     node, val);
                                                  node];
  else update(2 * ind + 2, mid + 1,
     high, node, val);
                                              push(node, 1, r);
  seg[ind] = seg[2 * ind + 1] + seg[2
     * ind + 2];
                                              int mid = (1 + r) >> 1;
}
                                              return query(node << 1, 1, mid, q1,</pre>
                                                  qr) + query(node << 1 | 1, mid +</pre>
4.2
     Segment Tree Lazy
                                                   1, r, ql, qr);
const int N = 1e5 + 5;
                                                 Fenwick Tree
int arr[N], seg[N << 2], lz[N << 2];</pre>
                                            4.3
void pull(int node, int 1, int r) {
                                            int fenwick[N];
  seg[node] = seg[node << 1] + seg[
     node << 1 | 1];
                                            void update(int ind, int val) {
                                              while (ind < N) {
                                                fenwick[ind] += val;
void push(int node, int 1, int r) {
                                                 ind += ind & -ind;
  int mid = (1 + r) >> 1;
                                              }
  lz[node << 1] += lz[node];</pre>
                                            }
  seg[node << 1] += lz[node] * (mid -
                                            int query(int ind) {
     1 + 1);
                                              int sum = 0;
  lz[node << 1 | 1] += lz[node];</pre>
                                              while (ind > 0) {
  seg[node << 1 | 1] += lz[node] * (r
                                                sum += fenwick[ind];
     - mid);
                                                ind -= ind & -ind;
  lz[node] = 0;
                                              }
                                              return sum;
                                            }
void build(int node, int 1, int r) {
                                            4.4 Disjoint Set
  if(1 == r) {
    seg[node] = arr[1];
    lz[node] = 0;
                                            class DisjointSet {
                                              vector < int > parent, sz;
    return;
  }
  int mid = (1 + r) >> 1;
                                             public:
  build(node << 1, 1, mid);</pre>
                                              DisjointSet(int n) {
  build(node << 1 | 1, mid + 1, r);
                                                sz.resize(n + 1);
  pull(node, 1, r);
                                                parent.resize(n + 2);
                                                for (int i = 1; i <= n; i++)</pre>
                                                    parent[i] = i, sz[i] = 1;
void update(int node, int 1, int r,
   int ql, int qr, int val) {
                                              int findUPar(int u) { return parent[
                                                  u] == u ? u : parent[u] =
  if(qr < 1 || r < q1) return;</pre>
                                                  findUPar(parent[u]); }
  if(ql <= l && r <= qr) {</pre>
                                              void unionBySize(int u, int v) {
    seg[node] += val;
    lz[node] += val;
                                                int a = findUPar(u);
                                                 int b = findUPar(v);
    return;
  }
                                                if (sz[a] < sz[b]) swap(a, b);
```

```
if (a != b) {
                                             }
      parent[b] = a;
      sz[a] += sz[b];
  }
};
     TRIE
4.5
const int N = 26;
class Node {
 public:
  int EoW;
  Node* child[N];
                                               }
  Node() {
    EoW = 0;
                                             }
    for (int i = 0; i < N; i++) child[</pre>
       i] = NULL;
  }
};
void insert(Node* node, string s) {
  for (size_t i = 0; i < s.size(); i</pre>
     ++) {
    int r = s[i] - 'A';
    if (node->child[r] == NULL) node->
       child[r] = new Node();
    node = node->child[r];
  }
  node \rightarrow EoW += 1;
                                             }
}
                                           }
int search(Node* node, string s) {
  for (size_t i = 0; i < s.size(); i</pre>
     ++) {
    int r = s[i] - A';
    if (node->child[r] == NULL) return
                                              hpp>
  }
  return node -> EoW;
void print(Node* node, string s = "")
  if (node->EoW) cout << s << "\n";</pre>
  for (int i = 0; i < N; i++) {</pre>
    if (node->child[i] != NULL) {
      char c = i + 'A';
      print(node->child[i], s + c);
    }
  }
}
bool isChild(Node* node) {
  for (int i = 0; i < N; i++)</pre>
    if (node->child[i] != NULL) return
        true;
  return false;
bool isJunc(Node* node) {
  int cnt = 0;
  for (int i = 0; i < N; i++) {</pre>
```

```
if (cnt > 1) return true;
  return false;
int trie_delete(Node* node, string s,
    int k = 0) {
  if (node == NULL) return 0;
  if (k == (int)s.size()) {
    if (node->EoW == 0) return 0;
    if (isChild(node)) {
      node -> EoW = 0;
      return 0;
    return 1;
  int r = s[k] - A';
  int d = trie_delete(node->child[r],
     s, k + 1);
  int j = isJunc(node);
  if (d) delete node->child[r];
  if (j) return 0;
  return d;
void delete_trie(Node* node) {
  for (int i = 0; i < 15; i++) {</pre>
    if (node->child[i] != NULL)
        delete_trie(node->child[i]);
  delete node;
4.6 Ordered Set
#include <ext/pb_ds/assoc_container.</pre>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
#define ordered_set tree<int,</pre>
    null_type,less<int>, rb_tree_tag,
    tree_order_statistics_node_update>
ordered_set o_set;
// insert function to insert in
// ordered set same as SET STL
o_set.insert(5);
o_set.insert(1);
o_set.insert(2);
// Finding the second smallest element
// in the set using * because
// find_by_order returns an iterator
*(o_set.find_by_order(1))
// Finding the number of elements
// strictly less than k=4
o_set.order_of_key(4)
4.7 Mo's Algorithm
void remove(int idx);
```

```
int get_answer();
int block_size = 700;
struct Query {
  int 1, r, idx;
  bool operator < (Query other) const {</pre>
    if (1 / block_size != other.1 /
        block_size) return 1 /
        block_size < other.1 /
        block_size;
    return r < other.r;</pre>
  }
};
vector<int> mo_s_algorithm(vector<
   Query > queries) {
  vector < int > answers(queries.size());
  sort(queries.begin(), queries.end())
  // TODO: initialize data structure
  int cur_l = 0;
  int cur_r = -1;
  // invariant: data structure will
      always reflect the range [cur_l,
       cur\_r]
  for (Query q : queries) {
    while (cur_1 > q.1) {
      cur_1--;
      add(cur_1);
    while (cur_r < q.r) {</pre>
      cur_r++;
      add(cur_r);
    while (cur_l < q.1) {</pre>
      remove(cur_1);
      cur_1++;
    }
    while (cur_r > q.r) {
      remove(cur_r);
      cur_r --;
    answers[q.idx] = get_answer();
  return answers;
    Algorithms
5
```

All about algorithms.

Monotonic Stack (Immediate 5.1Large)

```
for (int i = n - 1; i >= 0; i--) {
  while (!stk.empty() && v[i] >= v[stk
     .top()]) stk.pop();
  ind[i] = stk.empty() ? -1 : stk.top
     ();
```

```
stk.push(i);
}
// 3 1 5 4 10
// 2 2 4 4 -1
```

5.2 Kadane's Algorithm

```
int maxSubArraySum(vector<int> &a) {
 int size = a.size();
  int maxTill = INT_MIN, maxEnd = 0;
  for (int i = 0; i < size; i++) {</pre>
    maxEnd = maxEnd + a[i];
    if (maxTill < maxEnd) maxTill =</pre>
        maxEnd;
    if (maxEnd < 0) maxEnd = 0;</pre>
 }
  return maxTill;
```

2D Prefix Sum

```
pref[i][j] = a[i][j] + pref[i - 1][j]
   + pref[i][j - 1] - pref[i - 1][j -
    1];
Sum of region = pref[row2 + 1][col2 +
   1] - pref[row2 + 1][col1] - pref[
   row1][col2 + 1] + pref[row1][col1
```

5.4 2D Convex Hull

```
struct pt {
    double x, y;
    bool operator == (pt const& t)
       const {
        return x == t.x && y == t.y;
};
int orientation(pt a, pt b, pt c) {
    double v = a.x*(b.y-c.y)+b.x*(c.y-
       a.y)+c.x*(a.y-b.y);
    if (v < 0) return -1; // clockwise
    if (v > 0) return +1; // counter-
       clockwise
    return 0;
}
bool cw(pt a, pt b, pt c, bool
   include_collinear) {
    int o = orientation(a, b, c);
    return o < 0 || (include_collinear</pre>
        && o == 0);
}
bool collinear(pt a, pt b, pt c) {
   return orientation(a, b, c) == 0;
void convex_hull(vector<pt>& a, bool
   include_collinear = false) {
    pt p0 = *min_element(a.begin(), a.
        end(), [](pt a, pt b) {
        return make_pair(a.y, a.x) <</pre>
            make_pair(b.y, b.x);
```

}

}

return ((x * x) % m * base) % m;

```
});
                                             }
    sort(a.begin(), a.end(), [&p0](
                                             else {
        const pt& a, const pt& b) {
                                               int x = power(base, n / 2);
        int o = orientation(p0, a, b);
                                               return (x * x) % m;
        if (o == 0)
            return (p0.x-a.x)*(p0.x-a.
                                           }
                x) + (p0.y-a.y)*(p0.y-
                                           6.3
                                                 Miller Rabin
                a.y)
                < (p0.x-b.x)*(p0.x-b.x)
                                           using u64 = uint64_t;
                    ) + (p0.y-b.y)*(p0
                                           using u128 = __uint128_t;
                    .y-b.y);
        return o < 0;</pre>
                                           u64 binpower(u64 base, u64 e, u64 mod)
    });
    if (include_collinear) {
        int i = (int)a.size()-1;
                                             u64 result = 1;
                                             base %= mod;
        while (i >= 0 && collinear(p0,
                                             while (e) {
            a[i], a.back())) i--;
                                               if (e & 1) result = (u128)result *
        reverse(a.begin()+i+1, a.end()
                                                    base % mod;
           );
    }
                                               base = (u128)base * base % mod;
                                               e >>= 1;
                                             }
    vector <pt> st;
    for (int i = 0; i < (int)a.size();</pre>
                                             return result;
        i++) {
        while (st.size() > 1 && !cw(st
                                           bool check_composite(u64 n, u64 a, u64
            [st.size()-2], st.back(),
                                               d, int s) {
            a[i], include_collinear))
                                             u64 x = binpower(a, d, n);
            st.pop_back();
                                             if (x == 1 || x == n - 1) return
        st.push_back(a[i]);
                                                false;
                                             for (int r = 1; r < s; r++) {</pre>
                                               x = (u128)x * x % n;
    if (include_collinear == false &&
                                               if (x == n - 1) return false;
       st.size() == 2 && st[0] == st
        [1])
        st.pop_back();
                                             return true;
                                           };
    a = st;
                                           bool MillerRabin(u64 n, int iter = 5)
                                              { // returns true if n is
    Number Theory
                                               probably prime, else returns false
                                             if (n < 4) return n == 2 || n == 3;
All about math.
                                             int s = 0;
                                             u64 d = n - 1;
6.1 nCr
                                             while ((d & 1) == 0) {
                                               d >>= 1;
int inverseMod(int a, int m) { return
                                               s++;
   power(a, m - 2); }
int nCr(int n, int r, int m = mod){
                                             for (int i = 0; i < iter; i++) {</pre>
   if(r==0) return 1;
                                               int a = 2 + rand() % (n - 3);
   if(r>n) return 0;
                                               if (check_composite(n, a, d, s))
   return (fact[n] * inverseMod((fact[
                                                   return false;
      r] * fact[n-r]) % m , m)) % m;
                                             }
                                             return true;
6.2
    Power
                                                 Sieve
                                           6.4
int power(int base, int n, int m = mod
   ) {
                                           const int N = 1e7 + 3;
  if (n == 0) return 1;
                                           vector<int> primes;
  if (n & 1) {
                                           int notprime[N];
    int x = power(base, n / 2);
```

void sieve() {

```
primes.push_back(2);
  for (int i = 4; i < N; i += 2) {</pre>
    notprime[i] = true;
  for (int i = 3; i < N; i += 2) {</pre>
    if (!notprime[i]) {
      primes.push_back(i);
      for (int j = i * i; j < N; j +=
          2 * i) {
        notprime[j] = true;
    }
 }
}
     Bitset Sieve
```

6.5

```
const int sieve_size = 10000006;
bitset < sieve_size > sieve;
void Sieve() {
  sieve.flip();
  int finalBit = sqrt(sieve.size()) +
      1;
  for (int i = 2; i < finalBit; ++i) {</pre>
    if (sieve.test(i))
      for (int j = 2 * i; j <
          sieve\_size; j += i) sieve.
          reset(j);
}
```

6.6Segmented Sieve

Need Sieve() from 6.4

```
void segmentedSieve(int L, int R) {
  bool isPrime[R - L + 1];
  for (int i = 0; i <= R - L + 1; i++)
       isPrime[i] = true;
  if (L == 1) isPrime[0] = false;
  for (int i = 0; primes[i] * primes[i
     ] <= R; i++) {
    int curPrime = primes[i];
    int base = curPrime * curPrime;
    if (base < L) {</pre>
      base = ((L + curPrime - 1) /
          curPrime) * curPrime;
    for (int j = base; j \le R; j +=
       curPrime) isPrime[j - L] =
       false;
  for (int i = 0; i <= R - L; i++) {</pre>
    if (isPrime[i] == true) cout << L</pre>
       + i << '\n';
  cout << '\n';
```

6.7 Divisors

```
constexpr int N = 1000005;
```

```
int Prime[N + 4], kk;
bool notPrime[N + 5];
void SieveOf() {
  notPrime[1] = true;
  Prime[kk++] = 2;
  for (int i = 4; i <= N; i += 2)</pre>
     notPrime[i] = true;
  for (int i = 3; i <= N; i += 2) {</pre>
    if (!notPrime[i]) {
      Prime[kk++] = i;
      for (int j = i * i; j <= N; j +=</pre>
           2 * i) notPrime[j] = true;
  }
}
void Divisors(int n) {
  int sum = 1, total = 1;
  int mnP = INT_MAX, mxP = INT_MIN,
      cntP = 0, totalP = 0;
  for (int i = 0; i <= N && Prime[i] *</pre>
      Prime[i] <= n; i++) {
    if (n % Prime[i] == 0) {
      mnP = min(mnP, Prime[i]);
      mxP = max(mnP, Prime[i]);
      int k = 0;
      cntP++;
      while (n % Prime[i] == 0) {
        k++;
        n /= Prime[i];
      sum *= (k + 1); // NOD
      totalP += k;
      int s = 0, p = 1;
      while (k-->=0) {
        s += p;
        p *= Prime[i];
      };
      total *= s; // SOD
    }
 }
  if (n > 1) {
    cntP++, totalP++;
    sum *= 2;
    total *= (1 + n);
    mnP = min(mnP, n);
    mxP = max(mnP, n);
 }
  cout << mnP << " " << mxP << " " <<
     cntP << " " << totalP << " " <<
     sum << " " << total << "\n";
}
6.8
     Euler's Totient Phi Function
```

```
const int N = 5000005;
int phi[N];
unsigned long long phiSum[N];
void phiCalc() {
  for (int i = 2; i < N; i++) phi[i] =</pre>
  for (int i = 2; i < N; i++) {</pre>
```

(N + 30);

```
vector < int > primes = {2, 3, 5}; int
    if (phi[i] == i) {
      for (int j = i; j < N; j += i) {</pre>
                                                  psize = 3;
        phi[j] -= phi[j] / i;
                                               primes.resize(rsize);
    }
                                               vector <P> sprimes; size_t pbeg = 0;
  }
                                               int prod = 1;
  for (int i = 2; i < N; i++) {</pre>
                                               for (int p = 7; p \le v; ++p) {
    phiSum[i] = (unsigned long long)
                                                 if (!isp[p]) continue;
        phi[i] * (unsigned long long)
                                                 if (p <= Q) prod *= p, ++pbeg,</pre>
                                                    primes[psize++] = p;
        phi[i] + phiSum[i - 1];
                                                 auto pp = P(p);
  }
}
                                                 for (int t = 0; t < 8; ++t) {</pre>
                                                   int j = (p \le Q) ? p : p * p;
     Log a base b
6.9
                                                   while (j % 30 != rs[t]) j += p
                                                       << 1;
int logab (int a, int b){
                                                  pp.pos[t] = j / 30;
 return log2(a) / log2(b);
                                                 sprimes.push_back(pp);
                                               }
6.10
       Count 1's from 0 to n
                                               vector < unsigned char > pre(prod, 0xFF
int cntOnes(int n) {
                                                  );
  int cnt = 0;
                                               for (size_t pi = 0; pi < pbeg; ++pi)</pre>
  for(int i=1;i<=n;i<<=1) {</pre>
                                                   {
    int x = (n + 1) / (i << 1);</pre>
                                                 auto pp = sprimes[pi]; const int p
    cnt += x * i;
                                                      = pp.p;
    if((n + 1) % i && n & i) cnt += (n
                                                 for (int t = 0; t < 8; ++t) {</pre>
         + 1) % i;
                                                   const unsigned char m = ~(1 << t</pre>
                                                      );
  return cnt;
                                                   for (int i = pp.pos[t]; i < prod</pre>
                                                       ; i += p) pre[i] &= m;
6.11 Primes Upto 1e9
                                               }
// credit: min_25
                                               const int block_size = (L + prod -
// takes 0.5s for n = 1e9
                                                  1) / prod * prod;
vector < int > sieve(const int N, const
                                               vector < unsigned char > block(
   int Q = 17, const int L = 1 << 15)
                                                  block_size); unsigned char*
                                                  pblock = block.data();
  static const int rs[] = \{1, 7, 11,
                                               const int M = (N + 29) / 30;
     13, 17, 19, 23, 29};
  struct P {
                                               for (int beg = 0; beg < M; beg +=</pre>
    P(int p) : p(p) {}
                                                  block_size, pblock -= block_size
    int p; int pos[8];
                                                  ) {
                                                 int end = min(M, beg + block_size)
  auto approx_prime_count = [] (const
     int N) -> int {
                                                 for (int i = beg; i < end; i +=</pre>
    return N > 60184 ? N / (log(N) -
                                                    prod) {
        1.1)
                                                   copy(pre.begin(), pre.end(),
                      : max(1., N / (
                                                       pblock + i);
                          log(N) -
                          1.11)) + 1;
                                                 if (beg == 0) pblock[0] &= 0xFE;
  };
                                                 for (size_t pi = pbeg; pi <</pre>
                                                     sprimes.size(); ++pi) {
  const int v = sqrt(N), vv = sqrt(v);
                                                   auto& pp = sprimes[pi];
  vector < bool > isp(v + 1, true);
                                                   const int p = pp.p;
  for (int i = 2; i <= vv; ++i) if (
                                                   for (int t = 0; t < 8; ++t) {</pre>
      isp[i]) {
                                                     int i = pp.pos[t]; const
    for (int j = i * i; j \le v; j += i
                                                         unsigned char m = (1 \ll t)
        ) isp[j] = false;
                                                         );
                                                     for (; i < end; i += p) pblock</pre>
                                                         [i] &= m;
  const int rsize = approx_prime_count
                                                     pp.pos[t] = i;
```

```
}
    }
    for (int i = beg; i < end; ++i) {</pre>
      for (int m = pblock[i]; m > 0; m
           \&= m - 1) {
        primes[psize++] = i * 30 + rs[
            __builtin_ctz(m)];
      }
    }
  }
  assert(psize <= rsize);</pre>
  while (psize > 0 && primes[psize -
      1] > N) --psize;
  primes.resize(psize);
  return primes;
6.12
       Matrix Multiplication
```

```
const int mod = 1000000007;
vector < vector < i64>> multi(vector <
   vector < i64>> &m1, vector < vector <</pre>
   i64>> &m2) {
  int n = m1.size();
  vector < vector < i64>> res(n, vector <
     i64>(n, 0));
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++) {
      for (int k = 0; k < n; k++) {
        res[i][j] = (res[i][j] + (m1[i
            ][k] * m2[k][j]) % mod) %
            mod;
      }
    }
 }
  return res;
vector < vector < i64>> power (vector <</pre>
   vector < i64>> &base, int n, int m =
    mod) {
  int r = base.size();
  vector < vector < i64>> identity (r,
      vector < i64 > (r));
  for (int i = 0; i < r; i++) {</pre>
    identity[i][i] = 1;
  while (n > 0) {
    if (n % 2 == 1) {
      identity = multi(base, identity)
    base = multi(base, base);
    n /= 2;
  }
  return identity;
      Extended GCD
```

```
int exgcd(int a, int b, int &x, int &y
   ) {
   if (b == 0) {
```

```
x = 1;
        y = 0;
        return a;
    int g = exgcd(b, a \% b, y, x);
    y -= a / b * x;
    return g;
}
```

7 Dynamic Programming

7.1 LCS (Longest Common Subsequence)

```
string s, t;
vector<vector<int>> dp(3003, vector<
   int > (3003, -1));
vector < vector < int >> mark (3003, vector <
   int > (3003));
int f(int i, int j) {
  if (i < 0 || j < 0) return 0;</pre>
  if (dp[i][j] != -1) return dp[i][j];
  int res = 0;
  if (s[i] == t[j]) {
    mark[i][j] = 1;
    res = 1 + f(i - 1, j - 1);
  }
  else {
    int iC = f(i - 1, j);
    int jC = f(i, j - 1);
    if (iC > jC) mark[i][j] = 2;
    else mark[i][j] = 3;
    res = max(iC, jC);
 }
  return dp[i][j] = res;
}
void printWay(int i, int j) {
  if (i < 0 || j < 0) return;</pre>
  if (mark[i][j] == 1) printWay(i - 1,
      j - 1), cout << s[i];
  else if (mark[i][j] == 2) printWay(i
      - 1, j);
  else if (mark[i][j] == 3) printWay(i
      , j - 1);
}
```

7.2 LIS (Longest Increasing Subsequence)

```
void lis(vector<int> &v) {
  int n = v.size();
  vector < int > dp(n + 1, 1), hash(n);
  int mx = 1, lastInd = 0;
  for (int i = 0; i < n; i++) {</pre>
    hash[i] = i;
    for (int prev = 0; prev < i; prev</pre>
        ++) {
      if (v[i] > v[prev] && 1 + dp[
          prev] > dp[i]) {
        dp[i] = 1 + dp[prev];
        hash[i] = prev;
```

```
}
    }
    if (mx < dp[i]) {</pre>
      mx = dp[i];
      lastInd = i;
  }
  vector<int> printSeq;
  printSeq.push_back(v[lastInd]);
  while (hash[lastInd] != lastInd) {
    lastInd = hash[lastInd];
    printSeq.push_back(v[lastInd]);
  reverse(printSeq.begin(), printSeq.
     end());
  cout << mx << "\n";
  for (int i : printSeq) cout << i <<</pre>
     " ";
  cout << "\n";
}
```

7.3 SOS (Sum Of Subsets)

```
void SOS (vector<int> &v) {
  const int BITS = log2(*max_element(v
      .begin(), v.end())) + 1;
  vector<int> freq(1 << BITS, 0);</pre>
  for (int mask : v) freq[mask]++;
  vector < vector < int >> dp(BITS + 1,
      vector < int > (1 << BITS, 0));</pre>
  for (int mask = 0; mask < 1 << BITS;</pre>
       mask++) {
    dp[0][mask] = freq[mask];
  for (int bits = 1; bits <= BITS;</pre>
      bits++) {
    for (int mask = 0; mask < 1 <<</pre>
        BITS; mask++) {
      if ((mask & (1 << (bits - 1)))</pre>
          == 0) {
        dp[bits][mask] = dp[bits - 1][
            mask1:
      else {
        int other_mask = mask - (1 <<</pre>
            (bits - 1));
        dp[bits][mask] = dp[bits - 1][
            mask] + dp[bits - 1][
            other_mask];
      }
    }
  for (int mask : v) cout << dp[BITS][</pre>
      mask] << '\n';
// dp[bits][mask] means left most '
    bits' of submasks can be differ
// for dp[1][11] 01, 00 are now allow
    because leftmost 1 bit can be
    differ. 10 and 11 are allowed.
```

```
// for travarsing all the submask of a
    mask we can use
// submask = mask
// do {
// submask = (submask - 1) & mask
// } while (submask)
```

7.4 MCM (Matrix Chain Multiplication)

```
ans = min(ans, MCM(i, k) + MCM(k + 1,
j) + v[i - 1] * v[k] * v[j];
```

8 Strings

8.1 KMP

```
vector<int> prefix_function(string s)
    int n = (int)s.length();
    vector < int > pi(n);
    for (int i = 1; i < n; i++) {</pre>
        int j = pi[i - 1];
        while (j > 0 \&\& s[i] != s[j])
            j = pi[j - 1];
        if (s[i] == s[j]) j++;
        pi[i] = j;
    }
    return pi;
}
vector<int> find_matches(string text,
   string pat) {
    int n = pat.length(), m = text.
        length();
    string s = pat + "$" + text;
    vector<int> pi = prefix_function(s
       ), ans;
    for (int i = n; i <= n + m; i++) {</pre>
        if (pi[i] == n) {
            ans.push_back(i - 2 * n);
    }
    return ans;
}
```

8.2 Double Hashing

```
Need power() from 6.2
const int N = 1e6 + 9;
const int MOD1 = 127657753, MOD2 =
    987654319;
const int p1 = 137, p2 = 277;
int ip1, ip2;
pair<int, int> pw[N], ipw[N];
void prec() {
  pw[0] = {1, 1};
  for (int i = 1; i < N; i++) {
    pw[i].first = 1LL * pw[i - 1].
        first * p1 % MOD1;
    pw[i].second = 1LL * pw[i - 1].
        second * p2 % MOD2;
}</pre>
```

```
ip1 = power(p1, MOD1 - 2, MOD1);
  ip2 = power(p2, MOD2 - 2, MOD2);
  ipw[0] = \{1, 1\};
  for (int i = 1; i < N; i++) {</pre>
    ipw[i].first = 1LL * ipw[i - 1].
        first * ip1 % MOD1;
    ipw[i].second = 1LL * ipw[i - 1].
        second * ip2 % MOD2;
  }
}
struct Hashing {
 int n;
  string s; // 0 - indexed
  vector<pair<int, int>> hs; // 1 -
      indexed
  Hashing() {}
  Hashing(string _s) {
    n = _s.size();
    s = _s;
    hs.emplace_back(0, 0);
    for (int i = 0; i < n; i++) {</pre>
      pair < int , int > p;
      p.first = (hs[i].first + 1LL *
          pw[i].first * s[i] % MOD1) %
      p.second = (hs[i].second + 1LL *
           pw[i].second * s[i] % MOD2)
           % MOD2;
      hs.push_back(p);
    }
  }
  pair < int , int > get_hash(int 1, int r
     ) { // 1 - indexed
    assert(1 <= 1 && 1 <= r && r <= n)
    pair<int, int> ans;
    ans.first = (hs[r].first - hs[l -
        1].first + MOD1) * 1LL * ipw[1
         - 1].first % MOD1;
    ans.second = (hs[r].second - hs[l])
       - 1].second + MOD2) * 1LL *
       ipw[l - 1].second % MOD2;
    return ans;
  pair < int , int > get_hash() {
    return get_hash(1, n);
};
```

8.3 Manacher's Algorithm

```
p[i]++;
        }
        if(i + p[i] > r) {
            1 = i - p[i], r = i + p[i]
                ];
    }
    return vector<int>(begin(p) + 1,
        end(p) - 1);
}
vector<int> manacher(string s) {
    string t;
    for(auto c: s) t += string("#") +
       с;
    auto res = manacher_odd(t + "#");
    return vector<int>(begin(res) + 1,
         end(res) - 1);
}
// p[2 * i + (!odd)] - 1
// if r - l + 1 \le p[(r + l) / 2 + (r
   % 2 != l % 2)] is true then
   palindrome. Same parity means odd
   length.
```

9 Stress Testing

9.1 Bash Stress File

```
#!/usr/bin/env bash
g++ -g "$1".cpp -DONPC -o "$1"
g++ -g "$2".cpp -DONPC -o "$2"
g++ -g "$3".cpp -DONPC -o "$3"
for ((testNum=0;testNum<$4;testNum++))</pre>
do
    ./$3 > stdinput
    ./$2 < stdinput > outSlow
    ./$1 < stdinput > outWrong
    H1='md5sum outWrong'
    H2='md5sum outSlow'
    if !(cmp -s "outWrong" "outSlow")
        echo "Error found!"
        echo "Input:"
        cat stdinput
        echo "Wrong Output:"
        cat outWrong
        echo "Slow Output:"
        cat outSlow
        exit
done
echo Passed $4 tests
# ./stress.sh wrong correct gen times
9.2 C++ Generator File
```

#include <bits/stdc++.h>

using namespace std;

```
}
#define i64 long long
#define accuracy chrono::steady_clock
                                             int main() {
    ::now().time_since_epoch().count()
                                               srand(accuracy);
                                               int t = 1;
mt19937 rng(accuracy);
                                               t = rand(1, 10), cout << t << '\n';
                                               while (t--) {
                                                 // TODO
int rand(int 1, int r) {
  \verb"uniform_int_distribution<" int> | \verb"ludo" (1)|
                                             }
     , r);
  return ludo(rng);
```