CodeBook for Competitive Programming

$MIST_EaglesExpr$

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1 Useful Tips

Big Integer C++ __int128_t

C++ FastIO

ios::sync_with_stdio(false); cin.tie(nullptr);

Python FastIO

import sys; input = sys.stdin.readline

Integer - Binary Conversion in C++

bitset <size >(val).to_string();
(int) bitset <size >(val).to_ulong
();

Input From File

freopen("input.txt", "r", stdin);

Python Array Input

list(map(int, input().split()))

Vim Setup

"install xclip, vim-gtk3 set nocp ai bs=2 hls ic is lbr ls=2 mouse=a nu ru sc scs smd so=3 sw=4 ts=4 rnufiletype plugin indent on syn on map gA m'ggVG"+y'' inoremap ${< CR> {< CR>} < Esc>ko}$ nnoremap = mzgg=G'zvnoremap $\langle C-k \rangle : s / \langle s | zs / \rangle /$ /<CR>:nohl<CR> nnoremap <C $-k> : s/^ \ s * \ z s / \ / /$ /<CR>:nohl<CR> vnoremap <C-l $> : s/^\s*\zs\/\/$ //<CR>:nohl<CR> nnoremap $\langle C-l \rangle : s / \hat{s} \times zs / / /$ //<CR>:nohl<CR> autocmd FileType cpp map <F9>: w<CR> :! clear; g++ % -DONPC-o %:r && ./%:r<CR> "autocmd FileType cpp map <F9> :w<CR> :!clear; g++ % -o %&& gnome-terminal — ./%<<CR

2 Formula

2.1 Area Formula

 $\textbf{Rectangle} \ \textit{Area} = length*width$

Square Area = Side * Side

Triangle $Area = \frac{1}{2} * length * width$

Circle $Area = \pi * radius^2$

Parallelogram Area = base * height

Pyramid Base $Area = \frac{1}{2} * base * slantHeight$

Polygon

a $Area = \frac{1}{2} |\sum_{n=1}^{n-1} (x_i y_{i+1})|$

b (Pick's formula) $Area = a + \frac{b}{2} - 1$ (for int coordinates). Here a = int points inside polygon and b = int points outside polygon.

2.2 Perimeter Formulas

Rectangle Perimeter = 2 * (length + width)

Square Perimeter = 4 * side

Triangle Perimeter = 4 * side

Circle $Perimeter = 2 * \pi * radius$

2.3 Volume Formula

Cube $Volume = side^3$

Rect Prism Volume = length * width * height

Cylinder $Volume = \pi * radius^2 * height$

Sphere $Volume = \frac{4}{3} * \pi * radius^3$

Pyramid $Volume = \frac{1}{3} * baseArea * height$

2.4 Surface Area Formula

Cube $SurfaceArea = 6 * side^2$

Rectangle Prism SurfaceArea = 2 * (length * width + length * height + width * height)

Cylinder $SurfaceArea = 2 * \pi * radius * (radius + height)$

Sphere $SurfaceArea = 4 * \pi * radius^2$

Pyramid $SurfaceArea = basearea + \frac{1}{2} * perimeterOfBase*slantHeight$

2.5 Triangles

Side Lengths a, b, c

Semi Perimeter $p = \frac{a+b+c}{2}$

Area $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumstance $R = \frac{abc}{4A}$

In Radius $r = \frac{A}{n}$

2.6 Summation Of Series

- $c^k + c^{k+1} + \dots + c^n = c^{n+1} c^k$
- $1+2+3+\ldots+n=\frac{n(n+1)}{2}$
- $1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$
- $1^3 + 2^3 + 3^3 + \dots + n^3 = (\frac{n(n+1)}{2})^2$

2.7 Miscellaneous

- $2^{100} = 2^{50} * 2^{50}$
- $\bullet \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}$
- $log n! = log 1 + log 2 + \dots + log n$
- gcd(a,b) = gcd(a-b,b)
- Number of occurrence of a prime number p in n! is $\lfloor \frac{n}{p} \rfloor + \lfloor \frac{n}{p^2} \rfloor + \lfloor \frac{n}{p^3} \rfloor + \dots + 0$
- Formula of Catalan number is $\binom{2n}{n} \binom{2n}{n-1}$
- Number of divisors of $p^x q^y$ where p and q are prime is (x+1)*(y+1)
- Sum of divisors of $p^x q^y$ where p and q are prime is $(1+p+p^2+...+p^x)(1+q+q^2+...+q^y)$
- Divisibility rules for a number
 - **3** The sum of its digits is divisible by 3.
 - **4** Last two digits form a number that is divisible by 4.
 - **5** The number ends in 0 or 5.
 - 6 Divisible by both 2 and 3.
 - **7** A rule for 7 is to double the last digit, subtract it from the rest of the number, and check if the result is divisible by 7. Repeat if necessary.
 - 8 Last three digits form a number divisible by 8.
 - **9** Sum of its digits is divisible by 9
 - 11 Difference between the sum of digits of odd places and even places is divisible by 11.
- Derangement Number n person put their hat in a boy. How many way they can collect hat such that no one get their own hat $D_n = (n-1)D_{n-2} + (n-1)D_{n-1}$
- Golden Ratio $\Phi = (\frac{1+\sqrt{5}}{2}) \approx 1.618034$

- nth Fibonacci number $F_n = \frac{\Phi n (1 \Phi)n}{\sqrt{5}}$
- Stars and Bars Number of solutions of $x_1+x_2+x_3+...+x_k = n$ is $\binom{n-1}{k-1}$ where $x_i > 0$ and $\binom{n+k-1}{k-1}$ where $x_i > 0$

3 Graph Theory

All about graph.

3.1 BFS

```
void bfs(int start, int target = -1) {
   queue < int > q;
   q.push(start);
   vis[start] = true;
   while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int i : adj[u]) {
        if (!vis[i]) {
            vis[i] = true;
            q.push(i);
        }
    }
   }
}
```

3.2 DFS

```
map < int , vector < int >> adj;
map < int , int > visited , parent , level , color
   ;

void dfs(int start)
{
    visited[start]=1;
    for (auto child : adj[start])
    {
        if (!visited[child])
        {
            dfs(child);
        }
    }
    visited[start]=2;
}
```

3.3 Dijkstra Algorithm

```
void Dijkstra(int start) {
   // vector < pair < int, int >> adj[N];
   priority_queue < pair < int, int >>,
       vector < pair < int, int >>, greater <
       pair < int, int >>> pq;
   pq.push({0, start});
   while (!pq.empty()) {
       auto it = pq.top();
       pq.pop();
       int wt = it.first;
       int u = it.second;
```

```
if (vis[u])
    continue;
vis[u] = 1;
for (pair<int, int> i : adj[u]) {
    int adjWt = i.second;
    int adjNode = i.first;
    if (dist[adjNode] > wt + adjWt)
        {
        dist[adjNode] = wt + adjWt;
        pq.push({dist[adjNode],
            adjNode});
    }
}
```

3.4 Bellman Ford

```
vector < int > dist;
vector < int > parent;
vector < vector < pair < int , int >>> adj;
// resize the vectors from main
    function
void bellmanFord(int num_of_nd, int
    src) {
  dist[src] = 0;
  for (int step = 0; step < num_of_nd;</pre>
       step++) {
    for (int i = 1; i <= num_of_nd; i</pre>
        ++) {
      for (auto it : adj[i]) {
        int u = i;
         int v = it.first;
         int wt = it.second;
         if (dist[u] != inf &&
             ((dist[u] + wt) < dist[v])
                 ) {
           if (step == num_of_nd - 1) {
             cout << "Negative cycle</pre>
                 found\n ";
                 return:
           dist[v] = dist[u] + wt;
           parent[v] = u;
        }
      }
    }
  }
  for (int i = 1;i <= num_of_nd; i++)</pre>
    cout << dist[i] << " ";</pre>
  cout << endl;</pre>
```

3.5 Floyed Warshall Algorithm

```
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FloydWarshall(VVT &w, VVI &prev)
    {
  int n = w.size();
  prev = VVI(n, VI(n, -1));
```

3.6 Kruskal Algorithm (MST)

```
vector<pair<int, pair<int, int>>>
   Krushkal(vector<pair<int, pair<int</pre>
   , int>>> &edges, int n) {
  sort(edges.begin(), edges.end());
  vector<pair<int, pair<int, int>>>
     ans;
  DisjointSet D(n);
  for (auto it : edges) {
    if (D.findUPar(it.second.first) !=
        D.findUPar(it.second.second))
      ans.push_back({it.first, {it.
          second.first, it.second.
          second}});
      D.unionBySize(it.second.first,
          it.second.second);
    }
 }
 return ans;
```

3.7 Prims Algorithm (MST)

```
void Prims(int start) {
  // map <int, vector <pair <int, int>>>
      adj, ans;
  priority_queue < pair < int , pair < int ,</pre>
     int>>, vector<pair<int, pair<int</pre>
      , int>>>, greater<pair<int, pair
      <int, int>>>> pq;
  pq.push({0, {start, -1}});
  while (!pq.empty()) {
    auto it = pq.top();
    pq.pop();
    int wt = it.first;
    int u = it.second.first;
    int v = it.second.second;
    if (vis[u]) continue;
    vis[u] = 1;
    if (v != -1) ans[u].push_back({v,
    for (pair<int, int> i : adj[u]) {
      int adjWt = i.second;
      int adjNode = i.first;
```

```
if (!vis[adjNode]) pq.push({
          adjWt, {adjNode, u}});
}
```

3.8 Strongly Connected Components

```
vector < bool > visited; // keeps track
    of which vertices are already
   visited
// runs depth first search starting at
// each visited vertex is appended to
   the output vector when dfs leaves
   it.
void dfs(int v, vector<vector<int>>
   const &adj, vector<int> &output) {
  visited[v] = true;
  for (auto u : adj[v])
    if (!visited[u])
      dfs(u, adj, output);
  output.push_back(v);
// input: adj -- adjacency list of G
// output: components -- the strongy
    connected components in G
// output: adj\_cond -- adjacency list
   of G^SCC (by root vertices)
void scc(vector<vector<int>> const &
   adj, vector < vector < int >> &
   components, vector<vector<int>> &
   adj_cond) {
  int n = adj.size();
  components.clear(), adj_cond.clear()
  vector<int> order; // will be a
     sorted list of G's vertices by
      exit time
  visited.assign(n, false);
  // first series of depth first
     searches
  for (int i = 0; i < n; i++)</pre>
    if (!visited[i])
      dfs(i, adj, order);
  // create adjacency list of G^T
  vector < vector < int >> adj_rev(n);
  for (int v = 0; v < n; v++)</pre>
    for (int u : adj[v])
      adj_rev[u].push_back(v);
  visited.assign(n, false);
  reverse(order.begin(), order.end());
  vector<int> roots(n, 0); // gives
      the root vertex of a vertex's
```

```
SCC
  // second series of depth first
      searches
  for (auto v : order)
    if (!visited[v]) {
      std::vector<int> component;
      dfs(v, adj_rev, component);
      components.push_back(component);
      int root = *min_element(begin(
          component), end(component));
      for (auto u : component)
        roots[u] = root;
  // add edges to condensation graph
  adj_cond.assign(n, {});
  for (int v = 0; v < n; v++)
    for (auto u : adj[v])
      if (roots[v] != roots[u])
        adj_cond[roots[v]].push_back(
            roots[u]);
}
3.9 LCA
struct LCA {
  vector<int> height, euler, first,
      segtree, parent;
  vector < bool > visited;
  vector < int >> jump;
  int n;
  LCA(vector < vector < int >> &adj, int
     root = 0) {
    n = adj.size();
    height.resize(n);
    first.resize(n);
    parent.resize(n);
    euler.reserve(n * 2);
    visited.assign(n, false);
    dfs(adj, root);
    int m = euler.size();
    segtree.resize(m * 4);
    build(1, 0, m - 1);
    jump.resize(n, vector<int>(32, -1)
        );
    for(int i=0;i<n;i++) {</pre>
        jump[i][0] = parent[i];
    for(int j=1; j<20; j++) {</pre>
        for(int i=0;i<n;i++) {</pre>
            int mid = jump[i][j-1];
if(mid != -1) jump[i][j] =
                 jump[mid][j-1];
        }
    }
  }
```

```
void dfs(vector<vector<int>> &adj,
   int node, int h = 0) {
  visited[node] = true;
  height[node] = h;
  first[node] = euler.size();
  euler.push_back(node);
  for (auto to : adj[node]) {
    if (!visited[to]) {
      parent[to] = node;
      dfs(adj, to, h + 1);
      euler.push_back(node);
 }
}
void build(int node, int b, int e) {
  if (b == e) {
    segtree[node] = euler[b];
  } else {
    int mid = (b + e) / 2;
    build(node << 1, b, mid);</pre>
    \verb|build(node << 1 | 1, mid + 1, e)|
    int l = segtree[node << 1], r =</pre>
        segtree[node << 1 | 1];</pre>
    segtree[node] = (height[1] <
        height[r]) ? 1 : r;
}
int query(int node, int b, int e,
   int L, int R) {
  if (b > R || e < L)
    return -1;
  if (b >= L && e <= R)
   return segtree[node];
  int mid = (b + e) >> 1;
  int left = query(node << 1, b, mid</pre>
     , L, R);
  int right = query(node << 1 | 1,</pre>
     mid + 1, e, L, R);
  if (left == -1)
    return right;
  if (right == -1)
    return left;
  return height[left] < height[right</pre>
     ] ? left : right;
int lca(int u, int v) {
  int left = first[u], right = first
      [v];
  if (left > right)
    swap(left, right);
  return query(1, 0, euler.size() -
      1, left, right);
}
int kthParent(int u, int k) {
    for(int i=0;i<19;i++) {</pre>
        if(k & (1LL<<i)) u = jump[u</pre>
            ][i];
```

```
return u;
 }
};
3.10
     Max Flow
const int N = 505;
int capacity[N][N];
int vis[N], p[N];
int n, m;
int bfs(int s, int t) {
  memset(vis, 0, sizeof vis);
  queue < int > qu;
  qu.push(s);
  vis[s] = 1;
  while (!qu.empty()) {
    int u = qu.front();
    qu.pop();
    for (int i = 0; i \le n + m + 2; i
       ++) {
      if (capacity[u][i] > 0 && !vis[i
          ]) {
        p[i] = u;
        vis[i] = 1;
        qu.push(i);
    }
 }
  return vis[t] == 1;
int maxflow(int s, int t) {
  int cnt = 0;
  while (bfs(s, t)) {
    int cur = t;
    while (cur != s) {
     int prev = p[cur];
      capacity[prev][cur] -= 1;
      capacity[cur][prev] += 1;
      cur = prev;
    }
    cnt++;
  }
  return cnt;
    Data Structures
Different Data Structures.
4.1
     Segment Tree
constexpr int N = 100005;
int arr[N], seg[N];
void build(int ind, int low, int high)
  if (low == high) {
```

seg[ind] = arr[low];

return;

}

```
int mid = (low + high) / 2;
                                              pull(node, 1, r);
  build(2 * ind + 1, low, mid);
  build(2 * ind + 2, mid + 1, high);
  seg[ind] = seg[2 * ind + 1] + seg[2
                                            void update(int node, int 1, int r,
     * ind + 2];
                                                int ql, int qr, int val) {
                                              if(qr < 1 || r < q1) return;</pre>
int query(int ind, int low, int high,
                                              if (q1 <= 1 && r <= qr) {
   int 1, int r) {
                                                seg[node] += val;
  if (low >= 1 && high <= r) return</pre>
                                                lz[node] += val;
     seg[ind];
                                                return;
  if (low > r || high < 1) return 0;</pre>
                                              }
  int mid = (low + high) / 2;
                                              push(node, 1, r);
  int left = query(2 * ind + 1, low,
     mid, 1, r);
  int right = query(2 * ind + 2, mid +
                                              int mid = (1 + r) >> 1;
      1, high, l, r);
                                              update(node << 1, 1, mid, ql, qr,
  return left + right;
                                                  val);
}
                                              update(node << 1 | 1, mid + 1, r, ql
void update(int ind, int low, int high
                                                  , qr, val);
    , int node, int val) {
  if (low == high) {
                                              pull(node, 1, r);
    seg[ind] = val;
    return;
                                            int query(int node, int 1, int r, int
  int mid = (low + high) / 2;
                                                ql, int qr) {
  if (low <= node && node <= mid)</pre>
                                              if(qr < 1 || r < ql) return 0;</pre>
     update(2 * ind + 1, low, mid,
                                              if(ql <= l && r <= qr) return seg[</pre>
     node, val);
                                                  node];
  else update(2 * ind + 2, mid + 1,
     high, node, val);
                                              push(node, 1, r);
  seg[ind] = seg[2 * ind + 1] + seg[2
     * ind + 2];
                                              int mid = (1 + r) >> 1;
                                              return query(node << 1, 1, mid, ql,</pre>
                                                  qr) + query(node << 1 | 1, mid +
     Segment Tree Lazy
                                                   1, r, ql, qr);
const int N = 1e5 + 5;
                                            4.3 Fenwick Tree
int arr[N], seg[N << 2], lz[N << 2];</pre>
void pull(int node, int 1, int r) {
                                            int fenwick[N];
  seg[node] = seg[node << 1] + seg[</pre>
     node << 1 | 1];
                                            void update(int ind, int val) {
                                              while (ind < N) {
                                                fenwick[ind] += val;
void push(int node, int 1, int r) {
                                                 ind += ind & -ind;
  int mid = (1 + r) >> 1;
                                              }
  lz[node << 1] += lz[node];</pre>
                                            }
  seg[node << 1] += lz[node] * (mid -
                                            int query(int ind) {
     1 + 1);
                                              int sum = 0;
  lz[node << 1 | 1] += lz[node];</pre>
                                              while (ind > 0) {
  seg[node << 1 | 1] += lz[node] * (r
                                                sum += fenwick[ind];
     - mid);
                                                ind -= ind & -ind;
  lz[node] = 0;
                                              }
}
                                              return sum;
void build(int node, int 1, int r) {
                                            4.4 Disjoint Set
  if(1 == r) {
    seg[node] = arr[1];
                                            class DisjointSet {
    lz[node] = 0;
                                              vector < int > parent, sz;
    return;
  int mid = (1 + r) >> 1;
                                             public:
  build(node << 1, 1, mid);</pre>
                                              DisjointSet(int n) {
  build(node << 1 | 1, mid + 1, r);
                                                sz.resize(n + 1);
```

```
bool isChild(Node* node) {
    parent.resize(n + 2);
    for (int i = 1; i <= n; i++)</pre>
                                              for (int i = 0; i < N; i++)</pre>
       parent[i] = i, sz[i] = 1;
                                                if (node->child[i] != NULL) return
                                                     true;
  int findUPar(int u) { return parent[
                                              return false;
     u] == u ? u : parent[u] =
     findUPar(parent[u]); }
  void unionBySize(int u, int v) {
                                            bool isJunc(Node* node) {
    int a = findUPar(u);
                                              int cnt = 0;
    int b = findUPar(v);
                                              for (int i = 0; i < N; i++) {</pre>
    if (sz[a] < sz[b]) swap(a, b);</pre>
                                                if (node->child[i] != NULL) cnt++;
    if (a != b) {
      parent[b] = a;
                                              if (cnt > 1) return true;
      sz[a] += sz[b];
                                              return false;
 }
};
                                            int trie_delete(Node* node, string s,
                                                int k = 0) {
     TRIE
4.5
                                              if (node == NULL) return 0;
                                              if (k == (int)s.size()) {
const int N = 26;
                                                if (node->EoW == 0) return 0;
class Node {
                                                if (isChild(node)) {
public:
                                                   node -> EoW = 0;
 int EoW;
                                                   return 0;
  Node* child[N];
                                                }
  Node() {
                                                return 1;
    EoW = 0;
                                              }
    for (int i = 0; i < N; i++) child[</pre>
                                              int r = s[k] - A';
       i] = NULL;
                                               int d = trie_delete(node->child[r],
 }
                                                  s, k + 1);
};
                                              int j = isJunc(node);
                                              if (d) delete node->child[r];
void insert(Node* node, string s) {
                                              if (j) return 0;
  for (size_t i = 0; i < s.size(); i</pre>
                                              return d;
     ++) {
                                            }
    int r = s[i] - A';
    if (node->child[r] == NULL) node->
                                            void delete_trie(Node* node) {
       child[r] = new Node();
                                              for (int i = 0; i < 15; i++) {</pre>
    node = node->child[r];
                                                if (node->child[i] != NULL)
 }
                                                    delete_trie(node->child[i]);
  node \rightarrow EoW += 1;
                                              }
                                              delete node;
                                            }
int search(Node* node, string s) {
                                            4.6 Ordered Set
  for (size_t i = 0; i < s.size(); i</pre>
     ++) {
    int r = s[i] - A';
                                            #include <ext/pb_ds/assoc_container.</pre>
    if (node->child[r] == NULL) return
                                                hpp>
                                            #include <ext/pb_ds/tree_policy.hpp>
 }
                                            using namespace __gnu_pbds;
  return node -> EoW;
                                            #define ordered_set tree<int,</pre>
                                                null_type,less<int>, rb_tree_tag,
                                                tree_order_statistics_node_update>
void print(Node* node, string s = "")
                                            ordered_set o_set;
  if (node->EoW) cout << s << "\n";</pre>
  for (int i = 0; i < N; i++) {</pre>
                                            // insert function to insert in
                                            // ordered set same as SET STL
    if (node->child[i] != NULL) {
      char c = i + 'A';
                                            o_set.insert(5);
      print(node->child[i], s + c);
                                            o_set.insert(1);
    }
                                            o_set.insert(2);
 }
}
                                            // Finding the second smallest element
                                            // in the set using * because
```

```
Monotonic Stack (Immediate
                                            5.1
// find_by_order returns an iterator
*(o_set.find_by_order(1))
                                                  Large)
// Finding the number of elements
                                            for (int i = n - 1; i >= 0; i--) {
// strictly less than k=4
                                              while (!stk.empty() && v[i] >= v[stk
o_set.order_of_key(4)
                                                  .top()]) stk.pop();
     Mo's Algorithm
4.7
                                              ind[i] = stk.empty() ? -1 : stk.top
                                                  ();
                                              stk.push(i);
void remove(int idx);
void add(int idx);
                                            // 3 1 5 4 10
int get_answer();
                                            // 2 2 4 4 -1
int block_size = 700;
                                            5.2 Kadane's Algorithm
struct Query {
  int l, r, idx;
                                            int maxSubArraySum(vector<int> &a) {
  bool operator < (Query other) const {</pre>
                                              int size = a.size();
    if (1 / block_size != other.1 /
                                              int maxTill = INT_MIN, maxEnd = 0;
       block_size) return 1 /
                                              for (int i = 0; i < size; i++) {</pre>
        block_size < other.1 /
                                                maxEnd = maxEnd + a[i];
        block_size;
                                                if (maxTill < maxEnd) maxTill =</pre>
    return r < other.r;</pre>
                                                    maxEnd;
  }
                                                if (maxEnd < 0) maxEnd = 0;</pre>
};
                                              }
                                              return maxTill;
vector<int> mo_s_algorithm(vector<</pre>
                                            }
   Query > queries) {
  vector < int > answers (queries.size());
                                                 2D Prefix Sum
                                            5.3
  sort(queries.begin(), queries.end())
                                            pref[i][j] = a[i][j] + pref[i - 1][j]
                                                + pref[i][j - 1] - pref[i - 1][j -
  // TODO: initialize data structure
                                                1];
                                            Sum of region = pref[row2 + 1][col2 +
  int cur_1 = 0;
                                                1] - pref[row2 + 1][col1] - pref[
  int cur_r = -1;
                                                row1][col2 + 1] + pref[row1][col1
  // invariant: data structure will
                                                1:
      always reflect the range [cur_l,
       cur\_r]
                                            5.4 2D Convex Hull
  for (Query q : queries) {
    while (cur_1 > q.1) {
      cur_1 --;
                                            struct pt {
      add(cur_1);
                                                double x, y;
                                                bool operator == (pt const& t)
    while (cur_r < q.r) {</pre>
                                                    const {
      cur_r++;
                                                    return x == t.x && y == t.y;
      add(cur_r);
                                                }
    }
                                            };
    while (cur_1 < q.1) {</pre>
      remove(cur_1);
                                            int orientation(pt a, pt b, pt c) {
      cur_1++;
                                                double v = a.x*(b.y-c.y)+b.x*(c.y-
    }
                                                    a.y)+c.x*(a.y-b.y);
    while (cur_r > q.r) {
                                                if (v < 0) return -1; // clockwise
      remove(cur_r);
                                                if (v > 0) return +1; // counter-
      cur_r --;
                                                    clockwise
                                                return 0;
    answers[q.idx] = get_answer();
  return answers;
                                            bool cw(pt a, pt b, pt c, bool
}
                                                include_collinear) {
                                                int o = orientation(a, b, c);
    Algorithms
5
                                                return o < 0 || (include_collinear</pre>
                                                     \&\& o == 0);
All about algorithms.
                                            }
```

```
bool collinear(pt a, pt b, pt c) {
   return orientation(a, b, c) == 0;
void convex_hull(vector<pt>& a, bool
    include_collinear = false) {
    pt p0 = *min_element(a.begin(), a.
        end(), [](pt a, pt b) {
        return make_pair(a.y, a.x) <</pre>
            make_pair(b.y, b.x);
    });
    sort(a.begin(), a.end(), [&p0](
        const pt& a, const pt& b) {
        int o = orientation(p0, a, b);
        if (o == 0)
            return (p0.x-a.x)*(p0.x-a.
                x) + (p0.y-a.y)*(p0.y-
                a.y)
                 < (p0.x-b.x)*(p0.x-b.x)
                    ) + (p0.y-b.y)*(p0
                    .y-b.y);
        return o < 0;</pre>
    });
    if (include_collinear) {
        int i = (int)a.size()-1;
        while (i >= 0 && collinear(p0,
             a[i], a.back())) i--;
        reverse(a.begin()+i+1, a.end()
           );
    }
    vector <pt> st;
    for (int i = 0; i < (int)a.size();</pre>
        i++) {
        while (st.size() > 1 && !cw(st
            [st.size()-2], st.back(),
            a[i], include_collinear))
            st.pop_back();
        st.push_back(a[i]);
    }
    if (include_collinear == false &&
        st.size() == 2 && st[0] == st
        [1])
        st.pop_back();
    a = st;
}
    Number Theory
All about math.
6.1 nCr
```

```
int inverseMod(int a, int m) { return
   power(a, m - 2); }
int nCr(int n, int r, int m = mod){
   if(r==0) return 1;
   if(r>n) return 0;
   return (fact[n] * inverseMod((fact[
      r] * fact[n-r]) % m , m)) % m;
```

```
}
6.2 Power
int power(int base, int n, int m = mod
   ) {
  if (n == 0) return 1;
  if (n & 1) {
    int x = power(base, n / 2);
    return ((x * x) % m * base) % m;
  }
  else {
    int x = power(base, n / 2);
    return (x * x) % m;
  }
}
6.3
    Miller Rabin
using u64 = uint64_t;
using u128 = __uint128_t;
u64 binpower(u64 base, u64 e, u64 mod)
  u64 result = 1;
  base %= mod;
  while (e) {
    if (e & 1) result = (u128)result *
        base % mod;
    base = (u128)base * base % mod;
    e >>= 1;
 }
  return result;
bool check_composite(u64 n, u64 a, u64
    d, int s) {
  u64 x = binpower(a, d, n);
  if (x == 1 || x == n - 1) return
     false;
  for (int r = 1; r < s; r++) {</pre>
    x = (u128)x * x % n;
    if (x == n - 1) return false;
  }
 return true;
};
bool MillerRabin(u64 n, int iter = 5)
   { // returns true if n is
   probably prime, else returns false
  if (n < 4) return n == 2 || n == 3;</pre>
  int s = 0;
  u64 d = n - 1;
  while ((d & 1) == 0) {
    d >>= 1;
    s++;
  for (int i = 0; i < iter; i++) {</pre>
    int a = 2 + rand() % (n - 3);
    if (check_composite(n, a, d, s))
       return false;
  }
```

for (int i = 0; i <= R - L; i++) {</pre>

```
if (isPrime[i] == true) cout << L</pre>
  return true;
                                                     + i << '\n';
                                              }
     Sieve
6.4
                                               cout << '\n';
const int N = 1e7 + 3;
vector<int> primes;
int notprime[N];
                                                  Divisors
void sieve() {
                                             constexpr int N = 1000005;
  primes.push_back(2);
  for (int i = 4; i < N; i += 2) {</pre>
                                             int Prime[N + 4], kk;
    notprime[i] = true;
                                             bool notPrime[N + 5];
                                             void SieveOf() {
  for (int i = 3; i < N; i += 2) {</pre>
                                               notPrime[1] = true;
    if (!notprime[i]) {
                                               Prime[kk++] = 2;
      primes.push_back(i);
                                               for (int i = 4; i <= N; i += 2)</pre>
      for (int j = i * i; j < N; j +=</pre>
                                                  notPrime[i] = true;
          2 * i) {
                                               for (int i = 3; i <= N; i += 2) {</pre>
        notprime[j] = true;
                                                 if (!notPrime[i]) {
                                                   Prime[kk++] = i;
    }
                                                   for (int j = i * i; j <= N; j +=</pre>
  }
                                                        2 * i) notPrime[j] = true;
}
                                              }
6.5
     Bitset Sieve
                                            }
const int sieve_size = 10000006;
                                             void Divisors(int n) {
bitset < sieve_size > sieve;
                                               int sum = 1, total = 1;
                                               int mnP = INT_MAX, mxP = INT_MIN,
void Sieve() {
                                                  cntP = 0, totalP = 0;
  sieve.flip();
                                               for (int i = 0; i <= N && Prime[i] *</pre>
  int finalBit = sqrt(sieve.size()) +
                                                   Prime[i] <= n; i++) {
                                                 if (n % Prime[i] == 0) {
  for (int i = 2; i < finalBit; ++i) {</pre>
                                                   mnP = min(mnP, Prime[i]);
    if (sieve.test(i))
                                                   mxP = max(mnP, Prime[i]);
      for (int j = 2 * i; j <</pre>
                                                   int k = 0;
          sieve_size; j += i) sieve.
                                                   cntP++;
          reset(j);
                                                   while (n % Prime[i] == 0) {
                                                     k++;
}
                                                     n /= Prime[i];
6.6
     Segmented Sieve
                                                   sum *= (k + 1); // NOD
Need Sieve() from 6.4
                                                   totalP += k;
void segmentedSieve(int L, int R) {
                                                   int s = 0, p = 1;
  bool isPrime[R - L + 1];
                                                   while (k-- >= 0) {
  for (int i = 0; i <= R - L + 1; i++)
                                                     s += p;
       isPrime[i] = true;
                                                     p *= Prime[i];
  if (L == 1) isPrime[0] = false;
                                                   };
  for (int i = 0; primes[i] * primes[i
                                                   total *= s; // SOD
     ] <= R; i++) {
                                               }
    int curPrime = primes[i];
                                               if (n > 1) {
    int base = curPrime * curPrime;
    if (base < L) {</pre>
                                                 cntP++, totalP++;
      base = ((L + curPrime - 1) /
                                                 sum *= 2;
          curPrime) * curPrime;
                                                 total *= (1 + n);
                                                 mnP = min(mnP, n);
    for (int j = base; j <= R; j +=</pre>
                                                 mxP = max(mnP, n);
                                              }
        curPrime) isPrime[j - L] =
                                               cout << mnP << " " << mxP << " " <<
        false;
```

cntP << " " << totalP << " " <<
sum << " " << total << "\n";</pre>

```
}
                                              const int v = sqrt(N), vv = sqrt(v);
     Euler's Totient Phi Function
const int N = 5000005;
int phi[N];
unsigned long long phiSum[N];
void phiCalc() {
  for (int i = 2; i < N; i++) phi[i] =</pre>
       i;
  for (int i = 2; i < N; i++) {</pre>
    if (phi[i] == i) {
      for (int j = i; j < N; j += i) {</pre>
        phi[j] -= phi[j] / i;
    }
 }
  for (int i = 2; i < N; i++) {</pre>
    phiSum[i] = (unsigned long long)
       phi[i] * (unsigned long long)
       phi[i] + phiSum[i - 1];
}
     Log a base b
6.9
int logab (int a, int b){
  return log2(a) / log2(b);
                                                }
6.10 Count 1's from 0 to n
int cntOnes(int n) {
                                                 ):
 int cnt = 0;
  for(int i=1;i<=n;i<<=1) {</pre>
                                                  {
    int x = (n + 1) / (i << 1);
    cnt += x * i;
    if((n + 1) % i && n & i) cnt += (n
        + 1) % i;
  return cnt;
                                                }
       Primes Upto 1e9
6.11
                                              }
// credit: min_25
// takes 0.5s for n = 1e9
vector<int> sieve(const int N, const
   int Q = 17, const int L = 1 << 15)
    {
  static const int rs[] = \{1, 7, 11,
     13, 17, 19, 23, 29};
  struct P {
    P(int p) : p(p) {}
    int p; int pos[8];
  };
  auto approx_prime_count = [] (const
     int N) -> int {
    return N > 60184 ? N / (log(N) -
        1.1)
```

: max(1., N / (

log(N) -

};

1.11)) + 1;

```
vector < bool > isp(v + 1, true);
for (int i = 2; i <= vv; ++i) if (
   isp[i]) {
  for (int j = i * i; j <= v; j += i
     ) isp[j] = false;
const int rsize = approx_prime_count
   (N + 30);
vector < int > primes = {2, 3, 5}; int
   psize = 3;
primes.resize(rsize);
vector <P> sprimes; size_t pbeg = 0;
int prod = 1;
for (int p = 7; p \le v; ++p) {
  if (!isp[p]) continue;
  if (p <= Q) prod *= p, ++pbeg,</pre>
     primes[psize++] = p;
  auto pp = P(p);
  for (int t = 0; t < 8; ++t) {</pre>
    int j = (p \le Q) ? p : p * p;
    while (j \% 30 != rs[t]) j += p
        << 1;
   pp.pos[t] = j / 30;
  sprimes.push_back(pp);
vector<unsigned char> pre(prod, 0xFF
for (size_t pi = 0; pi < pbeg; ++pi)</pre>
  auto pp = sprimes[pi]; const int p
      = pp.p;
  for (int t = 0; t < 8; ++t) {</pre>
    const unsigned char m = ~(1 << t</pre>
    for (int i = pp.pos[t]; i < prod</pre>
        ; i += p) pre[i] &= m;
const int block_size = (L + prod -
   1) / prod * prod;
vector < unsigned char > block (
   block_size); unsigned char*
   pblock = block.data();
const int M = (N + 29) / 30;
for (int beg = 0; beg < M; beg +=</pre>
   block_size, pblock -= block_size
  int end = min(M, beg + block_size)
  for (int i = beg; i < end; i +=</pre>
     prod) {
    copy(pre.begin(), pre.end(),
        pblock + i);
  }
  if (beg == 0) pblock[0] &= 0xFE;
```

```
for (size_t pi = pbeg; pi <</pre>
      sprimes.size(); ++pi) {
    auto& pp = sprimes[pi];
    const int p = pp.p;
    for (int t = 0; t < 8; ++t) {</pre>
      int i = pp.pos[t]; const
          unsigned char m = (1 \ll t)
          );
      for (; i < end; i += p) pblock</pre>
          [i] &= m;
      pp.pos[t] = i;
  }
  for (int i = beg; i < end; ++i) {</pre>
    for (int m = pblock[i]; m > 0; m
         \&= m - 1) {
      primes[psize++] = i * 30 + rs[
          __builtin_ctz(m)];
    }
  }
}
assert(psize <= rsize);</pre>
while (psize > 0 && primes[psize -
   1] > N) --psize;
primes.resize(psize);
return primes;
     Matrix Multiplication
```

6.12

```
const int mod = 1000000007;
vector < vector < i64>> multi(vector <</pre>
   vector < i64>> &m1, vector < vector <</pre>
   i64>> &m2) {
  int n = m1.size();
  vector < vector < i64>> res(n, vector <
     i64>(n, 0));
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++) {
      for (int k = 0; k < n; k++) {
        res[i][j] = (res[i][j] + (m1[i
            ][k] * m2[k][j]) % mod) %
            mod:
      }
    }
 }
  return res;
vector < vector < i64>> power (vector <
   vector < i64>> &base, int n, int m =
    mod) {
  int r = base.size();
  vector < vector < i64 >> identity(r,
      vector < i64 > (r));
  for (int i = 0; i < r; i++) {</pre>
    identity[i][i] = 1;
  while (n > 0) {
    if (n % 2 == 1) {
      identity = multi(base, identity)
    }
```

```
base = multi(base, base);
  n /= 2;
}
return identity;
```

6.13 Extended GCD

```
int exgcd(int a, int b, int &x, int &y
   ) {
    if (b == 0) {
        x = 1;
        y = 0;
        return a;
    int g = exgcd(b, a % b, y, x);
    y -= a / b * x;
    return g;
}
```

Dynamic Programming

7.1 LCS (Longest Common Subsequence)

```
string s, t;
vector < vector < int >> dp(3003, vector <
   int > (3003, -1));
vector < vector < int >> mark(3003, vector <</pre>
   int > (3003));
int f(int i, int j) {
  if (i < 0 || j < 0) return 0;</pre>
  if (dp[i][j] != -1) return dp[i][j];
  int res = 0;
  if (s[i] == t[j]) {
    mark[i][j] = 1;
    res = 1 + f(i - 1, j - 1);
  }
  else {
    int iC = f(i - 1, j);
    int jC = f(i, j - 1);
    if (iC > jC) mark[i][j] = 2;
    else mark[i][j] = 3;
    res = max(iC, jC);
  }
  return dp[i][j] = res;
}
void printWay(int i, int j) {
  if (i < 0 || j < 0) return;
  if (mark[i][j] == 1) printWay(i - 1,
       j - 1), cout << s[i];</pre>
  else if (mark[i][j] == 2) printWay(i
       - 1, j);
  else if (mark[i][j] == 3) printWay(i
      , j - 1);
}
```

7.2 LIS (Longest Increasing Subsequence)

other_mask];

}

```
void lis(vector<int> &v) {
                                                }
  int n = v.size();
  vector < int > dp(n + 1, 1), hash(n);
  int mx = 1, lastInd = 0;
                                              for (int mask : v) cout << dp[BITS][</pre>
  for (int i = 0; i < n; i++) {</pre>
                                                 mask] << '\n';
    hash[i] = i;
    for (int prev = 0; prev < i; prev</pre>
                                            // dp[bits][mask] means left most '
                                               bits' of submasks can be differ
       ++) {
                                            // for dp[1][11] 01, 00 are now allow
      if (v[i] > v[prev] && 1 + dp[
                                               because leftmost 1 bit can be
         prev] > dp[i]) {
        dp[i] = 1 + dp[prev];
                                               differ. 10 and 11 are allowed.
        hash[i] = prev;
                                            // for travarsing all the submask of a
      }
                                                mask we can use
    }
                                            // submask = mask
                                            // do {
    if (mx < dp[i]) {</pre>
      mx = dp[i];
                                            lastInd = i;
                                            // } while (submask)
    }
                                            7.4 MCM (Matrix Chain Multipli-
  }
  vector < int > printSeq;
                                                 cation)
  printSeq.push_back(v[lastInd]);
  while (hash[lastInd] != lastInd) {
                                            ans = min(ans, MCM(i, k) + MCM(k + 1,
    lastInd = hash[lastInd];
                                               j) + v[i - 1] * v[k] * v[j]);
    printSeq.push_back(v[lastInd]);
                                                Strings
  reverse(printSeq.begin(), printSeq.
     end());
                                           8.1 KMP
  cout << mx << "\n";
  for (int i : printSeq) cout << i <<</pre>
                                            vector<int> prefix_function(string s)
     " ";
                                               {
  cout << "\n";
                                                int n = (int)s.length();
                                                vector < int > pi(n);
     SOS (Sum Of Subsets)
                                                for (int i = 1; i < n; i++) {</pre>
                                                    int j = pi[i - 1];
void SOS (vector<int> &v) {
                                                    while (j > 0 \&\& s[i] != s[j])
  const int BITS = log2(*max_element(v
                                                        j = pi[j - 1];
                                                    if (s[i] == s[j]) j++;
     .begin(), v.end())) + 1;
                                                    pi[i] = j;
  vector<int> freq(1 << BITS, 0);</pre>
  for (int mask : v) freq[mask]++;
                                                }
                                                return pi;
  vector < vector < int >> dp(BITS + 1,
     vector < int > (1 << BITS, 0));</pre>
  for (int mask = 0; mask < 1 << BITS;</pre>
                                            vector<int> find_matches(string text,
      mask++) {
                                               string pat) {
                                                int n = pat.length(), m = text.
    dp[0][mask] = freq[mask];
                                                    length();
                                                string s = pat + "$" + text;
  for (int bits = 1; bits <= BITS;</pre>
                                                vector<int> pi = prefix_function(s
     bits++) {
                                                   ), ans;
    for (int mask = 0; mask < 1 <<</pre>
                                                for (int i = n; i <= n + m; i++) {
       BITS; mask++) {
                                                    if (pi[i] == n) {
      if ((mask & (1 << (bits - 1)))</pre>
                                                        ans.push_back(i - 2 * n);
          == 0) {
                                                    }
        dp[bits][mask] = dp[bits - 1][
            mask];
                                                return ans;
                                            }
      else {
                                                Double Hashing
                                            8.2
        int other_mask = mask - (1 <<</pre>
            (bits - 1));
                                           Need power() from 6.2
        dp[bits][mask] = dp[bits - 1][
            mask] + dp[bits - 1][
                                            const int N = 1e6 + 9;
```

const int MOD1 = 127657753, MOD2 =

987654319;

```
const int p1 = 137, p2 = 277;
int ip1, ip2;
pair < int , int > pw[N] , ipw[N];
void prec() {
  pw[0] = {1, 1};
for (int i = 1; i < N; i++) {</pre>
    pw[i].first = 1LL * pw[i - 1].
        first * p1 % MOD1;
    pw[i].second = 1LL * pw[i - 1].
        second * p2 % MOD2;
  }
  ip1 = power(p1, MOD1 - 2, MOD1);
  ip2 = power(p2, MOD2 - 2, MOD2);
  ipw[0] = \{1, 1\};
  for (int i = 1; i < N; i++) {</pre>
    ipw[i].first = 1LL * ipw[i - 1].
        first * ip1 % MOD1;
    ipw[i].second = 1LL * ipw[i - 1].
        second * ip2 % MOD2;
  }
struct Hashing {
  int n;
  string s; // 0 - indexed
  vector<pair<int, int>> hs; // 1 -
      indexed
  Hashing() {}
  Hashing(string _s) {
    n = _s.size();
    s = _s;
    hs.emplace_back(0, 0);
    for (int i = 0; i < n; i++) {</pre>
      pair < int , int > p;
      p.first = (hs[i].first + 1LL *
          pw[i].first * s[i] % MOD1) %
      p.second = (hs[i].second + 1LL *
           pw[i].second * s[i] % MOD2)
           % MOD2;
      hs.push_back(p);
    }
  }
  pair < int , int > get_hash(int 1, int r
     ) { // 1 - indexed
    assert(1 <= 1 && 1 <= r && r <= n)
    pair<int, int> ans;
    ans.first = (hs[r].first - hs[l -
        1].first + MOD1) * 1LL * ipw[1
         - 1].first % MOD1;
    ans.second = (hs[r].second - hs[l])
        - 1].second + MOD2) * 1LL *
        ipw[l - 1].second % MOD2;
    return ans;
  }
  pair < int , int > get_hash() {
    return get_hash(1, n);
  }
};
```

8.3 Manacher's Algorithm

```
vector<int> manacher_odd(string s) {
    int n = s.size();
    s = "$" + s + "^";
    vector < int > p(n + 2);
    int 1 = 1, r = 1;
    for(int i = 1; i <= n; i++) {
        p[i] = max(0, min(r - i, p[1 +
             (r - i)]));
        while (s[i - p[i]] == s[i + p[i]]
            ]]) {
            p[i]++;
        if(i + p[i] > r) {
            1 = i - p[i], r = i + p[i]
                ];
    }
    return vector<int>(begin(p) + 1,
       end(p) - 1);
}
vector<int> manacher(string s) {
    string t;
    for(auto c: s) t += string("#") +
       с;
    auto res = manacher_odd(t + "#");
    return vector<int>(begin(res) + 1,
         end(res) - 1);
}
// p[2 * i + (!odd)] - 1
// if r - l + 1 \le p[(r + l) / 2 + (r + l)]
   % 2 != 1 % 2)] is true then
   palindrome. Same parity means odd
   length.
```

9 Stress Testing

9.1 Bash Stress File

```
#!/usr/bin/env bash
g++ -g "$1".cpp -DONPC -o "$1"
g++ -g "$2".cpp -DONPC -o "$2"
g++ -g "$3".cpp -DONPC -o "$3"
for ((testNum=0;testNum<$4;testNum++))</pre>
    ./$3 > stdinput
    ./$2 < stdinput > outSlow
    ./$1 < stdinput > outWrong
    H1='md5sum outWrong'
    H2='md5sum outSlow'
    if !(cmp -s "outWrong" "outSlow")
    then
        echo "Error found!"
        echo "Input:"
        cat stdinput
        echo "Wrong Output:"
        cat outWrong
        echo "Slow Output:"
        cat outSlow
        exit
```