Lab Report-05

Course Title: Artificial Intelligence

Course Code: CSE 366

Semester: Fall 2021

Section No: 01

Submitted By

Name: Syeda Tamanna Sheme

ID:2018-2-60-010

Submitted To

Md Al-Imran

Lecturer

Department of Computer Science & Engineering

Date of Submission: 1 December, 2021

Lab-05: A* Algorithm

Theory

Python is a general-purpose interpreted, interactive, object-oriented, and high-level programming language. It was created by Guido van Rossum during 1985- 1990. Like Perl, Python source code is also available under the GNU General Public License (GPL). Python is named after a TV Show called ëMonty Pythonís Flying Circusí and not after Python-the snake.

Python 3.0 was released in 2008. Although this version is supposed to be backward incompatibles, later on many of its important features have been backported to be compatible with version 2.7.

Python is a high-level, interpreted, interactive and object-oriented scripting language. Python is designed to be highly readable.

Characteristics of Python

Following are important characteristics of python –

- It supports functional and structured programming methods as well as OOP.
- It provides very high-level dynamic data types and supports dynamic type checking.
- It supports automatic garbage collection.
- It can be easily integrated with C, C++, COM, ActiveX, CORBA, and Java.

Python is one of the most widely used language over the web.

A* Search Algorithm

It is a searching algorithm that is used to find the shortest path between an initial and a final point. It is a handy algorithm that is often used for map traversal to find the shortest path to be taken. A* was initially designed as a graph traversal problem, to help build a robot that can find its own course. It still remains a widely popular algorithm for graph traversal. It searches for shorter paths first, thus making it an optimal and complete algorithm. An optimal algorithm will find the least cost outcome for a problem, while a complete algorithm finds all the possible outcomes of a problem. Another aspect that makes A* so powerful is the use of weighted graphs in its implementation. A weighted graph uses numbers to represent the cost of taking each path or course of action. This means that the algorithms can take the path with the least cost, and find the best route in terms of distance and time.

The Basic Concept of A* Algorithm: A heuristic algorithm sacrifices optimality, with precision and accuracy for speed, to solve problems faster and more efficiently. All graphs have different nodes or points which the algorithm has to take, to reach the final node. The paths between these nodes all have a numerical value, which is considered as the weight of the path. The total of all paths transverse gives you the cost of that route. Initially, the Algorithm calculates the cost to all its immediate neighboring nodes,n, and chooses the one incurring the least cost. This process repeats until no new nodes can be chosen and all paths have been traversed. Then, you should consider the best path among them. If f(n) represents the final cost, then it can be denoted as: f(n) = g(n) + h(n), where:

 $g(n) = \cos t$ of traversing from one node to another. This will vary from node to node

h(n) = heuristic approximation of the node's value. This is not a real value but an approximation cost

A* Search Algorithms Steps:

Step1: Place the starting node in the OPEN list.

Step 2: Check if the OPEN list is empty or not, if the list is empty then return failure and stops.

Step 3: Select the node from the OPEN list which has the smallest value of evaluation function (g+h), if node n is goal node then return success and stop, otherwise

Step 4: Expand node n and generate all of its successors, and put n into the closed list. For each successor n', check whether n' is already in the OPEN or CLOSED list, if not then compute evaluation function for n' and place into Open list.

Step 5: Else if node n' is already in OPEN and CLOSED, then it should be attached to the back pointer which reflects the lowest g(n') value.

Step 6: Return to Step 2.

The Pseudo-Code of the Algorithm goes like this:

let the openList equal empty list of nodes
let the closedList equal empty list of nodes
put the startNode on the openList (leave it's f at zero)
while the openList is not empty
let the currentNode equal the node with the least f value
remove the currentNode from the openList
add the currentNode to the closedList
if currentNode is the goal
You've found the end!

```
let the children of the currentNode equal the adjacent nodes
for each child in the children
if child is in the closedList
continue to beginning of for loop
child.g = currentNode.g + distance between child and current
child.h = distance from child to end
child.f = child.g + child.h
if child.position is in the openList's nodes positions
if the child.g is higher than the openList node's g
continue to beginning of for loop
add the child to the openList
```

A* Algorithm implementation in Python (Source Code)

###First Method:

```
def aStar(start, stop):
  open\_set = set(start)
  closed\_set = set()
  \mathbf{g} = \{\}
  parents = \{\}
  g[start] = 0
  parents[start] = start
  while len(open_set) > 0:
       n = None
       for v in open_set:
         if n == None \text{ or } g[v] + heuristic(v) < g[n] + heuristic(n):
            \mathbf{n} = \mathbf{v}
       if n == stop or graph_nodes[n] == None:
       else:
          for(m, weight) in get_neighbors(n):
              if m not in open_set and m not in closed_set:
                open_set.add(m)
                parents[m] = n
                g[m] = g[n] + weight
```

```
else:
             if g[m] > g[n] + weight:
               g[m] = g[n] + weight
               parents[m] = n
               if m in closed_set:
                  closed_set.remove(m)
                  open\_set.add(m)
      if n == None:
        print("Path does not exist")
        return None
      if n == stop:
        path = []
        while parents[n] != n:
           path.append(n)
           n = parents[n]
        path.append(start)
        path.reverse()
        print("Path found: {}".format(path))
        return path
      open_set.remove(n)
      closed\_set.add(n)
  print("Path does not exist")
  return None
def get_neighbors(v):
  if v in graph_nodes:
    return graph_nodes[v]
  else:
    return None
```

```
def heuristic(n):
  h_dist = {
     'A': 11,
     'B':6,
     'C':99,
     'D': 1,
     'E': 7,
     'G':0,
  return h_dist[n]
graph_nodes = {
   'A' : [('B', 2), ('E', 3)],
   'B': [('C', 1), ('G', 9)],
   'C': None,
   'E': [('D', 6)],
   'D': [('G', 1)],
 }
aStar('A', 'G')
Output:
<Path found: ['A', 'E', 'D', 'G']
###Second Method:
# This class represent a graph
class Graph:
  # Initialize the class
  def __init__(self, graph_dict=None, directed=True):
     self.graph_dict = graph_dict or {}
     self.directed = directed
    if not directed:
       self.make_undirected()
  # Create an undirected graph by adding symmetric edges
  def make_undirected(self):
     for a in list(self.graph_dict.keys()):
```

Add a link from A and B of given distance, and also add the inverse link if the

for (b, dist) in self.graph_dict[a].items():
 self.graph_dict.setdefault(b, {})[a] = dist

```
graph is undirected
  def connect(self, A, B, distance=1):
    self.graph_dict.setdefault(A, {})[B] = distance
    if not self.directed:
       self.graph_dict.setdefault(B, {})[A] = distance
  # Get neighbors or a neighbor
  def get(self, a, b=None):
    links = self.graph_dict.setdefault(a, {})
    if b is None:
       return links
    else:
       return links.get(b)
  # Return a list of nodes in the graph
  def nodes(self):
    s1 = set([k for k in self.graph_dict.keys()])
    s2 = set([k2 for v in self.graph_dict.values() for k2, v2 in v.items()])
    nodes = s1.union(s2)
    return list(nodes)
# This class represent a node
class Node:
  # Initialize the class
  def __init__(self, name:str, parent:str):
    self.name = name
    self.parent = parent
    self.g = 0 # Distance to start node
    self.h = 0 # Distance to goal node
    self.f = 0 # Total cost
  # Compare nodes
  def __eq__(self, other):
    return self.name == other.name
  # Sort nodes
  def __lt__(self, other):
     return self.f < other.f
  # Print node
  def __repr__(self):
    return ('({0},{1})'.format(self.name, self.f))
# A* search
def astar_search(graph, heuristics, start, end):
  # Create lists for open nodes and closed nodes
  open = []
  closed = []
  # Create a start node and an goal node
  start_node = Node(start, None)
```

```
goal_node = Node(end, None)
  # Add the start node
  open.append(start_node)
  # Loop until the open list is empty
  while len(open) > 0:
    # Sort the open list to get the node with the lowest cost first
    open.sort()
    # Get the node with the lowest cost
    current_node = open.pop(0)
    # Add the current node to the closed list
    closed.append(current_node)
    # Check if we have reached the goal, return the path
    if current_node == goal_node:
      path = []
      while current_node != start_node:
         path.append(current_node.name + ': ' + str(current_node.g))
         current_node = current_node.parent
      path.append(start_node.name + ': ' + str(start_node.g))
      # Return reversed path
      return path[::-1]
    # Get neighbours
    neighbors = graph.get(current_node.name)
    # Loop neighbors
    for key, value in neighbors.items():
      # Create a neighbor node
      neighbor = Node(key, current_node)
      # Check if the neighbor is in the closed list
      if(neighbor in closed):
         continue
      # Calculate full path cost
      neighbor.g = current_node.g + graph.get(current_node.name, neighbor.name)
      neighbor.h = heuristics.get(neighbor.name)
      neighbor.f = neighbor.g + neighbor.h
      # Check if neighbor is in open list and if it has a lower f value
      if(add_to_open(open, neighbor) == True):
         # Everything is green, add neighbor to open list
         open.append(neighbor)
  # Return None, no path is found
  return None
# Check if a neighbor should be added to open list
def add_to_open(open, neighbor):
  for node in open:
```

```
if (neighbor == node and neighbor.f > node.f):
       return False
  return True
# The main entry point for this module
def main():
  # Create a graph
  graph = Graph()
  # Create graph connections (Actual distance)
  graph.connect('Arad', 'Zerind', 75)
  graph.connect('Timisoara', 'Lugoj', 111)
  graph.connect('Zerind', 'Oradea', 71)
  graph.connect('Timisoara', 'Arad', 118)
  graph.connect('RimnicuVilcea', 'Sibiu', 80)
  graph.connect('Pitesti', 'RimnicuVilcea', 97)
  graph.connect('Fagaras', 'Sibiu', 99)
  graph.connect('Dobreta', 'Mehadia', 75)
  graph.connect('Bucharest', 'Fagaras', 211)
  graph.connect('Giurgiu', 'Bucharest', 90)
  graph.connect('Pitesti', 'Craiova', 70)
  graph.connect('Timisoara', 'Mehadia', 97)
  graph.connect('Zerind', 'Sibiu', 151)
  graph.connect('Fagaras', 'RimnicuVilcea', 80)
  graph.connect('Pitesti', 'Bucharest', 101)
  graph.connect('Dobreta', 'Giurgiu', 90)
  graph.connect('RimnicuVilcea', 'Pitesti', 138)
  graph.connect('Craiova', 'Bucharest',101)
  graph.connect('Sibiu', 'Craiova', 81)
  graph.connect('Mehadia', 'Craiova', 120)
  graph.connect('Fagarasu', 'Pitesti', 101)
  # Make graph undirected, create symmetric connections
  graph.make undirected()
  # Create heuristics (straight-line distance, air-travel distance)
  heuristics = {}
  heuristics['Giurgiu'] = 204
  heuristics['Pitesti'] = 101
  heuristics['Fagaras'] = 211
  heuristics['Craiova'] = 120
  heuristics['Sibiu'] = 99
  heuristics['RimnicuVilcea'] = 146
  heuristics['Mehadia'] = 138
  heuristics['Fagarasu'] = 47
  heuristics['Timisoara'] = 132
  heuristics['Passau'] = 257
  heuristics['Zerind'] = 168
```

```
heuristics['Dobreta'] = 70
heuristics['Oradea'] = 75
heuristics['Bucharest'] = 0
# Run the search algorithm
path = astar_search(graph, heuristics, 'RimnicuVilcea', 'Zerind')
print(path)
print()
# Tell python to run main method
if __name__ == "__main__": main()
```

Output:

```
['RimnicuVilcea: 0', 'Sibiu: 80', 'Zerind: 231']
```

Advantages:

- A* search algorithm is the best algorithm than other search algorithms.
- A* search algorithm is optimal and complete.
- This algorithm can solve very complex problems.

Time Complexity: The time complexity of A* search algorithm depends on heuristic function, and the number of nodes expanded is exponential to the depth of solution d. So the time complexity is O(b^d), where b is the branching factor.

Space Complexity: The space complexity of A* search algorithm is **O**(**b**^**d**)

Conclusion: A* is a very powerful algorithm with almost unlimited potential. However, it is only as good as its heuristic function, which can be highly variable considering the nature of a problem. It has found applications in many software systems, from Machine Learning and Search Optimization to game development where NPC characters navigate through complex terrain and obstacles to reach the player.

Results

After doing the lab task and lab work we learn -A* Search algorithm. Now we are able to do programs related to these topics

Discussion

Python is a general-purpose, versatile and popular programming language. It's great as a first language because it is concise and easy to read, and it is also a good language to have in any programmer's stack as it can be used for everything from web development to software development and data science applications.

This lab task is a great introduction to both fundamental programming concepts and the Python programming language. Python 3 is the most up-to-date version of the language with many improvements made to increase the efficiency and simplicity of the code that we write. Daybyday, python new version are realising and all the new versions have new features and these new features are better than previous one. So, finally I can say that python will be more user friendly in future.