

DATA STRUCTURE ALGORITHMS AND APPLICATIONS (CT- 159)

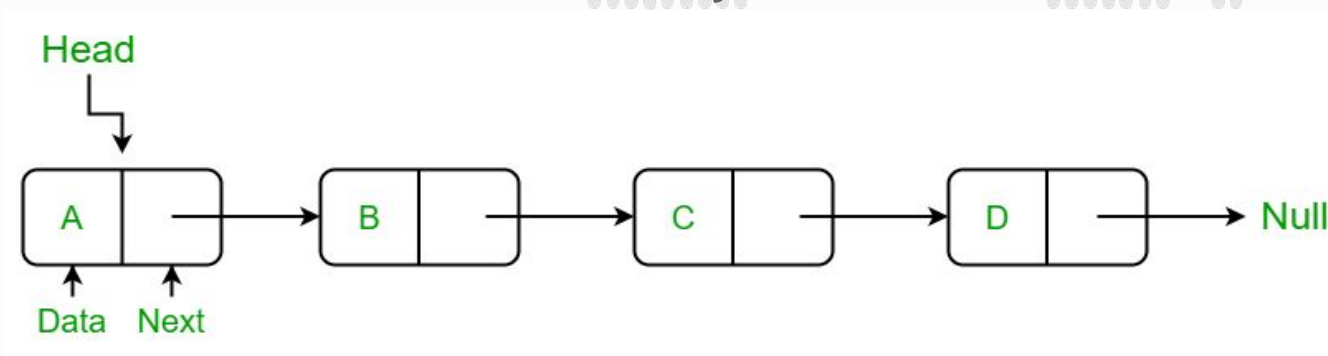
Lecture – 6, 7 Linked List, Doubly Linked List, Circular Linked List

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What is Linked List

- A **linked list** is a linear data structure where elements (called **nodes**) are stored in a sequence, but instead of using a contiguous block of memory (like an array), each node points to the next node in the sequence.
- Each node typically contains two things:
Data: The value stored in the node.
Pointer: A reference to the next node in the list.
- Unlike arrays, linked lists allow for efficient insertion and deletion of elements, as they don't require shifting elements around.
- However, searching for an element in a linked list is generally slower since you have to traverse the list node by node.



Pseudo Code for Creating a Linked List

```
struct Node {  
    int data;  
    Node* next;  
};
```

CREATING NODE

```
Node* createNode(int value) {  
    Node* newNode = new Node();  
    newNode->data = value;  
    newNode->next = nullptr; // No next  
node initially  
    return newNode;  
}
```

Basic Operations in Linked List

- **Insertion:** Adding an element to the list. Efficient when adding nodes at the beginning or end of the list.

INSERT AT BEGINNING

Function InsertAtBeginning(Head, Data)

Create new Node Node.Data = Data

Node.Prev = NULL Node.Next = Head

If Head != NULL Head.Prev = Node

End If

Head = Node

End Function

Basic Operations in Linked List

- **Insertion:** Adding an element to the list. Efficient when adding nodes at the beginning or end of the list.

INSERT IN THE MIDDLE

Function InsertInMiddle(Head, Data, Position)

Create new Node

Node.Data = Data

If Position == 1

 InsertAtBeginning(Head, Data)

Else

 Temp = Head

 Count = 1

 While Count < Position - 1 AND Temp != NULL

 Temp = Temp.Next

 Count = Count + 1

 End While

If Temp == NULL

 Print "Position out of bounds"

Else

 Node.Next = Temp.Next

 Node.Prev = Temp

 If Temp.Next != NULL

 Temp.Next.Prev = Node

 End If

 Temp.Next = Node

 End If

End If

End Function

Basic Operations in Linked List

- **Insertion:** Adding an element to the list. Efficient when adding nodes at the beginning or end of the list.

INSERT AT THE END

Function InsertAtEnd(Head, Data)

 Create new Node

 Node.Data = Data

 Node.Next = NULL

 If Head == NULL

 Node.Prev = NULL

 Head = Node

Else

 Temp = Head

 While Temp.Next != NULL

 Temp = Temp.Next

 End While

 Temp.Next = Node

 Node.Prev = Temp

 End If

End Function

Basic Operations in Linked List

- **Deletion:** Removing an element from the list. Deleting a node is fast if the position is known, but it may require traversal for arbitrary positions.

DELETE AT THE BEGINNING

Function DeleteAtBeginning(Head)

 If Head == NULL

 Print "List is empty"

 Else

 Temp = Head

 Head = Head.Next

 If Head != NULL

 Head.Prev = NULL

 End If

 Delete Temp

 End If

End Function

Basic Operations in Linked List

- **Deletion:** Removing an element from the list. Deleting a node is fast if the position is known, but it may require traversal for arbitrary positions.

DELETE AT THE END

```
Function DeleteAtEnd(Head)
```

```
    If Head == NULL
```

```
        Print "List is empty"
```

```
    Else If Head.Next == NULL
```

```
        Delete Head
```

```
        Head = NULL
```

```
    Else
```

```
        Temp = Head
```

```
        While Temp.Next != NULL
```

```
            Temp = Temp.Next
```

```
        End While
```

```
        Temp.Prev.Next = NULL
```

```
        Delete Temp
```

```
    End If
```

```
End Function
```


Basic Operations in Linked List

- **Searching:** Finding an element in the list. Searching in a linked list requires linear time ($O(n)$) as you must traverse node by node

SEARCHING A NODE/ TRAVERSING

Function Search(Head, Data)

Temp = Head

Position = 1

While Temp != NULL

 If Temp.Data == Data

 Print "Node found at position", Position

 Return

 End If

Temp = Temp.Next

Position = Position + 1

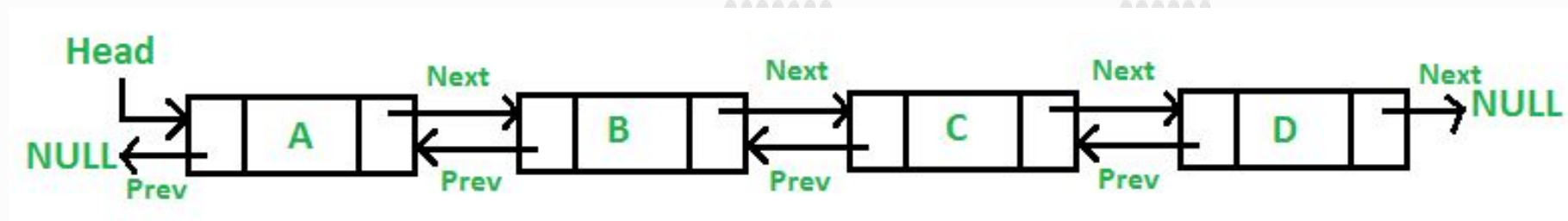
End While

Print "Node not found"

End Function

Doubly Linked List

- A doubly linked list is similar to a singly linked list, but each node contains an additional pointer that points to the previous node.
- This allows for traversal in both directions: forward and backward.
- **Advantages:**
- Easier to implement deletion of a node, as you can directly access the previous node.
- You can traverse the list in both directions.



Pseudo Code for Doubly Linked List

```
struct Node {  
    int data;  
    Node* next;  
    Node* prev;  
};
```

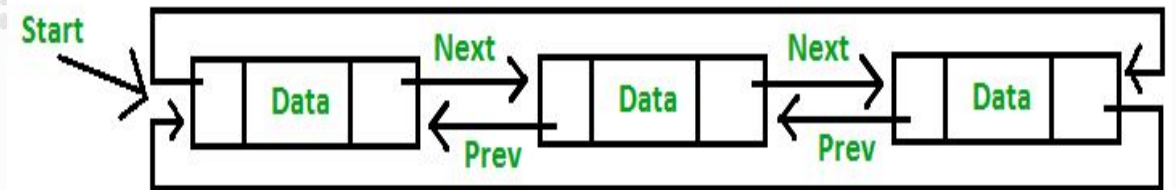
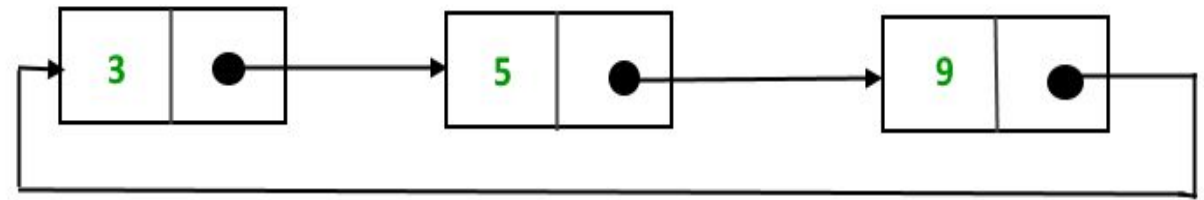
```
// Function to create a new node  
Node* createNode(int value) {  
    Node* newNode = new Node();  
    newNode->data = value;  
    newNode->next = nullptr;  
    newNode->prev = nullptr; // Points to the previous  
    node  
    return newNode;  
}
```

Circular Linked List

- In a circular linked list, the last node points back to the head, forming a circular structure.

There are two types:

- **Singly Circular Linked List:** The last node's next pointer points to the head.
- **Doubly Circular Linked List:** The last node's next pointer points to the head, and the head's previous pointer points to the last node.



Pseudo Code for Circular Linked List

```
// Node structure
struct Node {
    int data;
    Node* next;
};

// Function to create a new node
Node* createNode(int value) {
    Node* newNode = new Node();
    newNode->data = value;
    newNode->next = nullptr;
    return newNode;
}
```


Comparison of Linked Lists

- **Singly Linked List:** Simple and easy to implement. But limited traversal (only forward).
- **Doubly Linked List:** More versatile, allowing forward and backward traversal.
- Slightly more memory-intensive due to the additional pointer.
- **Circular Linked List:** Useful in applications where you need cyclic traversal (e.g., in round-robin scheduling).

Arrays VS Linked List

Feature	Linked List	Array
Size	Dynamic	Fixed (unless dynamically resized)
Memory Utilization	Efficient for variable sizes, but overhead for pointers	Efficient for fixed sizes, but wastes space if underused
Insertion/Deletion	Fast for middle and end operations	Slow due to shifting elements
Access Time	Sequential, $O(n)$	Random access, $O(1)$
Memory Allocation	Scattered memory, more flexible	Contiguous memory allocation

