

---

## 1. Process Management

**Concept:** A **Process** is a program in execution. A **Thread** is a lightweight unit of a process.

- **Deep Dive:**

- **Process Control Block (PCB):** Data structure storing process info (ID, State, Registers).
- **Context Switching:** Saving the state of the current process and loading the next one. **Cost:** Pure overhead (CPU does no useful work during this).
- **Process States:** New \$to\$ Ready \$to\$ Running \$to\$ Terminated (or Blocked/Wait).

### The `fork()` System Call (Crucial for C-CAT)

- **Explanation:** Creates a duplicate process. The child gets a copy of the parent's data/code.
- **Return Values:**
  - Returns 0 to the **Child**.
  - Returns **Child\_PID (>0)** to the **Parent**.
  - Returns <0 if it fails.

#### 💡 TRICK: Counting Processes

- **Total Processes:**  $2^n$  (where  $n$  is number of `fork()` calls).
- **Total Child Processes:**  $2^n - 1$ .
- Example:  
c `fork();` // 2 processes (1 parent, 1 child) `fork();` // 4 processes (2 parents, 2 children) `fork();` // 8 processes (4 parents, 4 children)

---

## 2. CPU Scheduling

**Concept:** Deciding which process runs on the CPU next.

Metrics:

- **Arrival Time (AT):** When process enters system.
- **Burst Time (BT):** Time required to execute.
- **Turnaround Time (TAT):** Completion Time - Arrival Time.
- **Waiting Time (WT):** TAT - Burst Time.

### Important Algorithms:

1. **FCFS (First Come First Serve):**

- **Type:** Non-preemptive.
- **Issue:** **Convoy Effect** (Short process stuck behind a long one).

2. **SJF (Shortest Job First):**
  - *Type:* Can be Preemptive (SRTF) or Non-preemptive.
  - *Advantage:* Gives **Minimum Average Waiting Time** (Optimal).
  - *Issue:* **Starvation** (Long jobs never run).
3. **Round Robin (RR):**
  - *Type:* Preemptive (uses **Time Quantum**).
  - *Best for:* Time-sharing systems (Response time).
  - *Trick:* If Time Quantum  $\rightarrow \infty$ , RR becomes FCFS.



Don't calculate mentally. Draw a horizontal bar (Gantt Chart) immediately.

- Write time on the bottom axis (\$0, 2, 5...\$).
- Cross out processes as they finish.

---

### 3. Deadlock

**Concept:** A situation where a set of processes are blocked because each is holding a resource and waiting for another resource held by someone else.

The 4 Necessary Conditions (Coffman Conditions):

(All 4 must happen for Deadlock)

1. **Mutual Exclusion:** Only one process can use a resource at a time.
2. **Hold and Wait:** A process holding at least one resource is waiting for others.
3. **No Preemption:** Resources cannot be forcibly taken away.
4. **Circular Wait:** P1 waits for P2, P2 waits for P3... Pn waits for P1.

**Handling Deadlock:**

- **Prevention:** Break one of the 4 conditions.
- **Avoidance:** **Banker's Algorithm.** (Check if system is in a "Safe State").
- **Detection:** Wait for deadlock, then kill processes.



- **Need Matrix = Max - Allocation.**
- If **Need <= Available**, the process can run.
- Once it finishes, **Available = Available + Allocation** (It returns resources).

---

### 4. Synchronization

**Concept:** Coordinating processes to avoid data inconsistency.

- **Race Condition:** When output depends on the order of execution (bad!).
- **Critical Section (CS):** Part of code where shared resources are accessed.
- **Semaphores:**
  - **Wait (P):** Decrements value. If  $\$ < 0$ , block.
  - **Signal (V):** Increments value. Wakes up a process.
  - **Binary Semaphore (Mutex):** Value is 0 or 1. Used for locks.
  - **Counting Semaphore:** Value can be  $\$ > 1$ . Used for managing multiple instances of a resource (e.g., 3 printers).

---

## 5. Memory Management

**Concept:** Managing RAM efficiently.

### Key Definitions:

- **Fragmentation:** Wasted space.
  - *Internal:* Wasted space *inside* a block (Fixed partitioning).
  - *External:* Wasted space scattered *outside* (Dynamic partitioning).
- **Paging:** Divides memory into fixed-size "Frames" and process into "Pages".
  - *Solves:* External Fragmentation.
  - *Hardware:* **MMU** (Memory Management Unit) maps Logical Address  $\$ \rightarrow \$$  Physical Address.

### Virtual Memory:

- Allows running programs **larger than RAM**.
- **Page Fault:** CPU asks for a page  $\$ \rightarrow \$$  Not in RAM  $\$ \rightarrow \$$  OS fetches from Disk.
- **Thrashing:** High paging activity (CPU is busy swapping pages, not working). *Cause:* Degree of Multi-programming is too high (too many processes).

#### 💡 TRICK: Address Calculation

- **Logical Address** = Page Number + Offset.
- **Physical Address** = Frame Number + Offset.
- *Formula:* Size of Page Table = (Total Logical Address Space) / (Page Size).

---

## 6. Important Linux Commands & Permissions

### File Permissions (`chmod`):

- **Read (r) = 4, Write (w) = 2, Execute (x) = 1.**
- **User Types:** Owner (u), Group (g), Others (o).
- **Example:** `chmod 751 file`

- **7** (Owner) = \$4+2+1\$ (rwx).
- **5** (Group) = \$4+0+1\$ (r-x).
- **1** (Others) = \$0+0+1\$ (–x).

### Must-Know Commands:

- `grep pattern file`: Search for text.
  - `ls -l`: List files with details (permissions, size).
  - `ps`: Display current processes.
  - `kill PID`: Terminate a process.
  - `top`: Real-time view of system (CPU/RAM usage).
- 

### Quick Recap for Exam Day

1. **Starvation** is solved by **Aging**.
2. **Threads** share Heap/Data/Code but have **own Stack/Registers**.
3. **Banker's Algorithm** is for **Avoidance**.
4. **SJF** is optimal for minimum waiting time.
5. **Kernel Mode** allows execution of privileged instructions; **User Mode** does not.

Would you like me to create a **visual comparison table** for "Process vs. Thread" or "Paging vs. Segmentation" to help you memorize the differences?

### Final "Exam Day" Checklist for OS

If you can answer these 3 questions in your head right now, you are done with OS.

1. **Math:** Can you calculate the number of processes for `fork(); fork(); fork();?` (*Ans: 8*)
2. **Logic:** Do you know that **SJF** is the best for average waiting time but causes starvation? (*Ans: Yes*)
3. **Command:** Do you know `chmod 777` means everyone has full permission? (*Ans: Yes*)