

# Tutorial 1

1. Determine whether each of the following sentences is a statement.
- i)  $y + 3$  is a positive integer. *Not a statement*
  - ii)  $128 = 2^6$ . *statement*
  - iii)  $x = 2^6$ . *Not a statement*
  - iv) Is 2 a positive number? *Not a statement*

2. Give the negation of the following statement.

i)  $2 + 7 \leq 11$        $2 + 7 > 11$

ii)  $2 + 1 = 3$        $2 + 1 \neq 3$

iii)  $p: (2 \text{ is an even integer}) \text{ and } (8 \text{ is an odd integer}) q$

iv) Today is Wednesday.

and :  $\wedge$

or :  $\vee$

iii)  $p \wedge q$   
 $\sim (p \wedge q) \equiv \sim p \vee \sim q$   
2 is not an even integer or 8 is not an odd integer.

iv) Today is not Wednesday.

$\vee$  : or  
 $\wedge$  : and  
 $\sim p$  : not  $p$

3. Let  $p, q, r$  be the propositions

$p$  : You have a flu.

$q$  : You miss the final exam.

$r$  : You pass the course.

Express each of these propositions as in English sentences.

i)  $p \vee q \vee r$     ii)  $(p \wedge q) \vee (\sim q \wedge r)$     iii)  $\sim p \wedge \sim q \wedge r$

i) You have a flu or you miss the final exam or you pass the course.

ii) Either you have a flu and miss the final exam or you didn't miss the final exam and pass the course.

iii) You did not have a flu and you did not miss the final exam and you pass the course.

but  $\equiv$  and

4. Let  $h$  : "John is healthy."  
 $w$  : "John is wealthy."  
 $s$  : "John is wise."

$\wedge, \vee, \dots$

Use the indicated letters and logical connectors to represent the following compound statements.

- i) John is healthy and wealthy.  $h \wedge w$
- ii) John is healthy and not wise.  $h \wedge \sim s$
- iii) John is healthy and wealthy but not wise.  $h \wedge w \wedge \sim s$
- iv) John is not wealthy but he is healthy and wise.  $\sim w \wedge (h \wedge s)$
- v) John is either wealthy or healthy, or both.  $w \vee h$
- vi) John is wealthy or he is healthy but not wealthy and healthy.  $w \vee h$
- vii) John is neither healthy nor wealthy.  $\sim (h \wedge w) \text{ or } h \wedge \sim w$
- viii) John is neither healthy, wealthy, nor wise.  $\sim (h \wedge w \wedge s)$

5. Determine the truth or falsity of each of the following statement.

T: True

F: False

i)  $(2 \geq 3)$  and  $(3 \text{ is a positive integer.})$  False

ii)  $(2 < 3)$  or  $(3 \text{ is not a positive integer.})$  True

iii)  $(2 \text{ is a prime})$  but  $(3 \text{ is not a prime.})$  False

iv) It is not true that  $(2 \text{ is not a prime})$  or  $(3 \text{ is prime.})$  False

v) It is false that  $(2 \text{ is prime})$  or  $(\text{multiple of } 4.)$  False

$$(i) \quad F \wedge T \equiv F \quad (T \wedge T \equiv T)$$

$$(ii) \quad T \vee F \equiv T \quad (F \vee F \equiv F)$$

$$(iii) \quad T \wedge F \equiv F$$

$$(iv) \quad \sim (F \vee T) \equiv \sim T \equiv F$$

$$(v) \quad \sim (T \vee F) \equiv \sim T \equiv F$$

prime:  $\checkmark 2, 3, 5, 7, 11, 13, \dots$

multiple of 4:  $\underline{4}, 8, 12, 16, 20, \dots$

$$p: T, q: F, r: T$$

6. Find the truth value of each proposition if  $p$  and  $r$  are true and  $q$  is false.

- |  |       |   |       |
|--|-------|---|-------|
| i) $\sim p \wedge (q \vee r)$              | False | ii) $\sim p \wedge (\sim(q \vee \sim r))$ | True  |
| iii) $(r \wedge \sim q) \vee (p \wedge r)$ | True  | iv) $(q \wedge r) \wedge (p \vee \sim r)$ | False |

$$i) F \wedge (F \vee T) \equiv F \wedge T \equiv F$$

$$ii) T \wedge (\sim(F \vee F)) \equiv T \wedge (\sim F) \equiv T \wedge T \equiv T$$

$$iii) (T \wedge T) \vee (T \wedge T) \equiv T \vee T \equiv T$$

$$iv) (F \wedge T) \wedge (T \vee F) \equiv F \wedge T \equiv F$$

7. Construct a truth table for the following compound statements.

i)  $(p \underline{\vee} q) \underline{\vee} (q \underline{\vee} r)$

✓ ✓ ✓ ✓

$\underline{\vee}$  : not both

$p$	$q$	$r$	$p \underline{\vee} q$	$q \underline{\vee} r$	$(p \underline{\vee} q) \underline{\vee} (q \underline{\vee} r)$
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	1	1	0
0	1	1	1	0	1
1	0	0	1	0	1
1	0	1	1	1	0
1	1	0	0	1	1
1	1	1	0	0	0



ii)  $(p \downarrow q) \wedge (q \downarrow r)$

$$p \downarrow q \equiv \neg(p \vee q)$$

$$0, 0 \Rightarrow 1$$

$\checkmark$ p	$\checkmark$ q	$\checkmark$ r	$\checkmark$ $p \downarrow q$	$\checkmark$ $q \downarrow r$	$(p \downarrow q) \wedge (q \downarrow r)$
0	0	0	1	1	1
0	0	1	1	0	0
0	1	0	0	0	0
0	1	1	0	0	0
1	0	0	0	1	0
1	0	1	0	0	0
1	1	0	0	0	0
1	1	1	0	0	0



iii)  $(p \mid q) \wedge r$

$p \mid q \equiv \sim (p \wedge q)$

$1, 1 \Rightarrow 0$

$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	
$p$	$q$	$r$	$p \mid q$	$(p \mid q) \wedge r$
0	0	0	1	0
0	0	1	1	1
0	1	0	1	0
0	1	1	1	1
1	0	0	1	0
1	0	1	1	1
1	1	0	0	0
1	1	1	0	0

iv)  $(p \mid q) \vee (p \mid r)$

$p \mid q \equiv \sim (p \wedge q)$

$1, 1 \Rightarrow 0$

$\checkmark$ $\checkmark$ p	$\checkmark$ q	$\checkmark$ r	$p \mid q$	$p \mid r$	$(p \mid q) \vee (p \mid r)$
0	0	0	1	1	1
0	0	1	1	1	1
0	1	0	1	1	1
0	1	1	1	1	1
1	0	0	1	1	1
1	0	1	1	0	1
1	1	0	0	1	1
1	1	1	0	0	0

8. Determine which of the pairs of statement forms are logically equivalent. Justify your answers using truth tables.

i)  $\sim(p \wedge (\sim(p \wedge q)))$  and  $p \vee (\sim q \wedge \sim p)$

$p$	$q$	$\sim p$	$\sim q$	$p \wedge q$	$\sim(p \wedge q)$	$p \wedge \sim(p \wedge q)$	$\sim(p \wedge (\sim(p \wedge q)))$ ✓	$\sim q \wedge \sim p$	$p \vee (\sim q \wedge \sim p)$ ✓
0	0	1	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1	0	0
1	0	0	1	0	1	1	0	0	1
1	1	0	0	1	0	0	1	0	1

$\uparrow$  not equal  $\uparrow$

$\therefore$  not logically equivalent

ii)

$$\overbrace{(p \downarrow p) \downarrow (q \downarrow q)}^{\text{LHS}} \text{ and } \overbrace{p \wedge q}^{\text{RHS}}$$

$$p \downarrow q \equiv \neg(p \vee q)$$

$$0, 0 : 1$$


p	q	$p \downarrow p$	$q \downarrow q$	$(p \downarrow p) \downarrow (q \downarrow q)$	$p \wedge q$
0	0	1	1	0	0
0	1	1	0	0	0
1	0	0	1	0	0
1	1	0	0	1	1

↑  
↑  
equal

∴ logically equivalent

iii)  $(p \underline{\vee} q) \wedge r$  and  $(p \wedge r) \underline{\vee} (q \wedge r)$

$p$	$q$	$r$	$p \vee q$	$(p \vee q) \wedge r$	$p \wedge r$	$q \wedge r$	$(p \wedge r) \vee (q \wedge r)$
0	0	0	0	0	0	0	0
0	0	1	0	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	1	0	1	1
1	0	0	1	0	0	0	0
1	0	1	1	1	1	0	1
1	1	0	0	0	0	0	0
1	1	1	0	0	1	1	0


  
 equal

$\therefore$  logically equivalent

9. Use <sup>0</sup>truth table to determine each of the statement forms below is a tautology, contradiction or contingency.

i)  $((\sim p \wedge q) \wedge (q \wedge r)) \wedge \sim q$

$p$	$q$	$r$	$\sim p$	$\sim q$	$\sim p \wedge q$	$q \wedge r$	$(\sim p \wedge q) \wedge (q \wedge r)$	$((\sim p \wedge q) \wedge (q \wedge r)) \wedge \sim q$
0	0	0	1	1	0	0	0	0
0	0	1	1	1	0	0	0	0
0	1	0	1	0	1	0	0	0
0	1	1	1	0	1	1	1	0
1	0	0	0	1	0	0	0	0
1	0	1	0	1	0	0	0	0
1	1	0	0	0	0	0	0	0
1	1	1	0	0	0	1	0	0

$\therefore$  contradiction

ii)  $(\sim p \vee q) \vee (p \wedge \sim q)$

$p$	$q$	$\sim p$	$\sim q$	$\sim p \vee q$	$p \wedge \sim q$	$(\sim p \vee q) \vee (p \wedge \sim q)$
0	0	1	1	1	0	1
0	1	1	0	1	0	1
1	0	0	1	0	1	1
1	1	0	0	1	0	1



$\therefore$  tautology



10. Simplify the following statements using law of logical equivalences.

i)  $(p \vee q) \wedge \sim(\sim p \wedge q)$

ii)  $((\underline{p} \vee q) \wedge (\underline{p} \vee \sim q)) \vee q$

$$\begin{aligned} & (\underline{2 \times 3}) + (\underline{2 \times 4}) \\ & = 2(3+4) \end{aligned}$$

$$i) (\underline{p \vee q}) \wedge (\underline{p \vee \sim q}) \equiv p \vee (q \wedge \sim q) \equiv p \vee c \equiv p \quad \#$$

$$ii) (p \vee (\overset{=c}{q \wedge \sim q})) \vee q \equiv p \vee q \quad \#$$

4) Identity laws:

■  $p \wedge t \equiv p$

■  $p \vee c \equiv p$

5) Negation/Inverse laws:

■  $p \vee \sim p = t$

■  $p \wedge \sim p = c$

11. Use the law of logical equivalences to verify the following logical equivalences.

LHS:  $(p \vee q) \vee (\sim p \wedge \sim q \wedge r) = p \vee q \vee r$  RHS

ii)  $\sim(p \downarrow q) = (\sim p \mid \sim q)$

$$\begin{aligned} \text{i) LHS: } & ((p \vee q) \vee \sim p) \wedge ((p \vee q) \vee \sim q) \wedge ((p \vee q) \vee r) \\ & \equiv ((p \vee \overset{=t}{\sim p}) \vee q) \wedge (p \vee (q \vee \overset{=t}{\sim q})) \wedge (p \vee q \vee r) \\ & \equiv (t \vee q) \wedge (p \vee t) \wedge (p \vee q \vee r) \\ & \equiv t \wedge t \wedge (p \vee q \vee r) \equiv p \vee q \vee r \quad \# \text{ (RHS)} \end{aligned}$$

$$\begin{aligned} \text{ii) LHS: } & \sim(p \downarrow q) \equiv \sim(\sim(p \vee q)) \\ & \equiv p \vee q \quad \checkmark \end{aligned}$$

$$\text{RHS: } \sim p \mid \sim q \equiv \sim(\sim p \wedge \sim q) \equiv p \vee q \quad \checkmark \# \text{ (LHS)}$$

9) Universal Bound / Domination laws:

■  $p \vee t \equiv t$

■  $p \wedge c \equiv c$