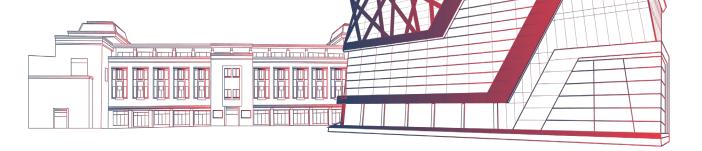




# **UNIT III**

**Relational Database Design** 









# **Equivalence of FD Sets**





## **FD Membership Test**





Let F be the FD set with FD X → Y

• FD X  $\rightarrow$  Y is implied in F iff

X+ determines Y in FD set F

i.e. FD X  $\rightarrow$  Y is implied in F iff X+ = {.....Y.....}





### **FD Membership Test Example**



Relation R (A, B, C, D)

Is AB 
$$\rightarrow$$
 D a member of F = { AB  $\rightarrow$  C, BC  $\rightarrow$  D}?

 $AB+ = ABCD : AB \rightarrow D$  is a member of F.

Relation R (A, B, C, D)

Is 
$$C \rightarrow AB$$
 a member of  $F = \{AB \rightarrow C, C \rightarrow D, D \rightarrow A\}$ ?

 $C+ = \{ C, D, A \} :: C \rightarrow A \text{ is implied but } C \rightarrow B \text{ is not implied}$ 

 $:: C \rightarrow AB$  is not a member of F





### **Equivalence of FD Sets**



- Two sets of FDs F and G are equivalent if:
  - every FD in F is implied in G, and
  - every FD in G is implied in F
- **Definition:** 
  - a) F covers G ( $F \supseteq G$ ) if every FD in G is implied in F.
  - b) Similarly, G covers F ( $G \supseteq F$ ) if every FD in F in implied in G.
- F and G are equivalent if F covers G and G covers F





## **Equivalence of FD Sets Example**



Consider the following two sets of FDs

$$F = \{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$$
  
 $G = \{A \rightarrow CD, E \rightarrow AH\}$ 

Check whether or not they are equivalent.

Solution

#### **Proof that G is covered by F:**

 $\{A\} + = \{A, C, D\}$  (with respect to F), which covers  $A \rightarrow CD$  in G

 $\{E\}$  + =  $\{E, A, D, H, C\}$  (with respect to F), which covers  $E \rightarrow AH$  in G

#### **Proof that F is covered by G:**

 $\{A\} + = \{A, C, D\}$  (with respect to G), which covers  $A \rightarrow C$  in F

 $\{A, C\} + = \{A, C, D\}$  (with respect to G), which covers  $AC \rightarrow D$  in F

 $\{E\}$  + =  $\{E, A, H, C, D\}$  (with respect to G), which covers  $E \rightarrow AD$  and  $E \rightarrow H$  in F





## **Equivalence of FD Sets Drill**



1. A relation R (A , C , D , E , H) is having two functional dependencies sets F and G as shown-

$$F = \{ A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H \}$$
  
 $G = \{ A \rightarrow CD, E \rightarrow H \}$ 

Check if F and G are equivalent?

2. Are these two FD sets equivalent or not?

FD1 = {AB 
$$\rightarrow$$
 C, D  $\rightarrow$  E, E  $\rightarrow$  C}  
FD2 = {AB  $\rightarrow$  C, D  $\rightarrow$  E, AB  $\rightarrow$  E, E  $\rightarrow$  C}







## Solution Equivalence of FD Sets Drill



#### 1. Determining whether F covers G-

$$(A)^{+} = \{A, C, D\}$$
 // using set F  
 $(E)^{+} = \{A, C, D, E, H\}$  // using set F

Thus, we conclude F covers G i.e.  $F \supseteq G$ 

#### **Determining whether G covers F-**

Thus, we conclude G does not cover F









## Solution Equivalence of FD Sets Drill



2. FD2 covers FD 1

FD1 does not cover FD 2 as AB  $\rightarrow$  E is not implied in FD 1.

Thus, FD1 ≠ FD2







### **Practice Drill**



**Ques:** Relation R has eight attributes ABCDEFGH. Fields of R contain only atomic values.

$$F = \{ CH \rightarrow G, A \rightarrow BC, B \rightarrow CFH, E \rightarrow A, F \rightarrow EG \}$$

is a set of functional dependencies (FDs)

How many candidate keys does the relation R have?



### **Solution**



A+ = ABCEFGH which is all attributes except D.

B+ = ABCEFGH which is all attributes except D.

E+ = ABCEFGH which is all attributes except D.

F+ = ABCEFGH which is all attributes except D.

So, there are total 4 candidate keys AD, BD, ED and FD





### **Practice Drill**



Consider a relation with schema R(A,B,C,D)

FDs are  $\{AB \rightarrow C, C \rightarrow D, D \rightarrow A\}$ .

- a. What are some of the nontrivial FDs that can be inferred from the given FDs?
- b. What are all candidate keys of R?



### **Solution**



#### a) Some of the non-trivial FDs:

 $C \rightarrow ACD$   $D \rightarrow AD$   $AB \rightarrow ABCD$ 

 $AC \rightarrow ACD$   $BC \rightarrow ABCD$   $BD \rightarrow ABCD$ 

 $CD \rightarrow ACD$   $ABC \rightarrow ABCD$   $ABD \rightarrow ABCD$ 

BCD → ABCD

#### b) Attribute Closure:

 $A \rightarrow A$   $B \rightarrow B$   $C \rightarrow ACD$   $D \rightarrow AD$ 

 $AB \rightarrow ABCD$   $AC \rightarrow ACD$   $AD \rightarrow AD$   $BC \rightarrow ABCD$ 

 $BD \rightarrow ABCD$   $CD \rightarrow ACD$   $ABC \rightarrow ABCD$ 

 $ABD \rightarrow ABCD$   $ACD \rightarrow ACD$   $BCD \rightarrow ABCD$ 

Thus, candidate keys are: AB, BC, and BD







# Thanks!!



