

Figure 5-15. 2-Mbyte PDE—PAE Paging Legacy-Mode

5.3 Long-Mode Page Translation

Long-mode page translation requires the use of physical-address extensions (PAE). Before activating long mode, PAE must be enabled by setting CR4.PAE to 1. Activating long mode before enabling PAE causes a general-protection exception (#GP) to occur.

The PAE-paging data structures support mapping of 64-bit virtual addresses into 52-bit physical addresses. PAE expands the size of legacy page-directory entries (PDEs) and page-table entries (PTEs) from 32 bits to 64 bits, allowing physical-address sizes of greater than 32 bits.

The AMD64 architecture enhances the page-directory-pointer entry (PDPE) by defining previously reserved bits for access and protection control. A new translation table is added to PAE paging, called the page-map level-4 (PML4). The PML4 table precedes the PDP table in the page-translation hierarchy.

Because PAE is always enabled in long mode, the PS bit in the page directory entry (PDE.PS) selects between 4-Kbyte and 2-Mbyte page sizes, and the CR4.PSE bit is ignored. When 1-Gbyte pages are supported, the PDPE. PS bit selects the 1-Gbyte page size.

5.3.1 Canonical Address Form

The AMD64 architecture requires implementations supporting fewer than the full 64-bit virtual address to ensure that those addresses are in canonical form. An address is in canonical form if the address bits from the most-significant implemented bit up to bit 63 are all ones or all zeros. If the addresses of all bytes in a virtual-memory reference are not in canonical form, the processor generates a general-protection exception (#GP) or a stack fault (#SS) as appropriate.

5.3.2 CR3

In long mode, the CR3 register is used to point to the PML4 base address. CR3 is expanded to 64 bits in long mode, allowing the PML4 table to be located anywhere in the 52-bit physical-address space. Figure 5-16 on page 131 shows the long-mode CR3 format.

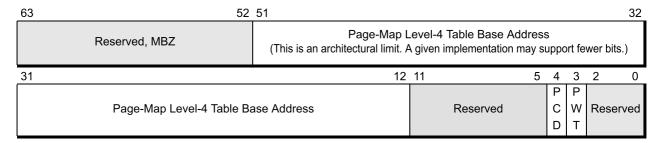


Figure 5-16. Control Register 3 (CR3)—Long Mode

The CR3 register fields for long mode are:

Table Base Address Field. Bits 51:12. This 40-bit field points to the PML4 base address. The PML4 table is aligned on a 4-Kbyte boundary with the low-order 12 address bits (11:0) assumed to be 0. This yields a total base-address size of 52 bits. System software running on processor implementations supporting less than the full 52-bit physical-address space must clear the unimplemented upper base-address bits to 0.

Page-Level Writethrough (PWT) Bit. Bit 3. Page-level writethrough indicates whether the highest-level page-translation table has a writeback or writethrough caching policy. When PWT=0, the table has a writeback caching policy. When PWT=1, the table has a writethrough caching policy.

Page-Level Cache Disable (PCD) Bit. Bit 4. Page-level cache disable indicates whether the highest-level page-translation table is cacheable. When PCD=0, the table is cacheable. When PCD=1, the table is not cacheable.

Reserved Bits. Reserved fields should be cleared to 0 by software when writing CR3.

5.3.3 4-Kbyte Page Translation

In long mode, 4-Kbyte physical-page translation is performed by dividing the virtual address into six fields. Four of the fields are used as indices into the level page-translation hierarchy. The virtual-address fields are described as follows, and are shown in Figure 5-17 on page 132:

- Bits 63:48 are a sign extension of bit 47, as required for canonical-address forms.
- Bits 47:39 index into the 512-entry page-map level-4 table.
- Bits 38:30 index into the 512-entry page-directory pointer table.
- Bits 29:21 index into the 512-entry page-directory table.
- Bits 20:12 index into the 512-entry page table.
- Bits 11:0 provide the byte offset into the physical page.

Note: The sizes of the sign extension and the PML4 fields depend on the number of virtual address bits supported by the implementation.

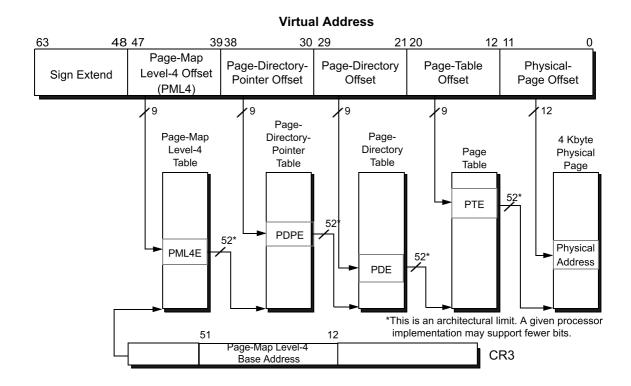


Figure 5-17. 4-Kbyte Page Translation—Long Mode

Figures 5-18 through 5-20 on page 133 and Figure 5-21 on page 133 show the long-mode 4-Kbyte translation-table formats:

- Figure 5-18 on page 133 shows the PML4E (page-map level-4 entry) format.
- Figure 5-19 on page 133 shows the PDPE (page-directory-pointer entry) format.
- Figure 5-20 on page 133 shows the PDE (page-directory entry) format.
- Figure 5-21 on page 133 shows the PTE (page-table entry) format.

The fields within these table entries are described in "Page-Translation-Table Entry Fields" on page 137.

Figure 5-20 on page 133 shows the PDE.PS bit (bit 7) cleared to 0, indicating a 4-Kbyte physical-page translation.

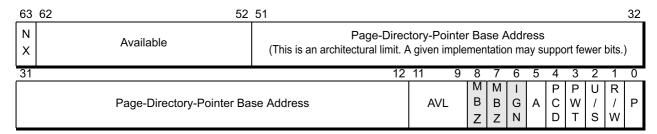


Figure 5-18. 4-Kbyte PML4E—Long Mode

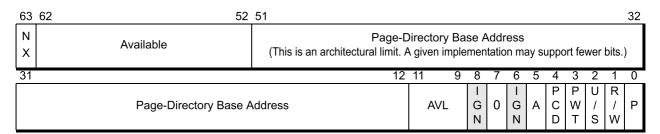


Figure 5-19. 4-Kbyte PDPE—Long Mode

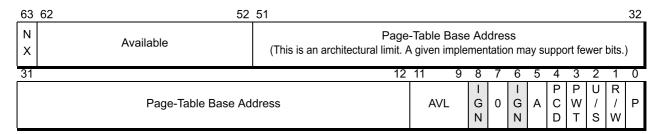


Figure 5-20. 4-Kbyte PDE—Long Mode

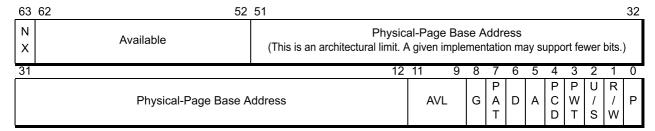


Figure 5-21. 4-Kbyte PTE—Long Mode

5.3.4 2-Mbyte Page Translation

In long mode, 2-Mbyte physical-page translation is performed by dividing the virtual address into five fields. Three of the fields are used as indices into the level page-translation hierarchy. The virtual-address fields are described as follows, and are shown in Figure 5-22:

- Bits 63:48 are a sign extension of bit 47 as required for canonical address forms.
- Bits 47:39 index into the 512-entry page-map level-4 table.
- Bits 38:30 index into the 512-entry page-directory-pointer table.
- Bits 29:21 index into the 512-entry page-directory table.
- Bits 20:0 provide the byte offset into the physical page.

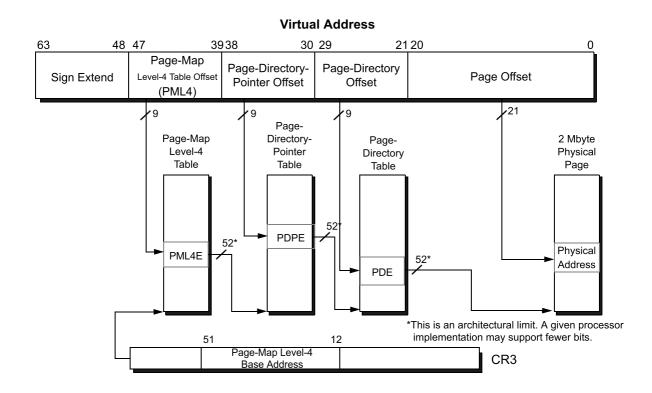


Figure 5-22. 2-Mbyte Page Translation—Long Mode

Figures 5-23 through 5-25 on page 135 show the long-mode 2-Mbyte translation-table formats (the PML4 and PDPT formats are identical to those used for 4-Kbyte page translations and are repeated here for clarity):

- Figure 5-23 on page 135 shows the PML4E (page-map level-4 entry) format.
- Figure 5-24 on page 135 shows the PDPE (page-directory-pointer entry) format.
- Figure 5-25 on page 135 shows the PDE (page-directory entry) format.

The fields within these table entries are described in "Page-Translation-Table Entry Fields" on page 137. PTEs are not used in 2-Mbyte page translations.

Figure 5-25 shows the PDE.PS bit (bit 7) set to 1, indicating a 2-Mbyte physical-page translation.

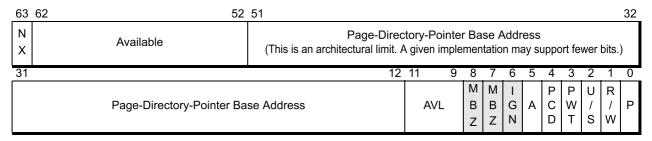


Figure 5-23. 2-Mbyte PML4E—Long Mode

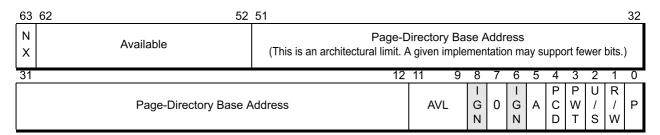


Figure 5-24. 2-Mbyte PDPE—Long Mode

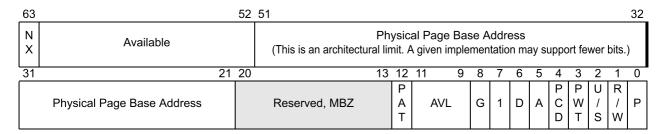


Figure 5-25. 2-Mbyte PDE—Long Mode

5.3.5 1-Gbyte Page Translation

In long mode, 1-Gbyte physical-page translation is performed by dividing the virtual address into four fields. Two of the fields are used as indices into the level page-translation hierarchy. The virtual-address fields are described as follows, and are shown in Figure 5-26 on page 136:

- Bits 63:48 are a sign extension of bit 47 as required for canonical address forms.
- Bits 47:39 index into the 512-entry page-map level-4 table.
- Bits 38:30 index into the 512-entry page-directory-pointer table.
- Bits 29:0 provide the byte offset into the physical page.