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Spec: (Same as 4.1)

- OS: AWS VM Ubuntu 18.04 (x2.2large)
- Kernal: Linux ip-172-31-37-178 5.4.0-1039-aws #41~18.04.1-Ubuntu SMP Fri Feb 26 11:20:14
 UTC 2021 x86_64 x86_64 x86_64 GNU/Linux
- Compiler: gcc version 7.5.0 (Ubuntu 7.5.0-3ubuntu1~18.04)
- Flag: -Ofast
- Details:

```
inux ip-172-31-37-178 5.4.0-1039-aws #41~18.04.1-Ubuntu SMP Fri Feb 26 11:20:14 UTC 2021 x86_64 x86_64 x86_64 GNU/Linux.
ubuntu@ip-172-31-37-178:~$ cat /proc/cpuinfo
processor
                 : 0
vendor_id
                 : GenuineIntel
cpu family
                 : 6
model
                 : 79
model name
                 : Intel(R) Xeon(R) CPU E5-2686 v4 @ 2.30GHz
stepping
                 : 0xb000038
microcode
                 : 2299.914
cpu MHz
cache size
                 : 46080 KB
physical id
                 : 0
siblings
                 : 8
                 : 0
core id
                 : 8
cpu cores
apicid
                 : 0
initial apicid
                : 0
fpu
                   yes
fpu_exception
                 : yes
cpuid level
                   13
                 : fpu vme de pse tsc msr pae mce cx8 apic sep mtrr pge mca cmov pat pse36 clfli
flaas
 x2apic movbe popcnt tsc_deadline_timer aes xsave avx f16c rdrand hypervisor lahf_lm abm cpuid
                 : cpu_meltdown spectre_v1 spectre_v2 spec_store_bypass l1tf mds swapgs itlb_mul
bugs
                 : 4600.10
bogomips
clflush size
                 : 64
cache_alignment : 64
address sizes
                 : 46 bits physical, 48 bits virtual
power management:
```

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Jubantualp-17-31-37-18-5 gcc ~V
Using built-in specs.

COLLECT_LTD_MRAPPER-/usr/lib/gcc/x86_64-linux-gnu/7/lto-wrapper

COLLECT_LTD_MRAPPER-/usr/lib/gcc/x86_64-linux-gnu/7/lto-wrapper

OFFLOWD_TAKET_MMES-mystx-nene

OFFLOWD_TAKET_MMSTAKET_MMSTAKET_MMSTAKET_NENE

OFFLOWD_TAKET_MMSTAKET_MMSTAKET_MMSTAKET_NENE

OFFLOWD_TAKET_MMSTAKET_MMSTAKET_MMSTAKET_NENE

OFFLOWD_TAKET_MMSTAKET_MMSTAKET_MMSTAKET_NENE

OFFLOWD_TAKET_MMSTAKET_MMSTAKET_MMSTAKET_NENE

OFFLOWD_TAKET_MMSTAKET_M

[ubuntu@ip-172-31-37-178:~\$ lscpu

Architecture: x86_64

CPU op-mode(s): 32-bit, 64-bit Byte Order: Little Endian

CPU(s): 8
On-line CPU(s) list: 0-7
Thread(s) per core: 1
Core(s) per socket: 8
Socket(s): 1
NUMA node(s): 1

Vendor ID: GenuineIntel

CPU family: 6
Model: 79

Model name: Intel(R) Xeon(R) CPU E5-2686 v4 @ 2.30GHz

Stepping: 1

CPU MHz: 2300.067 4600.10 BogoMIPS: Hypervisor vendor: Xen Virtualization type: full 32K L1d cache: L1i cache: 32K L2 cache: 256K L3 cache: 46080K NUMA node0 CPU(s): 0-7

4.2. Scheduling Policies (15 points)

The goal of this exercise is to illustrate the impact of the scheduling algorithm in performance. The code mandelbrot_omp.c generates an ASCII Portable Pixel Map (PPM) image of the Mandelbrot set using OpenMP.

Run the parallel code and take the execution time for a growing number of threads. Record the timing in a table. Try with the three scheduling algorithms. To change the schedule, you can either change the environment variable with export OMP_SCHEDULE=type where type can be any of static, dynamic, or guided. Alternatively, you can change the schedule in the source code as omp parallel for schedule (type).

Note: if you change the environment, then you don't have to change anything inside the code.

Submission

- P42.pdf: with replicability information (see Submission note above), Explain your results, what is the scheduling that provides the best performance? What is the reason for that?
- P42.c: Complete source code

(4.2). Table for improvements in elapsed execution time

² http://www.cs.cornell.edu/~bindel/class/cs5220-f11/hw2.html

(unit: seconds)

OMP_NUM_THREADS	2 cores	3 cores	4 cores	5 cores	6 cores	7 cores	8 cores
type = static	0.119479	0.213677	0.11835	0.151437	0.106952	0.117311	0.0903508
type = dynamic	0.118285	0.0788323	0.0592294	0.0474653	0.03963	0.0340479	0.0299404
type = guided	0.11817	0.153298	0.11719	0.0840822	0.0789541	0.0672873	0.0567206

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Discussion:

- The type=dynamics provides the best performance. We could observe the performance of dynamic > guided > static
 - It makes sense as static assigns each thread into a fixed-size chunk, and it is the default schedule version.
 - However, if using dynamic, work is assigned as a thread requests it.
 - For guided type, big chunks are allocated first and smaller and smaller chunks later.
 - For this task, we need dynamic to allocate tasks so that all threads could be parallelly and fully utilized, acheiving fastest execution time.
- Besides, we could also see even number of cores perform much better than odd number of
 cores, relatively. It may be due to task could be distributed evenly for all cores if it is an even
 number. And if it's an odd number, there are always core that is not fully utilized and not
 saturated, thus resulting in slower performance.
- We could also observe, as number of cores increases, the time it takes become less and less for odd number and even number of cores, respectively. This is as expected, as explained in 4.1, as more cores allow more concurrency when shared same memory, that could simultaneously execute statements in the parallel region.