

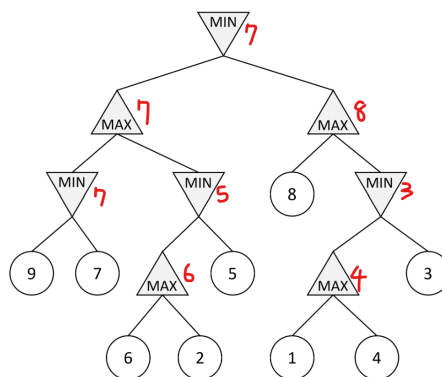
Foundations of Artificial Intelligence: Homework 2

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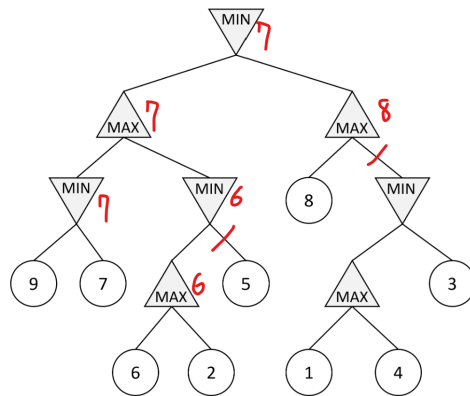
April 3, 2025

Problem 1

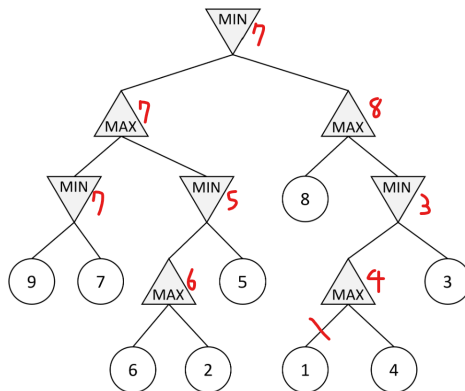
a)



b)



c)



Problem 2

a)

" 當官要掙錢，必須跪著"

$$(O \wedge K) \implies E$$

" 拿槍想掙錢，必須在山裡"

$$(G \wedge M) \implies E$$

" 我是想站著，還把錢掙了"

$$\neg K \wedge E$$

b)

For the left-hand side and right-hand side:

$$\begin{aligned}
 & ((O \wedge K) \implies E) \vee ((G \wedge M) \implies E) & (O \wedge G \wedge \neg K \wedge \neg M) \implies E \\
 & = (E \vee \neg(O \wedge K)) \vee (E \vee \neg(G \wedge M)) & = E \vee \neg(O \wedge G \wedge \neg K \wedge \neg M) \\
 & = E \vee \neg(O \wedge K) \vee \neg(G \wedge M) & = E \vee (\neg O \vee G \vee K \vee M) \\
 & = E \vee (\neg O \vee \neg K) \vee (\neg G \vee \neg M) & = E \vee \neg O \vee G \vee K \vee M \\
 & = E \vee \neg O \vee \neg K \vee \neg G \vee \neg M
 \end{aligned}$$

Then we have the final CNF:

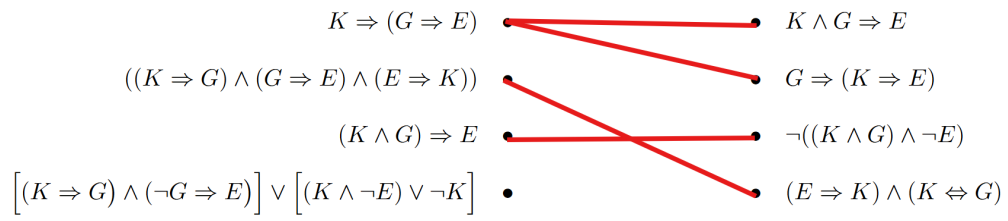
$$(E \vee \neg O \vee \neg K \vee \neg G \vee \neg M) \implies (E \vee \neg O \vee K \vee G \vee M)$$

Since there exist exactly one assignment let the sentence become false:

$$E = \text{false}, O = \text{true}, K = \text{false}, G = \text{false}, M = \text{false}$$

While the other assignments let the sentence remain true. Thus the sentence is **Satisfiable**.

c)



Problem 3

a)

$$\begin{aligned}
 \Pr(P_3 = 1) &= \sum_{p_1=0}^1 \sum_{p_2=0}^1 [\Pr(P_1 = p_1) \Pr(P_2 = p_2 \mid P_1 = p_1) \Pr(P_3 = 1 \mid P_2 = p_2)] \\
 &= \Pr(P_1 = 0) \Pr(P_2 = 0 \mid P_1 = 0) \Pr(P_3 = 1 \mid P_2 = 0) \\
 &\quad + \Pr(P_1 = 0) \Pr(P_2 = 1 \mid P_1 = 0) \Pr(P_3 = 1 \mid P_2 = 1) \\
 &\quad + \Pr(P_1 = 1) \Pr(P_2 = 0 \mid P_1 = 1) \Pr(P_3 = 1 \mid P_2 = 0) \\
 &\quad + \Pr(P_1 = 1) \Pr(P_2 = 1 \mid P_1 = 1) \Pr(P_3 = 1 \mid P_2 = 1) \\
 &= (0.4 \cdot 0.8 \cdot 0.3) + (0.4 \cdot 0.2 \cdot 0.8) \\
 &\quad + (0.6 \cdot 0.3 \cdot 0.3) + (0.6 \cdot 0.7 \cdot 0.8) \\
 &= 0.55
 \end{aligned}$$

b)

The additions are used to sum over all possible combinations of variables. Since the variable we want to compute (P_n) has a fixed value, we only need to consider the combinations of the $n - 1$ preceding variables, each of which has 2 possible values. Therefore, there are 2^{n-1} such combinations, and the number of additions required is $O(2^{n-1})$.

The multiplications are used to compute the probability of each variable for every combination. Since each combination involves n probability terms, we need $n - 1$ multiplications per combination. Given 2^{n-1} combinations, the total number of multiplications required is $2^{n-1} \cdot (n - 1)$, which implies a complexity of $O(n2^{n-1})$.

c)

Using variable elimination, we can substitute $\Pr(P_{n+1} = p_{n+1})$ with

$$\sum_{p_n} \Pr(P_n = p_n) \Pr(P_{n+1} = p_{n+1} \mid P_n = p_n)$$

This allows us to break down the problem involving n variables into $n - 1$ smaller problems, each involving only 2 variables.

For the case of $n = 2$, using naive enumeration, the number of additions and multiplications required is 2 and 4, respectively. Thus, for general n , the total number of operations follows:

- additions: $O(2 \cdot (n - 1)) = O(n)$
- multiplications: $O(4 \cdot (n - 1)) = O(n)$

Therefore, the computational complexity of variable elimination is $O(n)$ for both additions and multiplications.

Problem 4

a)

$$\begin{aligned} & \Pr(A) \Pr(B | A) \Pr(C | B) \Pr(D | B) \Pr(E) \\ & \Pr(F | C) \Pr(G | E) \Pr(H | D, F, G) \Pr(I) \\ & \Pr(J | F, H, I) \Pr(K) \Pr(L | J) \Pr(M | J, K) \Pr(N | M) \end{aligned}$$

b)

F 's Markov blankets: $[C, D, F, G, H, I, J]$

Variables that are d-separated from F : $[A, B, E, K, L, M, N]$

c)

- $D \perp\!\!\!\perp I \mid \{J\}$: false
- $H \perp\!\!\!\perp C \mid \{\}$: false
- $H \perp\!\!\!\perp K \mid \{\}$: true
- $B \perp\!\!\!\perp E \mid \{N\}$: false

d)

$[K, L, M, N]$

e)

$[I, N]$