



Recursion

C1

- Recursive functions & algorithms
- The call stack

- “...embedding an expression of some type within an expression of the same type”
(*linguistics*)



(Alf van Beem, CC0, via Wikimedia Commons)

- The body of a function can contain function calls, including *calls to the same function*.
 - This is known as **recursion**.
- The function must have a branching statement, such that a recursive call does not always take place (“base case”); otherwise, recursion never ends.
- Recursion is a way to think about solving a problem: how to reduce it to a simpler instance of itself?



Example: Factorial function

- Compute

$$\begin{aligned}f(n) &= n * (n - 1) * (n - 2) * \dots * 1 \\&= n * f(n - 1)\end{aligned}$$

- Base case:

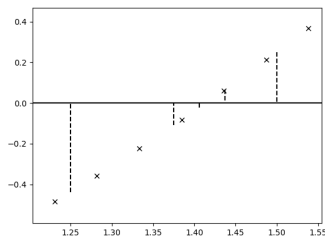
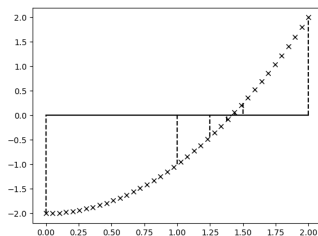
$$f(1) = 1$$

```
static int f(int n) {  
    if (n == 1)  
        return 1;  
    else  
        return n * f(n-1);  
}
```

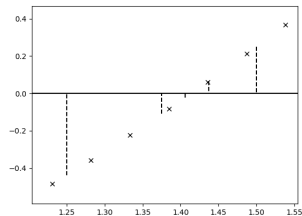
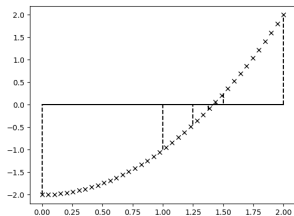


Example: Solving an equation

- Solve $g(x) = 0$.
- For example, find x such that $(x^2 - 2) = 0$.
- The interval-halving algorithm (a.k.a. binary search).

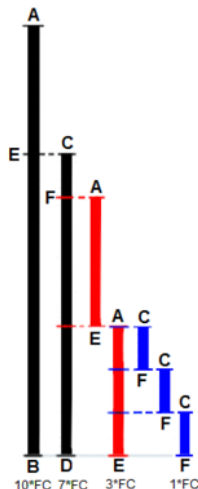


- Assumption: $g(x)$ is monotone increasing and crosses 0 in the interval $[l, u]$.
- Idea:
 - Find the middle of the interval, m :
 - if $g(m) \approx 0$, we're done;
 - if $g(m) < 0$, the solution lies in $[m, u]$;
 - if $g(m) > 0$, the solution lies in $[l, m]$.



Example: gcd

- The *greatest common divisor* (*gcd*) of numbers a and b is the greatest number c such that $a = ic$ and $b = jc$, for integers i and j .
- Euclid's algorithm.



- Assumption: a and b are positive, $a < b$.
- Euclid's idea:
 - Find m and r such that $b = ma + r$, $0 \leq r < a$.
 - $m = b / a$ (b/a rounded down)
 - $r = b \% a$ (remainder of b/a)
 - If $r = 0$, then $\gcd(a, b) = a$.
 - Else, $\gcd(a, b) = \gcd(r, a)$.



The call stack


- When a function call begins, the current instruction sequence is put “on hold” while the the function body executes.
- When the function ends, it returns to the next instruction after where the function was called.
- The “to-do list” of where to come back to after each current function call is called the **stack**.
- Variables declared in a function (including its parameters) are **local** to each call to that function.



```
static int f(int n) {
    if (n == 1)
        return 1;
    else
        return n * f(n-1);
}
```

1	a=f(3)			
	2	if (3 == 1)		
	3	return 3 * <u>f(3-1)</u>		
		4	if (2 == 1)	
		5	return 2 * <u>f(2-1)</u>	
			6	if (1 == 1)
			7	return 1
		8	return 2 * 1	
	9	return 3 * (2 * 1)		
10	a = 6			

stack depth




preview: Recursive data structures

- A recursive data structure is made up of parts that reference other parts of the same type.
- Example: A binary tree is
 - a leaf node; or
 - a node with two children, that are binary trees.

